

Loop Invariants

Insertion and Selection Sort

Goals:

- Be able to analyze an algorithm's correctness using a given loop invariant.
- Be able to describe in words, implement using a loop invariant, and analyze the runtime of:
 - Selection sort
 - Insertion sort

Tools for Reasoning about Algorithms

Precondition, Postcondition

```
/** return the max value in A
 * precondition: A is nonempty
 * postcondition: max value of A is returned */
public int findMax(int[] A) {
    int max = A[0];
    // invariant: max is the largest value in A[0..i]
    for (int i = 1; i < A.length; i++) {
        if (A[i] > max) {
            max = A[i];
        }
    }
    return max;
}
```

Interface, not implementation

The **precondition** is true **before** method execution.
The **postcondition** is true **after** method execution.

Loop Invariant

```
/** return the max value in A
 * precondition: A is nonempty
 * postcondition: max value of A is returned */
public int findMax(int[] A) {
    int currentMax = A[0];
    // invariant: currentMax is the max of A[0..i]
    for (int i = 1; i < A.length; i++) {
        if (A[i] > max) {
            max = A[i];
        }
    }
    return max;
}
```

currentMax is the largest value in:
A[0..1] A[0..i] A[0..a.length]

A **loop invariant** is true **before**, **during**, and **after** the loop.
(at the *end* of each iteration)

Loop Invariant

largest value
is currentMax

Precondition: A

$i=1$?
-------	---

Invariant: A

largest value is currentMax	i ?
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Postcondition: A

largest value is currentMax

$i=a.length$

$A[0..1]$

$A[0..i]$

$A[0..a.length]$

The **loop invariant** is true **before**, **during**, and **after** the loop.

Another Example

```
/** rearrange A so all negative values are to  
    * the left of all non-negative values */  
public void separateSign(int[] A);
```

Postcondition: A

< 0	≥ 0
-------	----------

```

/** rearrange A so all negative values are to
    * the left of all non-negative values */
public void separateSign(int[] A) {

```

Precondition: A

h	$?$	t
-----	-----	-----

Invariant: A

$h \rightarrow$	$?$	$\leftarrow t$
< 0	$?$	≥ 0

Postcondition: A

$t \quad h$	< 0	≥ 0
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Four concerns:

1. **Initialization:** Make the invariant true at the start.
2. **Termination:** Make the loop end when the postcondition is true.
3. **Progress:** Make progress towards the postcondition.
4. **Maintenance:** Make the invariant true after each iteration.

Insertion Sort

Insert $A[i]$ into the sorted sublist $A[0..i-1]$.

Selection Sort

Find the smallest element in $A[i..n]$ and place it at $A[i]$.

<https://visualgo.net/bn/sorting>

Insertion Sort

Insert $A[i]$ into the sorted sublist $A[0..i-1]$.

Invariant: A

sorted	$?$
-----------------	-----

Selection Sort

Find the smallest element in $A[i..n]$ and place it at $A[i]$.

Invariant: A

$\text{sorted}, \leq A[i..n]$	$?$
-------------------------------	-----

<https://visualgo.net/bn/sorting>

```
insertionSort(A):
```

```
    i = 0;
```

```
    while i < A.length:
```

```
        // push A[i] to its sorted position
```

```
        // increment i
```

```
selectionSort(A):
```

```
    i = 0;
```

```
    while i < A.length:
```

```
        // find min of A[i..A.length]
```

```
        // swap it with A[i]
```

```
        // increment i
```

Developing InsertionSort

Precondition: A

?

Invariant: A

i
sorted

?

Postcondition: A

sorted

Four tasks:

1. **Initialization:** Make the invariant true at the start.
2. **Termination:** Make the loop end when the postcondition is true.
3. **Progress:** Make progress towards the postcondition.
4. **Maintenance:** Make the invariant true after each iteration.

Developing SelectionSort

Precondition: A

?

Invariant: A

$sorted, \leq A[i..n]$

i

?

Postcondition: A

$sorted$

Four concerns:

1. **Initialization:** Make the invariant true at the start.
2. **Termination:** Make the loop end when the postcondition is true.
3. **Progress:** Make progress towards the postcondition.
4. **Maintenance:** Make the invariant true after each iteration.

```

insertionSort(A):
    i = 0;
    while i < A.length:
        j = i;
        while j > 0 and A[j] > A[j-1]:
            swap(A[j], A[j-1])
            j--
        i++

```

ABCD: What's the best and worst-case asymptotic runtime complexity of insertionSort?

	Best	Worst
A	$O(n)$	$O(n)$
B	$O(n^2)$	$O(n)$
C	$O(n)$	$O(n^2)$
D	$O(n^2)$	$O(n^2)$

Why is this best-case runtime interesting?

```

insertionSort1(A):
    i = 0;
    while i < A.length:
        j = i;
        while j > 0 and A[j] < A[j-1]:
            swap(A[j], A[j-1])
            j--
        i++

```

```

insertionSort2(A):
    i = 0;
    while i < A.length:
        j = i;
        tmp = A[i];
        while j > 0 and tmp < A[j-1]:
            A[j] = A[j-1]
            j--
        i++

```

ABCD: What's the best and worst-case asymptotic runtime complexity of insertionSort2?

	Best	Worst
A	$O(n)$	$O(n)$
B	$O(n^2)$	$O(n)$
C	$O(n)$	$O(n^2)$
D	$O(n^2)$	$O(n^2)$

Practice problems

Write a precondition, postcondition, and loop invariant for the *inner* loop of InsertionSort.

Develop SelectionSort using the precondition, loop invariant, and postcondition by completing the four tasks:

- Initialization

- Termination

- Progress

- Maintenance