

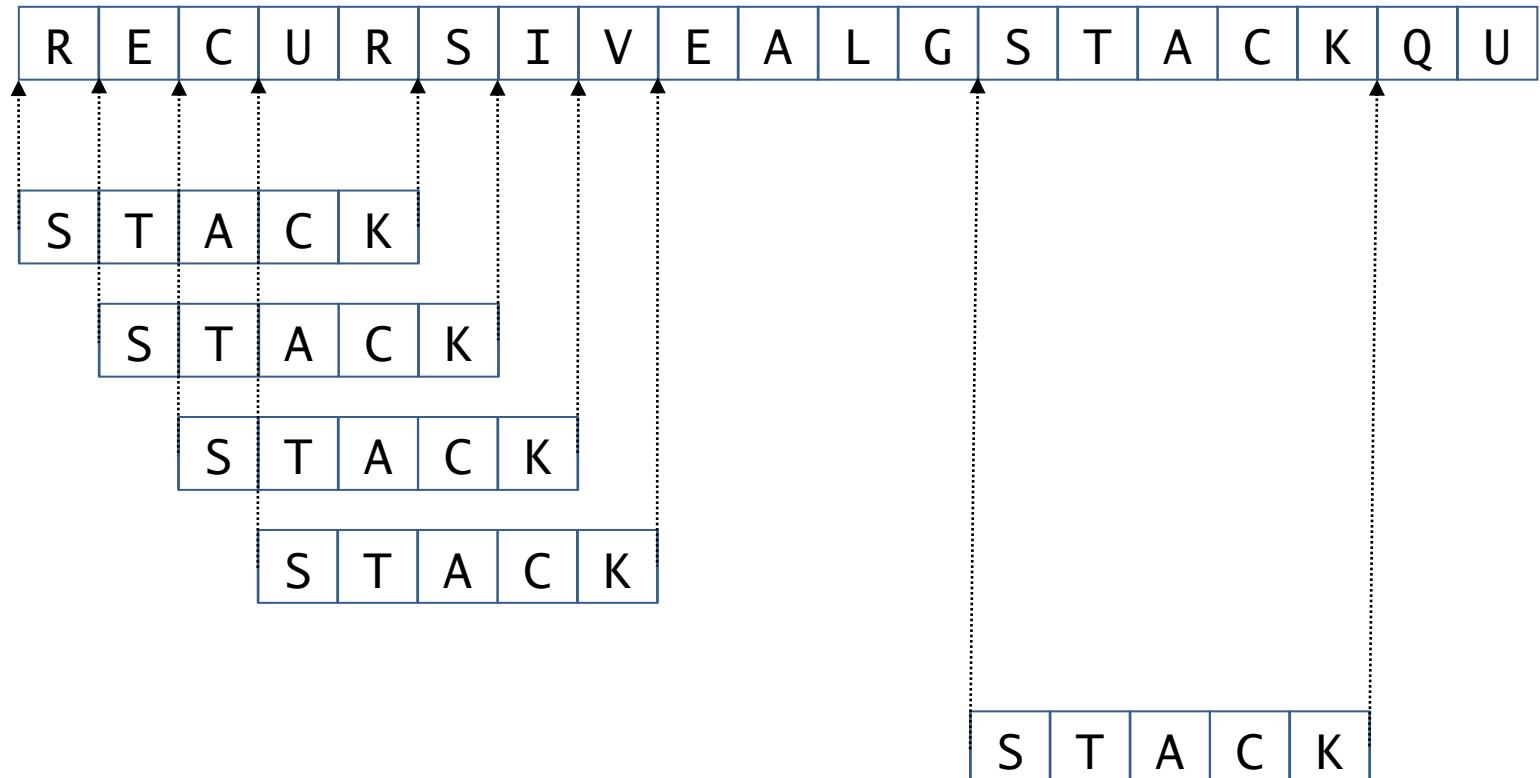
Ch. 32 String Matching : KMP algorithm



국민대학교 컴퓨터공학부 최준수
Modified by 임은진

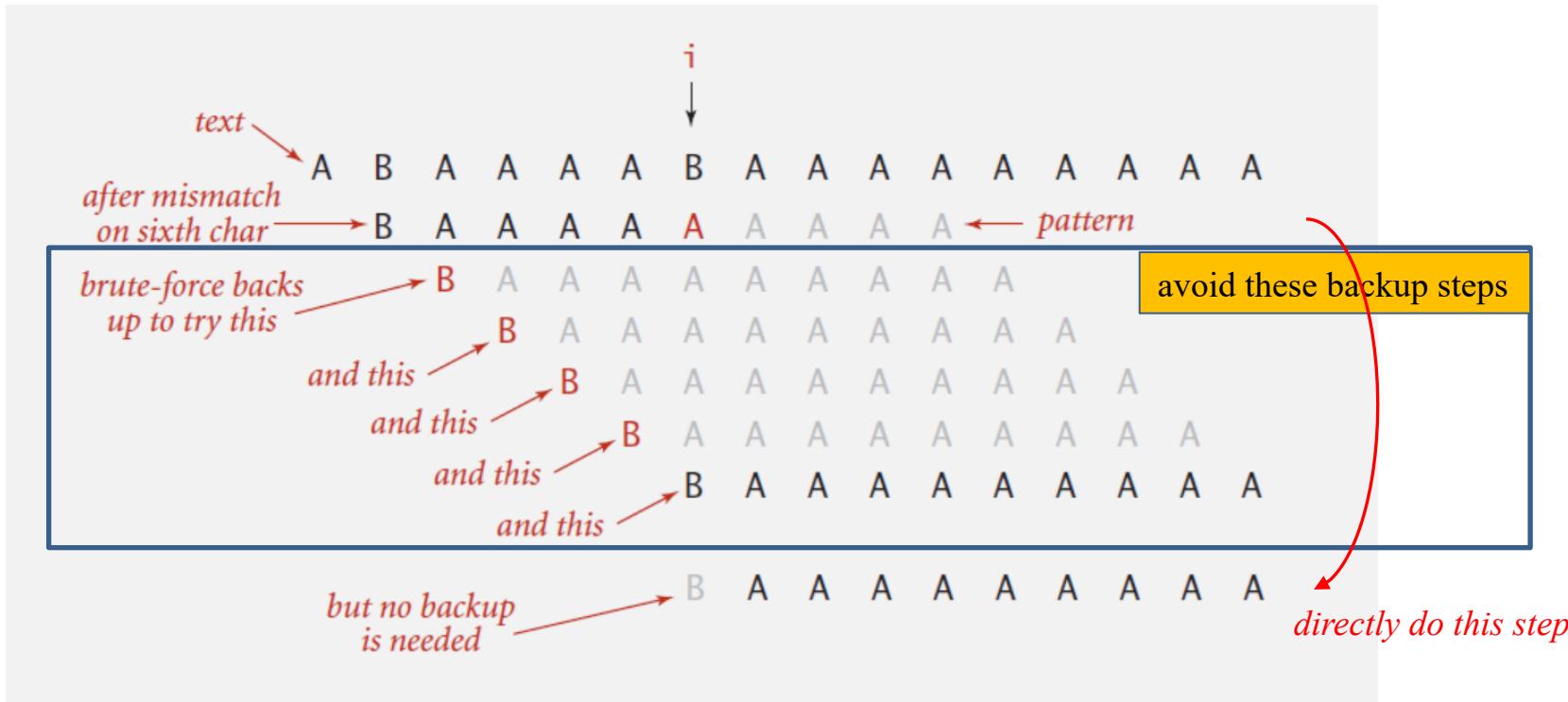
Brute-Force String Matching $O(MN)$

- Naïve Algorithm
 - Check for pattern starting at each text position



String Matching with DFA $O(N)$

- Clever method to avoid **backup** problem.



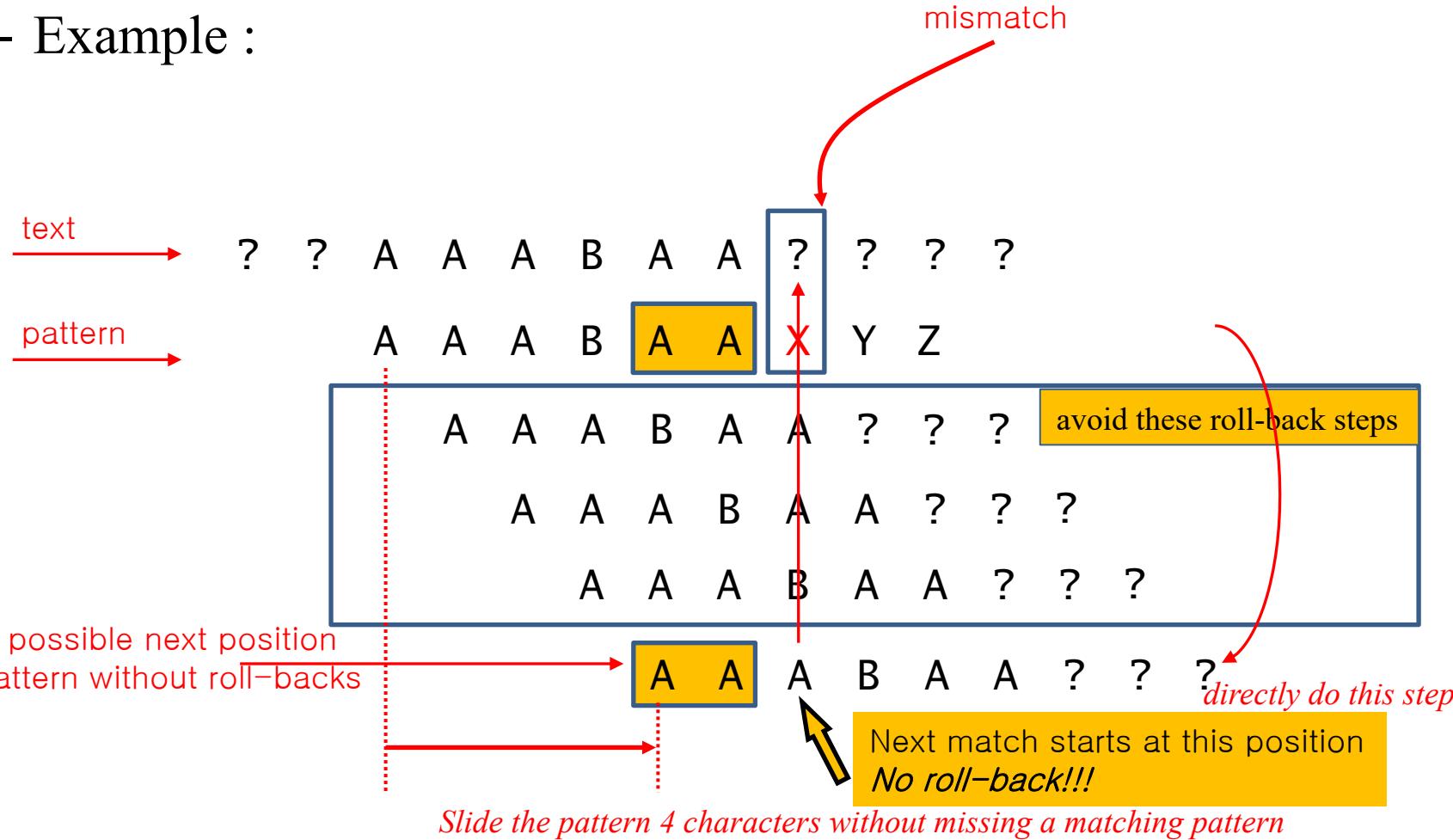
String Matching algorithms

Algorithm	Preprocessing time	Matching time
Naive	0	$O((n - m + 1)m)$
Rabin-Karp	$\Theta(m)$	$O((n - m + 1)m)$
Finite automaton	$O(m \Sigma)$	$\Theta(n)$
Knuth-Morris-Pratt	$\Theta(m)$	$\Theta(n)$

Knuth-Morris-Pratt(KMP) Algorithm

- Avoid roll-backs in naïve algorithm:

- Example :



Prefix and Proper Suffix of the Prefix

pattern

A A A B A A X Y Z

Prefix

Proper Suffix of the Prefix “AAABAA”

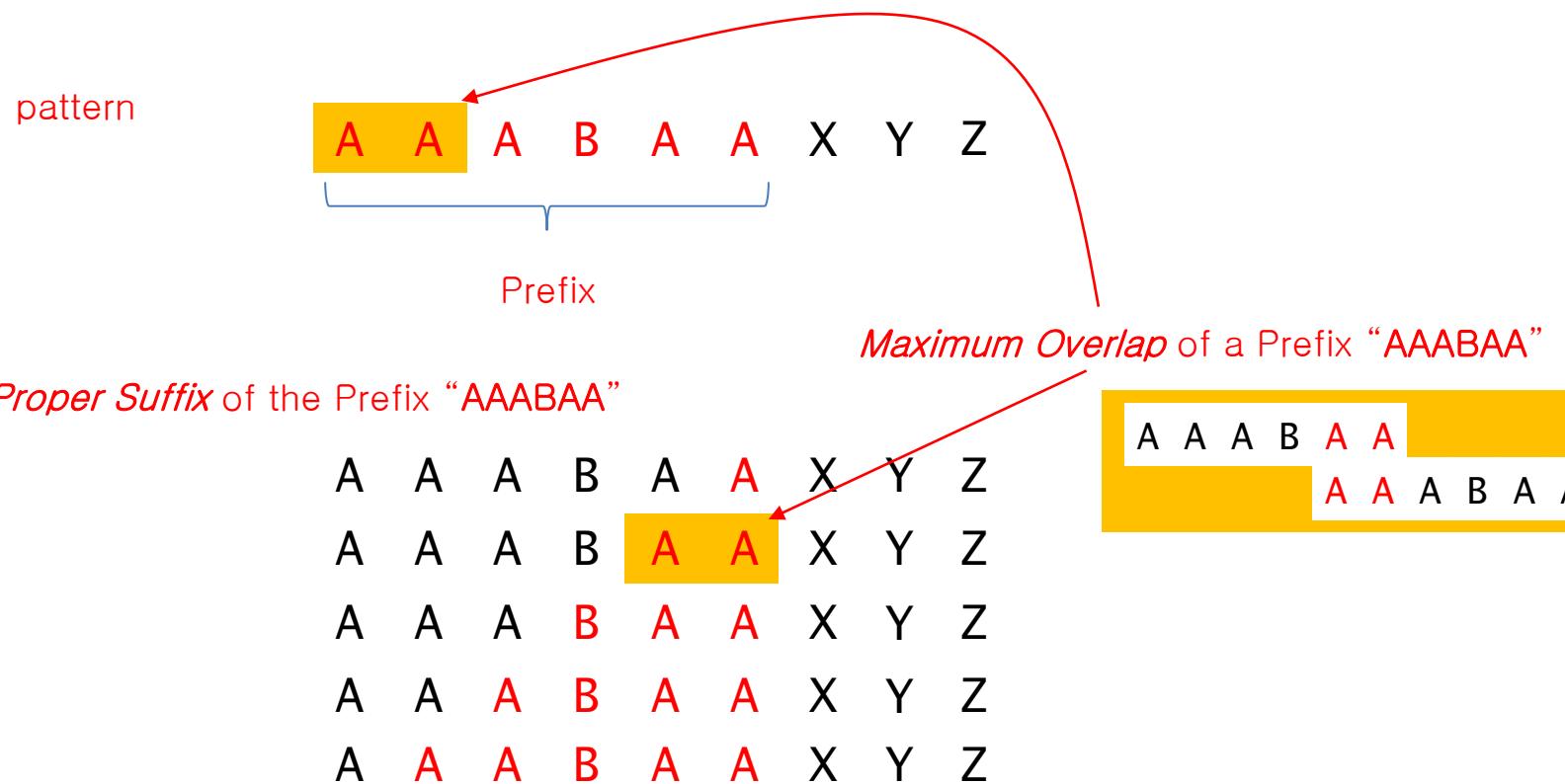
A A A B A A X Y Z
A A A B A A X Y Z
A A A B A A X Y Z
A A A B A A X Y Z
A A A B A A X Y Z

A A A B A A X Y Z

←
Not a *proper* suffix of the prefix
(the same as the prefix)

Maximum Overlap of a Prefix

- the longest proper suffix that is equal to prefix of the prefix



Examples : Maximum Overlap of a Prefix

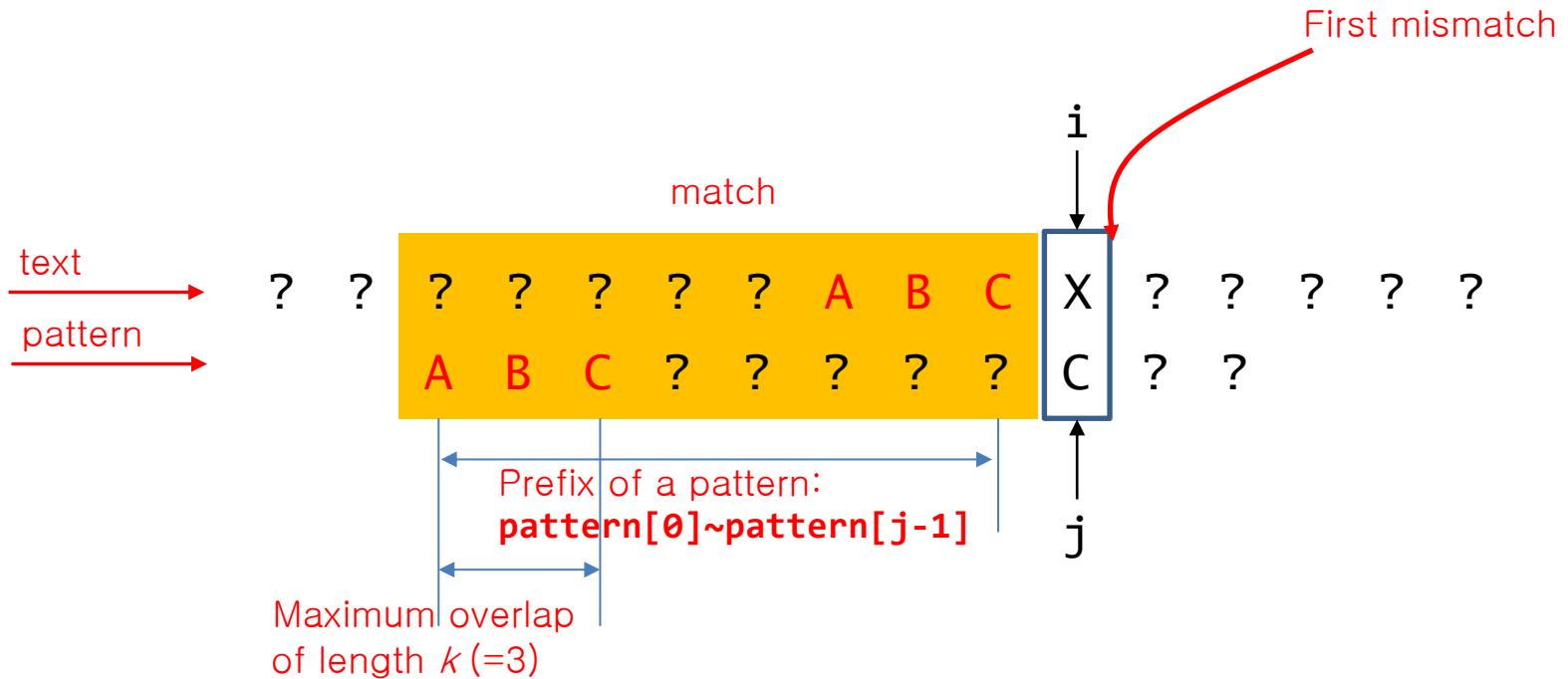
Prefix	Maximum Overlap
AAAAAA	AAAA
AABA	A
AAAB	
ABABABAB	ABABAB

not AAAAA

NULL String

maximum overlap of a prefix in KMP Algorithm

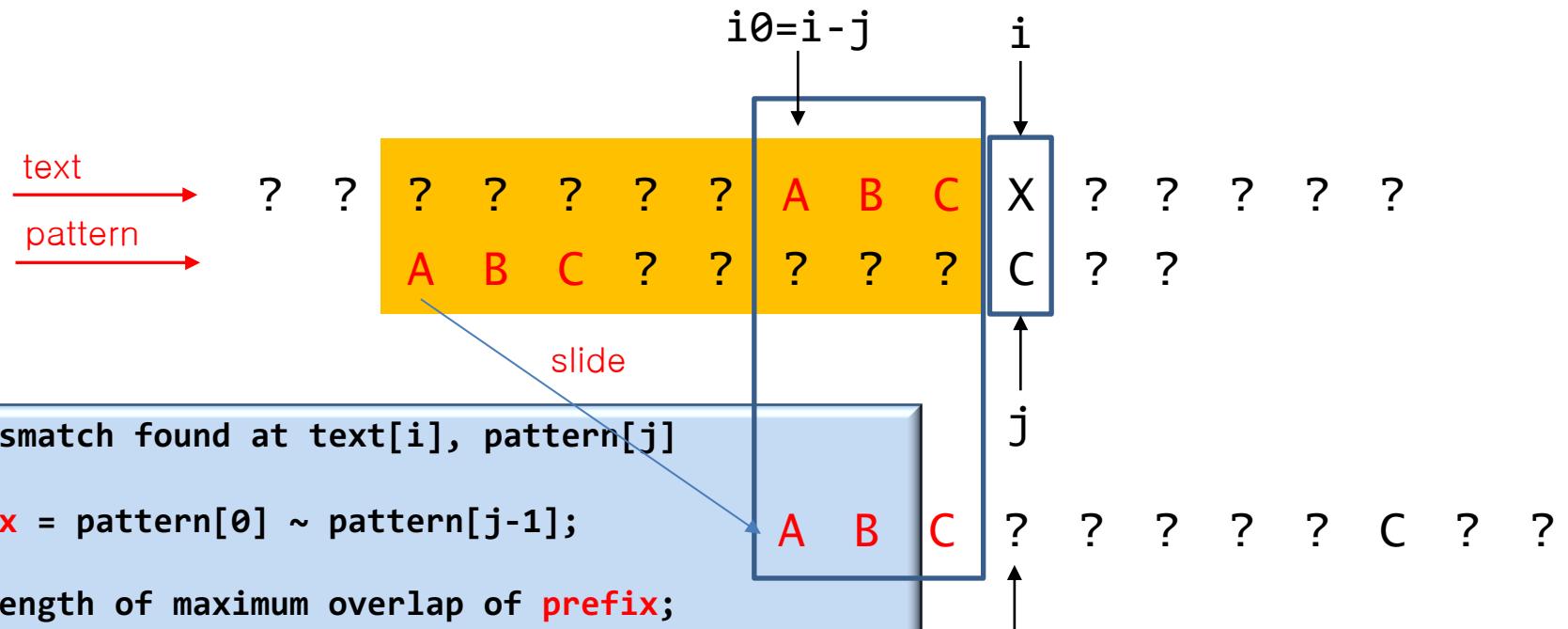
- Reuse of prefix information when there is a mismatch
 - Mismatch at **text[i]** and **pattern[j]**



Note that if the mismatched *location* is **pattern[j]**, then *prefix* is:
pattern[0]~pattern[j-1]

KMP Algorithm

- Then we can slide the pattern so that the *suffix and prefix aligns without missing out on a match*:

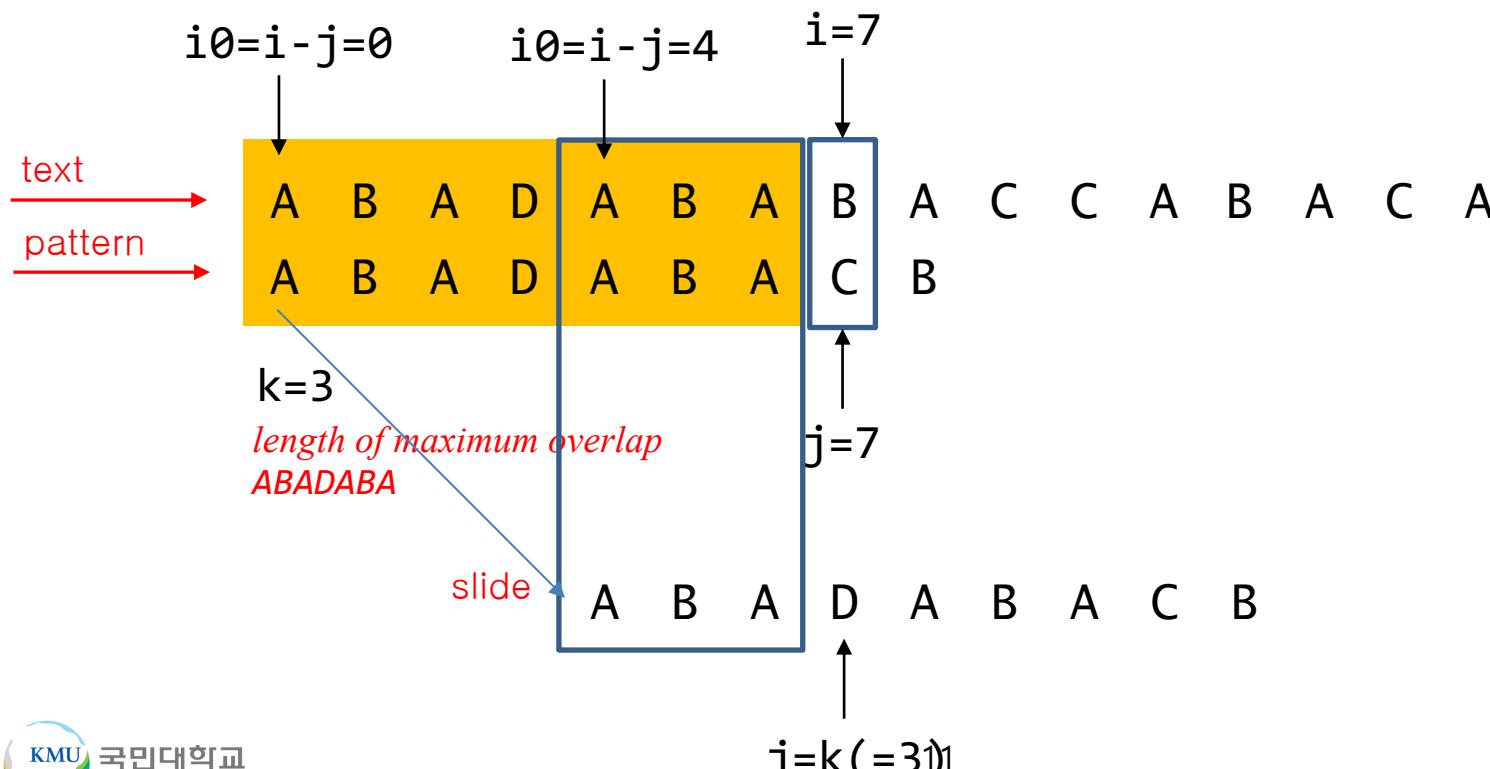


*length of maximum overlap
ABC?????*

KMP Algorithm

- Fast sliding algorithm:

```
// mismatch found at text[i], pattern[j]  
  
prefix = pattern[0] ~ pattern[j-1];  
  
k = Length of maximum overlap of prefix;  
  
j = k;  
// i is unchanged !  
  
// Matched position i0 in text starts from (i - j);  
i0 = i - j;
```



failure function in KMP Algorithm

- Failure function:
 - M : the length of a pattern
 - For $0 < k < M$, the **failure function** $fail(k)$ is the length of maximum overlap of a prefix $\text{pattern}[0] \sim \text{pattern}[k]$
 - Note that $fail(0) = 0$

banabana	k	prefix	$fail(k)$
	0	b	0
	1	ba	0
	2	ban	0
	3	bana	0
	4	banab	1
	5	banaba	2
	6	banab	3
	7	banabana	4

KMP Algorithm

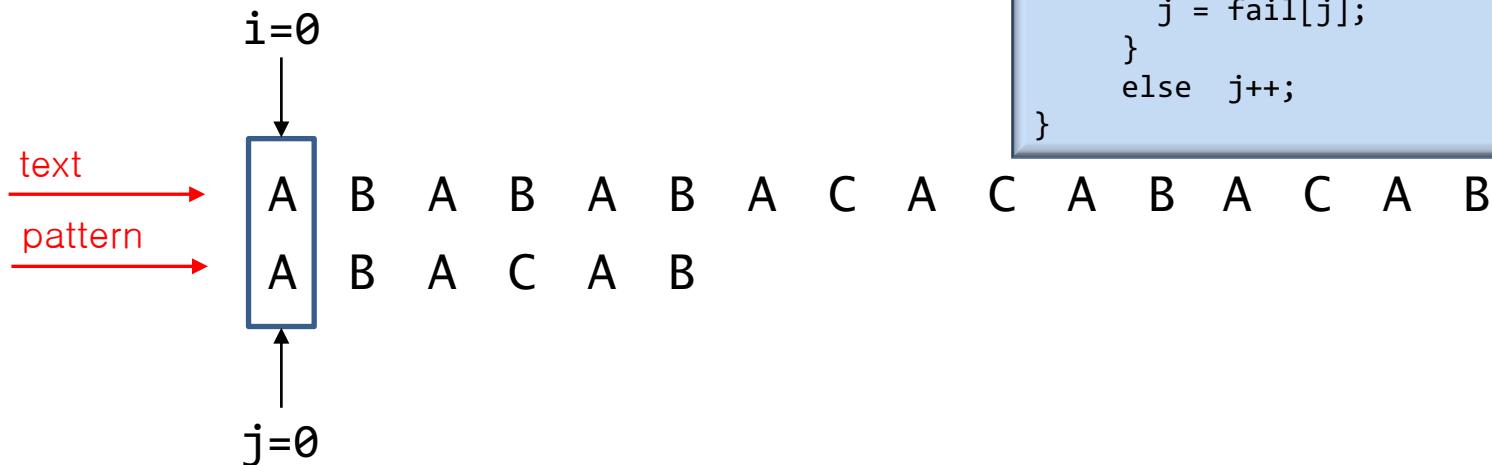
- Knuth-Morris-Pratt(KMP) Algorithm

```
vector<int> kmp(string text, string pattern)
{
    vector<int> ans;
    fail = getFail(pattern);          // failure function
    int n = (int) text.size(), m = (int) pattern.size();
    int j =0;                         // j : index of pattern

    for(int i = 0 ; i < n ; i++)     // i : index of text
    {
        while(j>0 && text[i] != pattern[j])
            j = fail[j-1];
        if(text[i] == pattern[j])
            if(j==m-1)                // pattern matching is found
            {
                ans.push_back(i-j); // save the matched position
                j = fail[j];
            }
        else
            j++;
    }
    return ans;
}
```

KMP Algorithm

- Example:



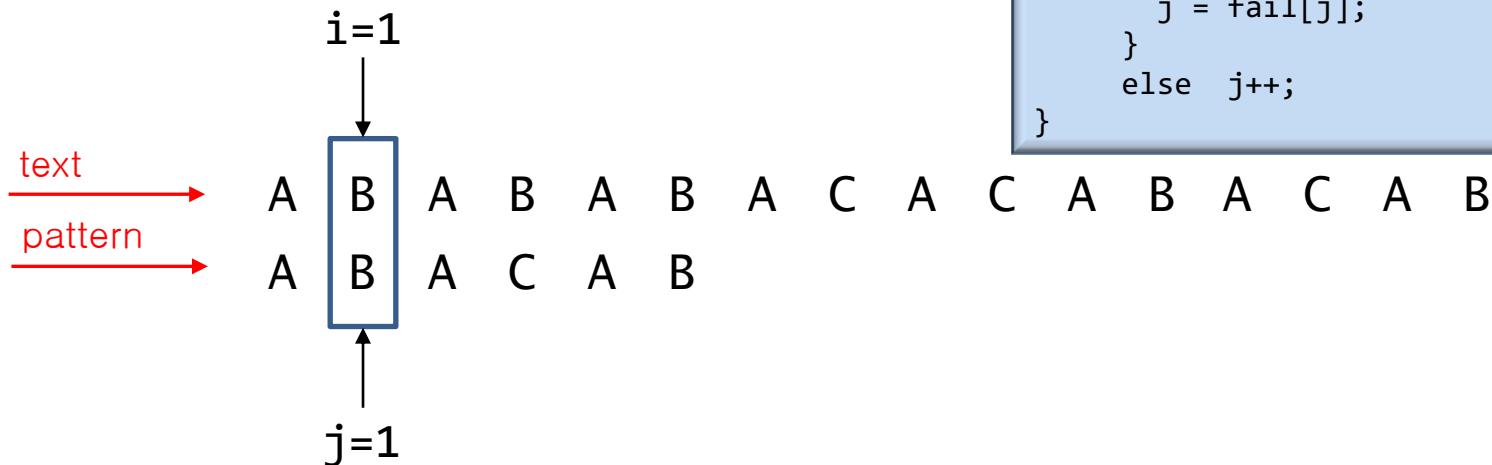
```
for(int i = 0 ; i < n ; i++){
    while(j>0 && text[i] != pattern[j])
        j = fail[j-1];
    if(text[i] == pattern[j])
        if(j==m-1){
            ans.push_back(i-j);
            j = fail[j];
        }
        else j++;
}
```

text[i] == pattern[j]
→ i++, j++

i	0	1	2	3	4	5
fail(i)	0	0	1	0	1	2

KMP Algorithm

- Example:



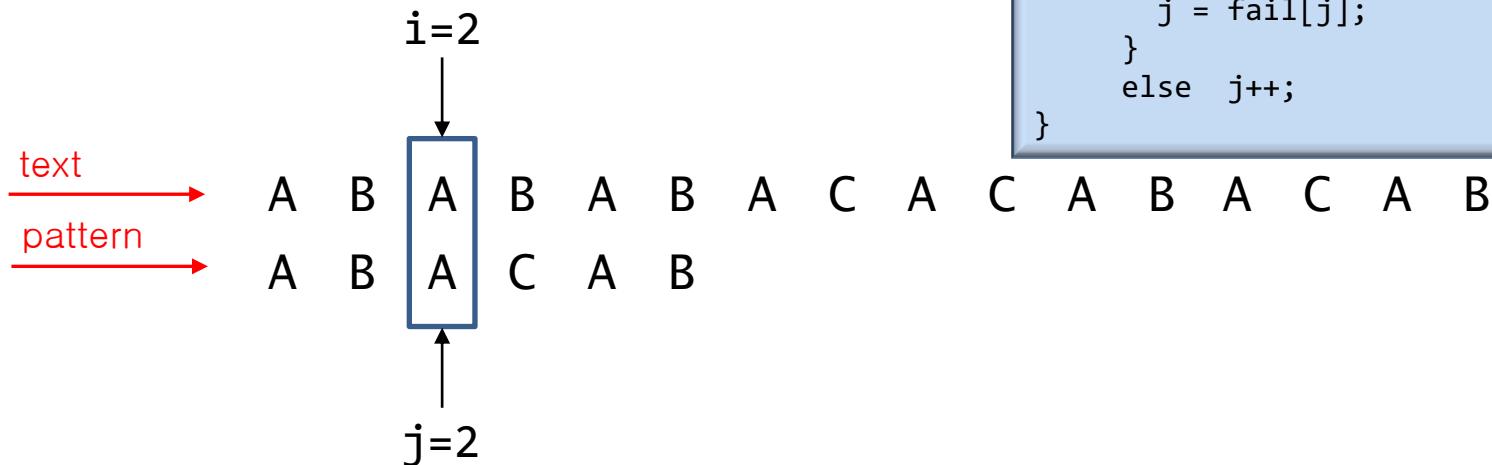
```
for(int i = 0 ; i < n ; i++){
    while(j>0 && text[i] != pattern[j])
        j = fail[j-1];
    if(text[i] == pattern[j])
        if(j==m-1){
            ans.push_back(i-j);
            j = fail[j];
        }
        else j++;
}
```

text[i] == pattern[j]
→ i++, j++

i	0	1	2	3	4	5
fail(i)	0	0	1	0	1	2

KMP Algorithm

- Example:



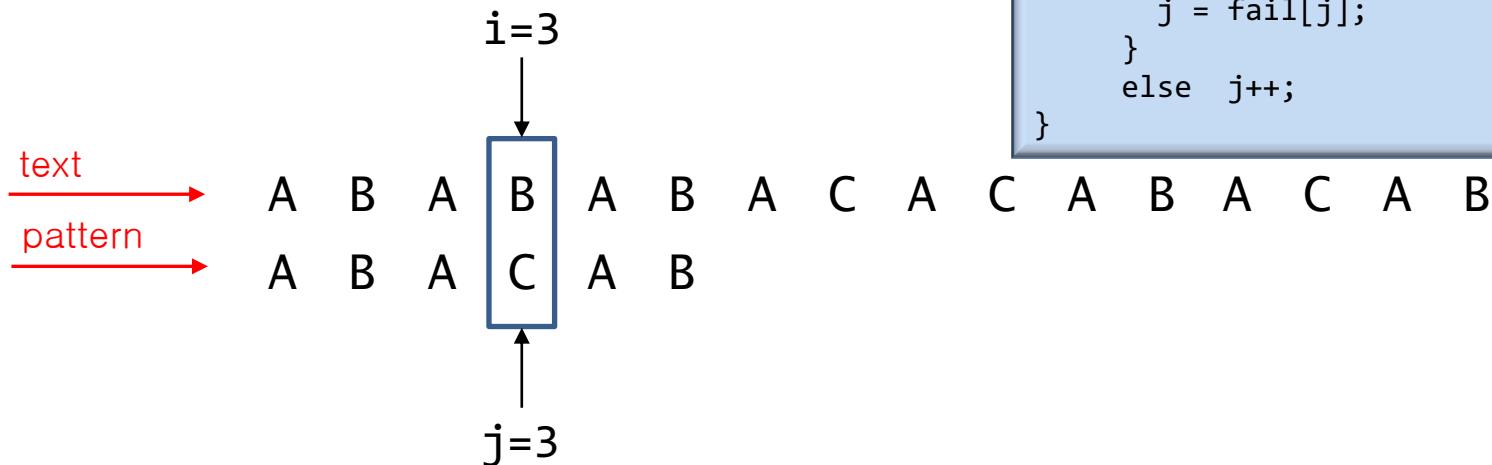
```
for(int i = 0 ; i < n ; i++){
    while(j>0 && text[i] != pattern[j])
        j = fail[j-1];
    if(text[i] == pattern[j])
        if(j==m-1){
            ans.push_back(i-j);
            j = fail[j];
        }
        else j++;
}
```

$\text{text}[i] == \text{pattern}[j]$
→ $i++$, $j++$

i	0	1	2	3	4	5
fail(i)	0	0	1	0	1	2

KMP Algorithm

- Example:



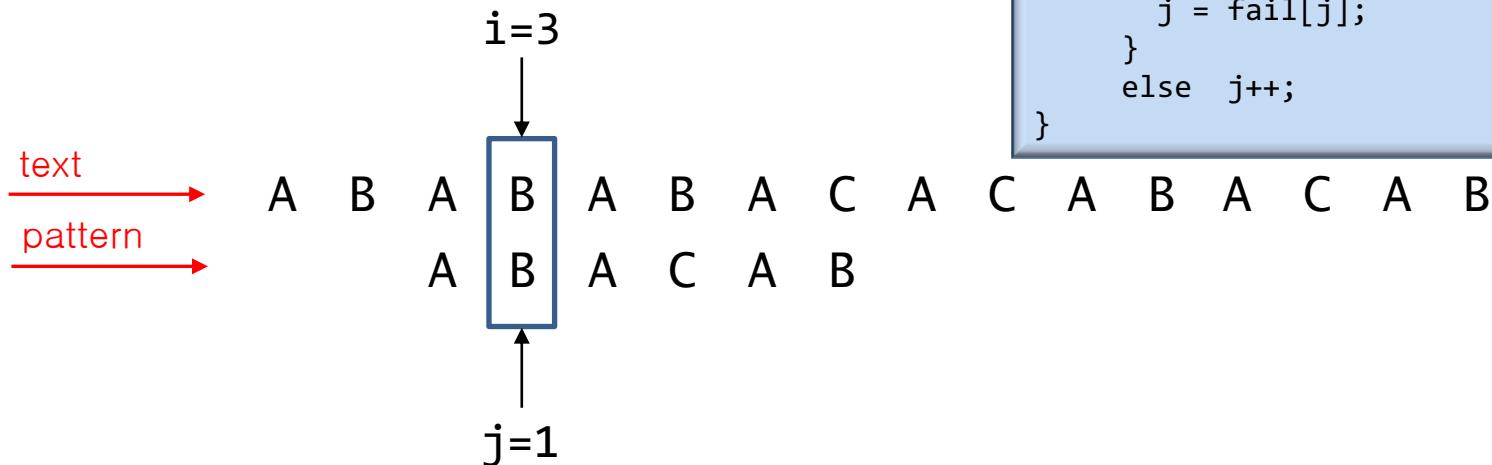
```
for(int i = 0 ; i < n ; i++){
    while(j>0 && text[i] != pattern[j])
        j = fail[j-1];
    if(text[i] == pattern[j])
        if(j==m-1){
            ans.push_back(i-j);
            j = fail[j];
        }
        else j++;
}
```

$\text{text}[i] \neq \text{pattern}[j]$
→ $j = \text{fail}[j-1]$

i	0	1	2	3	4	5
fail(i)	0	0	1	0	1	2

KMP Algorithm

- Example:



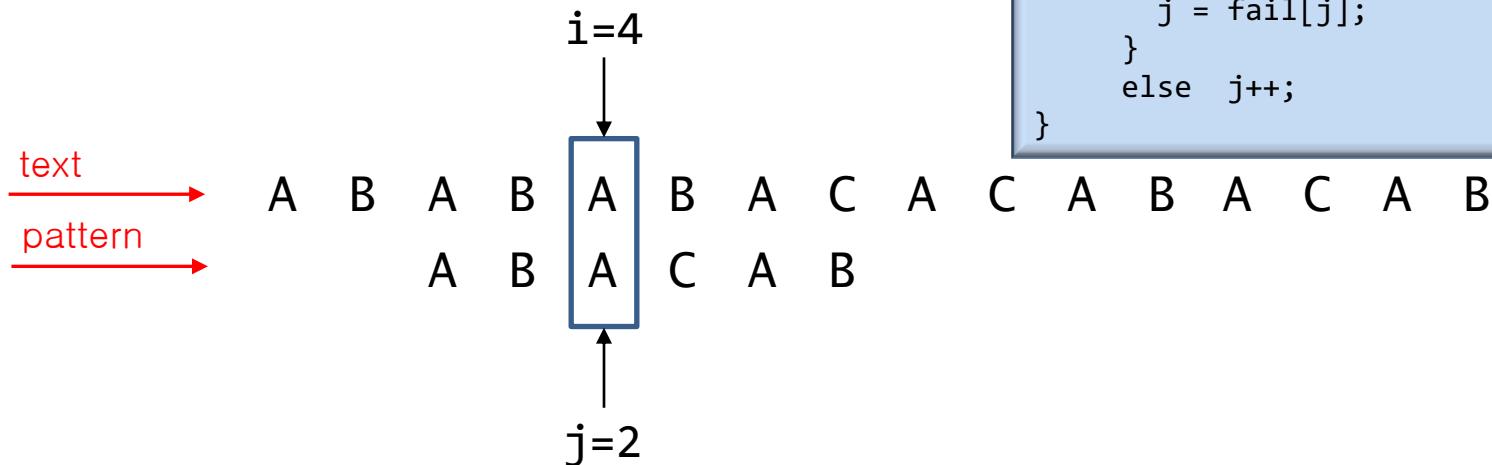
```
for(int i = 0 ; i < n ; i++){  
    while(j>0 && text[i] != pattern[j])  
        j = fail[j-1];  
    if(text[i] == pattern[j])  
        if(j==m-1){  
            ans.push_back(i-j);  
            j = fail[j];  
        }  
        else j++;  
}
```

text[i] == pattern[j]
→ i++, j++

i	0	1	2	3	4	5
fail(i)	0	0	1	0	1	2

KMP Algorithm

- Example:



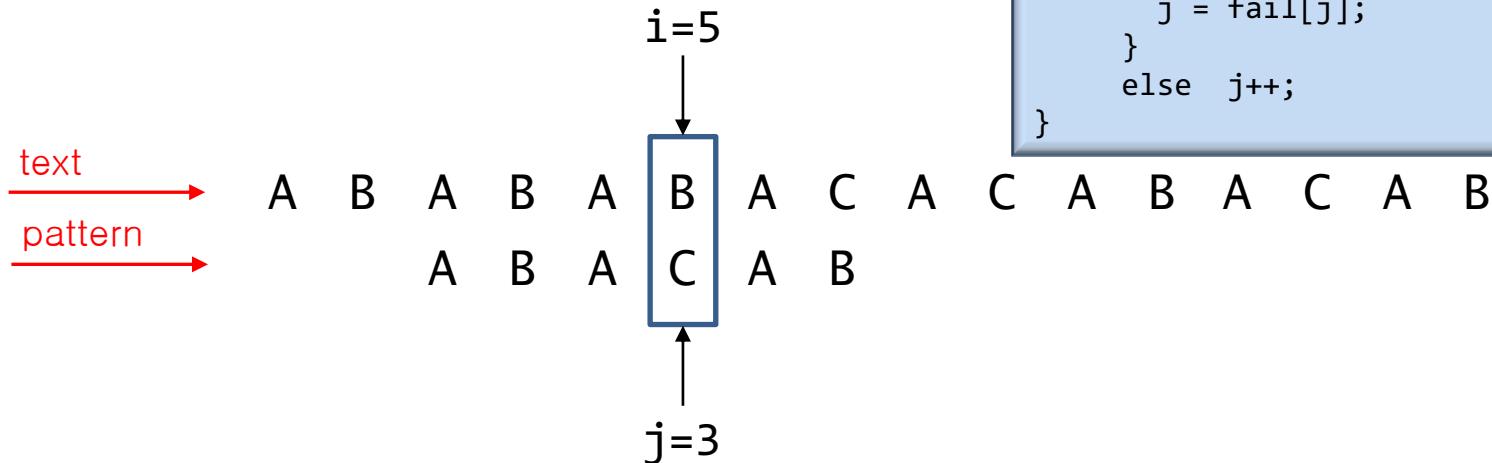
```
for(int i = 0 ; i < n ; i++){
    while(j>0 && text[i] != pattern[j])
        j = fail[j-1];
    if(text[i] == pattern[j])
        if(j==m-1){
            ans.push_back(i-j);
            j = fail[j];
        }
        else j++;
}
```

$\text{text}[i] == \text{pattern}[j]$
→ $i++$, $j++$

i	0	1	2	3	4	5
fail(i)	0	0	1	0	1	2

KMP Algorithm

- Example:



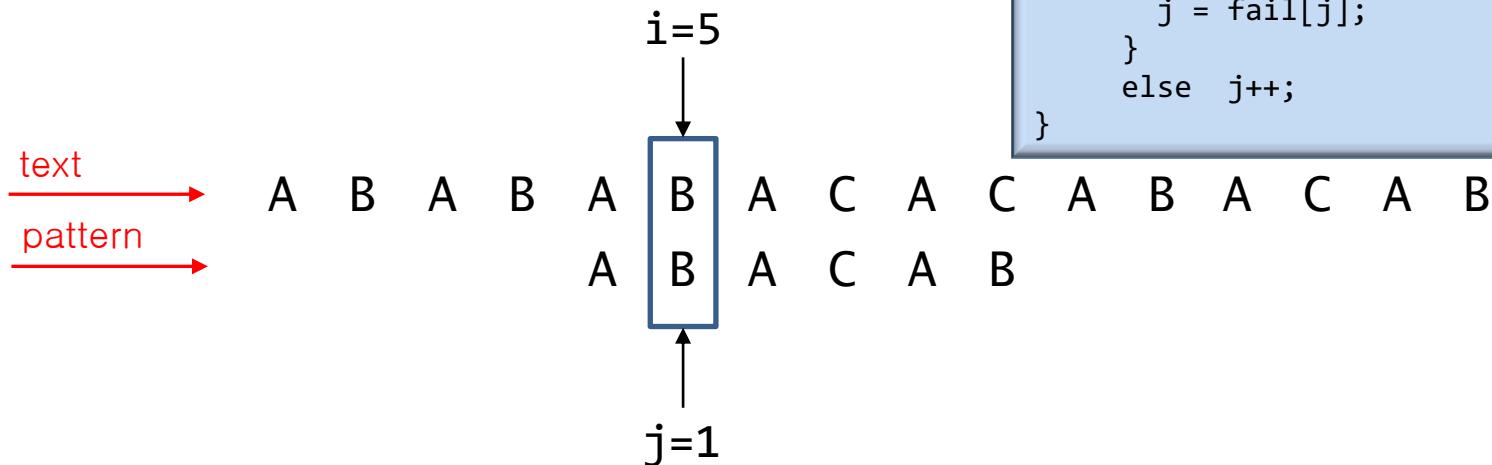
```
for(int i = 0 ; i < n ; i++){
    while(j>0 && text[i] != pattern[j])
        j = fail[j-1];
    if(text[i] == pattern[j])
        if(j==m-1){
            ans.push_back(i-j);
            j = fail[j];
        }
        else j++;
}
```

$\text{text}[i] \neq \text{pattern}[j]$
→ $j = \text{fail}[j-1]$

i	0	1	2	3	4	5
fail(i)	0	0	1	0	1	2

KMP Algorithm

- Example:



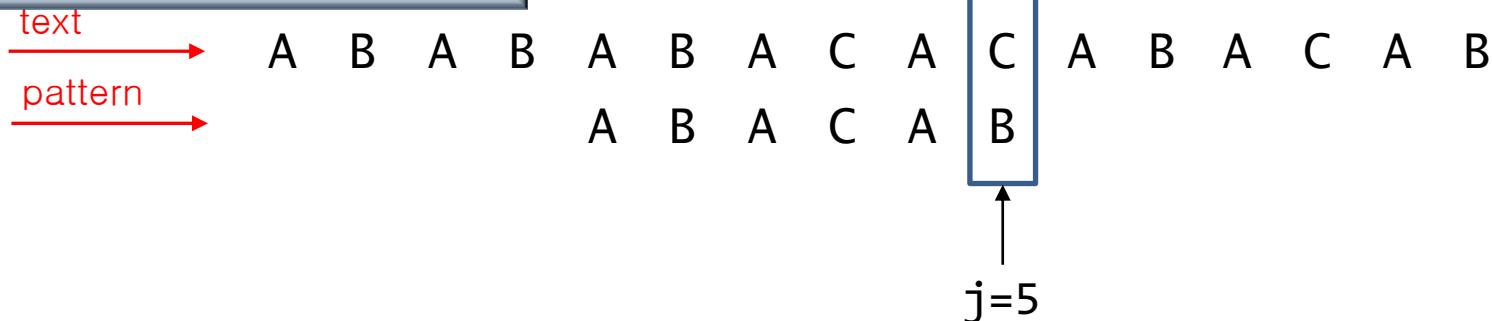
```
for(int i = 0 ; i < n ; i++){
    while(j>0 && text[i] != pattern[j])
        j = fail[j-1];
    if(text[i] == pattern[j])
        if(j==m-1){
            ans.push_back(i-j);
            j = fail[j];
        }
        else j++;
}
```

$\text{text}[i] == \text{pattern}[j]$
→ $i++$, $j++$

i	0	1	2	3	4	5
fail(i)	0	0	1	0	1	2

KMP Algorithm

```
for(int i = 0 ; i < n ; i++){
    while(j>0 && text[i] != pattern[j])
        j = fail[j-1];
    if(text[i] == pattern[j])
        if(j==m-1){
            ans.push_back(i-j);
            j = fail[j];
        }
        else j++;
}
```



`text[i] != pattern[j]`
→ `j = fail[j-1]`

i	0	1	2	3	4	5
fail(i)	0	0	1	0	1	2

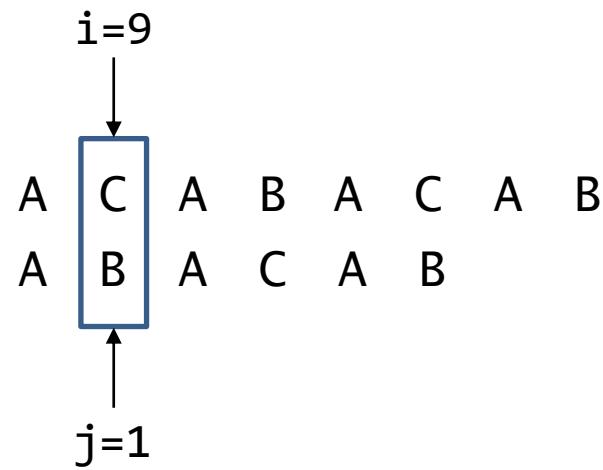
KMP Algorithm

```
for(int i = 0 ; i < n ; i++){
    while(j>0 && text[i] != pattern[j])
        j = fail[j-1];
    if(text[i] == pattern[j])
        if(j==m-1){
            ans.push_back(i-j);
            j = fail[j];
        }
        else    j++;
}
```

text
pattern

A B A B A B A C A

A B A C A B



text[i] != pattern[j]

→ j = fail[j-1]

i	0	1	2	3	4	5
fail(i)	0	0	1	0	1	2

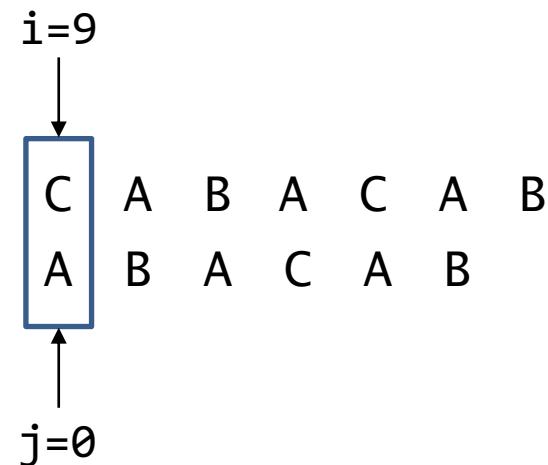
KMP Algorithm

```
for(int i = 0 ; i < n ; i++){
    while(j>0 && text[i] != pattern[j])
        j = fail[j-1];
    if(text[i] == pattern[j])
        if(j==m-1){
            ans.push_back(i-j);
            j = fail[j];
        }
        else    j++;
}
```

text
_____→
pattern
_____→

A B A B A B A C A

B A C A B



j==0 && text[i] != pattern[j]

→ i++

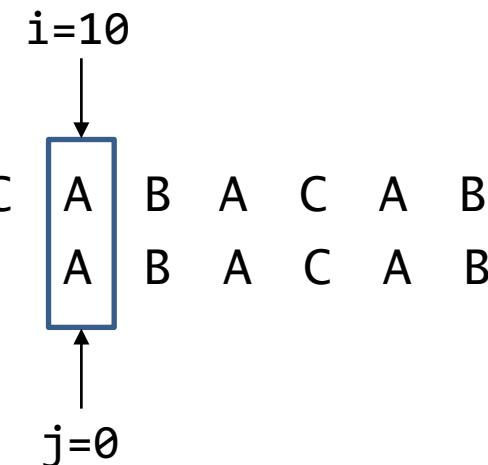
i	0	1	2	3	4	5
fail(i)	0	0	1	0	1	2

KMP Algorithm

```
for(int i = 0 ; i < n ; i++){
    while(j>0 && text[i] != pattern[j])
        j = fail[j-1];
    if(text[i] == pattern[j])
        if(j==m-1){
            ans.push_back(i-j);
            j = fail[j];
        }
        else j++;
}
```

text
_____→
pattern
_____→

A B A B A B A C A C A B



text[i] == pattern[j]
→ i++, j++

i	0	1	2	3	4	5
fail(i)	0	0	1	0	1	2

KMP Algorithm

```
for(int i = 0 ; i < n ; i++){
    while(j>0 && text[i] != pattern[j])
        j = fail[j-1];
    if(text[i] == pattern[j])
        if(j==m-1){
            ans.push_back(i-j);
            j = fail[j];
        }
        else j++;
}
```

text

pattern

A B A B

A B A C A

C A B A C A

B
B

$$i_0 = i - j = 10$$

$$i = 15$$

$$j = 5$$

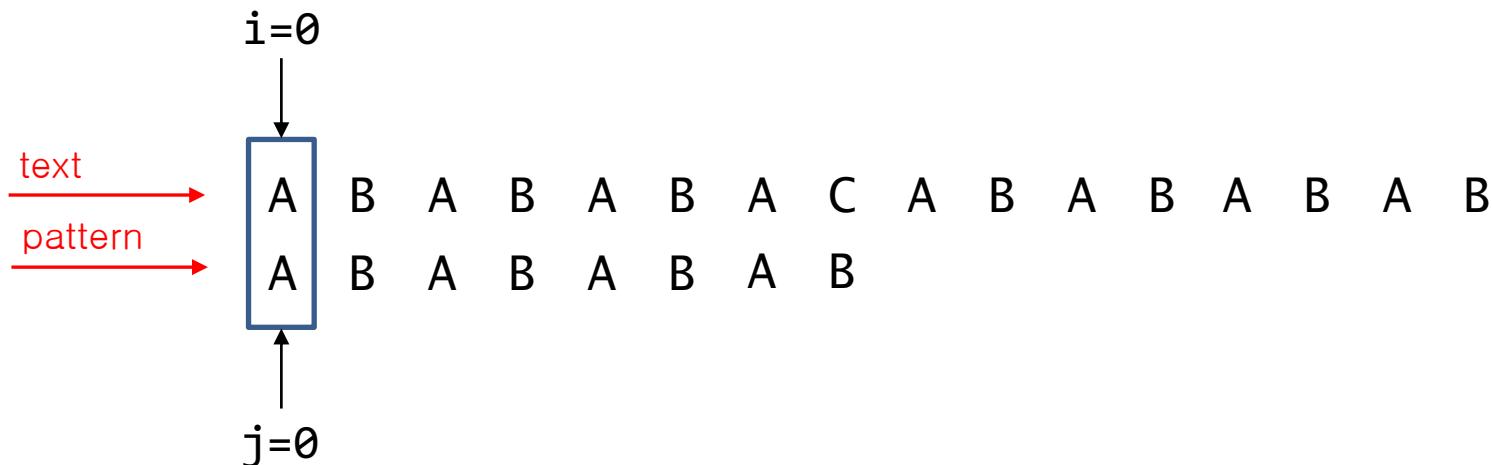
i	0	1	2	3	4	5
fail(i)	0	0	1	0	1	2

another example

KMP Algorithm

- Knuth-Morris-Pratt(KMP) Algorithm
 - Why?

```
while(j>0 && text[i] != pattern[j])
    j = fail[j-1];
```



$\text{text}[i] == \text{pattern}[j]$

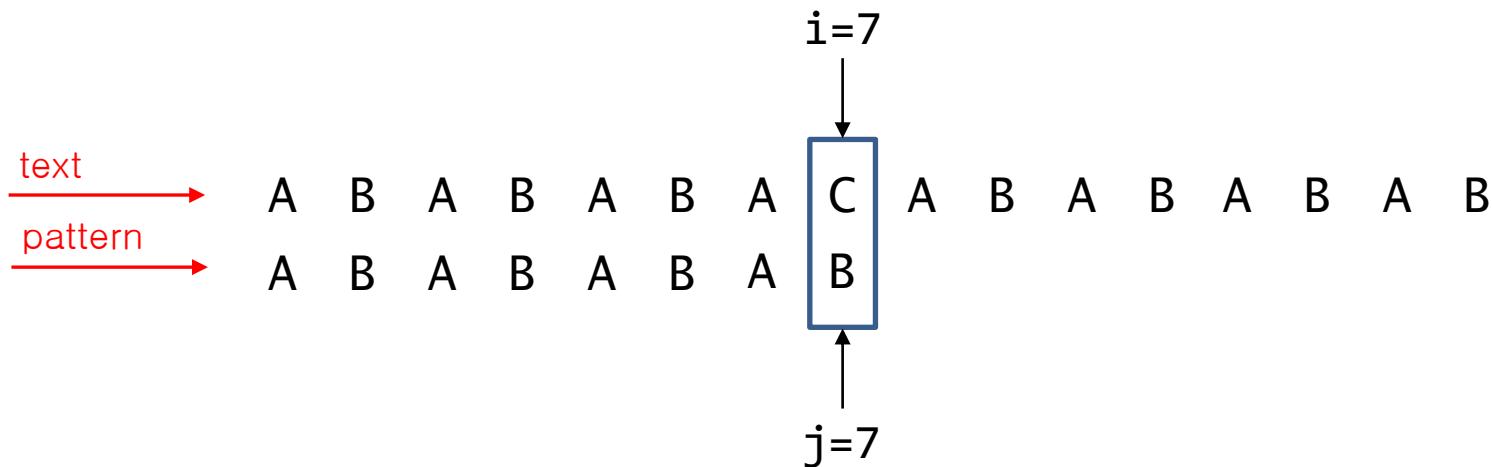
→ $i++$, $j++$

i	0	1	2	3	4	5	6	7
fail(i)	0	0	1	2	3	4	5	6

KMP Algorithm

- Knuth-Morris-Pratt(KMP) Algorithm
 - Why?

```
while(j>0 && text[i] != pattern[j])
    j = fail[j-1];
```



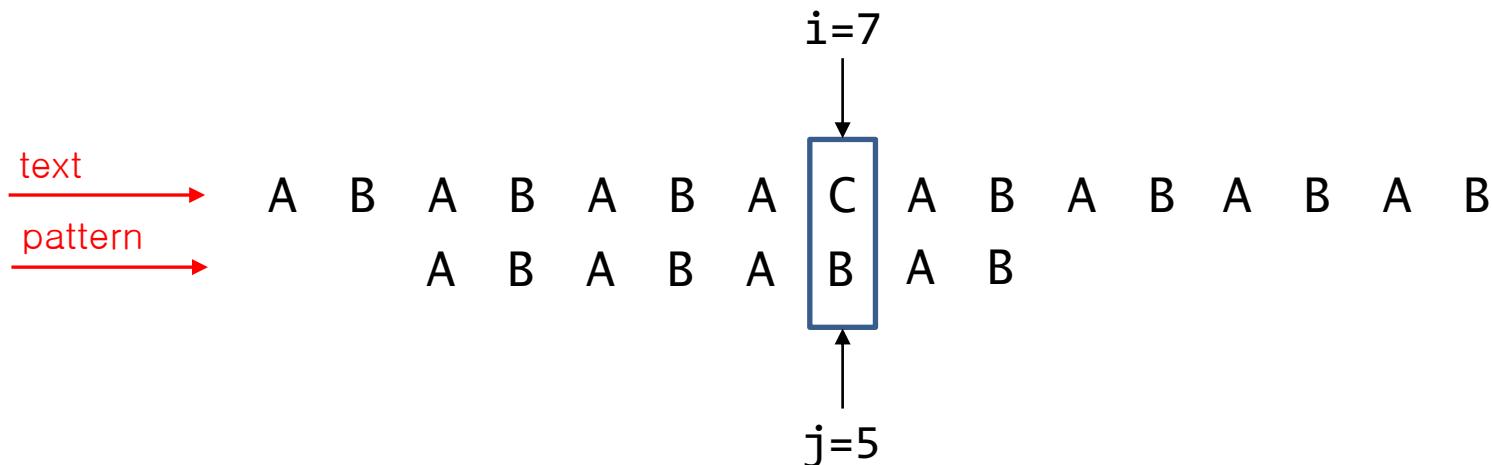
text[i] != pattern[j]
→ j = fail[j-1]

i	0	1	2	3	4	5	6	7
fail(i)	0	0	1	2	3	4	5	6

KMP Algorithm

- Knuth-Morris-Pratt(KMP) Algorithm
 - Why?

```
while(j>0 && text[i] != pattern[j])
    j = fail[j-1];
```



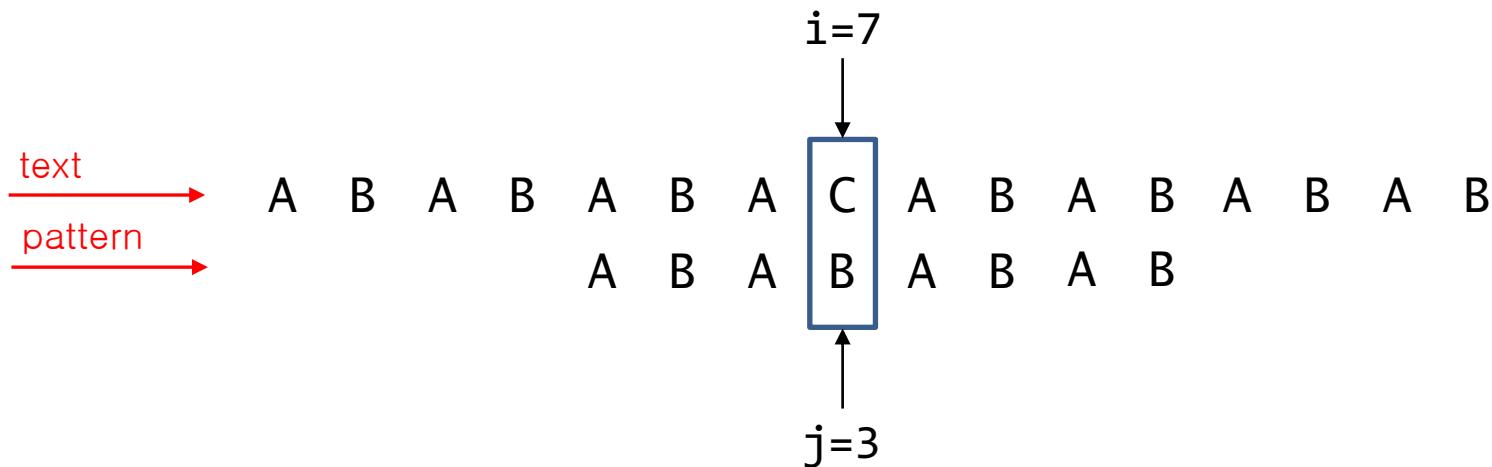
`text[i] != pattern[j]`
⇒ `j = fail[j-1]`

i	0	1	2	3	4	5	6	7
fail(i)	0	0	1	2	3	4	5	6

KMP Algorithm

- Knuth-Morris-Pratt(KMP) Algorithm
 - Why?

```
while(j>0 && text[i] != pattern[j])
    j = fail[j-1];
```



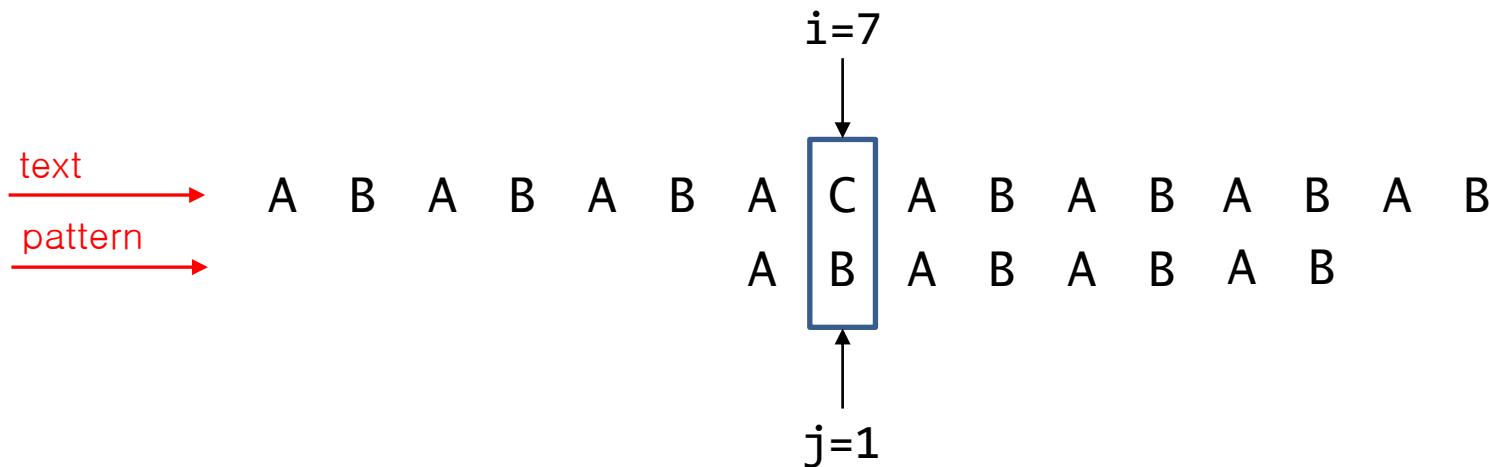
text[i] != pattern[j]
→ j = fail[j-1]

i	0	1	2	3	4	5	6	7
fail(i)	0	0	1	2	3	4	5	6

KMP Algorithm

- Knuth-Morris-Pratt(KMP) Algorithm
 - Why?

```
while(j>0 && text[i] != pattern[j])
    j = fail[j-1];
```



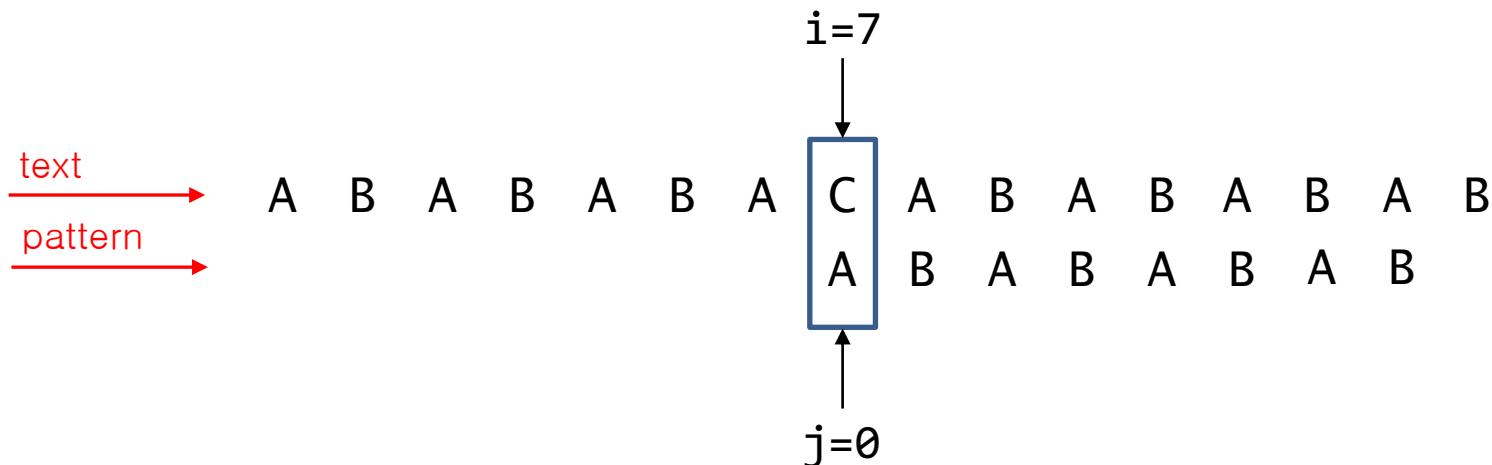
text[i] != pattern[j]
→ j = fail[j-1]

i	0	1	2	3	4	5	6	7
fail(i)	0	0	1	2	3	4	5	6

KMP Algorithm

- Knuth-Morris-Pratt(KMP) Algorithm
 - Why?

```
while(j>0 && text[i] != pattern[j])
    j = fail[j-1];
```

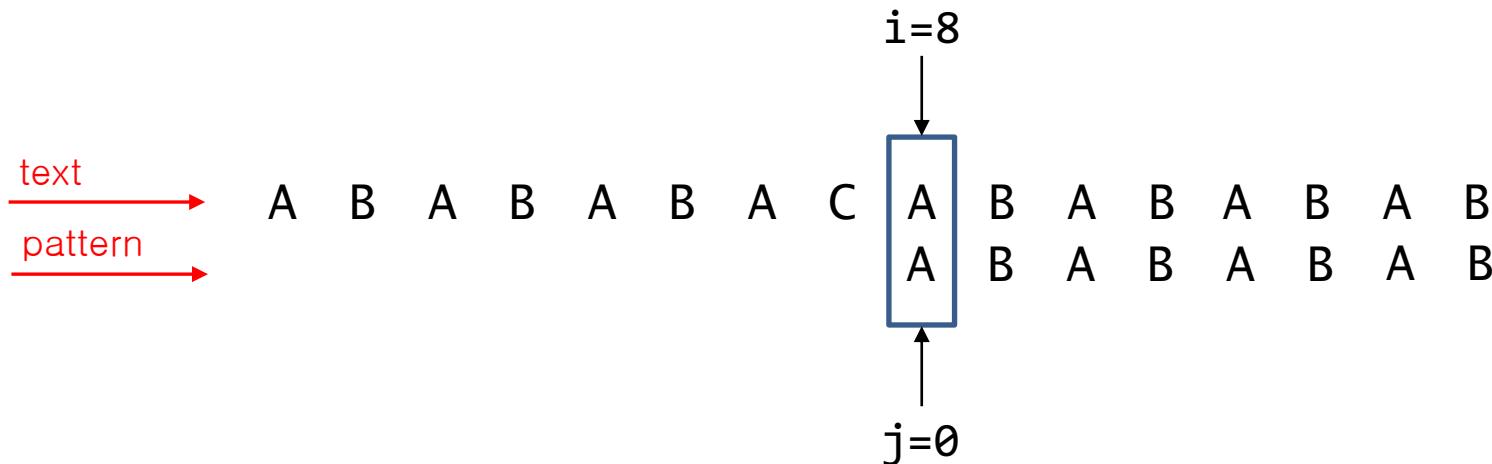


text[i] != pattern[j]
→ j = fail[j-1]

KMP Algorithm

- Knuth-Morris-Pratt(KMP) Algorithm
 - Why?

```
while(j>0 && text[i] != pattern[j])
    j = fail[j-1];
```



$j==0 \&\& \text{text}[i] \neq \text{pattern}[j]$

→ $i++$

i	0	1	2	3	4	5	6	7
fail(i)	0	0	1	2	3	4	5	6

computation of fail[] (= $\pi[]$)

- **getFail()** function
 - Simple **fail[]** computation needs $O(M^3)$ time.
 - $O(M)$ time algorithm: very similar to KMP algorithm itself

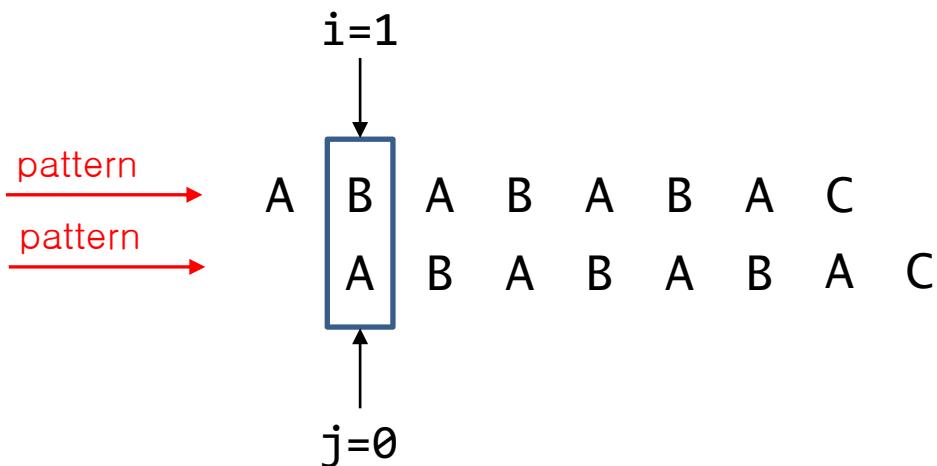
```
vector<int> getFail(string pattern)
{
    int m = (int) pattern.size();
    int j=0;
    vector<int> fail(m, 0);

    for(int i = 1; i< m ; i++)
    {
        while(j > 0 && pattern[i] != pattern[j])
            j = fail[j-1];
        if(pattern[i] == pattern[j])
            fail[i] = ++j;
        else fail[i] = j;
    }
    return fail;
}
```

computation of fail[] (= $\pi[]$)

- **getFail()** function
 - Example:

i	0	1	2	3	4	5	6	7
fail[i]	0	0						



```
vector<int> getFail(string pattern)
{
    int m = (int) pattern.size();
    int j=0;
    vector<int> fail(m, 0);

    for(int i = 1; i< m ; i++)
    {
        while(j > 0 &&
              pattern[i] != pattern[j])
            j = fail[j-1];
        if(pattern[i] == pattern[j])
            fail[i] = ++j;
        else
            fail[i] = j;
    }
    return fail;
}
```

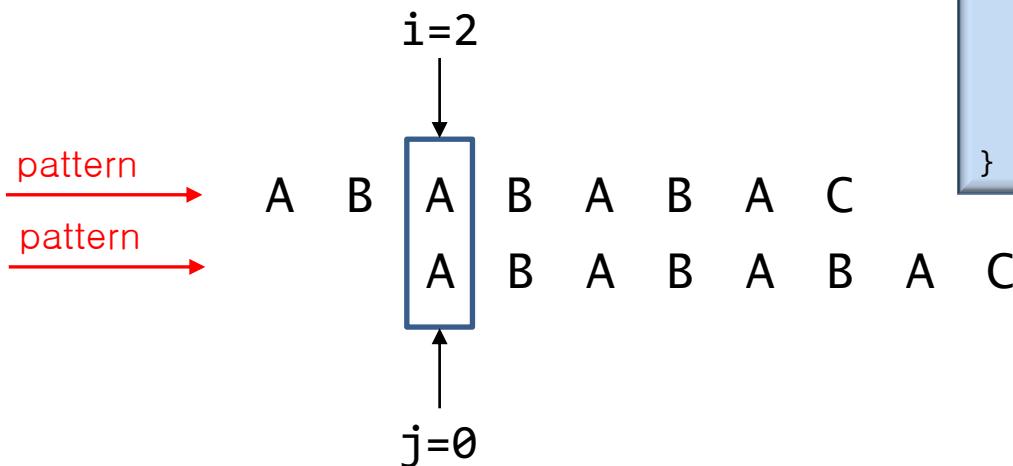
pattern[i] != pattern[j]

→ fail[i] = j
i++

computation of fail[] (= $\pi[]$)

- **getFail()** function
 - Example:

i	0	1	2	3	4	5	6	7
fail[i]	0	0	1					



$pattern[i] == pattern[j]$

→ $fail[i] = ++j$
 $i++$

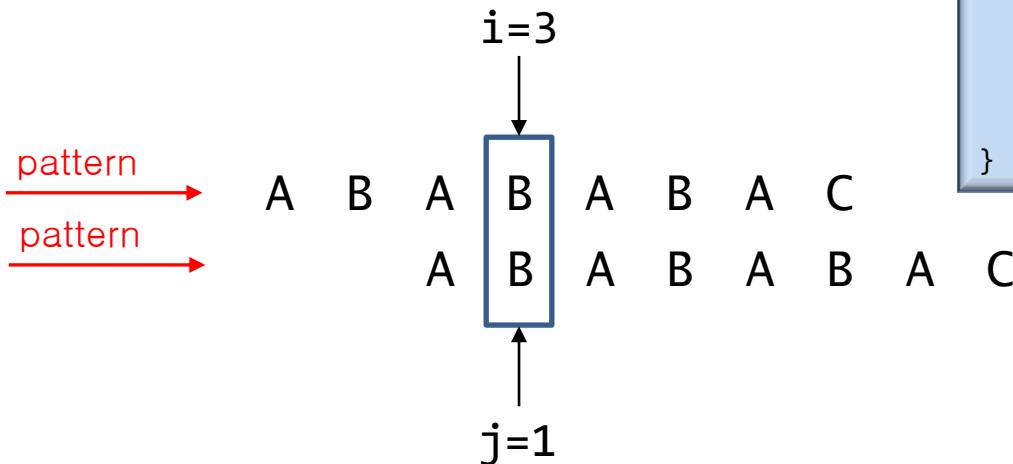
```
vector<int> getFail(string pattern)
{
    int m = (int) pattern.size();
    int j=0;
    vector<int> fail(m, 0);

    for(int i = 1; i< m ; i++)
    {
        while(j > 0 &&
              pattern[i] != pattern[j])
            j = fail[j-1];
        if(pattern[i] == pattern[j])
            fail[i] = ++j;
        else fail[i] = j;
    }
    return fail;
}
```

computation of fail[] (= $\pi[]$)

- **getFail()** function
 - Example:

i	0	1	2	3	4	5	6	7
fail[i]	0	0	1	2				



$\text{pattern}[i] == \text{pattern}[j]$

→ $\text{fail}[i] = ++j$
i++

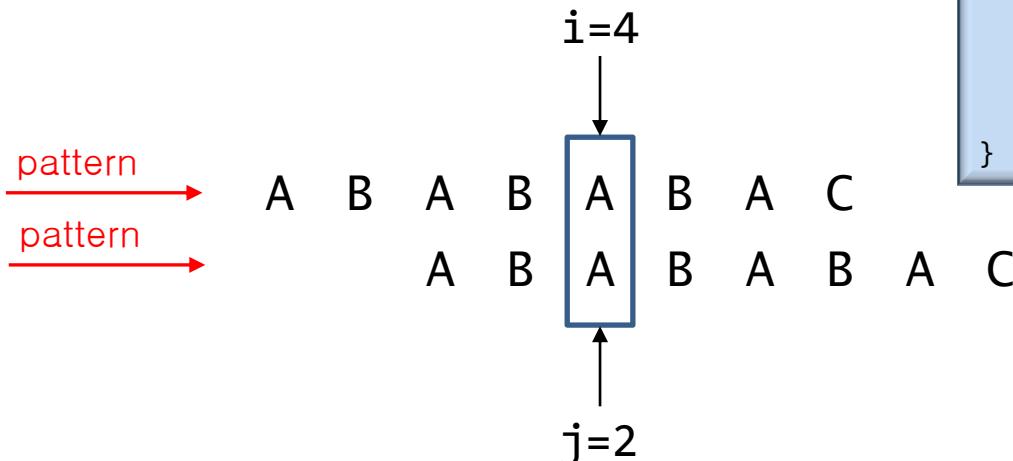
```
vector<int> getFail(string pattern)
{
    int m = (int) pattern.size();
    int j=0;
    vector<int> fail(m, 0);

    for(int i = 1; i< m ; i++)
    {
        while(j > 0 &&
              pattern[i] != pattern[j])
            j = fail[j-1];
        if(pattern[i] == pattern[j])
            fail[i] = ++j;
        else fail[i] = j;
    }
    return fail;
}
```

computation of fail[] (= $\pi[]$)

- **getFail()** function
 - Example:

i	0	1	2	3	4	5	6	7
fail[i]	0	0	1	2	3			



$\text{pattern}[i] == \text{pattern}[j]$

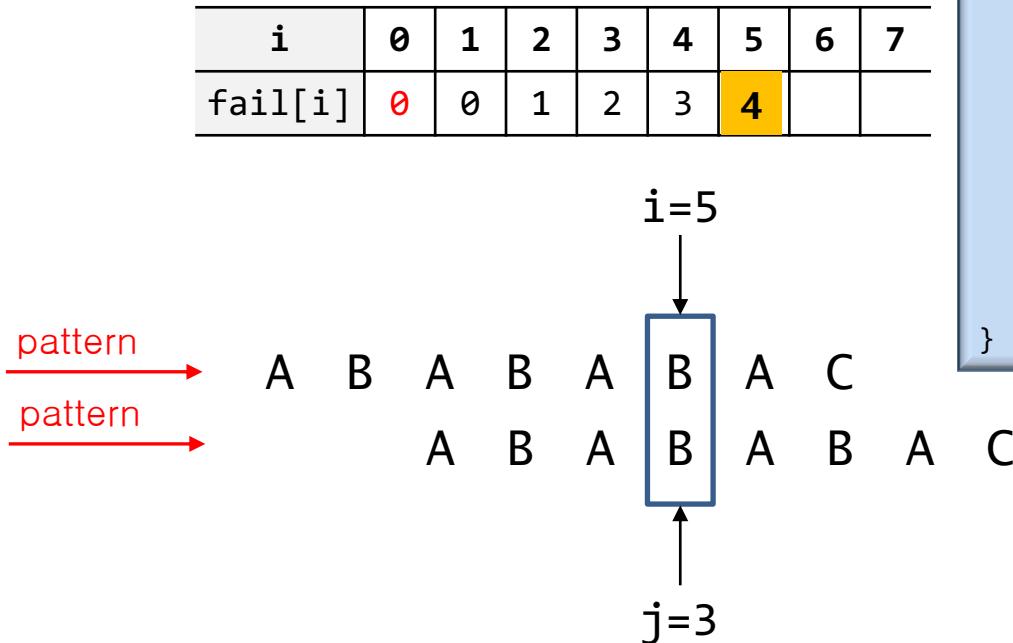
→ $\text{fail}[i] = ++j$
 $i++$

```
vector<int> getFail(string pattern)
{
    int m = (int) pattern.size();
    int j=0;
    vector<int> fail(m, 0);

    for(int i = 1; i< m ; i++)
    {
        while(j > 0 &&
              pattern[i] != pattern[j])
            j = fail[j-1];
        if(pattern[i] == pattern[j])
            fail[i] = ++j;
        else
            fail[i] = j;
    }
    return fail;
}
```

computation of fail[] (= $\pi[]$)

- **getFail()** function
 - Example:



$pattern[i] == pattern[j]$

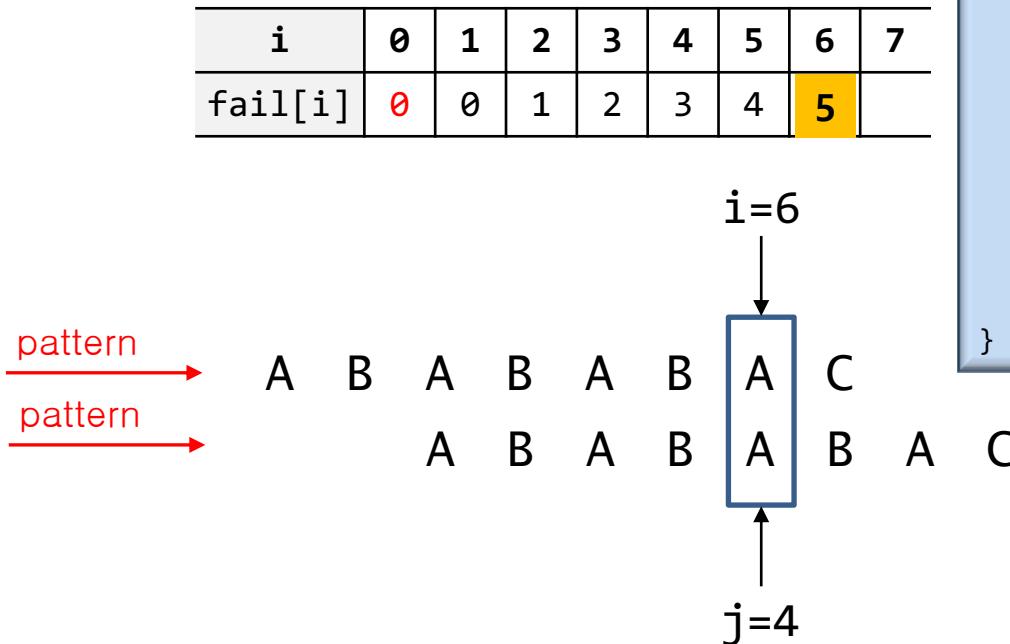
→ $fail[i] = ++j$
 $i++$

```
vector<int> getFail(string pattern)
{
    int m = (int) pattern.size();
    int j=0;
    vector<int> fail(m, 0);

    for(int i = 1; i< m ; i++)
    {
        while(j > 0 &&
              pattern[i] != pattern[j])
            j = fail[j-1];
        if(pattern[i] == pattern[j])
            fail[i] = ++j;
        else
            fail[i] = j;
    }
    return fail;
}
```

computation of fail[] (= $\pi[]$)

- **getFail()** function
 - Example:



$pattern[i] == pattern[j]$

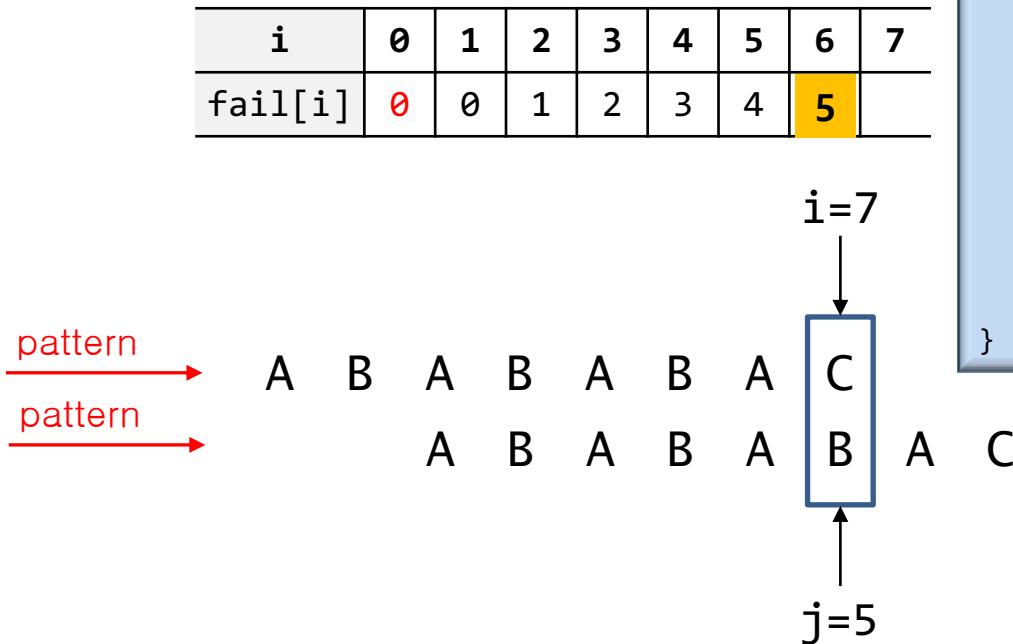
→ $fail[i] = ++j$
 $i++$

```
vector<int> getFail(string pattern)
{
    int m = (int) pattern.size();
    int j=0;
    vector<int> fail(m, 0);

    for(int i = 1; i< m ; i++)
    {
        while(j > 0 &&
              pattern[i] != pattern[j])
            j = fail[j-1];
        if(pattern[i] == pattern[j])
            fail[i] = ++j;
        else fail[i] = j;
    }
    return fail;
}
```

computation of fail[] (= $\pi[]$)

- **getFail()** function
 - Example:



$\text{pattern}[i] \neq \text{pattern}[j]$

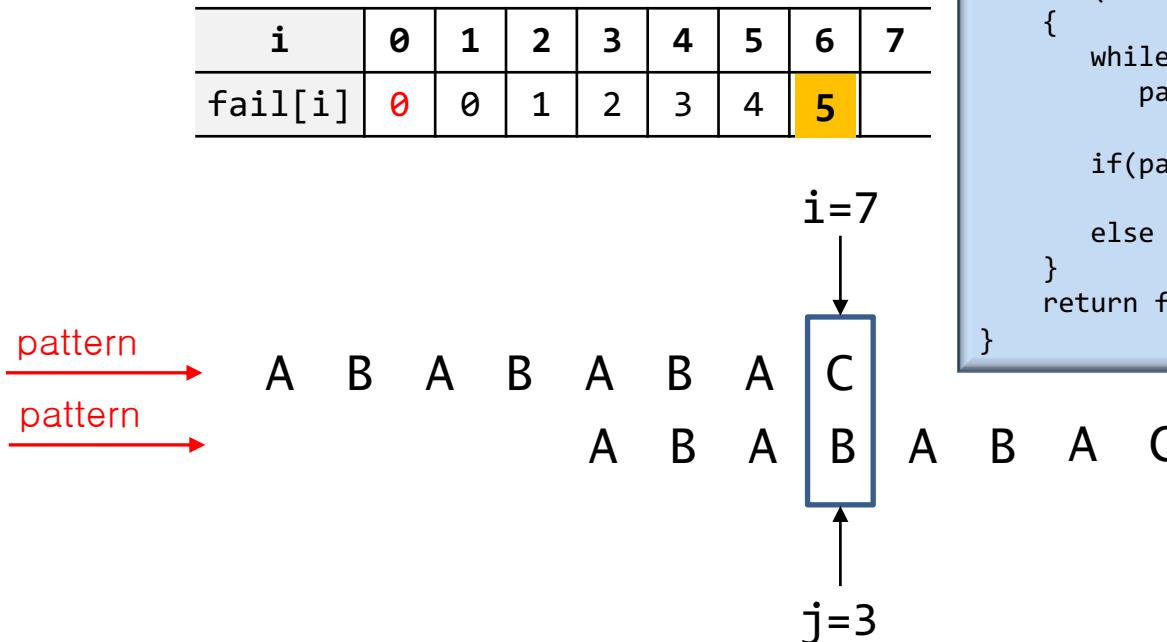
→ $j = \text{fail}[j-1]$ (=3)

```
vector<int> getFail(string pattern)
{
    int m = (int) pattern.size();
    int j=0;
    vector<int> fail(m, 0);

    for(int i = 1; i< m ; i++)
    {
        while(j > 0 &&
              pattern[i] != pattern[j])
            j = fail[j-1];
        if(pattern[i] == pattern[j])
            fail[i] = ++j;
        else fail[i] = j;
    }
    return fail;
}
```

computation of fail[] (= $\pi[]$)

- **getFail()** function
 - Example:



$\text{pattern}[i] \neq \text{pattern}[j]$

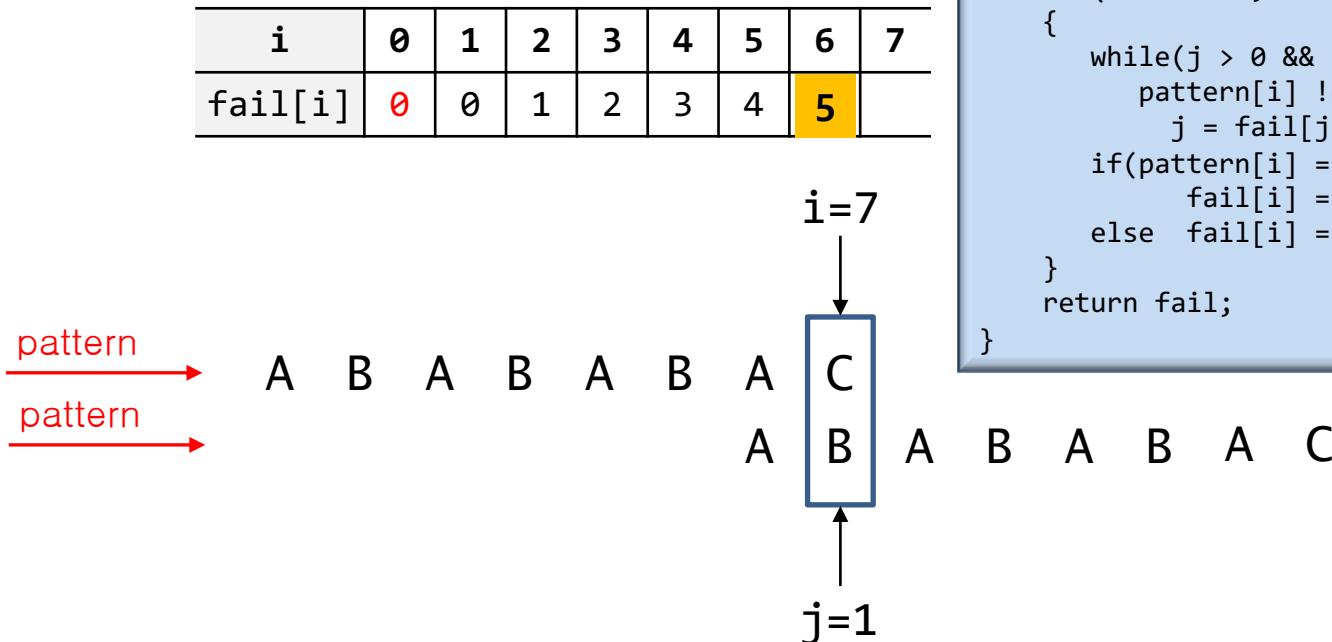
→ $j = \text{fail}[j-1]$ (=1)

```
vector<int> getFail(string pattern)
{
    int m = (int) pattern.size();
    int j=0;
    vector<int> fail(m, 0);

    for(int i = 1; i< m ; i++)
    {
        while(j > 0 &&
              pattern[i] != pattern[j])
            j = fail[j-1];
        if(pattern[i] == pattern[j])
            fail[i] = ++j;
        else fail[i] = j;
    }
    return fail;
}
```

computation of fail[] (= $\pi[]$)

- **getFail()** function
 - Example:



$\text{pattern}[i] \neq \text{pattern}[j]$

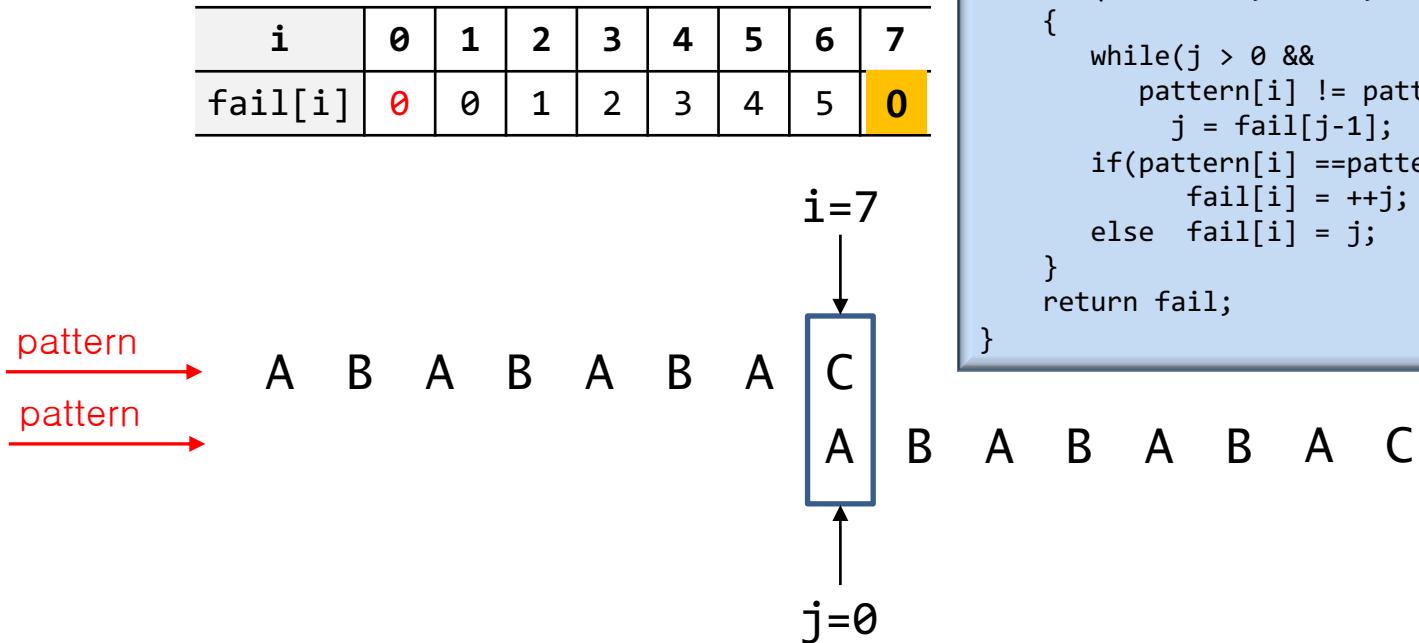
→ $j = \text{fail}[j-1]$ (=1)

```
vector<int> getFail(string pattern)
{
    int m = (int) pattern.size();
    int j=0;
    vector<int> fail(m, 0);

    for(int i = 1; i< m ; i++)
    {
        while(j > 0 &&
              pattern[i] != pattern[j])
            j = fail[j-1];
        if(pattern[i] == pattern[j])
            fail[i] = ++j;
        else fail[i] = j;
    }
    return fail;
}
```

computation of fail[] (= $\pi[]$)

- **getFail()** function
 - Example:



```
vector<int> getFail(string pattern)
{
    int m = (int) pattern.size();
    int j=0;
    vector<int> fail(m, 0);

    for(int i = 1; i< m ; i++)
    {
        while(j > 0 &&
              pattern[i] != pattern[j])
            j = fail[j-1];
        if(pattern[i] == pattern[j])
            fail[i] = ++j;
        else
            fail[i] = j;
    }
    return fail;
}
```

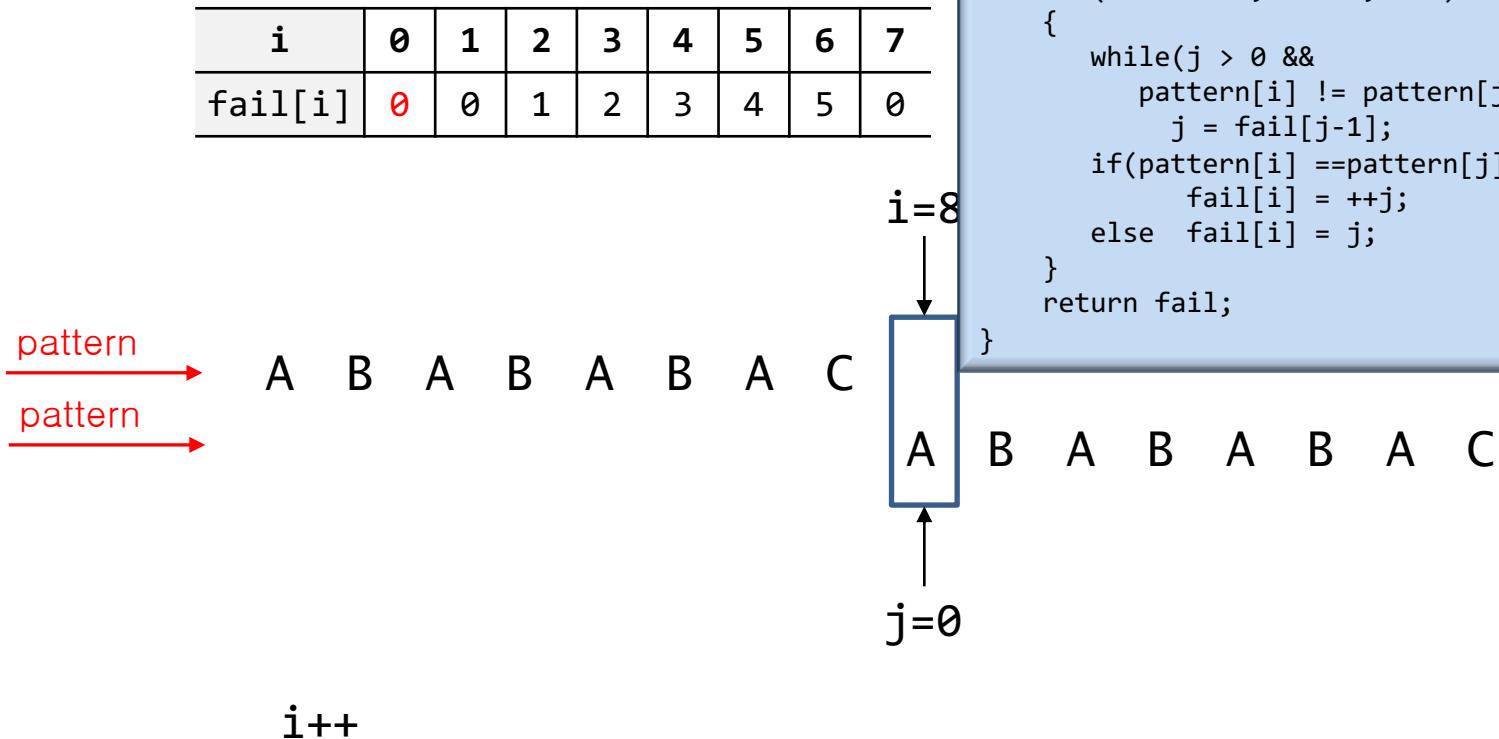
$pattern[i] \neq pattern[j]$

→ $j = fail[j-1] (=0)$
 $fail[i] = j$

45

computation of fail[] (= $\pi[]$)

- **getFail()** function
 - Example:



- `getFail()` function
 - Example: `aabaabac`

<code>i</code>	0	1	2	3	4	5	6	7
<code>fail[i]</code>	0	1	0	1	2	3	4	0

```
vector<int> getFail(string pattern)
{
    int m = (int) pattern.size();
    int j=0;
    vector<int> fail(m, 0);

    for(int i = 1; i< m ; i++)
    {
        while(j > 0 &&
              pattern[i] != pattern[j])
            j = fail[j-1];
        if(pattern[i] ==pattern[j])
            fail[i] = ++j;
        else fail[i] = j;
    }
    return fail;
}
```

Analysis of KMP Algorithm

- `getFail()` function
 - Why $O(M)$?
 - Index i increases from 1 to $M-1$
 - Index j *increases* maximally as many as i increases
 - Also index j *decreases* maximally as many as j increases
- KMP algorithm
 - Why $O(N)$ algorithm?
 - Similar logic to get the time complexity of `getFail()` function