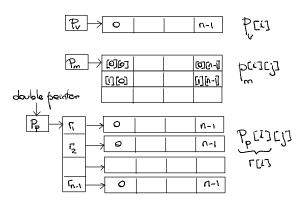
Data Structures

- arrays
- linked lists
- stack
- queues
- binary trees
- forests
- graphs
- hash tables

Arrays



C++ Example

```
#include <cstdio>
const int dim = 6:
1 void afunction()
2 {
3 int a[dim];
4 int m[dim][dim/2];
5 int* pv = a;
  int(*pm)[dim/2] = m; // NOTICE THE SYNTAX!!!
 int* r[dim]:
7
8 int** pp = r;
9 for (int i=0; i<dim; i++)
10 r[i] = &m[i][0]; // m+i*dim/2
11 // initializing matrix
12 for (int i=0; i< dim; i++)
13
     for (int j=0; j<dim/2; j++)
       m[i][j] = (i+1)*10+(j+1);
14
```

C++ Example

```
10
   /*
    int(*pm)[dim/2] = m; // NOTICE THE SYNTAX!!!
6
   int* r[dim];
7
   int**pp = r;
    for (int i=0; i<dim; i++)
      r[i] = &m[i][0]; // m+i*dim/2;
10
    */
    printf("printing matrix using pointer to matrix\n");
    for (int i=0; i< dim; i++)
16
17
18
      for (int j=0; j < dim/2; j++)
        printf("%d ", pm[i][j]);
19
20
     printf("\n");
21
22
    printf("printing matrix using array of pointers to rows\n");
23
    for (int i=0; i< dim; i++)
24
25
      for (int j=0; j<dim/2; j++)
        printf("%d ", pp[i][j]);
26
27
      printf("\n");
28
29
```

Running the program

```
int main(int argc, char* argv[])
  afunction();;
  return 0;
printing matrix using pointer to matrix: int(*pm)[dim/2] = m;
11 12 13
21 22 23
31 32 33
41 42 43
51 52 53
61 62 63
printing matrix using array of pointers to rows: int* r[dim]; int** pp = r;
11 12 13
21 22 23
31 32 33
41 42 43
51 52 53
61 62 63
```

Resizing

```
#include <cstring>
const double factor = 1.5;

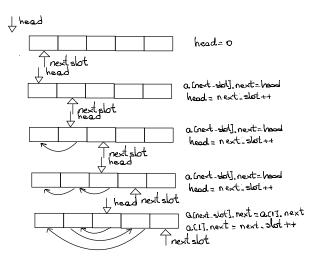
int* resize(const int* array, const unsigned int size)
{
   const unsigned int newSize = size*factor;
   int* newP = new int[newSize];
   memcpy(newP, array, size*sizeof(int));
   return newP;
}
```

Linked List

Linked List

```
typedef unsigned int nodedata;
struct Node
 const nodedata
                  cargo;
 const unsigned int key;
 Node*
                     link:
 Node(const nodedata mycargo, const unsigned int akey) :
    cargo(mycargo), key(akey), link(NULL) { }
};
class LinkedList
public:
 LinkedList():
 ~LinkedList():
 unsigned int insert(nodedata cargo)
 { unsigned int nodekey = key; Node* newNode = new Node(cargo, nodekey);
   newNode->link = head; head = newNode; key++; count++; return nodekey; }
 unsigned int insertafterkey(nodedata, unsigned int key);
 void remove(unsigned int);
 nodedata removeafterkey(unsigned int key);
private:
 Node*
              head:
 unsigned int count;
 unsigned int key;
};
                                                 4□ → 4□ → 4 □ → □ ● 900
```

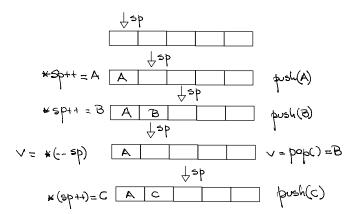
Array-based Linked List



Array-based Linked List Implementation

```
struct ArrayNode
  const nodedata
                  cargo;
  const unsigned int key;
 unsigned int
                    index;
 ArrayNode* link;
  ArrayNode(const nodedata mycargo, const unsigned int akey) :
    cargo(mycargo), key(akey), index(0), link(NULL) { }
};
class ArrayLinkedList
public:
  ArrayLinkedList();
 ~ArrayLinkedList();
 unsigned int insert(nodedata);
 unsigned int insertafterkey(nodedata, unsigned int key);
 void remove(unsigned int);
 nodedata removeafterkey(unsigned int key);
private:
 ArrayNode*
             array;
 unsigned int arraySize;
 unsigned int headIndex;
 unsigned int nextSlot;
 unsigned int emptySlotHead;
 unsigned int key;
                                                4□ → 4□ → 4 □ → □ ● 900
```

Stack



Stack - Array Implementation

```
struct StackNode
  nodedata cargo;
 StackNode() { }:
};
class Stack
public:
  Stack(unsigned int maxdepth) : array(new StackNode[maxdepth])
  { sp=array; endStack = sp+maxdepth; }
  void push(const StackNode& n) { *sp++ = n; } // check sp < endStack</pre>
  StackNode pop() { return *(--sp); } // check sp>array
private:
  StackNode* sp;
  StackNode* array;
  StackNode* endStack;
};
```

Stack - Linked List Implementation

```
struct LLStackNode
 nodedata
           cargo;
 LLStackNode* next:
 LLStackNode() { };
class LLStack
public:
 LLStack() : head(NULL) { }
 void push(const LLStackNode& n)
 { StackNode* sn = new LLStackNode(n); sn->next = head; head=sn; }
 // check head !=NULL
 LLStackNode pop() { LLStackNode sn = *(head); head=head->next; return sn; }
private:
 LLStackNode* head:
};
```

Example

```
a * (((b+c) * (d*e)) + f)
struct Item
 union
 int var;
 char oper;
 } token;
  int tokenType; // tokeType=0 for number, tokenType=1 for operator
};
     a * (((b+c) * (d*e)) + f)
    a (((b+c) * (d*e)) + f) *
     a (((b c +) (d e *) *) f +) *
Item * 1 = \{a b c + d e * * f + *\}
     (a) b c + d e * * f + *
push(a) | a |
     a (b) c + d e * * f + *
push(b) | b | a |
     a b (c) + d e * * f + *
push(c) | c | b | a |
```

Example

```
abc(+)de * * f + *
c = pop() | b | a |
b = pop() | a |
h = b+c
push(h) | h | a |
    a b c + (d) e * * f + *
push(d) | d | h | a |
    a b c + d (e) * * f + *
push(e) | e | d | h | a |
    a b c + d e (*) * f + *
e = pop() | d | h | a |
d = pop() | h | a |
i = d*e
push(i) | i | h | a |
```

Example

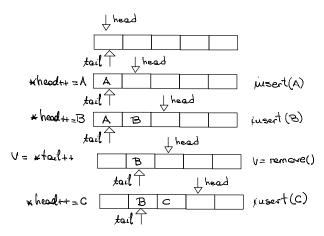
```
a b c + d e * (*) f + *
i = pop() | h | a |
h = pop() | a |
j = i*h
push(j) | j | a |
    a b c + d e * * (f) + *
push(f) | f | j | a |
    a b c + d e * * f (+) *
f=pop() | j | a |
j=pop() | a |
k = f + j
push(k) | k | a |
    abc+de** f+(*)
k = pop()
a = pop()
result = k*a
```

Queues

Queue

```
struct QueueNode
 nodedata
           cargo;
 QueueNode* next:
 QueueNode() : next(NULL) { }:
};
class Queue
public:
 Queue() { QueueNode* dummy = new QueueNode; head=tail = dummy; };
 void insert(QueueNode& node)
                                            head->next=nn
 { QueueNode* nn = new QueueNode(node); nn->next = head->next; head=nn; }
 QueueNode remove() {
    QueueNode nn(*tail->next); QueueNode* tn = tail->next;
    delete tail: tail=tn: return nn: }
private:
 QueueNode* head:
 QueueNode* tail:
};
```

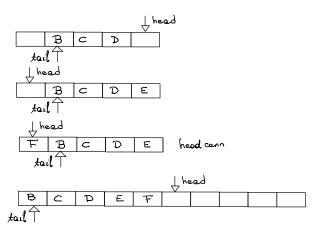
Array-based Implementation



Array-based Queue

```
struct ArrayQueueNode
 nodedata
            cargo;
 unsigned int next;
 ArrayQueueNode() : next(0) { };
};
class ArrayQueue
public:
 ArrayQueue(unsigned int maxSize) : maxQueueSize(maxSize)
 { head=tail= 0; array = new ArrayQueueNode[maxSize]; };
 void insert(ArrayQueueNode& node)
 { ArrayQueueNode nn = node;
   nn.next = head:
    array[head++] = nn; }
  ArrayQueueNode remove() {
    ArrayQueueNode nn = array[tail++]; return nn; }
private:
 unsigned int head;
 unsigned int tail;
 unsigned int maxQueueSize;
 ArrayQueueNode* array;
};
```

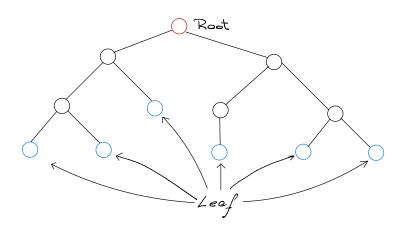
Resizing the Array



Wrap-around

```
class ArrayQueue2
public:
  ArrayQueue2(unsigned int maxSize) : maxQueueSize(maxSize)
  { head=tail= 0; array = new ArrayQueueNode[maxSize]; };
  void insert(ArrayQueueNode& node)
  { ArrayQueueNode nn = node;
    nn.next = head:
    array[head++] = nn;
    if (head == maxQueueSize ) head = 0;
    if (head == tail)
      { /* queue is full - raise exception
handle the execption by reallocating array and transferring items */
  ArrayQueueNode remove() {
    if (tail != head)
      { ArrayQueueNode nn = array[tail++]; return nn; }
    /* queue is empty - raise exception */ }
private:
  unsigned int head;
  unsigned int tail;
  unsigned int maxQueueSize;
  ArrayQueueNode* array;
};
```

Binary Trees



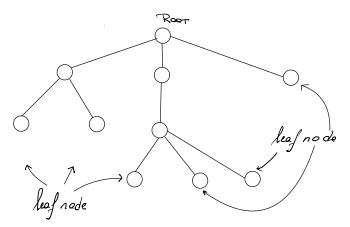
Properties of trees

- There is exactly one path connecting any two nodes in a tree
- A tree with N nodes hase N-1 edges
- A binary tree with N internal nodes has N+1 leaves
- ullet The hight of a full binary tree with N internal nodes is about log_2N

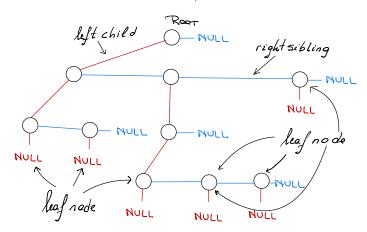
Binary Tree Representation

```
struct BinaryTreeNode
  BinaryTreeNode() : leftChild(NULL), rightChild(NULL) { }
  nodedata
                  cargo;
  BinaryTreeNode* leftChild;
  BinaryTreeNode* rightChild;
};
class BinaryTree
public:
  BinaryTree(BinaryTreeNode& root);
private:
  BinaryTree* rootNode;
};
```

General Tree Structure



General Tree Representation



Forest Representation

```
struct ForestNode
  ForestNode() : leftChild(NULL), rightChild(NULL) { }
  nodedata cargo;
  ForestNode* leftChild;
  ForestNode* rightSibling;
};
class Forest
public:
  Forest(ForestNode& root);
private:
  ForestNode* rootNode;
};
```

Recursion

- Function calls itself
- 2 Termination Condition

Mathematical recurrence \iff Recursive programs Factorial

```
N! = N \times (N-1)! \quad \text{for } N \geq 1 \qquad 0! = 1 int factorial(const int n) \{ if (n==0) return 1; return n*factorial(n-1); }
```

Recursion

Fibonacci Sequence

```
F_n = F_{n-1} + F_{n-2} for N > 2 F_0 = F_1 = 1
int fibonacci(const int n)
  if (n<2) return 1:
  return fibonacci(n-1)+fibonacci(n-2);
}
How many recursive calls are needed?
For F_0 or F_1: 1
For F_2: calls for F_0 + calls for F_1: 2
For F_3: calls for F_2 + calls for F_1: 3
For F_4: calls for F_3 + calls for F_2: 5
For F_N: \phi^N where \phi \approx 1.62 - golden ratio
```

Divide and Conquer

- Split input into 2 parts
- Two recursive calls, each operating on approximately half of the input
- Merge the results of processing 2 independent portions of the input
- Input is divided without overlap
- There may be code before, after, or in between ther recursive calls
- Top-down orientation

Tree Traversal

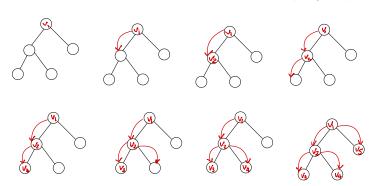
```
struct BinaryTreeNode
  BinaryTreeNode() : leftChild(NULL), rightChild(NULL) { }
  nodedata
                  cargo;
  BinaryTreeNode* leftChild;
  BinaryTreeNode* rightChild;
  void visit();
}:
void inorder_traverse(BinaryTreeNode* t)
  inorder traverse(t->leftChild):
  t->visit():
  inorder_traverse(t->rightChild);
void preorder_traverse(BinaryTreeNode* t)
  t->visit():
  preorder_traverse(t->leftChild);
  preorder_traverse(t->rightChild);
```

Inorder Traversal

traverse(t->leftChild) t->visit() traverse(t->rightChild)

Preorder Traversal

t->visit() traverse(t->leftChild) traverse(t->rightChild)



```
void preorder_traverse(BinaryTreeNode* t)
  if (t)
    t->visit();
    preorder_traverse(t->leftChild);
    preorder_traverse(t->rightChild);
void preorder_traverse(BinaryTreeNode* t)
    do
      t->visit();
      if (t->leftChild) preorder_traverse(t->leftChild);
      t = t->rightChild;
      } while(t);
```

On entry

- push values of the local variables on the stack
- push the address of the next instruction (return) on the stack
- set the values of the parameters to the procedure
- go to the beginning of the procedure

On return

- pop the return address
- reset local variables
- go to the return address

```
void preorder_traverse(BinaryTreeNode* t)
 if (t)
    t->visit():
   preorder_traverse(t->leftChild);
   preorder_traverse(t->rightChild);
void preorder_traverse(BinaryTreeNode* t)
 dο
   while (t)
      t->visit():
      stack.push(t); // push the local variable into the stack
      t = t->leftChild; // set the new value for the function parameter
                        // returns to the begining of the while loop
// we are here when t->leftChild == 0
// meaning we have to return to the parent node
    if (stack.isempty()) break; // no more nodes on the stack. exit
    t = stack.pop(); // retrieve the pointer to the parent node
   t = t->rightChild; // traverse the right subtree
    } while(1):
                      // reenter the function
                                                4 D > 4 A > 4 B > 4 B > B 9 9 0
```

```
void preorder_traverse(BinaryTreeNode* t)
 dο
   while (t)
     t->visit();
     stack.push(t);
     t = t->leftChild:
   if (stack.isempty()) break; -----
   t = stack.pop(); -----
   t = t->rightChild;
   } while(1);
void preorder_traverse(BinaryTreeNode* t)
 stack.push(t);
 while (!stack.isempty()) <-----
   t = stack.pop(); <-----
   if (t->rightChild) stack.push(t->rightChild);
   t->visit();
   if (t->leftChild) stack.push(t->leftChild);
```

Algorithms and their Classification of Algorithms

- Constant Time
- ullet log~N: logaritmic running time
- N: linear running time
- N log N: N log N running time
- N^2 : quadratic running time
- N^3 : cubic running time
- ullet 2^N : exponential running time

Computational Complexity - Big O notation

Definition: A function g(N) is said to be $\mathrm{O}\big(f(N)\big)$ if there exist constants c_0 and N_0 such that g(N) is less that $c_0f(N)$ for all $N>N_0$