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Big-O-Notation Analysis
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   Part A:
    void fillint n)
     \mathcal{E}
        int i=2;
         while (icn) &
              1100)
              i=ixi; -> exponential growth
         3
     3
                                                                     Oruntime = O (log(log(n)))
                           where x is the # of loops
                 io=2, so, ... log(n)=i**
                                                                thus log(log(n)
                                 log( log(n)) =x
Part B:
      void f2(intn) {
               for (inti=1; i<=n; i++) &
                   if( (i % (int) sqrt(n))==0) }
                         for (int k=0; k < pow (i,3); k++) }
                                110(1) run-time
                  3
           3
          outer loop runtime: icn;
           if statement run-time: Only runs if i is divisible by sartn. so combined
             \Theta(\overline{n} \cdot n^3) = \Theta(n^{\frac{1}{2}} \cdot n^{\frac{6}{2}}) = \Theta(n^{\frac{3}{2}})
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Part C:
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Solution: E @ routine = @ (n·n·1·logen)·1) (O(n*·logen))

for (int i=1/i<=n;i++)? > O(n)for (int k=1;k<=n;k++)? > O(n)if (A[k] ==i)? > O(n)for (int m=1;m<=n; m=m+m)? $> O(\log(n))$ // do something for O(1) runtime > O(1)

Outer Loop: i=n, i+t. So O(n)

2nd outmost loop: k<=n, k+t. So O(n)

if-statement: Only runs if A[k]=i
inner-loop: m<=n; m+m; so O(log(n))

 $O = O(n^2 \log cn)$ $\Omega = \Omega(n^2)$ Thus, θ run-time θ

Part D:

outer loop: $\sum_{i=0}^{n} \theta(i) = \Theta(N)$ if statement: runs if $i=size:(6ize)(\frac{3}{2})^{n}.so^{i}\theta(\log(n))$ for $i=\log(n)$ inner-loop: occurs for $\theta(n(\frac{3}{2})^{\log(n)})$

 $\Omega(n) \leftarrow \text{continue for } n < 10$ $O(n \cdot \log(n)) = (n \cdot \log(n))$ $Comight \text{ be } O(n \cdot \log(n))$

Final Answer

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Fun-time = 0(n)

is O (modra) is