

Implications on Specification and Verification

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- \* ESC/Java2 is an extended static checker
  - \* based upon DEC/Compaq SRC ESC/Java
  - operates on JML-annotated Java code
  - behaves like a compiler
    - error messages similar to javac & gcc
    - completely automated
    - hides enormous complexity from user



# What is Extended - L, L Static Checking?

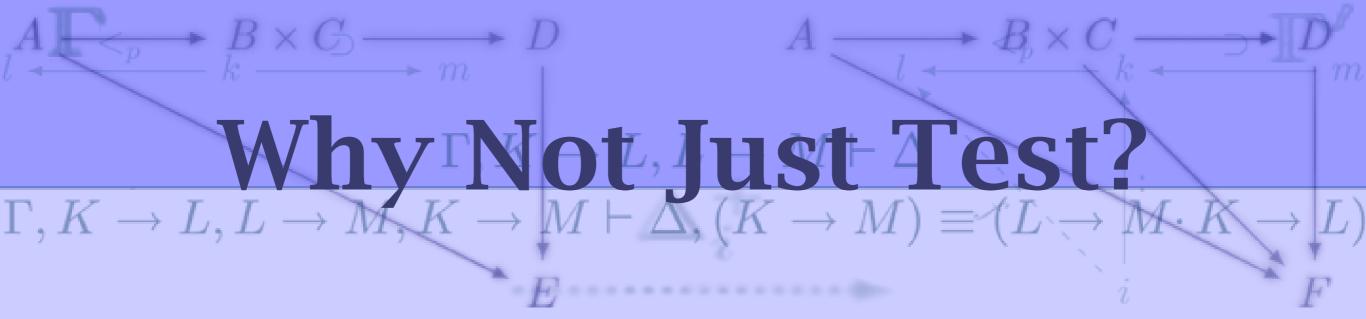
annotated source

static checker

→ Error: ...

- \* type systems
  - Error: wrong number of arguments to call
- lint & modern compilers
  - \* Error: unreachable code
- full program verification
  - Error: qsort does not yield a sorted array





- \* testing is essential, but
  - \* expensive to write and maintain
  - finds errors late
  - \* misses many errors

static checking and testing are complementary techniques

### Comparison of Static

Sheckers) =(L

quality

100%

ESC/Java2

ESC/Modula III

SRC ESC/Java

type systems

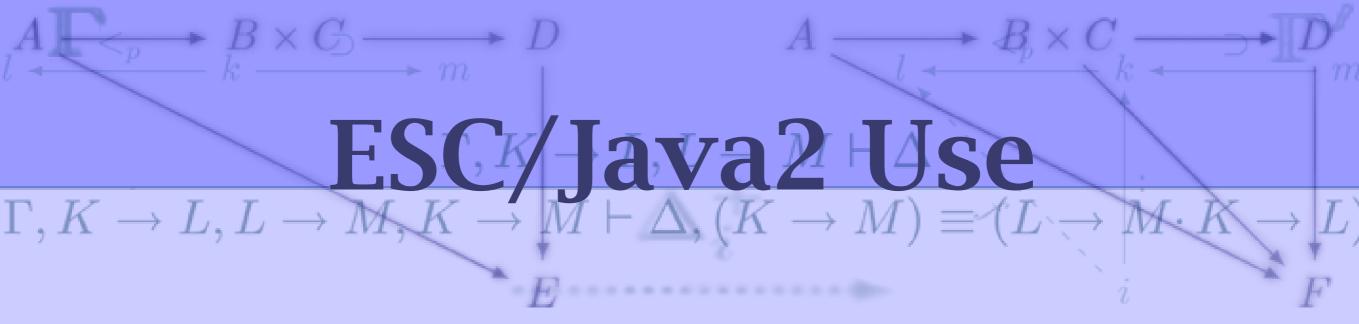
full program ion Loop

lint

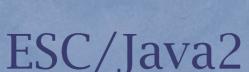
Note: graph is not to scale

effort





JMLannotaated program



"null-dereference error on line 486"

- Modularly checks for:
  - null-dereference errors
  - array bounds errors
  - \* type cast errors
  - \* specification violations
  - \* race conditions & deadlocks
  - ... dozens of other errors



# Soundness and Completeness

- a sound and complete prover is nonautomated, very complex, and expensive
  - modular checking
  - properties of arithmetic and floats
  - complex invariants and data structures
- instead, design and build an unsound and incomplete verification tool
  - \* trade soundness and completeness for automation and usability

# IML: The Java I, K - L Modeling Language

- a behavioral interface specification language
- syntax and semantics are very close to Java
- annotations written as comments in code
- - \* standard constructs include preconditions, postconditions, invariants, etc.
- \* one language used for documentation, runtime checking, and formal verification



## A JML Example and $\Gamma, K \rightarrow L, L$ ESC/Java2 Demo

```
class Bag {
   int[] a;
   int n;
   Bag(int[] input) {
     n = input.length;
     a = new int[n];
     System.arraycopy(input, 0,
                       a, 0, n);
   boolean isEmpty() {
     return n == 0;
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```

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```
int extractMin() {
  int m = Integer.MAX_VALUE;
  int mindex = 0;
  for (int i = 1; i <= n; i++) {
    if (a[i] < m) {
      mindex = i;
      m = a[i];
  a[mindex] = a[n];
  return m;
```

### The Annotated Class

```
class Bag {
 /*@ non_null */ int[] a;
 int n;
 //@ invariant 0 <= n && n <= a.length;</pre>
 //@ ghost public boolean empty;
 //@ invariant empty == (n == 0);
 //@ requires input != null;
  //@ ensures this.empty == (input.length == 0);
  public Bag(int[] input) {
    n = input.length;
   a = new int[n];
    System.arraycopy(input, 0, a, 0, n);
    //@ set empty = n == 0;
 //@ ensures \result == empty;
  public boolean isEmpty() {
    return n == 0;
```

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 $\Gamma, K \to L, L \to M, K \to M \vdash \Delta, (K)$ 

```
//@ requires !empty;
//@ modifies empty;
//@ modifies n, a[*];
public int extractMin() {
  int m = Integer.MAX_VALUE;
  int mindex = 0;
  for (int i = 0; i < n; i++) {
    if (a[i] < m) {
      mindex = i;
      m = a[i];
  n--;
 //@ set empty = n == 0;
  //@ assert empty == (n == 0);
  a[mindex] = a[n];
  return m;
```



JMLannotated -> program

translator

automatic theorem prover verification conditions

counterexamples

post-processor

warning messages



- partial semantics for Java and JML
- written in unsorted first-order logic
- \* highly tuned to current theorem prover's capabilities and quirks
  - Nelson's Simplify prover circa mid-80s
- originally consisted of 81 axioms
- \* extended by 20 axioms in ESC/Java2

## Example Java Mype-Axions

```
(DEFPRED (<: t0 t1))
(BG_PUSH (<: |T_java.lang.Object|
             IT_java.lang.Object())
; <: reflexive
(BG_PUSH
 (FORALL (t)
    (<: t t)))
; <: transitive
(BG_PUSH
 (FORALL (t0 t1 t2)
    (IMPLIES (AND (<: t0 t1) (<: t1 t2))
      (<: t0 t2))))
```

# Java Incompleteness

## Examples of Java & L - IML Semanties -

```
(DEFPRED (is x t))
(BG_PUSH (FORALL (x t)
                 (is (cast x t) t)))
(BG_PUSH (FORALL (x t)
                 (IMPLIES (is x t) (EQ (cast x t) x)))
(BG_PUSH
 (FORALL (e a i)
        (is (select (select (asElems e) a) i)
             (elemtype (typeof a)))))
(DEFPRED (nonnullelements x e)
  (AND (NEQ x null)
       (FORALL (i)
                (IMPLIES (AND (<= 0 i) (< i (arrayLength x)))
                         (NEQ (select (select e x) i) null)))))
```



- \* used for verification condition generation in Dijkstra wp/wlp style
- easy for small/research languages
- much harder for "real world" languages
  - \* typed concurrent object-oriented language
  - dynamic memory allocation and GC
  - \* exceptions
  - aliasing



## VC Generation

 $\Gamma, K \to L, L \to M, K + OF \Delta aVa_M) \equiv (L - C)$ 

### annotated source



guarded commands



verification condition

$$x = a[i++];$$

```
assume preconditions assume invariants
```

```
i0 = i;
i = i + 1;
assert (LABEL null@218: a != null);
assert (LABEL IndexNeg@218: 0 <= i0);
assert (LABEL IndexTooBig@218: i0 < a.length);
x = elems[a][i0];
...
assert postconditions</pre>
```

$$\forall i_0.(i_0=i\implies\ldots)$$

assert invariants

# $Verification \\ \Gamma, K \to L, L \to M, K Condition \\ E$

- \* formula in unsorted, first-order predicate calculus
  - equality and function symbols
  - \* quantifiers
  - arithmetic operators
  - \* select and store operations
  - \* e.g.,  $\forall x. \forall y. \exists z. (x > y \implies z \times 2 == \dots$

# Example Verification $L, L \rightarrow M$ Condition $\equiv (L \rightarrow M)$

#### verification condition large & unstructured

(EXPLIES (LBLNEG lvc.Bag.isEmpty.11.2| (IMPLIES (AND (EQ In@pre:3.6| In: 3.6|) (EQ In:3.6| (asField In:3.6| T\_int)) (EQ Ia@pre:2.8| Ia:2.8|) (EQ Ia: 2.8| (asField Ia:2.8| (array T\_int))) (< (fClosedTime Ia:2.8|) alloc) (EQ IMAX\_VALUE@pre:10..| IMAX\_VALUE:10..|) (EQ I@truel (is IMAX\_VALUE:10..| T\_int)) (EQ Ilength@pre:unknown| Ilength:unknown|) (EQ Ilength:unknown| (asField Ilength:unknown| T\_int)) (EQ Ielems@prel elems) (EQ elems (asElems elems)) (< (eClosedTime elems) alloc) (EQ LS (asLockSet LS)) (EQ Ialloc@prel alloc) (EQ Istate@prel state)) (NOT (AND (EQ I@truel (is this T\_Bag)) (EQ I@truel (isAllocated this alloc)) (NEQ this null) (EQ RES (integralEQ (select In:3.6| this) 0)) (LBLPOS Itrace.Return^0,12.4| (EQ I@truel I@truel)) (NOT (LBLNEG IException@13.2| (EQ IecReturn| IecReturn|))))))) (AND (DISTINCT IecReturn|) (< 10000000 pos2147483647)))

# Problems with Current Logic

- unsorted
  - no mental model, no type checking
- \* tightly coupled to Simplify prover
  - unmaintained, two generations old
- very incomplete
  - \* want to verify new properties and some functional specifications
- never checked for soundness

# $\begin{array}{c} \stackrel{\mathbb{P}}{\longrightarrow} \stackrel{\mathbb{R}}{\longrightarrow} \stackrel{\mathbb{R}}\longrightarrow \stackrel{\mathbb{R}}{\longrightarrow} \stackrel{\mathbb{R}}\longrightarrow \stackrel{\mathbb{R}}{$

- partial semantics for Java and JML
- \* written in *sorted* first-order logic
- independent of any particular prover
- \* written in SMT-LIB
  - supported by many new provers
- ~100 axioms thus far
  - no reduction because no subsorts or overloading



## Sorts of 4New-Logic $\Gamma, K \to L, L \to M, K \to M \vdash \Delta, (K \to M) \equiv L \to L$

```
:sorts ( # sort that represents *values* of Java's boolean base type
         Boolean
        # sort that represents *values* of all Java's base types
        # but for Boolean
        Number
        # sort that represents all Java non-base types
         ReferenceType
         # ... represents object references
         Reference
        # ... represents object values
         Object
        # Boolean, Number, Object fields
         BooleanField
         NumberField
         ReferenceField
         # ... represents the heap
        Memory )
```

## Example Axioms $\Gamma, K \to L, L \to M, K \to M \vdash \Delta, (K \to M) \equiv (L \to M)$

#### Benefits of New Logic $L, L \rightarrow M, K \rightarrow M \vdash \Delta, (K \rightarrow M) \equiv (L \rightarrow M)$

- \* ESC/Java2 will support multiple provers
  - use multiple provers concurrently
  - \* choose prover(s) based upon context
- proof of soundness
  - new logic being encoded in PVS and Coq
- increase ESC/Java2 soundness, completeness, and performance
  - \* able to verify larger, more complex programs than ever before



# $\begin{array}{c} A \\ \hline \\ Current Work \\ \hline \\ \Gamma, K \to L, L \to M, K \to M \vdash \Delta, (K \to M) \equiv (L \to M, K \to L) \\ E \end{array}$

- \* initial version of new logic sketched out
  - \* found several type errors in original logic
  - dramatically more understandable
- beginning to use new provers
  - Sami from Tinelli and Harvey from Ranise
- incorporate new provers into ESC/Java2
  - increase independence of tool from Simplify prover

### Open Questions $\rightarrow M, K \rightarrow M \vdash \Delta, (K \rightarrow M) \equiv (L \rightarrow M)$

- how to "factor out" calculus and logic from implementation of ESC/Java2
- would like to prove that the new logic subsumes the old logic
- how to integrate with other logics
  - \* of particular interest is how to integrate with full verification Loop
  - \* how to perform proof reuse when moving from first-order to higher-order semantics

### Acknowledgements

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