

Using ESC/Java2

Architecture, Hints, and Tricks

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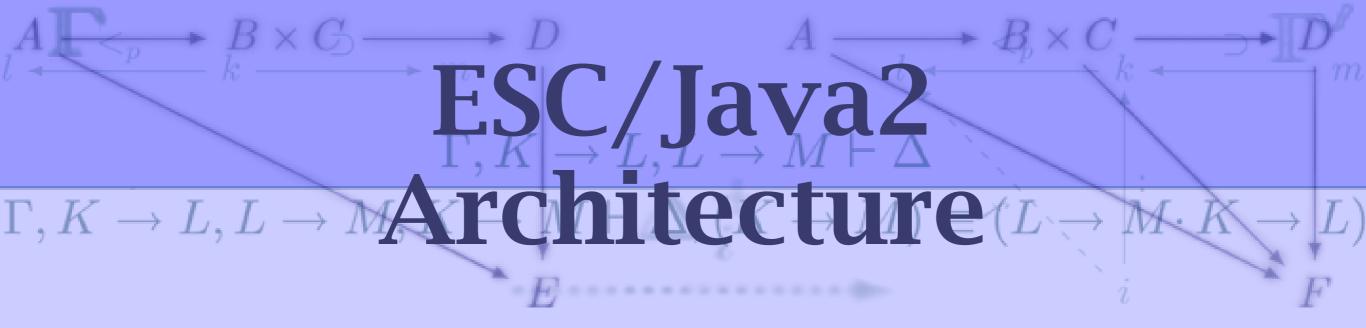
Using ESC/Java2 — M. Effectively (L)

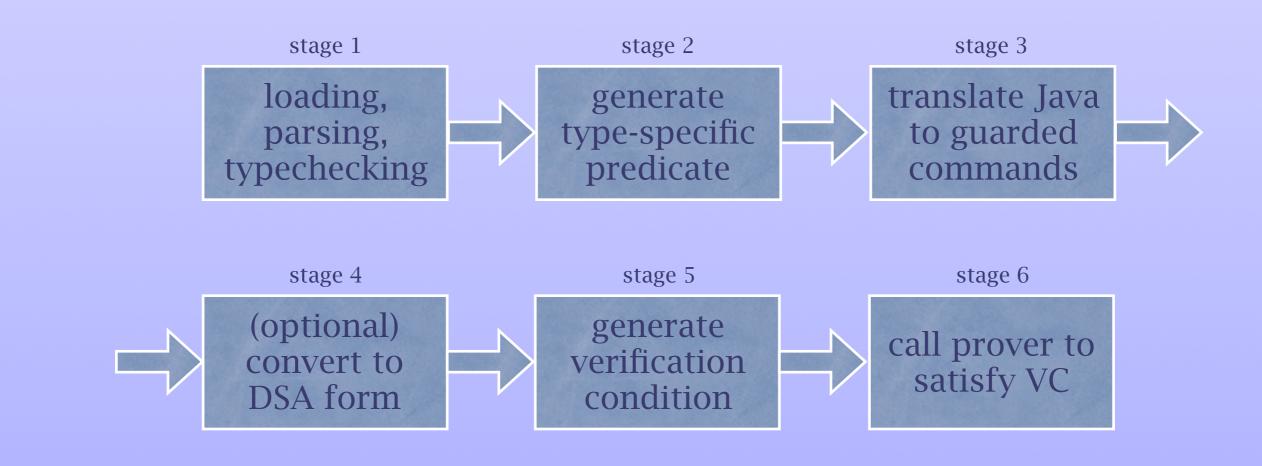
- * basic familiarity with ESC/Java2 is easy
 - * it is automatic and behaves like a compiler
- * but any non-trivial use quickly becomes very difficult and time-consuming
 - complexity of Java and JML semantics
 - limitations of logic
 - designed limitations of tool
 - limitations of Simplify theorem prover

Thinking, not $\Gamma, K \to L, L \to M, K + \mathbf{lacking}) \equiv (L \to M, K \to L)$

- * successful application of tool requires hard thought and very little labor
- recognizing that specific misbehavior implies particular errors in specifications or program code is key to effective use
- understanding theoretical underpinnings of extended static checking is very helpful
- a problem solving process for verification is needed for successful adoption

- * the *key aspects* of the verification process
 - * small steps in specification refinement and program development
 - * iterative and continuous application of ESC/Java2 to method or class of focus
 - use a standardized problem-solving technique for resolving warnings
 - * think before you type







How ESC/Java2 $\rightarrow MK \rightarrow MCKSM) \equiv (L$

* find, load, parse, and typecheck all relevant files and dependencies

- * this includes all refinements, models, etc.
- for each class being checked, generate a type-specific background predicate
 - * type and subtype information about classes and fields
 - * non_null-ness of references
 - size of constants

- * translate each routine to be check into a verification condition (VC)
 - intermediate step in this translation is to translate Java into a (Dijkstra-like) guarded-command language
 - * translation is accomplished by generating strongest-postconditions or weakestpreconditions for method body

$\begin{array}{c|c} & B \times C & D \\ \hline & How ESC/Jaya2 \\ \hline & K \rightarrow L, L \rightarrow M, KWorks (3) \equiv (L \rightarrow M, K \rightarrow L) \end{array}$

- * ask theorem prover to prove VC
 - * background predicate for Java expressed as a set of axioms
 - * type-specific background predicate generated in second step is assumed true
 - assert VC is true
- * if proof fails and prover finds counterexample, translate result back to warning message and Java, if possible

Examining the Track Results of Each Stage

- -v alone to print information on loading, parsing, refinement, etc.
- -showDesugaredSpecs to see heavyweight specifications desugared to lightweight ones (will also be -sds in next release)
- -pgc to print guarded command
- -ppvc to pretty-print verification condition
- -pxLog to print predicate sent to prover

A Stage-Driven M, K-Process M,

- ensure that the proper source and bytecode files are being loaded
 - * this is particularly important when initially setting up a verification problem and when using refinement
- * make sure that your specs mean what you think they do be examining the desugared specs
 - multiple heavyweight specs sometimes have unintuitive meaning for the beginner

A Stage-Driven M. Process (2)

- * check size of "local contributors"
 - e.g., 35 types 99 invariants 62 fields
- examine the generated VC
 - it must has a reasonable structure
 - * type-specific background predicate, followed by translated specification and program code
 - it is reasonably sized
 - [®] ~1MB is ok, multiple MB is a problem

Datagroups and Chost Fields

- datagroups are used to specify sets of fields that are interrelated
 - * the primary datagroup used by non-expert ESC/Java2 users is Object.objectState
- * ghost fields are specification-only fields that can be assigned using the set keyword
 - * the primary ghost field used by nonexperts is *Object.owner*

The Datagroup $\Gamma, K \to L, IObjectrobjectState$

```
/** A data group for the state of this object. This is used to
  * allow side effects on unknown variables in methods such as
  * equals, clone, and toString. It also provides a convenient way
  * to talk about "the state" of an object in assignable
  * clauses.
  */
//@ public non_null model JMLDataGroup objectState;
//@ represents objectState <- JMLDataGroup.IT;</pre>
```



public class Object {

The Ghost Field Object owner



$\begin{array}{c|c} \hline & B \times G \\ \hline & Dealing with \\ & \to L, L \to M, Complexity (L \to M, K \to L) \end{array}$

- * specification and code complexity are the primary factors in verification complexity
- if performing "Design by Contract" then one can "Design for Verification" also
- * if performing "Contract the Design" then verification is sometimes only possible with refinement if code modification is not permitted

Managing Spec M.Complexity

- * write and verify specs iteratively using very small steps
- * use independent heavyweight specification blocks to specify independent behaviors
- ensure your specs are sound
 - assert a false predicate to check
 - eliminate suspect predicates iteratively to determine source of unsoundness

Managing Spec Complexity (2)

- * use ghost variables or model fields to factor out complex specification subexpressions
 - helps with comprehension, not verification
- avoid universally quantified expressions
- * use the *objectState* datagroup as much as reasonable for your frame axioms
- use the owner field to disambiguate objects

- * track cyclic complexity of method bodies
 - each branch, switch case, loop, and exception block doubles complexity
- decompose methods into smallest reasonable units
 - Smalltalk and Eiffel method size rule-ofthumb applies (e.g., all methods <15 LOC)</p>
- avoid constructors that make calls

Managing Code Complexity (2)

- focus on methods that make no calls first
- work from low to high cyclic complexity
- use assertions to check
- * recognize sources of incompleteness of Java semantics
 - complex arithmetic
 - bit-level operations
 - String manipulations

Refinement for Complex Verification

- if you have a method with high cyclic complexity that you cannot refactor
 - inherit and override
 - implement and verify separate private methods for each branch of original method
 - implement overridden version as composition of verified new methods