Homework 3

1. Using the Theorem 3 in the lecture note for Chapter 3 could be a way to obtain a bound for the absolute error. Note $f(x) = x^{-2}$ is smooth on [1, 2], so, for the Lagrange interpolating polynomial P, we have

$$|f(x) - P(x)| = \left| \frac{f^{(3)}(\xi(x))}{3!} (x - 1) \left(x - \frac{11}{8} \right) (x - 2) \right|$$
$$= \left| -4\xi(x)^{-5} (x - 1) \left(x - \frac{11}{8} \right) (x - 2) \right|$$
$$\le 4 \cdot 1^{-5} \cdot \frac{9}{128} = \frac{9}{32} \approx 0.28125$$

because $\xi(x) \in [1, 2]$ and (x - 1)(x - 11/8)(x - 2) has two extrema on [1, 2], which are at x = 7/4 and x = 7/6, where the former one gives the maximum absolute value of (x - 1)(x - 11/8)(x - 2) on [1, 2].

(Idea for an alternative solution) We also can directly calculate the error term. Since (f(x) - P(x))' = 0 is a quartic equation, which can be always solved with an exact solution using radicals, the maximum of |f(x) - P(x)| also can be calculated exactly, though it is difficult for a human. The maximum absolute error is approximately 0.041300, which is quite better than the answer above.

2. Suppose $x \neq x_0$, otherwise we have $f(x_0) - P_n(x_0) = 0$. Denote $R_n(x) = f(x) - P_n(x)$, and let

$$g(t) = R_n(t) - R_n(x) \frac{(t - x_0)^{n+1}}{(x - x_0)^{n+1}}.$$

I will use somewhat different generalization of the Rolle's theorem than the Theorem 1.10 in the textbook:

Suppose $f \in C^n[a,b]$ and $f^{(n+1)}$ exists at x=a, with $f(a)=\cdots=f^{(n)}(a)=0$ and f(b)=0. Then there is a number $c\in(a,b)$ with $f^{(n+1)}(c)=0$.

This can be proved easily: find $\xi_1 \in (a,b)$ satisfying $f'(\xi_1) = 0$ using

Rolle's to f, find $\xi_2 \in (a, \xi_1)$ satisfying $f''(\xi_1) = 0$ using Rolle's to f', and so on.

Now the problem became easy: since

$$g^{(j)}(t) = R_n^{(j)}(t) - (n+1)\cdots(n+2-j)\frac{(t-x_0)^{n+1-j}}{(x-x_0)^{n+1}}$$

so that

$$g^{(j)}(x_0) = f^{(j)}(x_0) - P_n^{(j)}(x_0) = 0.$$

Also, g(x) = 0. Thus, applying the generalization of the Rolle's theorem above, we have a constant $c =: \xi(x)$ between x_0 and x, depending only on x. Since both x_0 and x are on [a, b], we have $\xi(x) \in (a, b)$.

- 3. See the code below for the implementation.
 - (a) $\sqrt{3} = f(0.5) \approx 1.70833333$.
 - (b) $\sqrt{3} = f(3) \approx 1.73138614$.

% Neville's Method

% Define f(x)

```
f = 0(x) 3^x;
                      % for (a)
% f = @(x) \ sqrt(x); % for (b)
X = [-2, -1, 0, 1, 2];
N = length(X);
% Q\{i, j\} = Q_{1} \{i - 1, j - 1\}
Q = cell(N, N);
for i = 1:N
    Q\{i, 1\} = [f(X(i))];
% Calculating Q{i, j}
for j = 2:N
    for i = j:N
        % (x - x_{i-j})Q_{i, j-1}
        first_term = sum_poly( ...
            [Q{i, j-1} 0],
            -X(i-j+1) * Q{i, j-1} ...
        % -(x - x_{i})Q_{i-1, j-1}
        second_term = sum_poly(
            -1 * [Q{i-1, j-1} 0], ...
```

```
X(i) * Q{i-1, j-1} \dots
        );
        Q\{i, j\} =
            sum_poly(first_term, second_term) / ...
            (X(i) - X(i-j+1));
    end
end
% Evaluate the polynomial
           % for (a)
x = 0.5;
% x = 3;
            % for (b)
fprintf('f(\%f) = \%.8f\n', x, polyval(Q{N, N}, x))
% x * P = [P \ 0]
% a * P = a * P
% P + Q
function s = sum_poly(a, b)
    N = max(length(a), length(b));
    pa = [zeros(1, N - length(a)) a];
    pb = [zeros(1, N - length(b)) b];
    s = pa + pb;
    return
end
```

4. (a) Note $P_{1,2} = Q_{2,1}$ and $P_2 = Q_{2,0}$. With x = 0.5,

$$\frac{27}{7} = Q_{2,2} = \frac{(x - x_0)Q_{2,1} - (x - x_2)Q_{1,1}}{x_2 - x_0} = \frac{5}{7}Q_{2,1} + 1, \qquad Q_{2,1} = 4;$$

$$4 = Q_{2,1} = \frac{(x - x_1)Q_{2,0} - (x - x_2)Q_{1,0}}{x_2 - x_1} = \frac{Q_{2,0} + 5.6}{3}; \qquad Q_{2,0} = 6.4.$$

(b)

$$P_{0,1,2}(1.5) = Q_{2,2}(1.5) = \frac{(x - x_0)Q_{2,1}(x) - (x - x_2)Q_{1,1}(x)}{x_2 - x_0} \bigg|_{x=1.5} = 3.25;$$

$$P_{0,1,2,3}(1.5) = Q_{3,3}(1.5) = \frac{(x - x_0)Q_{3,2}(x) - (x - x_3)Q_{2,2}(x)}{x_3 - x_0} \bigg|_{x=1.5} = 3.625.$$

(b: alternative way) Since $P_{1,2,3}(x)$ is a quadratic polynomial with $P_{1,2,3}(1) = P_{1,2}(1) = 2$, $P_{1,2,3}(1.5) = 4$, and $P_{1,2,3}(2) = P_{1,2}(2) = 5$, we have $P_{1,2,3}(x) = -2x^2 + 9x - 5$, having $P_{0,1,2,3}(3) = 4$. So the graph of $P_{0,1,2,3}$ passes (0,1), (1,2), (2,5), and (3,4), we have

$$P_{0,1,2,3}(x) = -x^3 + 4x^2 - 2x + 1.$$

Thus $P_{0,1,2,3}(1.5) = 3.625$.

5. Let P be a polynomial of least degree, and show $P = H_{2n+1}$. In this way, we can simultaneously show that H_{2n+1} is such a polynomial of least degree and is actually unique. Note H_{2n+1} is a polynomial agreeing with f and f' at given points, though we do not yet know that H_{2n+1} is of least degree. Thus, by the minimality of the degree of P, we have deg $P \le \deg H_{2n+1}$. Consequently, $\deg(H_{2n+1} - P) = \deg D \le \deg H_{2n+1} \le 2n + 1$. Then, $\deg D \le 2n$. However, by the Rolle's theorem, we have $D'(\xi_i) = 0$ with some $x_i < \xi_i < x_{i+1}$ $(i = 0, \ldots, n-1)$ so that D'(x) = 0 for $x = x_0, \ldots, x_n, \xi_0, \ldots, \xi_{n-1}$. According to the fundamental theorem of algebra, $D' \equiv 0$. Hence D is a constant, which should be 0, because, say, $D(x_0) = 0$.

(Alternative solution) $D(x_i) = D'(x_i) = 0$ for i = 0, ..., n. Letting $D(x) = \sum_{i=0}^{2n+1} d_i x^i$ (d_i could be zero),

$$d_0 + d_1 x_0 + \dots + d_{2n+1} x_0^{2n+1} = 0, \quad d_1 + 2d_2 x_0 + \dots + (2n+1)d_{2n+1} x_0^{2n} = 0,$$

$$\vdots \qquad \qquad \vdots$$

$$d_0 + d_1 x_0 + \dots + d_{2n+1} x_n^{2n+1} = 0, \quad d_1 + 2d_2 x_n + \dots + (2n+1)d_{2n+1} x_n^{2n} = 0,$$

i.e.,

$$\begin{pmatrix} 1 & x_0 & \cdots & x_0^{2n+1} \\ \vdots & \vdots & \ddots & \vdots \\ 1 & x_n & \cdots & x_n^{2n+1} \\ 0 & 1 & \cdots & (2n+1)x_0^{2n} \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 1 & \cdots & (2n+1)x_n^{2n} \end{pmatrix} \begin{pmatrix} d_0 \\ \vdots \\ d_n \\ d_{n+1} \\ \vdots \\ d_{2n+1} \end{pmatrix} = \mathbf{0}.$$

Redefine n to be n-1 in the formulae above for the sake of simplicity. The determinant of the $2n \times 2n$ square matrix on the left side can be calculated by the induction on the number ℓ of rows with leading zero in the following

form of matrix determinant: $(0 \le \ell \le m := 2n - \ell)$

$$\begin{vmatrix} 1 & x_0 & \cdots & x_0^{2n-1} \\ \vdots & \vdots & \ddots & \vdots \\ 1 & x_{m-1} & \cdots & x_{m-1}^{2n-1} \\ 0 & 1 & \cdots & (2n-1)x_0^{2n-2} \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 1 & \cdots & (2n-1)x_{\ell-1}^{2n-2} \end{vmatrix} = (-1)^{\ell(\ell-1)/2} \prod_{0 \le i < j \le m-1} (x_j - x_i)^{\alpha_i \alpha_j}$$

where $\alpha_j = 2$ if $j < \ell$, and $\alpha_j = 1$ otherwise. When $\ell = 0$, we have nothing to do due to the Vandermonde determinant. Observe that expanding the determinant along m-th row, the other parts does not depend on x_{m-1} . As an induction step (assuming $0 \le \ell < n$ so $\alpha_{m-1} = 1$),

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 so $\alpha_{m-1} = 1$),
$$\frac{\partial}{\partial x_{m-1}} \begin{vmatrix}
1 & x_0 & \cdots & x_0^{2n-1} \\
\vdots & \vdots & \ddots & \vdots \\
1 & x_{m-1} & \cdots & x_{m-1}^{2n-1} \\
0 & 1 & \cdots & (2n-1)x_0^{2n-2} \\
\vdots & \vdots & \ddots & \vdots \\
0 & 1 & \cdots & (2n-1)x_{\ell-1}^{2n-2}
\end{vmatrix}$$

$$= \begin{vmatrix}
1 & x_0 & \cdots & x_0^{2n-1} \\
\vdots & \vdots & \ddots & \vdots \\
0 & 1 & \cdots & (2n-1)x_{m-1}^{2n-2} \\
0 & 1 & \cdots & (2n-1)x_0^{2n-2} \\
\vdots & \vdots & \ddots & \vdots \\
0 & 1 & \cdots & (2n-1)x_{\ell-1}^{2n-2}
\end{vmatrix}$$

$$= (-1)^{\ell(\ell-1)/2} \prod_{i < j < m-1} (x_j - x_i)^{\alpha_i \alpha_j} \cdot \frac{\partial}{\partial x_{m-1}} \prod_{i=0}^{m-2} (x_{m-1} - x_i)^{\alpha_i}.$$

By replacing x_{m-1} by x_{ℓ} , we have

$$\begin{vmatrix} 1 & x_0 & \cdots & x_0^{2n-1} \\ \vdots & \vdots & \ddots & \vdots \\ 1 & x_{m-2} & \cdots & x_{m-2}^{2n-1} \\ 0 & 1 & \cdots & (2n-1)x_\ell^{2n-2} \\ 0 & 1 & \cdots & (2n-1)x_\ell^{2n-2} \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 1 & \cdots & (2n-1)x_{\ell-1}^{2n-2} \end{vmatrix}$$

$$= (-1)^{\ell(\ell-1)/2} \prod_{i < j < m-1} (x_j - x_i)^{\alpha_i \alpha_j} \cdot \left(\sum_{k=0}^{m-2} \alpha_k (x_{m-1} - x_i)^{\alpha_i} \prod_{\substack{i=0 \ i \neq k}}^{m-2} (x_{m-1} - x_i)^{\alpha_i} \prod_{\substack{x_{m-1} = x_\ell \\ i \neq k}}^{m-2} (x_{m-1} - x_i)^{\alpha_i} \prod_{\substack{x_{m-1} = x_\ell \\ i \neq k}}^{m-2} (x_{m-1} - x_i)^{\alpha_i} \prod_{\substack{x_{m-1} = x_\ell \\ i \neq \ell}}^{m-2} (x_{m-1} - x_i)^{\alpha_i} \prod_{\substack{x_{m-1} = x_\ell \\ i \neq \ell}}^{m-2} (x_{m-1} - x_i)^{\alpha_i} \prod_{\substack{x_{m-1} = x_\ell \\ i \neq \ell}}^{m-2} (x_{m-1} - x_i)^{\alpha_i} \prod_{\substack{x_{m-1} = x_\ell \\ i \neq \ell}}^{m-2} (x_\ell - x_i)^{\alpha_i} \prod_{\substack{x_{m-1} = x_\ell \\ i \neq \ell}}^{m-2} (x_\ell - x_i)^{\alpha_i} \prod_{\substack{x_{m-1} = x_\ell \\ i \neq \ell}}^{m-2} (x_\ell - x_i)^{\alpha_i} \prod_{\substack{x_{m-1} = x_\ell \\ i \neq \ell}}^{m-2} (x_\ell - x_i)^{\alpha_i} \prod_{\substack{x_{m-1} = x_\ell \\ i \neq \ell}}^{m-2} (x_\ell - x_i)^{\alpha_i} \prod_{\substack{x_{m-1} = x_\ell \\ i \neq \ell}}^{m-2} (x_\ell - x_i)^{\alpha_i} \prod_{\substack{x_{m-1} = x_\ell \\ i \neq \ell}}^{m-2} (x_\ell - x_i)^{\alpha_i} \prod_{\substack{x_{m-1} = x_\ell \\ i \neq \ell}}^{m-2} (x_\ell - x_i)^{\alpha_i} \prod_{\substack{x_{m-1} = x_\ell \\ i \neq \ell}}^{m-2} (x_\ell - x_i)^{\alpha_i} \prod_{\substack{x_{m-1} = x_\ell \\ i \neq \ell}}^{m-2} (x_\ell - x_i)^{\alpha_i} \prod_{\substack{x_{m-1} = x_\ell \\ i \neq \ell}}^{m-2} (x_\ell - x_i)^{\alpha_i} \prod_{\substack{x_{m-1} = x_\ell \\ i \neq \ell}}^{m-2} (x_\ell - x_i)^{\alpha_i} \prod_{\substack{x_{m-1} = x_\ell \\ i \neq \ell}}^{m-2} (x_\ell - x_i)^{\alpha_i} \prod_{\substack{x_{m-1} = x_\ell \\ i \neq \ell}}^{m-2} (x_\ell - x_i)^{\alpha_i} \prod_{\substack{x_{m-1} = x_\ell \\ i \neq \ell}}^{m-2} (x_\ell - x_i)^{\alpha_i} \prod_{\substack{x_{m-1} = x_\ell \\ i \neq \ell}}^{m-2} (x_\ell - x_i)^{\alpha_i} \prod_{\substack{x_{m-1} = x_\ell \\ i \neq \ell}}^{m-2} (x_\ell - x_i)^{\alpha_i} \prod_{\substack{x_{m-1} = x_\ell \\ i \neq \ell}}^{m-2} (x_\ell - x_i)^{\alpha_i} \prod_{\substack{x_{m-1} = x_\ell \\ i \neq \ell}}^{m-2} (x_\ell - x_i)^{\alpha_i} \prod_{\substack{x_{m-1} = x_\ell \\ i \neq \ell}}^{m-2} (x_\ell - x_i)^{\alpha_i} \prod_{\substack{x_{m-1} = x_\ell \\ i \neq \ell}}^{m-2} (x_\ell - x_i)^{\alpha_i} \prod_{\substack{x_{m-1} = x_\ell \\ i \neq \ell}}^{m-2} (x_\ell - x_i)^{\alpha_i} \prod_{\substack{x_{m-1} = x_\ell \\ i \neq \ell}}^{m-2} (x_\ell - x_i)^{\alpha_i} \prod_{\substack{x_{m-1} = x_\ell \\ i \neq \ell}}^{m-2} (x_\ell - x_i)^{\alpha_i} \prod_{\substack{x_{m-1} = x_\ell \\ i \neq \ell}}^{m-2} (x_\ell - x_i)^{\alpha_i} \prod_{\substack{x_{m-1} = x_\ell \\ i \neq \ell}}^{m-2} (x_\ell - x_i)^{\alpha_i} \prod_{\substack{x_{m-1} = x_\ell \\ i \neq \ell}}^{m-2}$$

where $\tilde{\alpha}_i = \alpha_i$ if $i \neq \ell$ and $\tilde{\alpha}_\ell = 2$. (Note that $m - 1 - (\ell + 1) \equiv 0 \pmod{2}$.)

Therefore,

$$\begin{vmatrix} 1 & x_0 & \cdots & x_0^{2n-1} \\ \vdots & \vdots & \ddots & \vdots \\ 1 & x_{m-2} & \cdots & x_{m-2}^{2n-1} \\ 0 & 1 & \cdots & (2n-1)x_0^{2n-2} \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 1 & \cdots & (2n-1)x_{\ell-1}^{2n-2} \\ 0 & 1 & \cdots & (2n-1)x_{\ell}^{2n-2} \end{vmatrix}$$

$$\begin{vmatrix} 1 & x_0 & \cdots & x_0^{2n-1} \\ \vdots & \vdots & \ddots & \vdots \\ 1 & x_{m-2} & \cdots & x_{m-2}^{2n-1} \\ 0 & 1 & \cdots & (2n-1)x_{\ell}^{2n-2} \\ 0 & 1 & \cdots & (2n-1)x_0^{2n-2} \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 1 & \cdots & (2n-1)x_{\ell-1}^{2n-2} \end{vmatrix}$$

$$= (-1)^{\ell} (-1)^{\ell(\ell-1)/2} \prod_{i < j < m-1} (x_j - x_i)^{\tilde{\alpha}_i \tilde{\alpha}_j}$$
The determinant is zero iff there are duplication.

as desired. Thus, the determinant is zero iff there are duplicated points, but the problem asserted that all x_i 's are distinct. So the linear equation above has only the trivial solution.