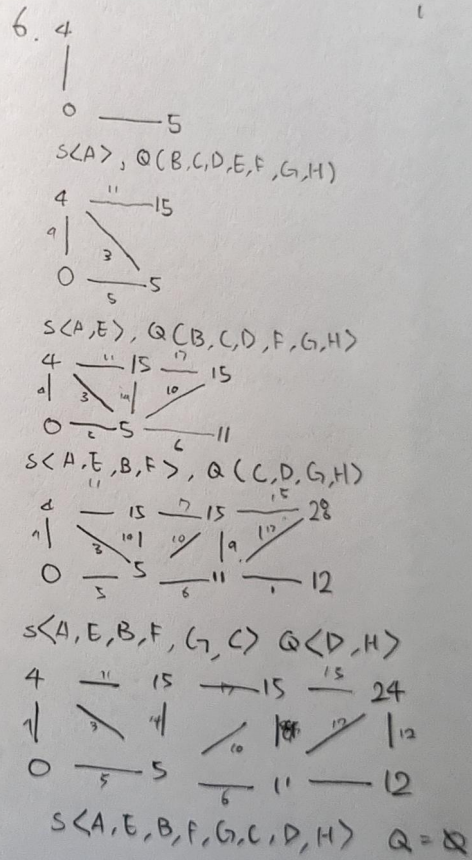
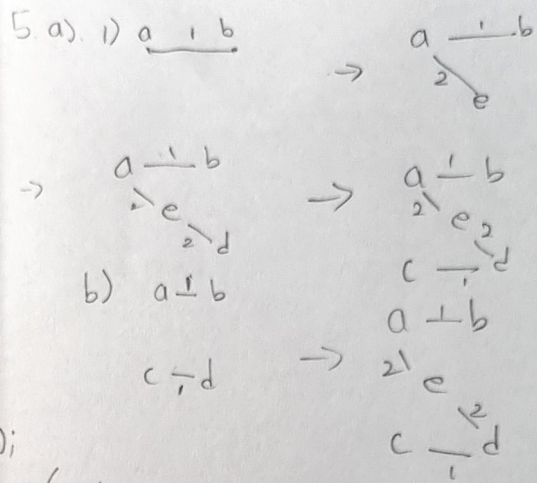


1. a) $1 + (4 \times 3) + 2 + (3 \times 4) + 2 = 29$
- b) $OPT[i] = \begin{cases} \max(OPT[i-1] + V_i, OPT[i-2] + V_i \times V_{i-1}) & (i \geq 2) \\ V_1 & (i = 1) \\ 0 & (i = 0) \end{cases}$
- c) $Prod = \text{Sum}(inc[], n)$
 if $(n == 0)$ return 0;
 int $OPT[] = \text{new int}[n+1];$
 $OPT[0] = 0;$
 $OPT[1] = V[1];$
 for int $j = 2$ to n
 $OPT[j] = \max(OPT[j-1] + V[j], OPT[j-2] + V[j] \times V[j-1]);$
 return $OPT[n];$

2. Least-Cents (int n)
 int $coin[8] = \{1, 10, 25\};$ $\leftarrow DP[0] = 0$
 for (int $i = 0; i < 8; i++)$
 for (int $j = coin[i]; j \leq n; j++)$
 $DP[j] = \min(DP[j], DP[j - coin[i]] + 1);$
 return $DP[n];$

3. When it's 14 cents $\frac{1}{11}$
 In greedy algorithm $\{11, 1, 1, 1\} \rightarrow 4 \text{ coins}$
 But the fewest coins is $\{7, 7\} \rightarrow 2 \text{ coins}$
 So, It does not always find fewest coin

4.
 $DP[0][0] = M[0][0]$
 for (int $i = 1; i \leq M; i++)$
 $DP[0][i] = DP[0][i-1] + M[0][i]$
 $DP[i][0] = DP[i-1][0] + M[i][0]$
 for (int $a = 1; a \leq M; a++)$
 for (int $b = 1; b \leq M; b++)$
 $DP[a][b] = \max(DP[a-1][b], M[a][b], DP[a][b-1] + M[a][b])$
 return $DP[N-1][M-1]$



7. $G_d C' G_b$ is true

G_b is works with negative-weight edges and detects if there is a negative-weight cycle but G_d requires all edge weights to be nonnegative

So $G_d C G_b$