

# 64-bit PowerPC Microprocessors

## Memory Management

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origin: [http://www.cs.cmu.edu/~412/lectures/2009-09-18\\_ppc64vm.pdf](http://www.cs.cmu.edu/~412/lectures/2009-09-18_ppc64vm.pdf)

# Abbreviations

EA	Effective Address	64-bit
VA	Virtual Address	80-bit
PA	Physical Address	62-bit
SDR1	Storage Description Register 1	
MSR	Machine State Register	

PR[49]: Privilege level

If set, run in user mode

IR[58] / DR [59]: Instruction address translation / Data address translation

If set, Instruction/Data address translation is enabled

Feature Category	64-Bit Imple
	Conventional
Address ranges	$2^{64}$ bytes of effective address
	$2^n$ where $65 \leq n \leq 80$ bytes of virtual address
	$\leq 2^m$ ( $m \leq 62$ ) bytes of physical address
Page size	4 Kbytes Some large page sizes ( $2^p$ where $13 \leq p \leq 28$ )
Segment size	256 Mbytes
Block address translation	Not applicable
	Not applicable
Memory protection	Segments selectable as no-execute
	Pages selectable as user/supervisor and read-only
	Blocks selectable as user/supervisor and read-only
Page history	Referenced and changed bits defined and maintained
Page address translation	Translations stored as PTEs in hashed page tables in memory
	Page table size determined by size programmed into SDR1 register
TLBs	Instructions for maintaining optional TLBs
Segment descriptors	Stored as STEs (explicit or implicit segment tables)
	Instructions for maintaining SLBs

# Memory Addressing Overview

- PowerPC processors support two types of address translation:

## **Real addressing mode**

$$EA = PA$$

## **Page address translation**

translates the EA to VA then to PA.

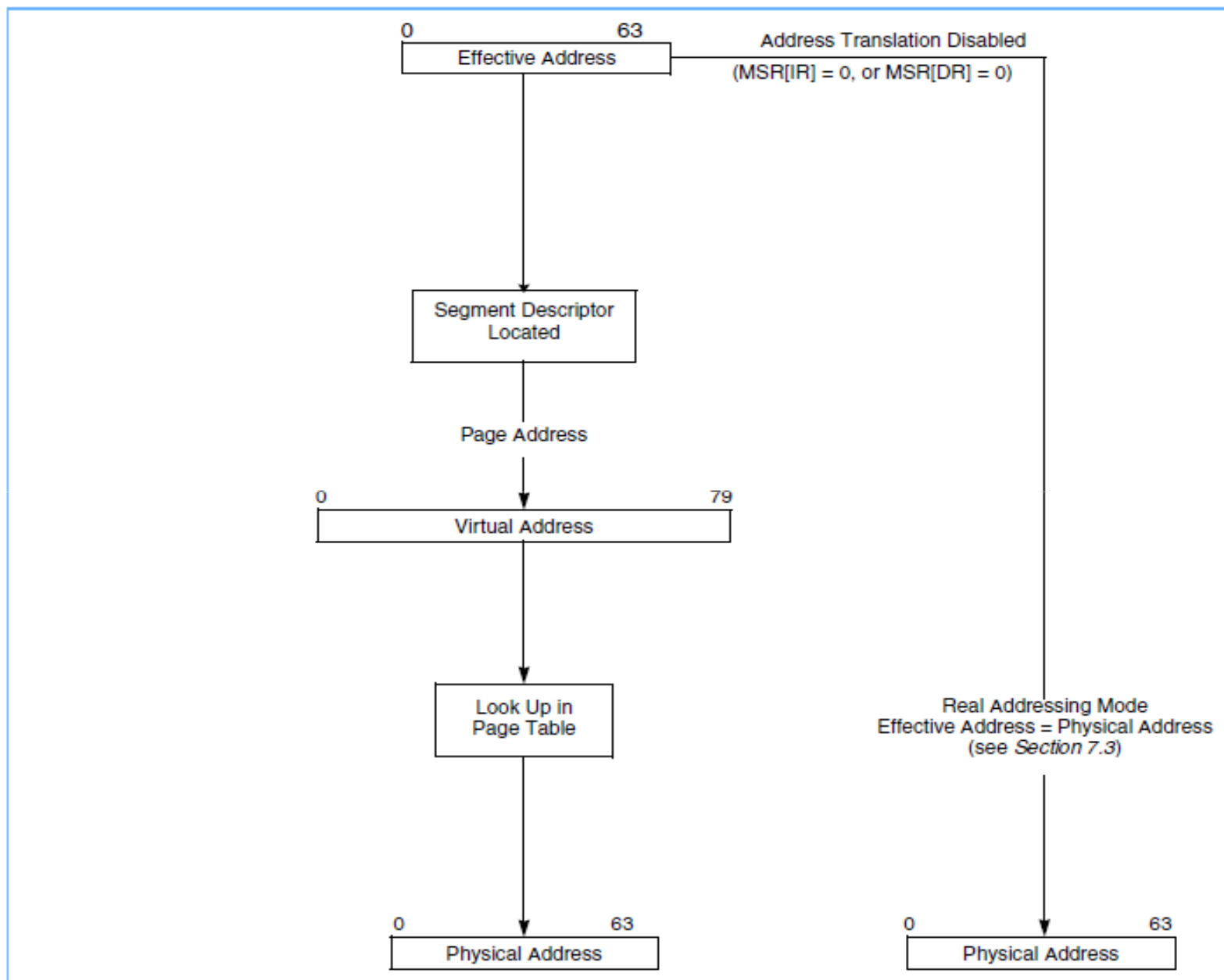
# Real Addressing Mode

MSR[IR] = 0 or MSR[DR] = 0

- EA is treated as PA.
- Bypass all memory protection checks.
- Processor would set MSR[IR] = 0 and MSR[DR] = 0 after exceptions occur. Must explicitly enable translation if needed.
- Attempts to access PA not physically present will cause a machine check exception

# Page address translation

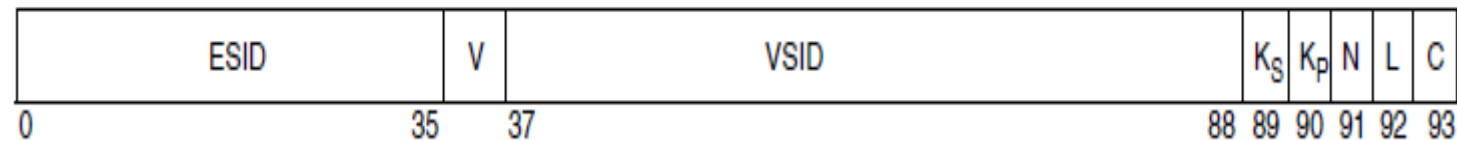
- The page translation proceeds in the following two steps:
- 1. From effective address to the virtual address, and
- 2. From virtual address to physical address.



# EA to VA

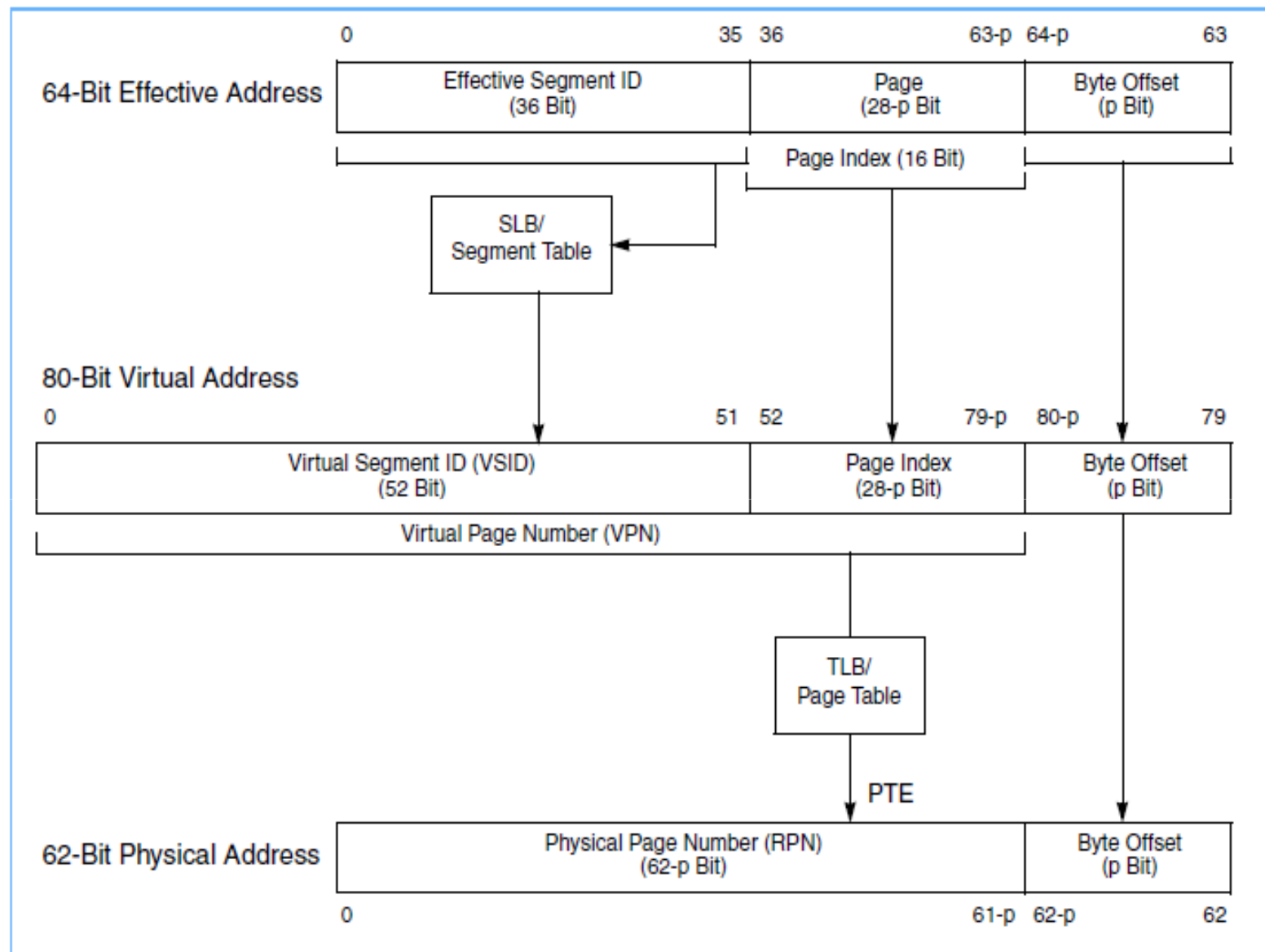
- Segment Lookaside Buffer (SLB)  
Each SLB entry maps one ESID (Effective Segment ID) to one VSID (Virtual Segment ID).

SLB entry format



- SLB Search  
Success if  
 $SLBE[V] = 1$  and  $SLBE[ESID] = EA[0-35]$



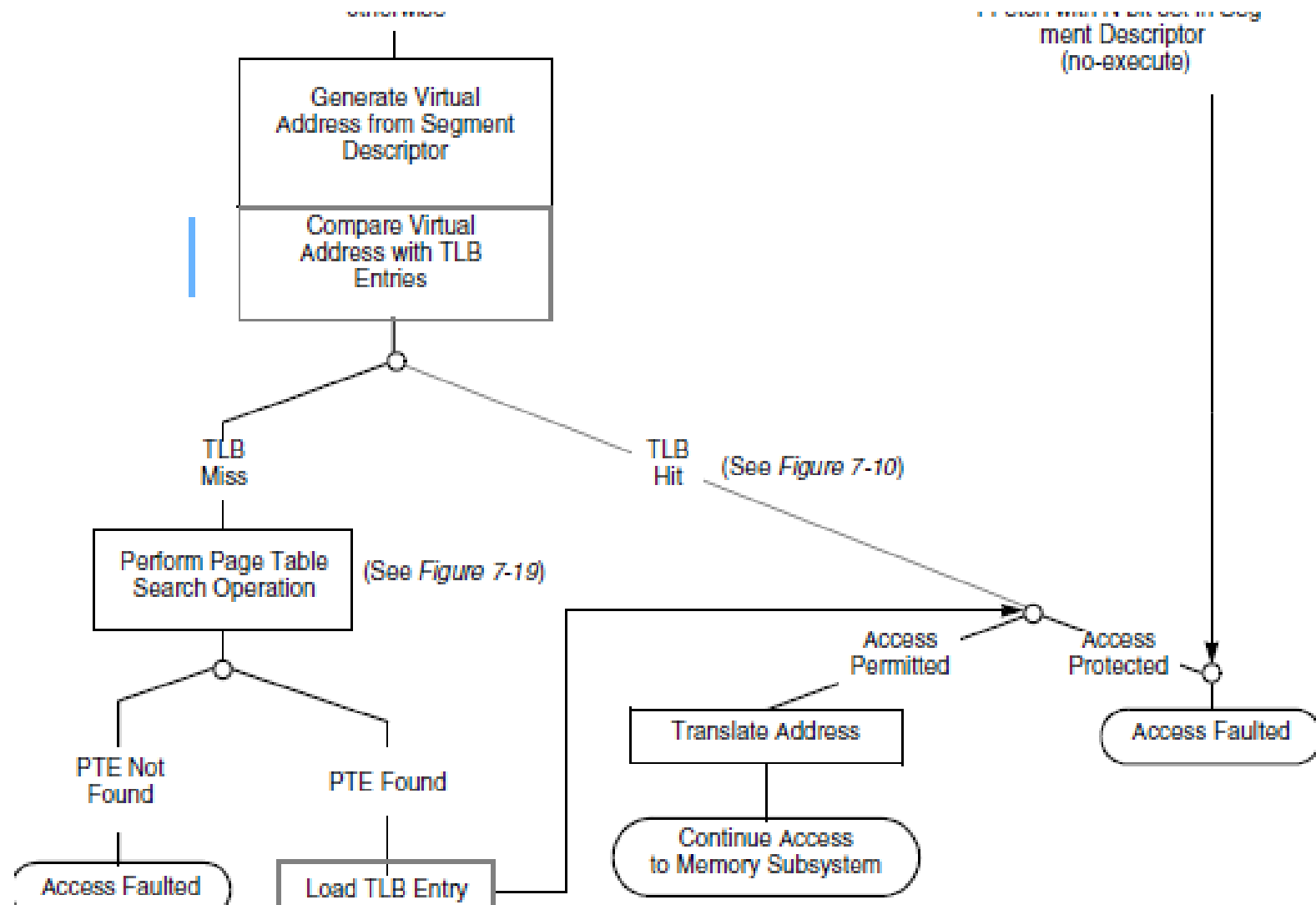


# VA to PA

## Translation Lookaside Buffer (TLB)

- Holds PTEs that have recently been used.
- Prior to searching the page table.
- Needs to keep consistency with the page table.
- Not required in 64-bit PowerPC architecture.

# VA to PA cont.



## VA to PA cont.

- Hashed Page Tables

If the PTE we search is not resident in a TLB.

- Page table entry groups (PTEG)

The page table contains a number of PTEGs.

A PTEG contains eight page table entries (PTE).

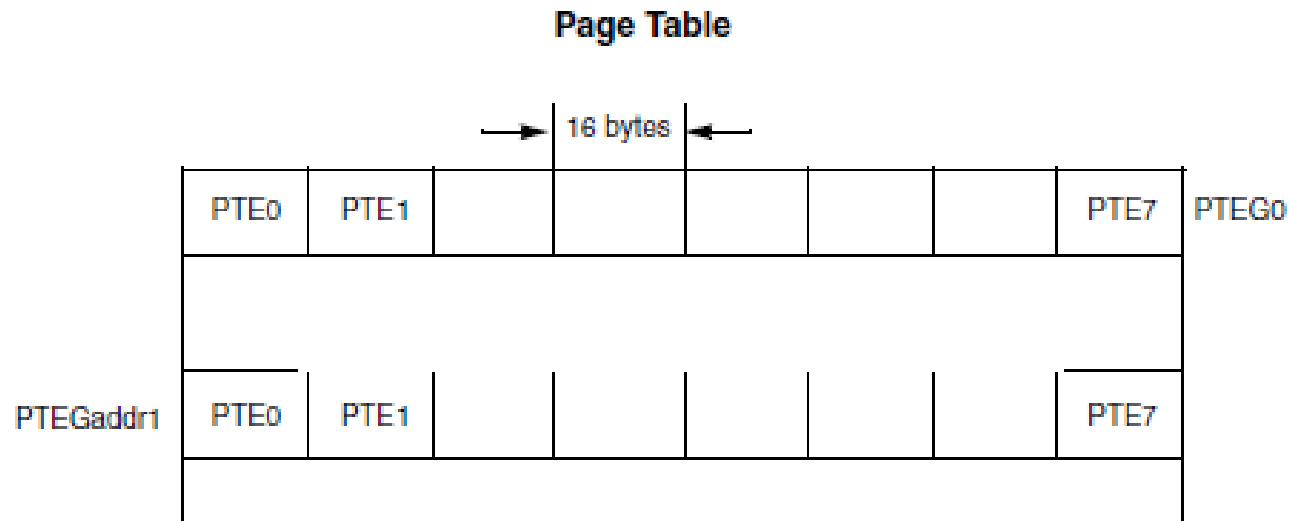
- Page table Entry (PTE)

A PTE can reside in two possible PTEGs, primary and secondary.

# VA to PA cont.

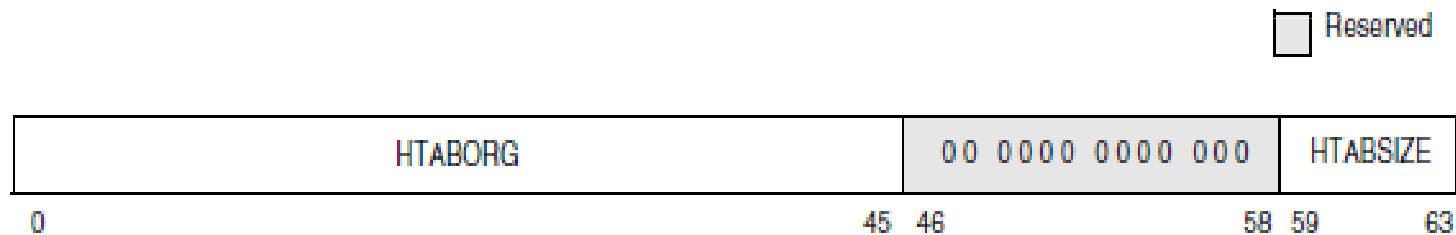
- Locate PTEG

The address of PTEG is derived from the **HTABORG** field of the SDR1 register and **output of hashing function of VA.**



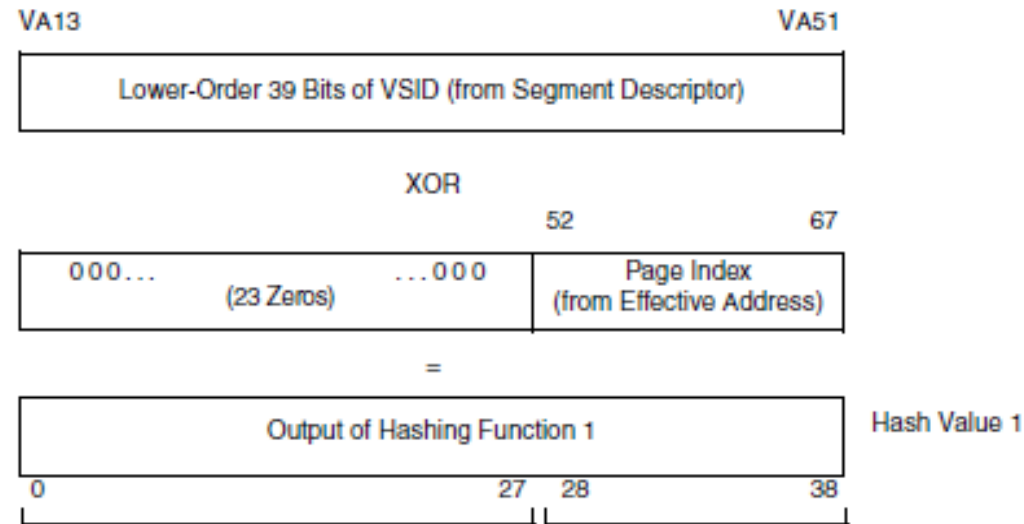
# VA to PA cont.

- Storage Description Register (SDR1)  
SDR1 defines the **physical base address** of page table and the size of the table.

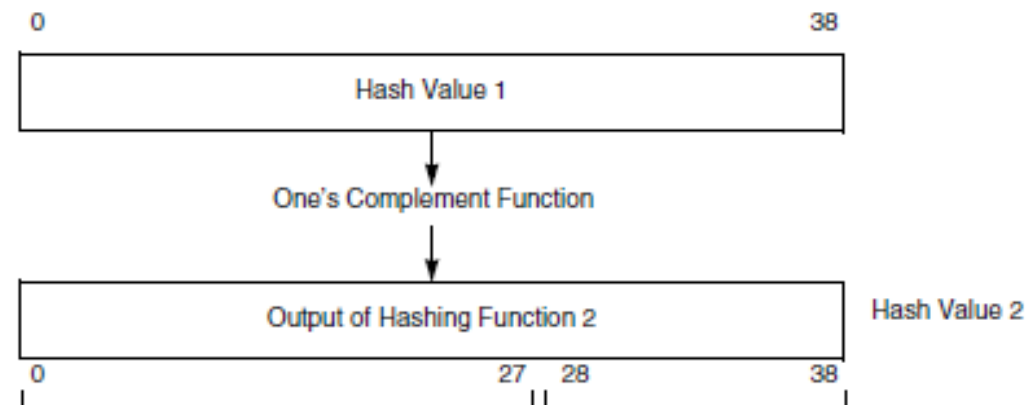


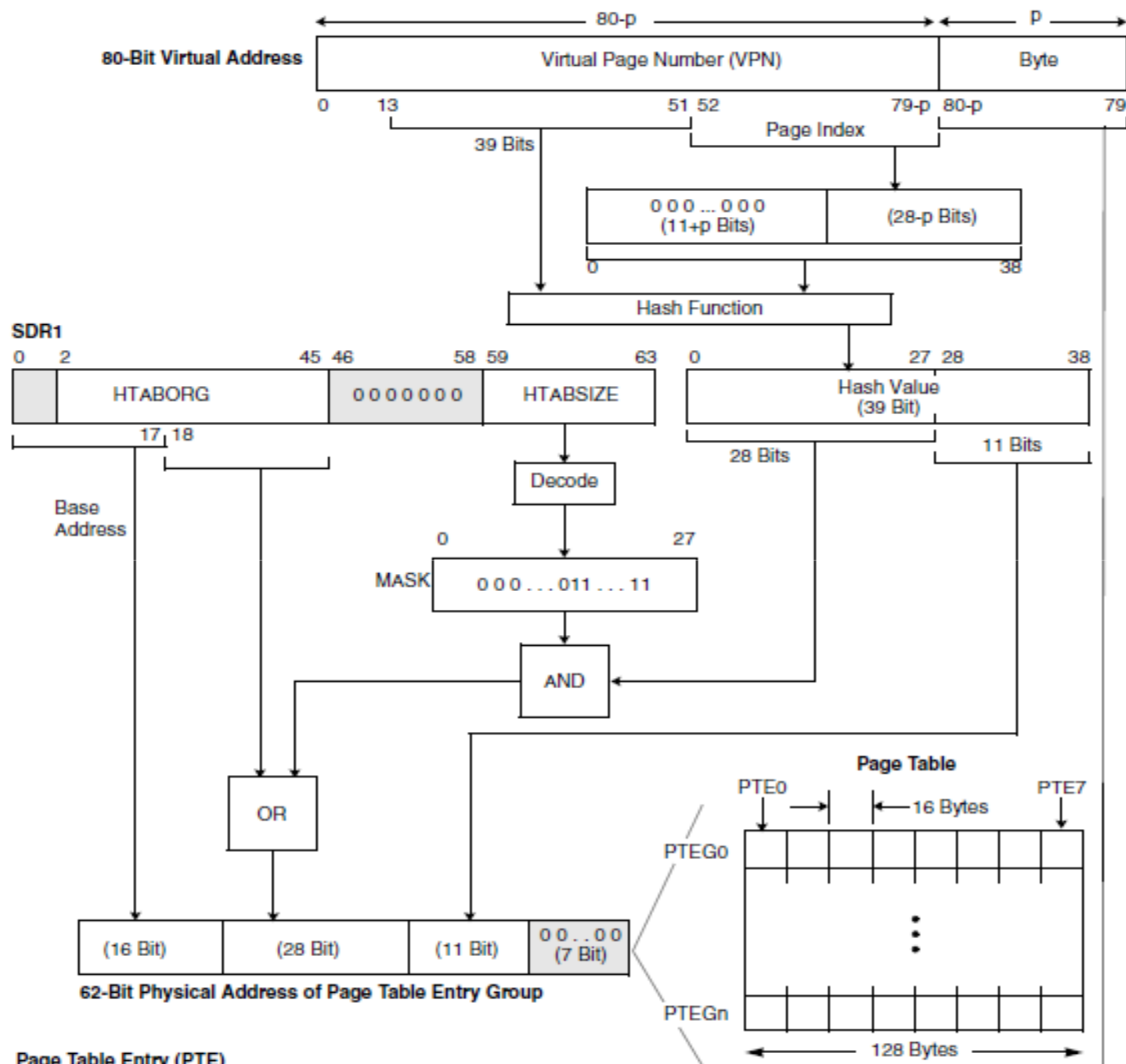
# VA to PA cont. (Hashing function)

Primary Hash:

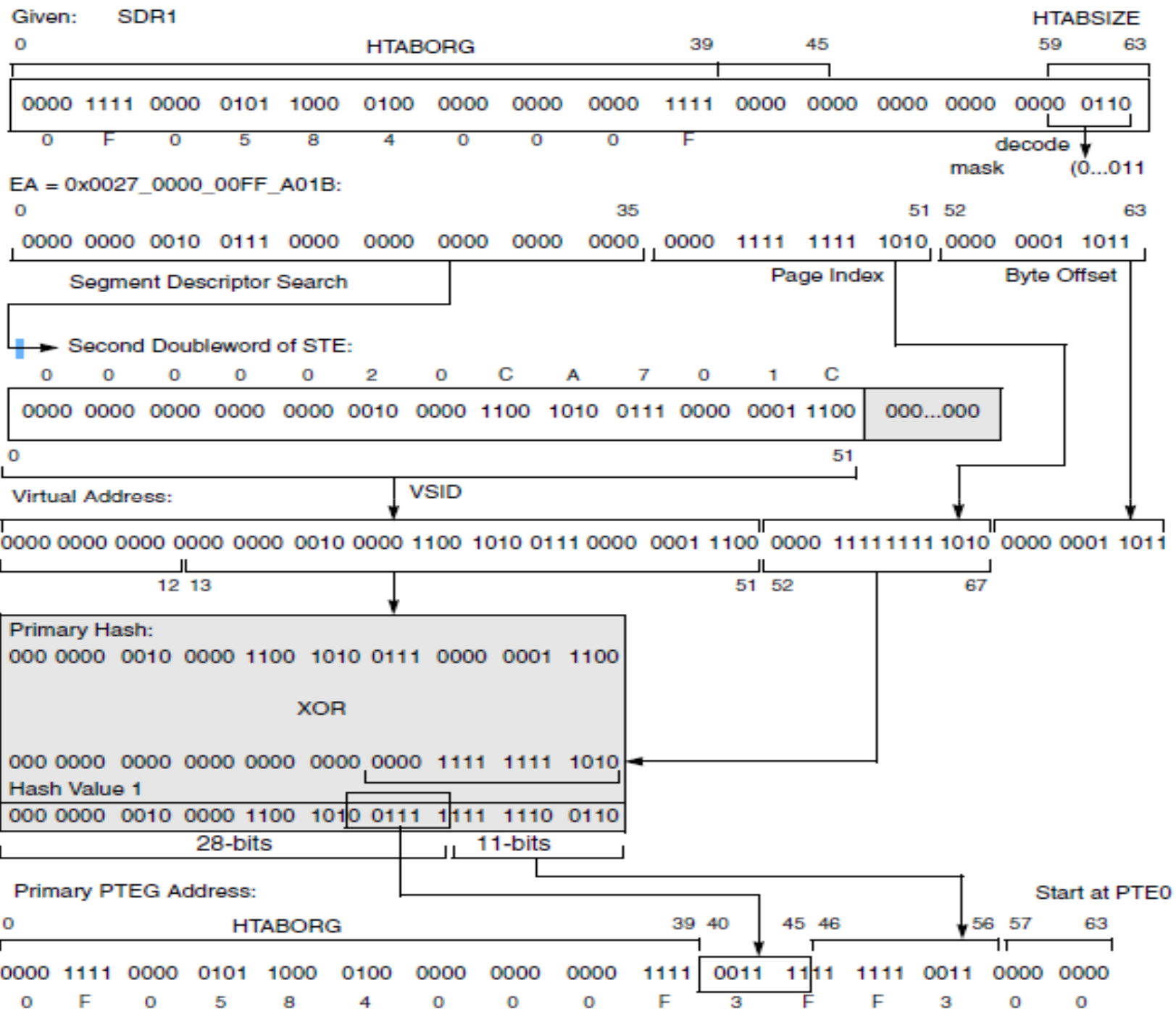


Secondary Hash:

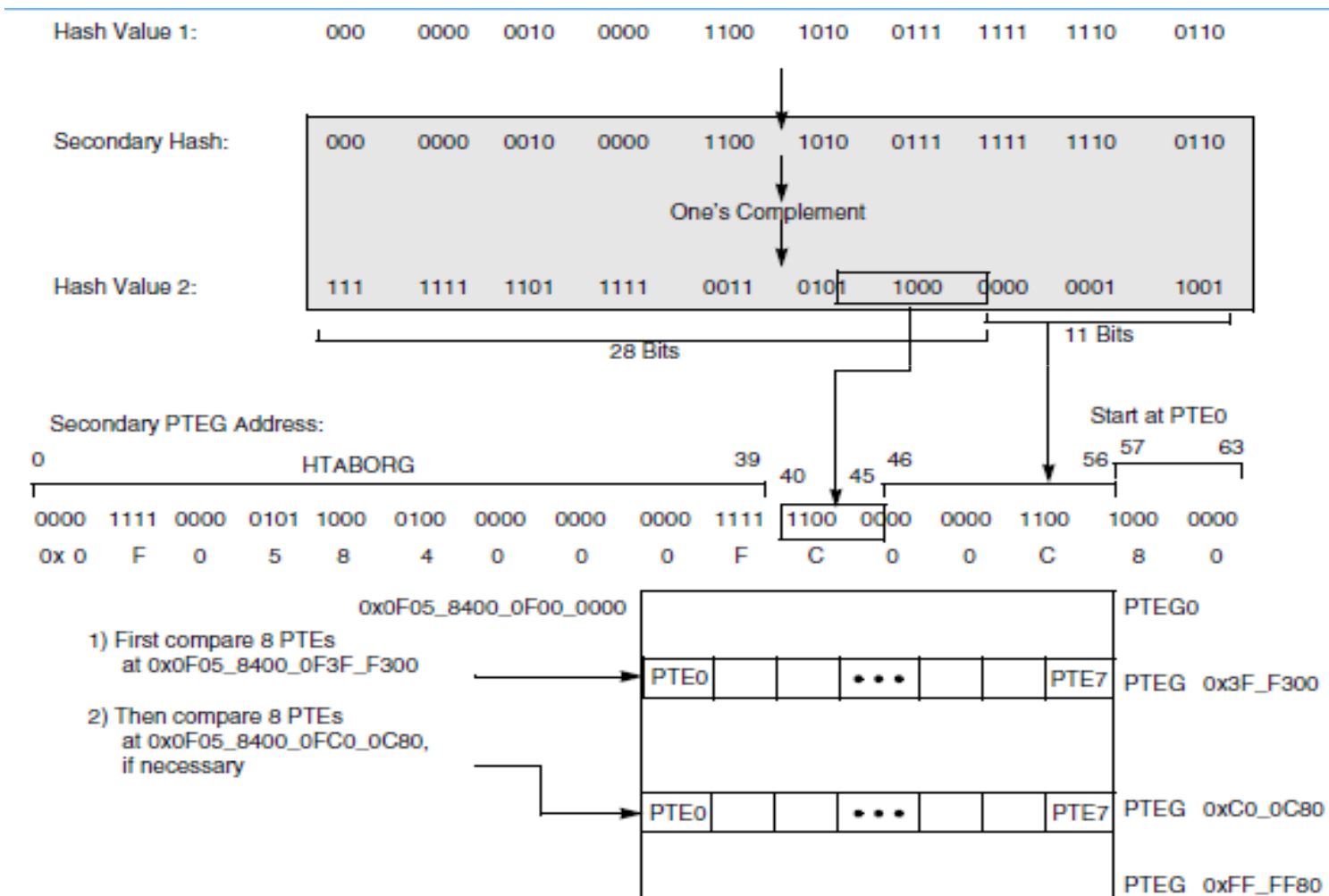






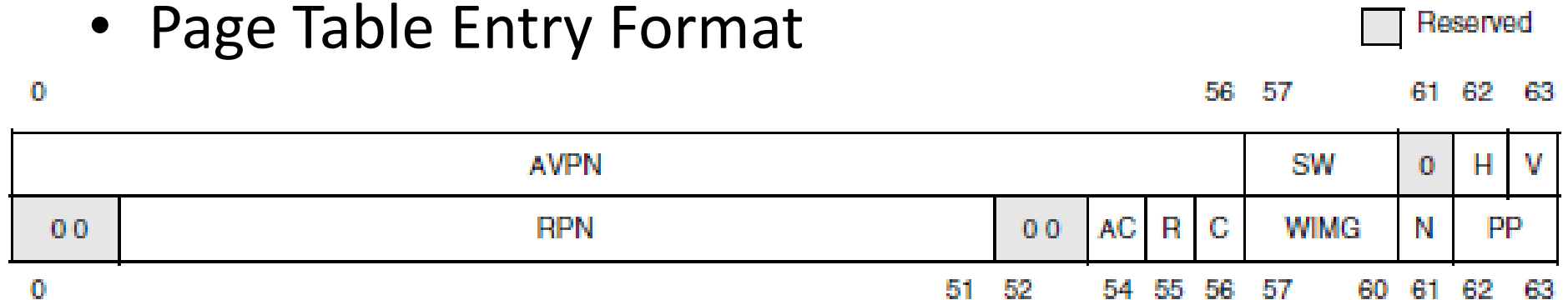


# VA to PA cont. (Secondary Hash fun.)



# VA to PA cont. (PTE)

- Page Table Entry Format



- Page Table Entry Search

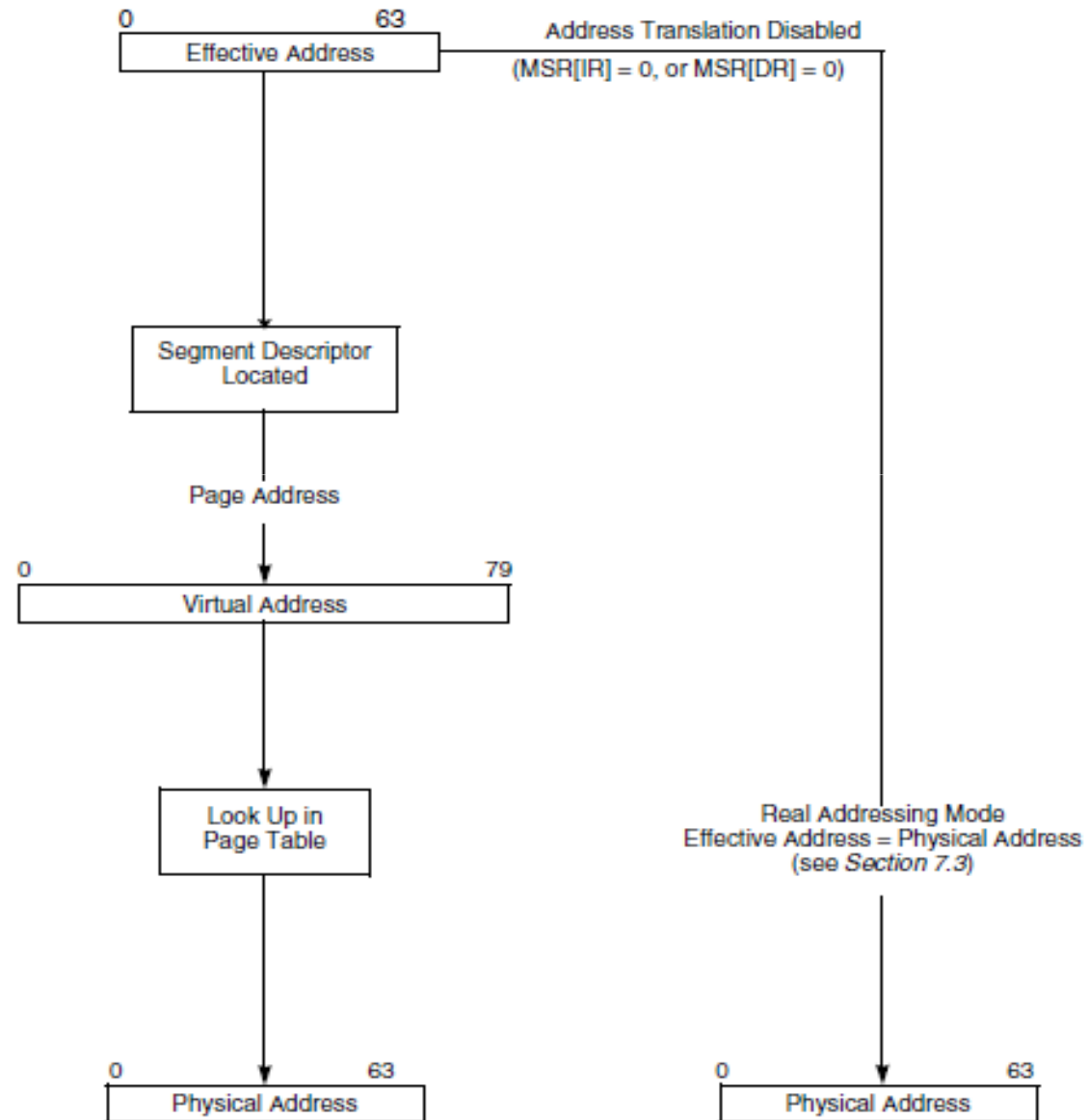
PTE[H] = 0 for primary PTEG; PTE[H] = 1 for secondary PTEG

$$\text{PTE}[V] = 1$$

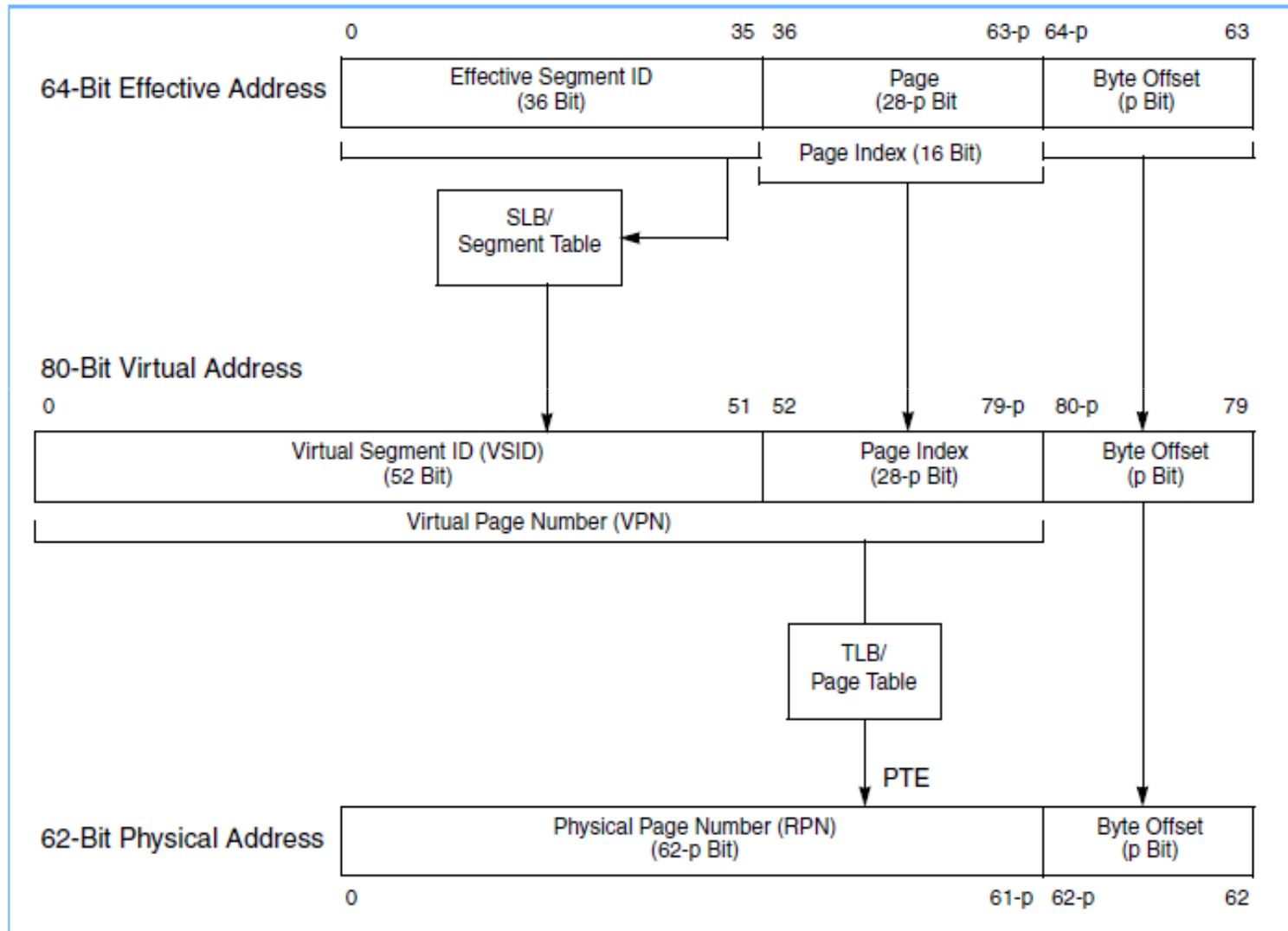
PTE[VSID] = VA[0-51]

PTE[API] = VA[52-56]

# Summary



# Summary



# Reference

- “Programming Environments Manual for 64-Bit Microprocessors” Chapter 7.

# Thanks!