

16.3 A 28nm 16.9-300TOPS/W Computing-in-Memory Processor Supporting Floating-Point NN Inference/Training with Intensive-CIM Sparse-Digital Architecture

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Computing-in-memory (CIM) has shown high energy efficiency on low-precision integer multiply-accumulate (MAC) [1-3]. However, implementing floating-point (FP) operations using CIM has not been thoroughly explored. Previous FP CIM chips [4-5] require either complex in-memory FP logic or have lengthy alignment-cycle latencies arising from converting FP data having different exponents into integer data. The challenges for an energy-efficient and accurate FP CIM processor are shown in Fig. 16.3.1. Firstly, aligning an FP vector onto a CIM module requires a long bit-serial sequence due to infrequent but long tail values, incurring many CIM cycles. In this work, we observe that most exponents of FP data are clustered in a small range, which motivates dividing FP operations into high-efficiency intensive-CIM and flexible sparse-digital parts. Secondly, to implement the intensive-CIM + sparse-digital FP workflow, a sparse digital core is required for flexible intensive/sparse processing. Thirdly, the FP alignment brings more random sparsity. Though analog CIM can utilize random sparsity with a low-resolution ADC, the corresponding sparse strategy for digital CIM has not been explored.

To overcome the above challenges, this work proposes an accurate and energy-efficient FP CIM processor for neural network (NN) inference/training applications. The main innovations include: 1) An FP-to-INT CIM workflow for the intensive FP operations with reduced execution cycles and high efficiency. 2) A flexible sparse-digital core for the remaining FP operations with encoded sparse activation/weight, which assists the CIM core to achieve both high energy efficiency and high accuracy. 3) An energy-efficient low-MAC-value (MACV) CIM macro that adopts a two-stage adder tree to utilize random sparsity and eliminate high-bit-position circuits with no accuracy loss.

Figure 16.3.2 shows the overall FP CIM processor, illustrating the intensive-CIM sparse-digital workload division. The CIM processor consists of a RISC-V CPU, 512KB global SRAM, a CIM core with four low-MACV CIM macros and a sparse digital core. A convolution or matrix-vector multiplication between FP weight (W) and activation (A) data is divided into intensive and sparse parts, and then re-organized into one intensive and two sparse parts. The intensive part $W_{intensive} \cdot A_{intensive}$ is considered heavy computation and implemented on the CIM core for high efficiency. $W_{all} \cdot A_{sparse}$ and $W_{sparse} \cdot A_{intensive}$ are comparatively sparse computations, so they are implemented on the flexible sparse-digital core, which also performs the final sum. Note that $W_{intensive}$ can be extracted on-chip from W_{all} , while W_{sparse} and A_{sparse} consume extra storage. The CIM core supports both FP and INT operations. Together with the RISC-V CPU and FP digital core, the back-propagation, error calculation and weight update operations can be executed on-chip to support FP/INT inference and training applications.

Figure 16.3.3 presents the overall FP-to-INT CIM architecture for intensive FP operations. It fetches and selects the intensive part ($W_{intensive}$) from the total weight (W_{all}), and then stores it into the CIM macro in INT8/INT16 format. The exponents (Exp) of the intensive FP activations reside in a small range $[E_{min}, E_{max}]$ to reduce the CIM execution cycles. A concise bit-serial FP-to-INT transfer unit converts the FP data as multi-cycle bit vectors into the CIM macro. When Exp is not aligned (1st step), it aligns Exp towards E_{max} with a self-increasing operator, and outputs the sign bit. Once Exp is aligned (2nd step), the MSB of the mantissa is sent to the CIM macro bit-serially with a left-shift operator. An example of the FP-to-INT conversion/transfer is presented, requiring $(E_{max} - E_{min} + 12)$ cycles.

Following the FP-to-INT conversion, the CIM macro runs FP MAC operations in integer format. The accumulation circuits support FP/INT activation and INT16/8/4 weight configurations, and convert the INT results back to FP format. The long bit-width INT results for various configurations are translated to an FP21 format with 5 additional mantissa bits (LSB) to avoid intermediate precision loss in the accumulation step, while the accumulation result is stored in a self-defined FP16 format. The self-defined FP16 format omits the Exp shift operation in the standard FP16, since it can be moved to the write-back step and be calculated together with the activation/weight-incurred Exp shift. By omitting the Exp shift in accumulation circuits, the critical path is shortened and power consumption is reduced. The result can also be written back as INT8/INT4 for a low-bit-precision NN. By only computing the intensive FP operations in the CIM module, this work significantly reduces the CIM execution cycles by 2.7 \times compared with the long-tail full FP16 execution on the CIM module. Compared with direct CIM FP alignment and digital FP circuits, this work shows 2.7 \times and 1.5 \times energy efficiency improvement.

Figure 16.3.4 shows the flexible sparse digital core for the remaining sparse FP operations, which implements $W_{all} \cdot A_{sparse}$, $W_{sparse} \cdot A_{intensive}$ and the final sum. Each 32b quantity in the digital core can represent either two 16b dense FP data items, or one 16b sparse FP data item with a 16b index. To simplify the circuit design with the limited index bitwidth, the row/column of the activation/weight are represented as absolute indices, while the channels of activation/weight are represented as an incremental index, which ensures large representation range with simple decoding circuits. Two examples are presented for the convolution of sparse weight/activation. In each cycle, the sparse digital core receives one sparse activation (weight) and eight dense weights (activations) for the MAC operation. Each MAC unit can be configured as BF16/FP16 by tuning the exponent/mantissa mask for bit[9:7]. Accumulating the results of the intensive-CIM and sparse-digital cores, the final accuracy matches with the FP16 baseline. The sparse digital FP execution only takes 4.59% of the total system energy. Layer-wise evaluation on ImageNet, ResNet50 shows that the sparse digital core incurs no performance loss by parallel CIM/digital execution.

Figure 16.3.5 shows the low-MACV CIM macro, consisting of a wordline decoder, activation input, normal IO, 16 128 \times 2b SRAM banks and 16 low-MACV adder trees. The 128 \times 2b SRAM bank adopts a ping-pong structure, with one 6T cell in each ping/pong part. The weight data can be directly selected from Q/QB with a fixed path instead of from the bitline (BL/BLB) with precharge/discharge in each cycle. Observe that with the ping-pong structure and weight update technique, the CIM macro can utilize a large external SRAM as its own storage with slight power overhead. Therefore, there is no need to enlarge CIM capacity with multi-SRAM cells in each CIM unit. The ping-pong structure with the fixed-weight path provides either 1.47 \times energy reduction by eliminating precharge/discharge similar to [6], or 2 \times transistor-count reduction at the same power compared with 12T SRAM cells [2].

The low-MACV adder tree in Fig. 16.3.5 contains two stages with the high-bit-position computation circuits discarded. We have observed that most of the CIM MAC results (out of a maximum of M) reside in a small region ($<N/2$ or $<N/4$). The FP-to-INT alignment further contributes to this phenomenon, which motivates random sparsity optimization for the digital CIM. The 1st stage of the 128 \times 1b adder tree accumulates 16 \times 1b to a 4b result, while the 2nd stage accumulates 8 \times 4b to a 6b result. Following the MACV result distribution, the 1st/2nd stage can discard the computation for the high 1b/2b (or high 2b/3b) in the conservative (or aggressive) mode. This work implements the conservative solution to avoid the accuracy loss. Compared with a normal adder tree, the conservative/aggressive low-MACV adder tree offers 9.0%/13.9% power reduction with no accuracy loss.

The measurement results of the fabricated 28nm CIM processor are shown in Fig. 16.3.6. This chip can work at 10-400MHz with 0.469-0.9V digital voltage and 0.397-0.9V CIM voltage, while the best system energy efficiency is at 50MHz with 0.485V (digital), 0.458V (CIM). This chip is verified on several NN models with INT8/4/FP16 bit-precision on the Cifar-10 and ImageNet datasets. From dense models to the average of sparse models, the chip achieves 275-1615TOPS/W@INT4 and 17.2-91.3TOPS/W@FP16 macro energy efficiency for inference tasks. In summary, this work introduces an intensive-CIM sparse-digital architecture, and demonstrates an accurate and energy-efficient FP CIM chip. The INT4/FP16 macro energy efficiency is 6.99 \times /1.97 \times compared with the state-of-the-art CIM processor [5]. Fig. 16.3.7 shows the die photo and chip summary.

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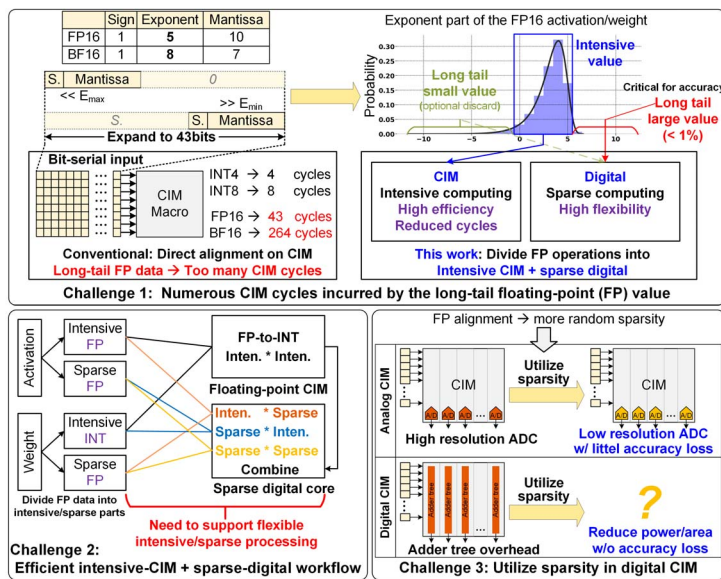


Figure 16.3.1: Challenges for energy-efficient and accurate floating-point (FP) CIM.

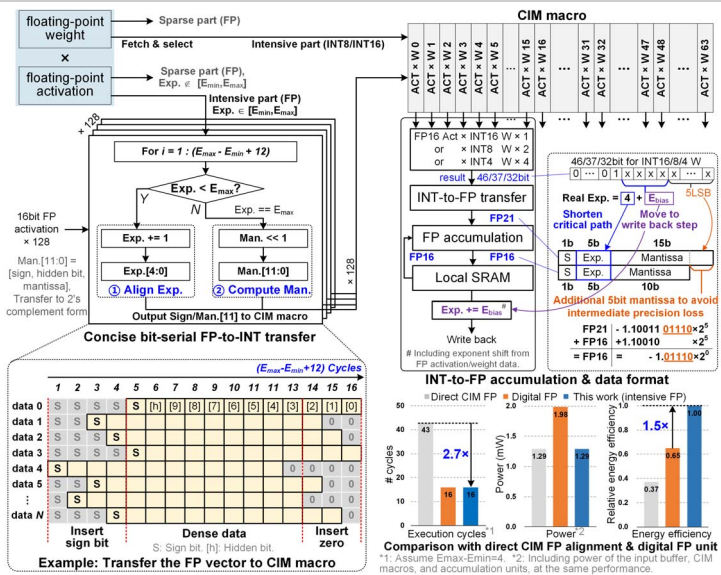


Figure 16.3.3: The FP-to-INT CIM workflow for intensive FP operations with cycle count reduction.

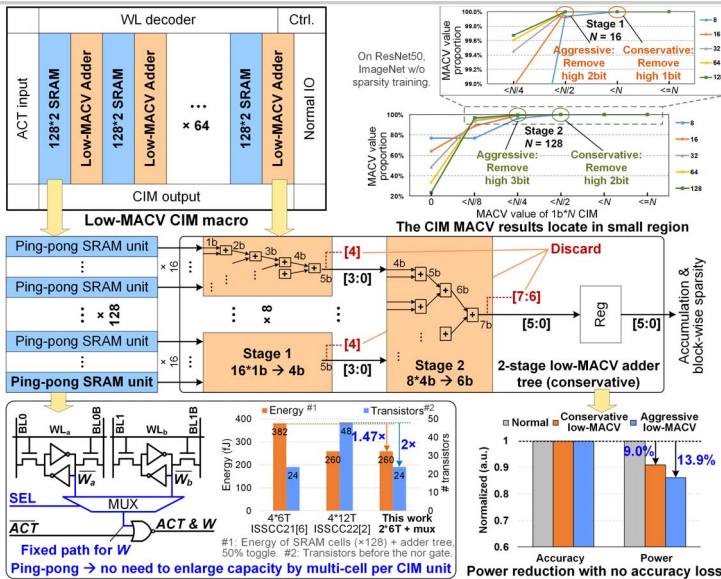


Figure 16.3.5: Energy-efficient low-MACV CIM macro to utilize random sparsity without accuracy loss.

$$W * A = (W_{intensive} + W_{sparse}) * (A_{intensive} + A_{sparse})$$

$$= W_{intensive} * A_{intensive} + (W_{intensive} + W_{sparse}) * A_{sparse} + W_{sparse} * A_{intensive}$$

$$= W_{intensive} * A_{intensive} + W_{all} * A_{sparse} + W_{sparse} * A_{intensive}$$

Heavy computation **Slight sparse computation**
Intensive CIM: High efficiency **Sparse digital: High flexibility**

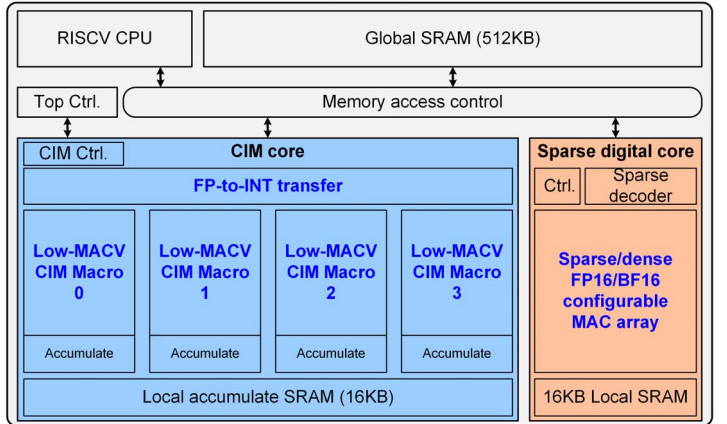


Figure 16.3.2: Overall architecture of the FP CIM processor with intensive-CIM sparse-digital division.

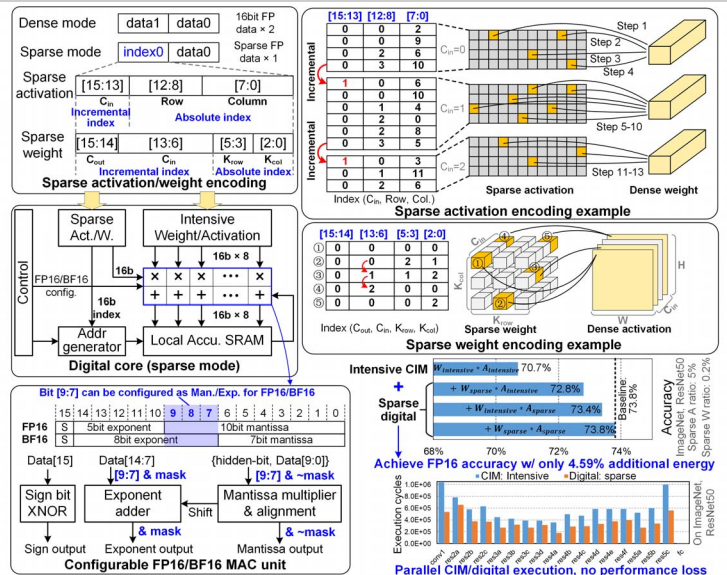


Figure 16.3.4: Flexible sparse-digital core for various sparse FP operations to ensure high accuracy.

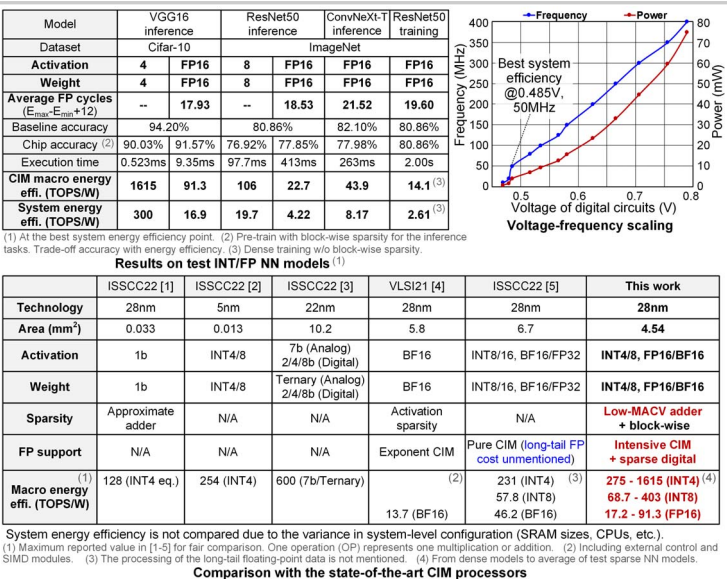


Figure 16.3.6: Measurement results and comparison with the state-of-the-art CIM processors.

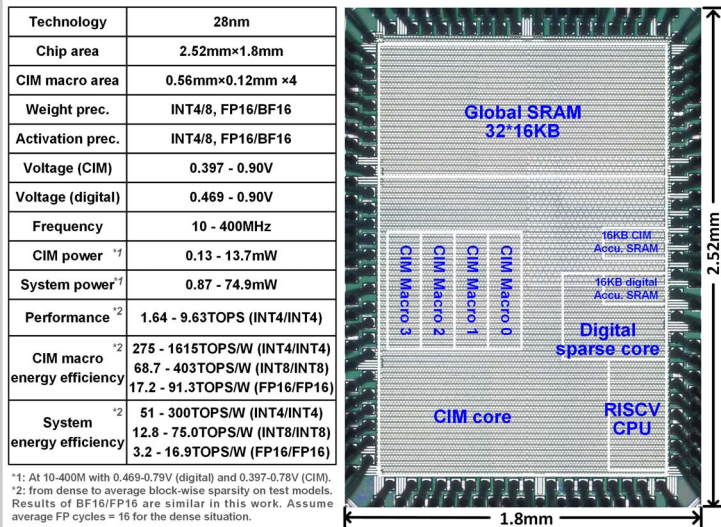


Figure 16.3.7: The chip metrics and die photo.