THE UNIVERSITY of ADELAIDE

Mock Examination Do Not Use!

001956 / 102691	**Computer Systems - UG / PG** **COMP SCI 2000 / 7081**

Official Reading Time:	10 mins
Writing Time:	120 mins
Total Duration:	130 mins

Questions	**Time**	**Marks**
Answer all 12 questions	120 mins	120 marks
		120 Total

Instructions for Candidates

- This is a Closed-book examination.
- Begin each answer on a new page.
- Examination material must not be removed from the examination room.

Materials

- Foreign Language Dictionaries are Permitted for Translation Only

DO NOT COMMENCE WRITING UNTIL INSTRUCTED TO DO SO

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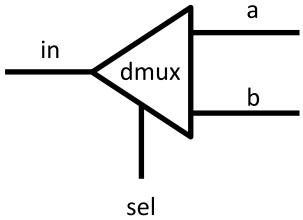
Basic Gates and Boolean Logic

Question 1

(a) **Published in Mock Exam**
The following diagram shows a one bit de-multiplexor (dmux) chip.

sel 0 1 / 11 / 10

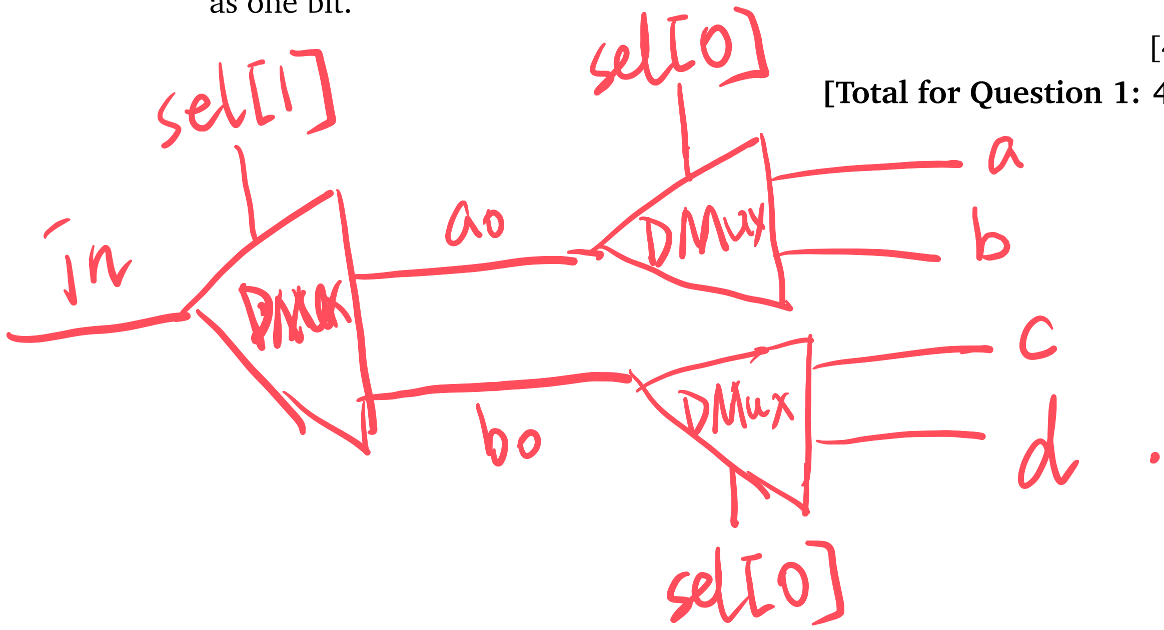


This chip directs the signal from `in` to either `a` or `b` depending on the value of `sel`. The non selected ouput is zero.

Now, given the 1-bit dmux above, draw an implementation for a dmux with four outputs and a two-bit selector. In your diagram assume that `in` remains as one bit.

[4 marks]

[Total for Question 1: 4 marks]



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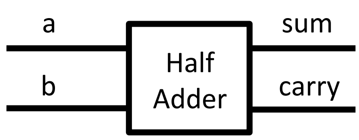
Boolean Arithmetic and ALU design

Question 2

For the following questions you may find the information in Figures 1 and 2 in the appendix of this paper useful.

(a) **Published in Mock Exam**

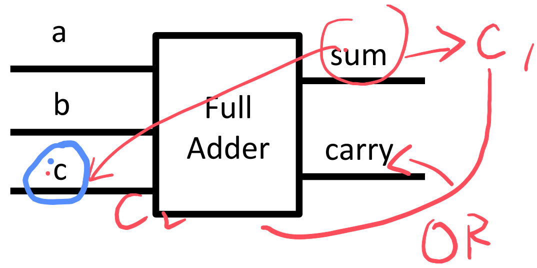
The following is a diagram the interface of a 1 bit half-adder:



a half-adder sums its two input bits to produce a sum bit and a carry bit.

Answer the following:

i. Draw an implementation of a full-adder chip composed from half-adder chips and/or other gates. Recall that the interface for a full adder is:

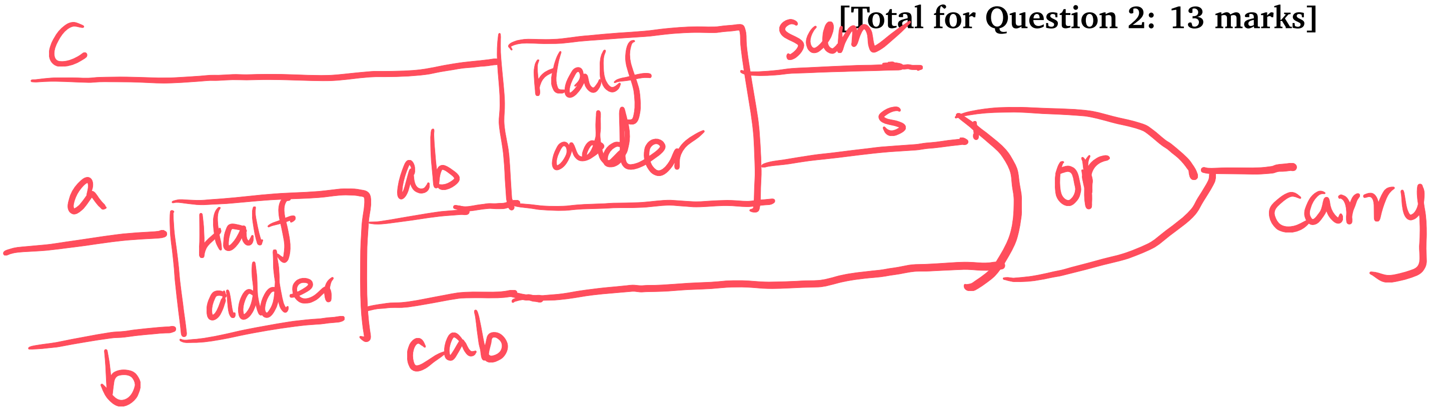


[7 marks]

ii. Write the code in the PARTS section of a HDL file describing the full-adder you defined in your answer to part (i) above. In your code you must assume that the inputs to the full adder are as labelled in the diagram above.

[6 marks]

[Total for Question 2: 13 marks]



HDL:
Halfadder(a=a,b=b,sum=ab,carry=cab)
Halfadder(a=c,b=ab, sum = sum, carry = s)
or(a= s, b= cab, out = carry)

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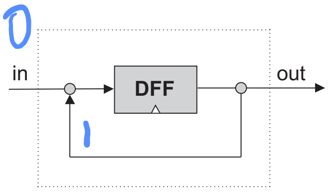
Sequential Logic

Question 3

(a) **Published in Mock Exam**
Look at the following diagram for an invalid design for a 1-bit register from figure 3.1 of the textbook.



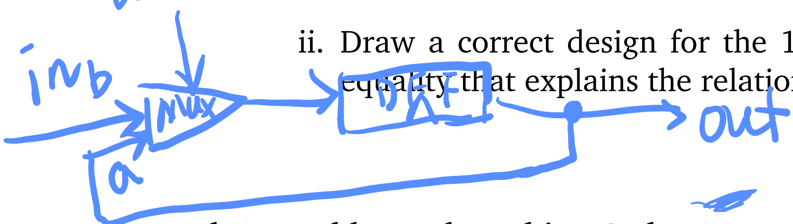
sel=0 ⇒ a
sel=1 ⇒ b

Because there is no reliable way to choose from the value in the input and the value in the output if they are different.

Answer the following.

i. Briefly explain what is wrong with the design of the register above. [2 marks]

ii. Draw a correct design for the 1-bit register above and write down the equality that explains the relationship between the `in` and `out` wires. [4 marks]



[Total for Question 3: 6 marks]

Hack Assembler and Machine Code

ROM → intro

Question 4

For the following questions you may find the information in Figures 3, 4, 5, 6, 7 and 8 in the appendix of this paper useful.

(a) **Published in Mock Exam**
Look at the following Hack machine code:

```
1 0000000000010000
2 1111110010001000
3 1111110000010000
4 0000000000000000
5 1110001100000001
6 0000000000000101
7 1110101010000111
```

Kaifeng, Peng
Clare, Heruin
Oc, Neo
xintong, Shengbin
Studenthelpand, Wai

1. @16
2. M=M-1
3. D=M
4. @0
5. D;JGT
6. @5
7. 0;JMP

A 0

Answer the following:

i. Using the instruction formats in Figures 3, 4, 5, 6, and 7 as a guide, write down the Hack assember instructions that are equivalent to this code. [7 marks]

bin → Hack

ii. Describe what the machine code does. [3 marks]

[Total for Question 4: 10 marks]

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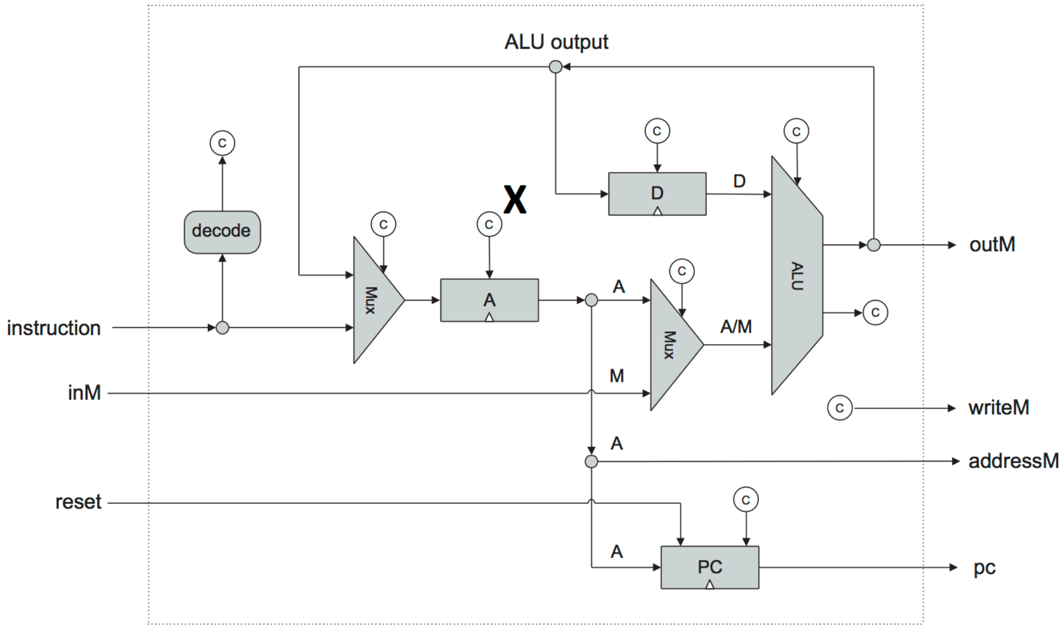
Computer Architecture

Question 5

For the following questions you may find the information in Figures 1, 2, 3, 4, 5, 6, 7 and 8 in the appendix of this paper useful.

(a) **Published in Mock Exam**

Look at the following partial diagram of a Hack CPU taken from Figure 5.9 of the textbook:



Some of the control logic is missing from this diagram. These missing gates and wires are marked with a © symbol. In the diagram one such section of missing control logic is marked with a large **X**. Given what you know about Hack instruction formats and ALU design, describe in detail what this missing control logic is.

Hint: feel free to use the figures in the appendix for some of the information you need.

[6 marks]

[Total for Question 5: 6 marks]

1. when A instructions --> !instr[15]
2. when C instruction has A register as the destination --> instr[5]

==>
!(instr[15]) | instr[5]

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ADM
010

① PASS: END, 6
② PASS: X, 16

RA
MEM
[16]

Assembler

Question 6

(a) **Published in Mock Exam**
Look at the following Hack assembler code:

```
    @X      0
    D=M     1
    @END    2
    D;JGE   3
    @X      4
    M=-M    5
(END)
    @END    6
    0;JMP   7
```

1. @16 --> 0000 0000 0001 0000
2. 111 1 110000 010 000
3. 0000000000000110 (@6)
4. 111 0 001100 000 011
5. 0000 0000 0001 0000
6. 111 1 110011 001 000
7. 0000000000000110
8. 111 0 101010 000 111

Hand-assemble this code by writing out the binary machine code the assembler would produce. For this question you may find the information in Figures 3, 4, 5, 6, and 7 useful.

[9 marks]

[Total for Question 6: 9 marks]

Virtual Machine - Expressions

Question 7

(a) **Published in Mock Exam**
Translate the following Jack let statement into Hack Virtual Machine language:

loc5 loc7

`let d = ((2 - x) * y) + 5`

loc2

The variables d, x and y are in memory segment *local* at indexes 2,5 and 7 respectively. Assume there is a function named *multiply* that will take two arguments and return the result of multiplying the two numbers together.

[8 marks]

[Total for Question 7: 8 marks]

push constant 2
push local 5
sub
push local 7
call Math.multiply 2
push constant 5
add
pop local 2

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Virtual Machine - Subroutines

Question 8

(a) **Published in Mock Exam**
The Hack Virtual Machine language provides three function related commands:

- call f m
- function f n
- return

i. Briefly describe what the function command does during program execution.

[2 marks]

ii. Briefly describe what the call command does during program execution.
[7 marks]

iii. Briefly describe what the return command does during program execution.

[8 marks]

arg 0
arg 1
...
arg N
return address
lcl
arg
this
that
local 0
local 1
.
.
.
.
.
local M-1

(b) **Published in Mock Exam**
The Hack Virtual Machine allocates an area of the stack for each active function call. Briefly describe the structure of one of these stack frames immediately after the execution of the function command in a **Jack** method that is declared with N parameters and M local variables.

[9 marks]

[Total for Question 8: 26 marks]

Jack

Question 9

(a) **Published in Mock Exam**
Write a **Jack** program that calls a recursive function to calculate the 7th fibonacci number. The result must be placed in an int variable x.

[8 marks]

[Total for Question 9: 8 marks]

Main.main

```
Class Main
{
  function void main()
   {
       var int x;
       let x = Foo.Fibonacci(7);
       return;
   }
}
```

```
Class Foo
{
   function int Fibonacci(int i)
   {
      if ( I=0)
      {
         return 0
      }
      if (I = 1)
      {
        return 1;
      }
       return Fibonacci(I-1) + Fibonacci(I-2);
   }
```

PLEASE SEE NEXT PAGE

```
}
```

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Parsing

Question 10

(a) **Published in Mock Exam**
Show the two symbol tables for the following code just ater the last variable declaration in the method has been parsed.

```
class BankAccount
{
        // Class variables
        static string key ;
        static int nAccounts ;

        // Instance variables ;
        field string owner ;
        field int balance ;

        method void transfer(int sum, BankAccount b2)
        {
                var Date due ;
                var int i,j ;

                let i = sum ;
        }
}
```

class table:
name, type, kind, segment, index
key, string,static, static, 0
nAccounts, int, static,static, 1
owner, string, field,this, 0
balance, int, field, this,1

method table:
this, BankAccount, argument, arg, 0
sum, int, argument, arg, 1
b2, BankAccount, argument, arg, 2

due, Date, local, local, 0
I, int, local, local, 1
j, int, local, loca, 2

[10 marks]

[Total for Question 10: 10 marks]

Code Generation

Question 11

(a) **Published in Mock Exam**
Consider the following **Jack** method:

```
method int useless(String x, String y)
{
        var Array local1 ;
        var int local0 ;
        var string local3 ;
}
```

set
point0 → this → obj
point1 → that → array

method table:

x, string, arg, 1
y, string, arg 2

local1, Array, local 0
local0, int, local 1
local3, string, local 2

What Hack Virtual Machine language code would implement the following Jack program fragments if they were in the body of the method useless?

i. `let local1[7] = x ;` → MEM[6000] [6 marks]

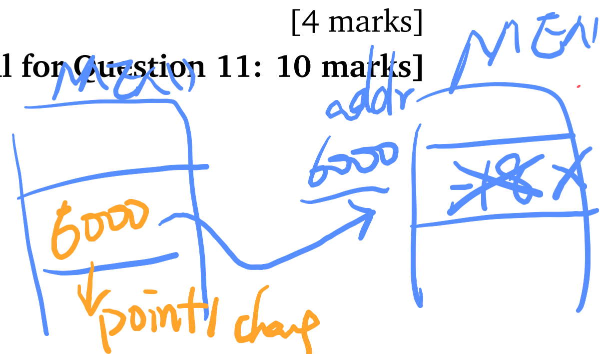
ii. `return local0 + 1 ;` [4 marks]

[Total for Question 11: 10 marks]

I.
push local 0
push constant 7
add 6000
pop pointer 1
push arg 1 → x
pop that 0

MEM[4] that

ii.
push local 1
push constant 1
add
return



pointer | chang value in "that"

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"that 0" chang the value in the location pointed by "that"

Jack OS, Optimisation

Question 12

(a) **Published in Mock Exam**
How do caches take advantage of temporal and spatial locality to improve the performance of a computer?

arr[0]
arr[1]

[4 marks]

(b) **Published in Mock Exam**
What determines the minimum length of a clock cycle in a processor?

[2 marks]

(c) **Published in Mock Exam**
The Jack Operating System provides a small number of libraries that extend the functionality of the Jack programming language. Excluding support for graphical user interfaces, identify two operating system services that are not provided by the Jack OS but are provided by Linux. In each case explain why the service is important.

[4 marks]

[Total for Question 12: 10 marks]

Security

Network

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APPENDICES

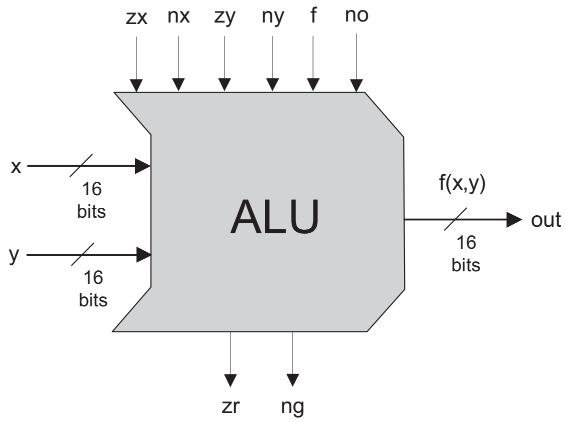

Figure 1: An interface diagram for the ALU. From figure 2.5 of the textbook.

These bits instruct how to preset the `x` input		These bits instruct how to preset the `y` input		This bit selects between + / And	This bit inst. how to postset `out`	Resulting ALU output
zx	nx	zy	ny	f	no	out=
if zx then x=0	if nx then x=!x	if zy then y=0	if ny then y=!y	if f then out=x+y else out=x&y	if no then out=!out	f(x,y)=
1	0	1	0	1	0	0
1	1	1	1	1	1	1
1	1	1	0	1	0	-1
0	0	1	1	0	0	x
1	1	0	0	0	0	y
0	0	1	1	0	1	!x
1	1	0	0	0	1	!y
0	0	1	1	1	1	-x
1	1	0	0	1	1	-y
0	1	1	1	1	1	x+1
1	1	0	1	1	1	y+1
0	0	1	1	1	0	x-1
1	1	0	0	1	0	y-1
0	0	0	0	1	0	x+y
0	1	0	0	1	1	x-y
0	0	0	1	1	1	y-x
0	0	0	0	0	0	x&y
0	1	0	1	0	1	x|y

Figure 2: The Hack ALU truth table. From figure 2.6 of the textbook.

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A-instruction: *@value* // Where *value* is either a non-negative decimal number
// or a symbol referring to such number.

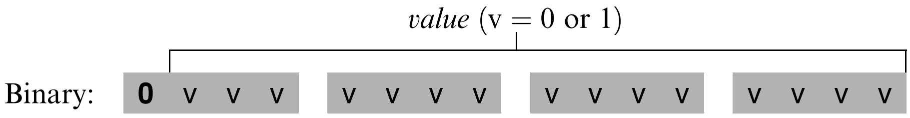

Figure 3: The format of an A-instruction. From page 64 of the text book.

C-instruction: *dest=comp;jump* // Either the *dest* or *jump* fields may be empty.
// If *dest* is empty, the "=" is omitted;
// If *jump* is empty, the ";" is omitted.

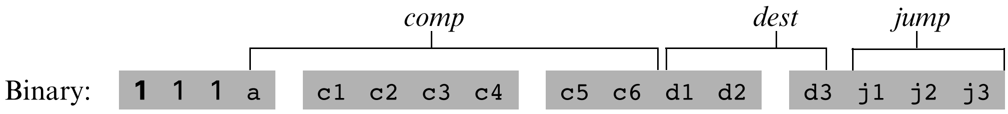

Figure 4: The format of an C-instruction. From page 66 of the text book.

(when a=0) ***comp*** *mnemonic*	c1	c2	c3	c4	c5	c6	(when a=1) ***comp*** *mnemonic*
0	1	0	1	0	1	0	
1	1	1	1	1	1	1	
-1	1	1	1	0	1	0	
D	0	0	1	1	0	0	
A	1	1	0	0	0	0	M
!D	0	0	1	1	0	1	
!A	1	1	0	0	0	1	!M
-D	0	0	1	1	1	1	
-A	1	1	0	0	1	1	-M
D+1	0	1	1	1	1	1	
A+1	1	1	0	1	1	1	M+1
D-1	0	0	1	1	1	0	
A-1	1	1	0	0	1	0	M-1
D+A	0	0	0	0	1	0	D+M
D-A	0	1	0	0	1	1	D-M
A-D	0	0	0	1	1	1	M-D
D&A	0	0	0	0	0	0	D&M
D|A	0	1	0	1	0	1	D|M

Figure 5: The meaning of C-instruction Fields. From figure 4.3 of the textbook.

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d1	d2	d3	*Mnemonic*	*Destination (where to store the computed value)*
0	0	0	null	The value is not stored anywhere
0	0	1	M	Memory[A] (memory register addressed by A)
0	1	0	D	D register
0	1	1	MD	Memory[A] and D register
1	0	0	A	A register
1	0	1	AM	A register and Memory[A]
1	1	0	AD	A register and D register
1	1	1	AMD	A register, Memory[A], and D register

Figure 6: The meaning of the destination bits of the C-instruction From figure 4.4 of the textbook.

j1 ($out < 0$)	j2 ($out = 0$)	j3 ($out > 0$)	**Mnemonic**	**Effect**
0	0	0	null	No jump
0	0	1	JGT	If $out > 0$ jump
0	1	0	JEQ	If $out = 0$ jump
0	1	1	JGE	If $out \geq 0$ jump
1	0	0	JLT	If $out < 0$ jump
1	0	1	JNE	If $out \neq 0$ jump
1	1	0	JLE	If $out \leq 0$ jump
1	1	1	JMP	Jump

Figure 4.5 The *jump* field of the *C*-instruction. *Out* refers to the ALU output (resulting from the instruction's *comp* part), and *jump* implies "continue execution with the instruction addressed by the A register."

Figure 7: The meaning of the jump bits of the C-instruction From figure 4.5 of the textbook.

Label	*RAM address*
SP	0
LCL	1
ARG	2
THIS	3
THAT	4
R0-R15	0-15
SCREEN	16384
KBD	24576

Figure 8: The predefined symbols in Hack Assembly language. From page 110 of the text book.

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Lexical Elements

```
keyword         ::= ’class’ | ’constructor’ | ’function’ | ’method’ | \
                   ’field’ | ’static’ | ’var’ | ’int’ | ’char’ | \
                   ’boolean’ | ’void’ | ’true’ | ’false’ | ’null’ | \
                   ’this’ | ’let’ | ’do’ | ’if’ | ’else’ | ’while’ | \
                   ’return’
symbol          ::= ’{’ | ’}’ | ’(’ | ’)’ | ’[’ | ’]’ | ’.’ | \
                   ’,’ | ’;’ | ’+’ | ’-’ | ’*’ | ’/’ | ’&’ | \
                   ’|’ | ’<’ | ’>’ | ’=’ | ’~’ | ’
integerConstant ::= A decimal number in the range 0 .. 32767
stringConstant  ::= ’"’ A sequence of Unicode characters not including
                       double quote or newline ’"’
identifier      ::= A sequence of letters, digits and underscore (’_’)
                   not starting with a digit.
```

Statements

```
statements      ::= statement*
statement       ::= letStatement | ifStatement | whileStatement | \
                   doStatement | returnStatement}
letStatement    ::= ’let’ varName (’[’ expression ’]’)? ’=’ expression ’;’
ifStatement     ::= ’if’ ’(’ expression ’)’ ’{’ statements ’}’ \
                  (’else’ ’{’ statements ’}’)?
whileStatement  ::= ’while’ ’(’ expression ’)’ ’{’ statements ’}’
doStatement     ::= ’do’ subroutineCall ’;’
returnStatement ::= ’return’ expression? ’;’
```

Expressions

```
expression      ::= term (op term)*
term            ::= integerConstant | stringConstant | \
                   keywordConstant | varName | \
                   varName ’[’ expression ’]’ | subroutineCall | \
                   ’(’ expression ’)’ |  unaryOp term
subroutineCall  ::= subroutineName ’(’ expressionList ’)’ |  \
                   (className | varName) ’.’ subroutineName ’(’ expressionList ’)’
expressionList  ::= (expression (’,’ expression)*)?
op              ::= ’+’ | ’-’ | ’*’ | ’/’ | ’&’ | ’|’ | ’<’ | ’>’ | ’=’
unaryOp         ::= ’-’ | ’~’
keywordConstant ::= ’true’ | ’false’ | ’null’ | ’this’
varName         ::= identifier
```

Figure 9: The Jack grammar. From figure 10.5 of the textbook.

END OF EXAMINATION PAPER