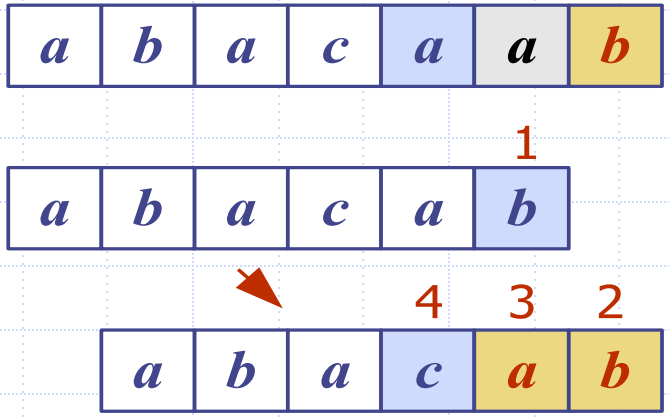
Pattern Matching



Strings

- A string is a sequence of characters
- Examples of strings:
  - Python program
  - HTML document
  - DNA sequence
  - Digitized image
- An alphabet Σ is the set of possible characters for a family of strings
- Example of alphabets:
  - ASCII
  - Unicode
  - {0, 1}
  - {A, C, G, T}
- Let P be a string of size m
  - A substring $P[i \,..\, j]$ of P is the subsequence of P consisting of the characters with ranks between i and j
  - A prefix of P is a substring of the type $P[0 \,..\, i]$
  - A suffix of P is a substring of the type $P[i \,..\, m-1]$
- Given strings T (text) and P (pattern), the pattern matching problem consists of finding a substring of T equal to P
- Applications:
  - Text editors
  - Search engines
  - Biological research

Brute-Force Pattern Matching

- The brute-force pattern matching algorithm compares the pattern P with the text T for each possible shift of P relative to T, until either
  - a match is found, or
  - all placements of the pattern have been tried
- Brute-force pattern matching runs in time $O(nm)$
- Example of worst case:
  - $T = aaa \ldots ah$
  - $P = aaah$
  - may occur in images and DNA sequences
  - unlikely in English text

Algorithm *BruteForceMatch*(T, P)

```
Input text T of size n and pattern
    P of size m
Output starting index of a
    substring of T equal to P or −1
    if no such substring exists
for i ← 0 to n − m
    { test shift i of the pattern }
    j ← 0
    while j < m ∧ T[i + j] = P[j]
        j ← j + 1
    if j = m
        return i {match at i}
    else
        break while loop {mismatch}
return -1 {no match anywhere}
```

Right to Left Matching

- Matching the pattern from right to left
- For a pattern abc:

T: `bbacdcbaabcddcdaddaaabcbcb`

P: `abc`

- Worst case is still O(n m)

Bad Character Rule (BCR) (1)

- On a mismatch between the pattern and the text, we can shift the pattern by more than one place.

```
ddbbacdcbaabcddcdaddaaabcbcb
acabc
```

Bad Character Rule (BCR) (2)

- Preprocessing –
  - A table, for each position in the pattern and a character, last occurrence of the mismatched character in P preceding the mismatch . O(n |Σ|) space. O(1) access time.

a c a b c

1 2 3 4 5

	1	2	3	4	5
a	1	1	3	3	3
b				4	4
c		2	2	2	5