



# Network Layer

Note: The slides are adapted from the materials from Prof. Richard Han at CU Boulder and Profs. Jennifer Rexford and Mike Freedman at Princeton University, and the networking book (Computer Networking: A Top Down Approach) from Kurose and Ross.

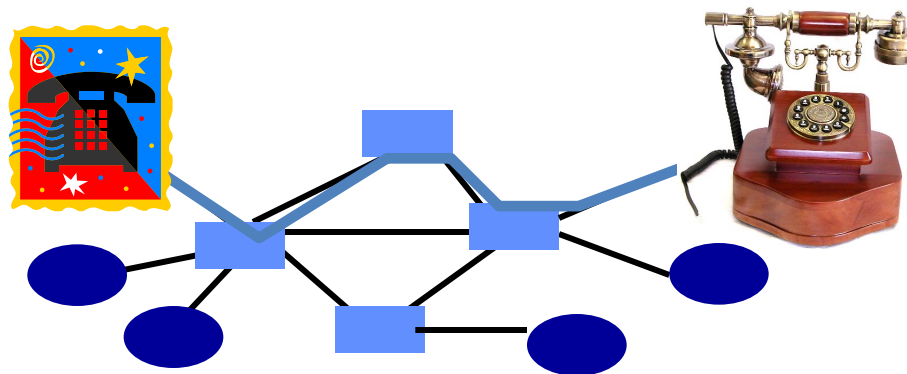
# IP Protocol Stack: Key Abstractions

<b>Application</b>	Applications	
<b>Transport</b>	Reliable streams	Messages
<b>Network</b>	Best-effort <i>global</i> packet delivery	
<b>Link</b>	Best-effort <i>local</i> packet delivery	

# Best-Effort Global Packet Delivery

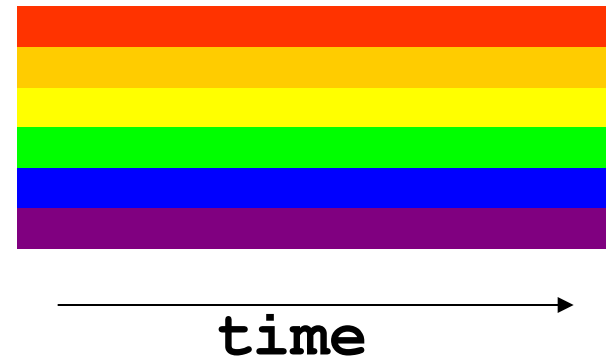
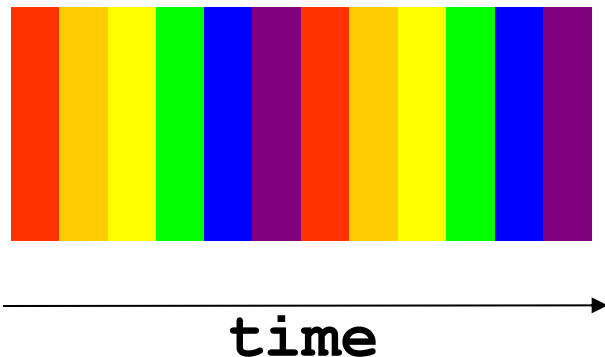
# Circuit Switching (e.g., Phone Network)

- **Source establishes connection**
  - Reserve resources along hops in the path
- **Source sends data**
  - Transmit data over the established connection
- **Source tears down connection**
  - Free the resources for future connections



# Circuit Switching: Static Allocation

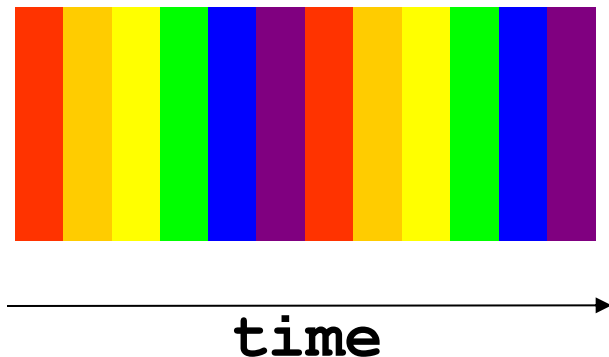
Q: Frequency-Division vs. Time-Division



# Circuit Switching: Static Allocation

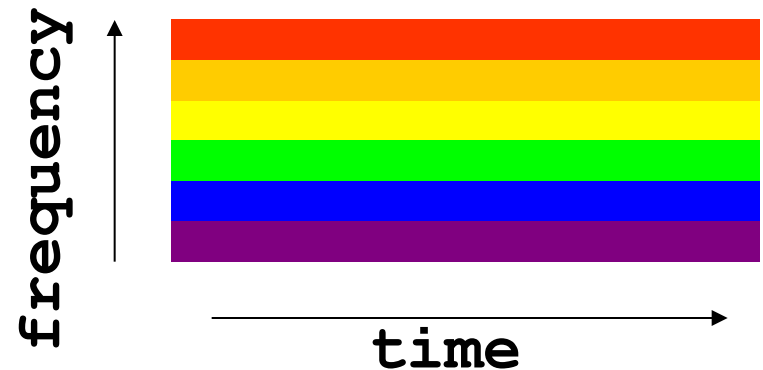
- Time-division

- Each circuit allocated certain time slots



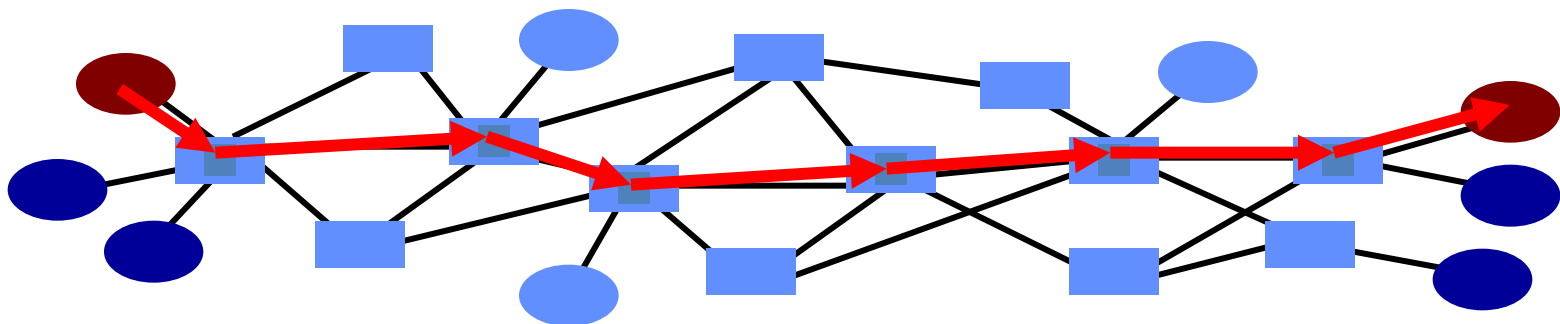
- Frequency-division

- Each circuit allocated certain frequencies

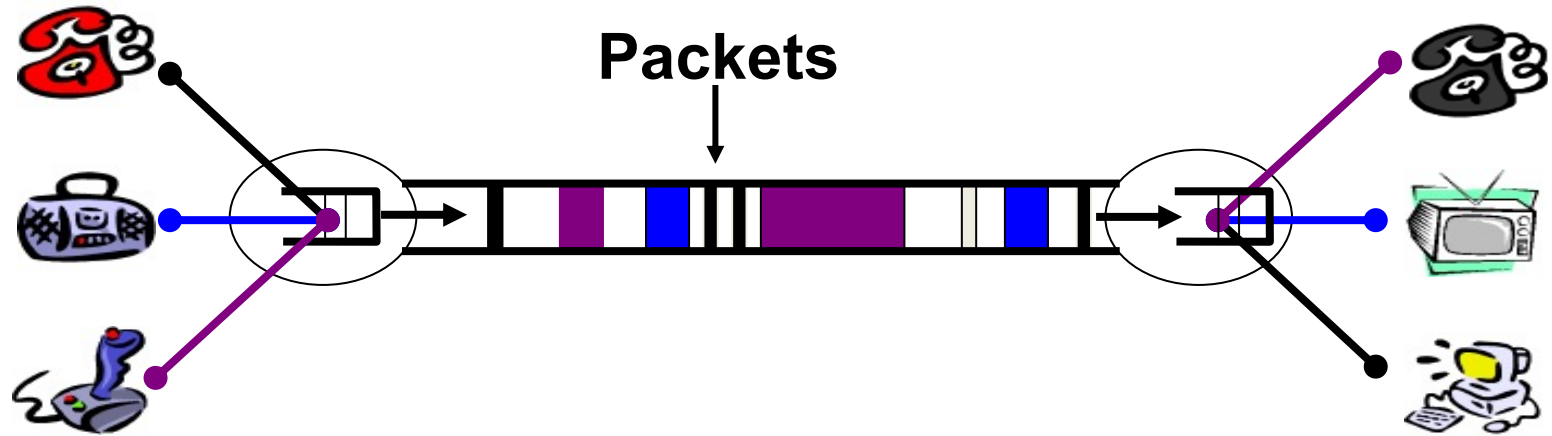


# Packet Switching (e.g., Internet)

1. Data traffic divided into packets
  - Each packet contains header (with src and dst addr)
2. Packets travel separately through network
  - Packet forwarding based on the header
  - Network nodes may store packets temporarily
  - Best effort: Packets may be loss, corrupted, reordered
3. Destination reconstructs the message



# Packet Switching: Statistical (Time Division) Multiplexing



- **Intuition: Traffic by computer end-points is bursty!**
  - Versus: Telephone traffic not bursty (e.g., constant 56 kbps)
  - One can use network while others idle
- **Packet queuing in network: tradeoff space for time**
  - Handle short periods when outgoing link demand  $>$  link speed



# Best Effort: Celebrating Simplicity

- Packets may be lost, corrupted, reordered
- Never having to say you're sorry...
  - Don't reserve bandwidth and memory
  - Don't do error detection and correction
  - Don't remember from one packet to next
- Easier to survive failures
  - Transient disruptions are okay during failover
- Easier to support on many kinds of links
  - Important for *interconnecting* different networks

# Is Best Effort Good Enough?

- Packet loss and delay
  - Sender can resend
- Packet corruption
  - Receiver can detect, and sender can resend
- Out-of-order delivery
  - Receiver can put the data back in order
- Packets follow different paths
  - Doesn't matter
- Network failure
  - Drop the packet
- Network congestion
  - Drop the packet

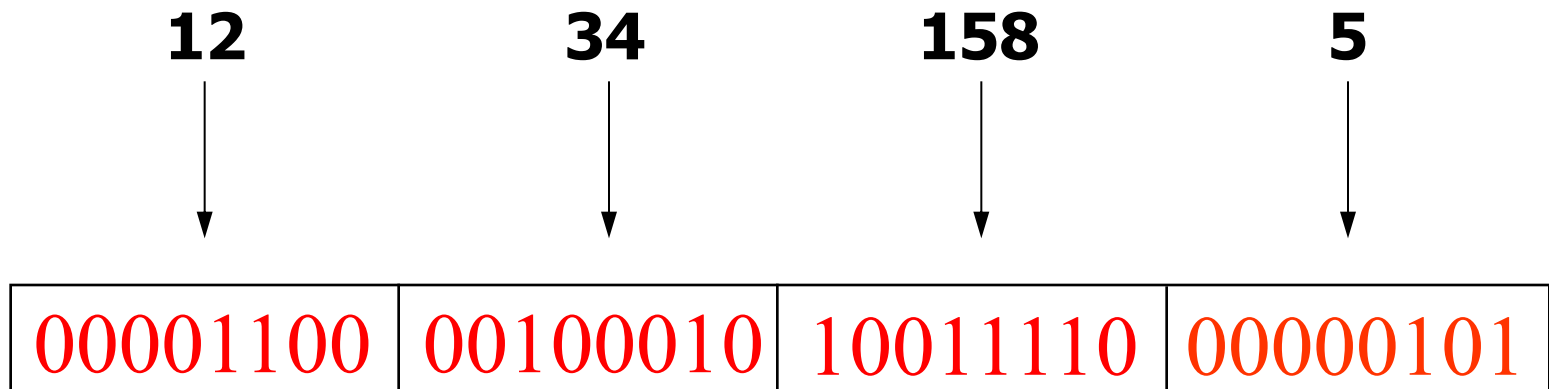
# Packet vs. Circuit Switching?

- Predictable performance Circuit
- Network never blocks senders Packet
- Reliable, in-order delivery Circuit
- Low delay to send data Packet
- Simple forwarding Circuit
- No overhead for packet headers Circuit
- High utilization under most workloads Packet
- No per-connection network state Packet

# Network Addresses

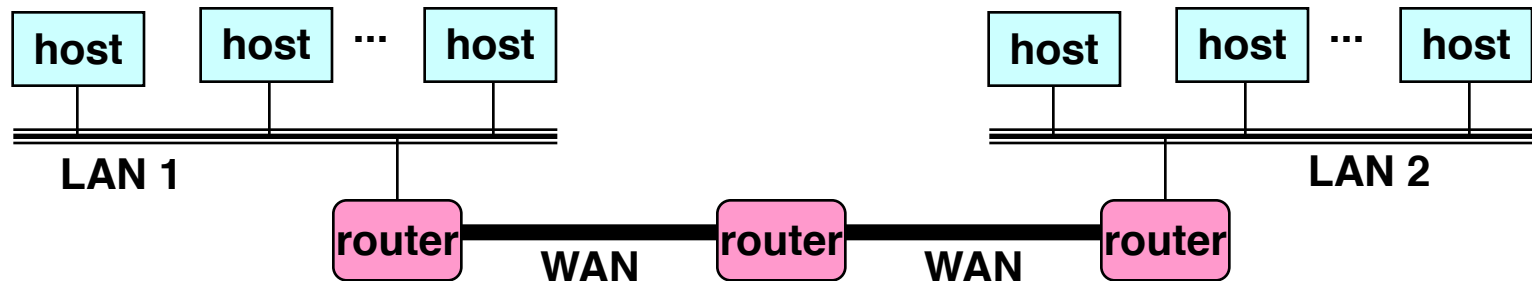
# IP Address (IPv4)

- A unique 32-bit number
- Identifies an interface (on a host, on a router, ...)
- Represented in dotted-quad notation



# Grouping Related Hosts

- The Internet is an “inter-network”
  - Used to connect networks together, not hosts
  - Need to address a network (i.e., group of hosts)

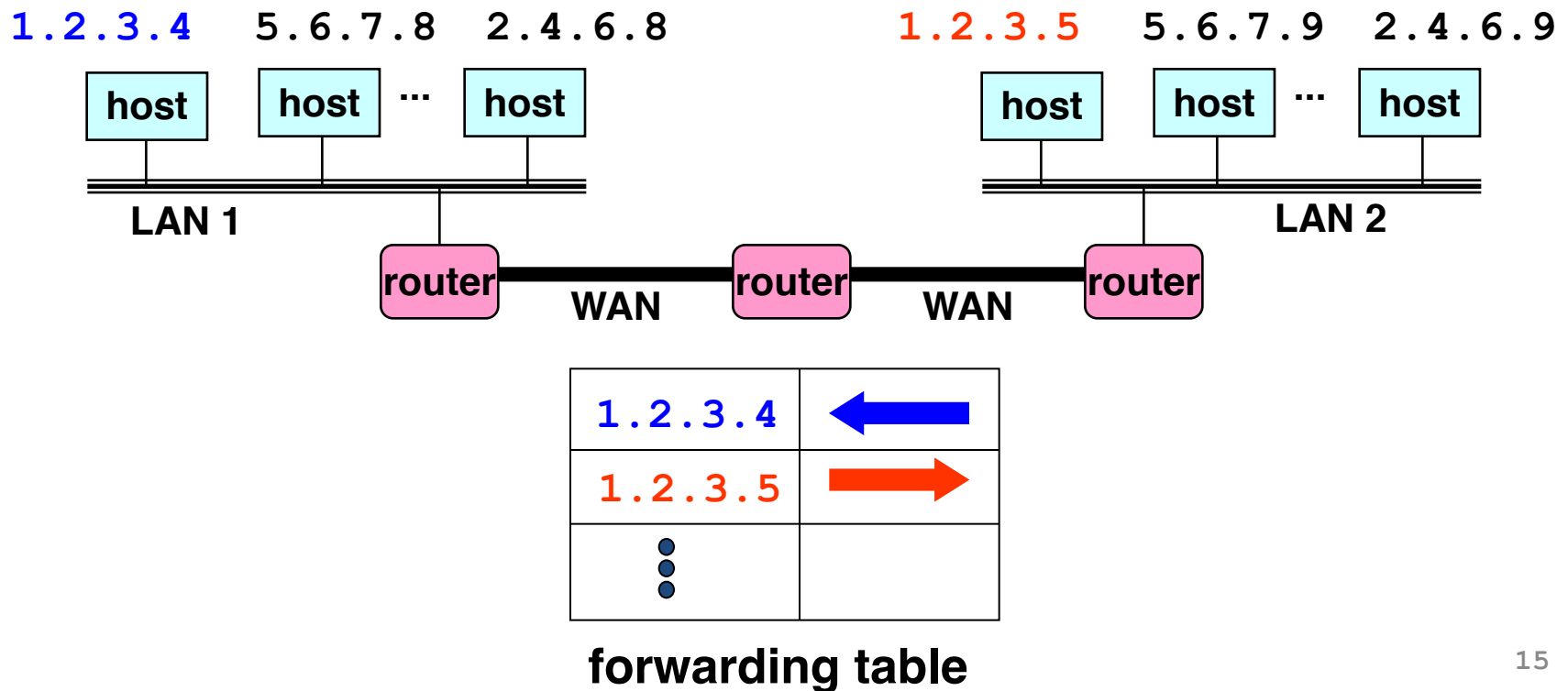


LAN = Local Area Network

WAN = Wide Area Network

# Scalability Challenge

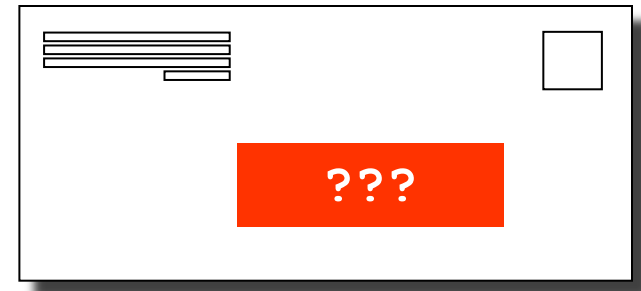
- Suppose hosts had arbitrary addresses
  - Then every router would need a lot of information
  - ...to know how to direct packets toward every host



# Hierarchical Addressing in U.S. Mail

- Addressing in the U.S. mail

- Zip code: 80309
- Building: 1045 Regent Drive
- Room in building: ECCR 1B14
- Name of occupant: Sangtae Ha



- Forwarding the U.S. mail

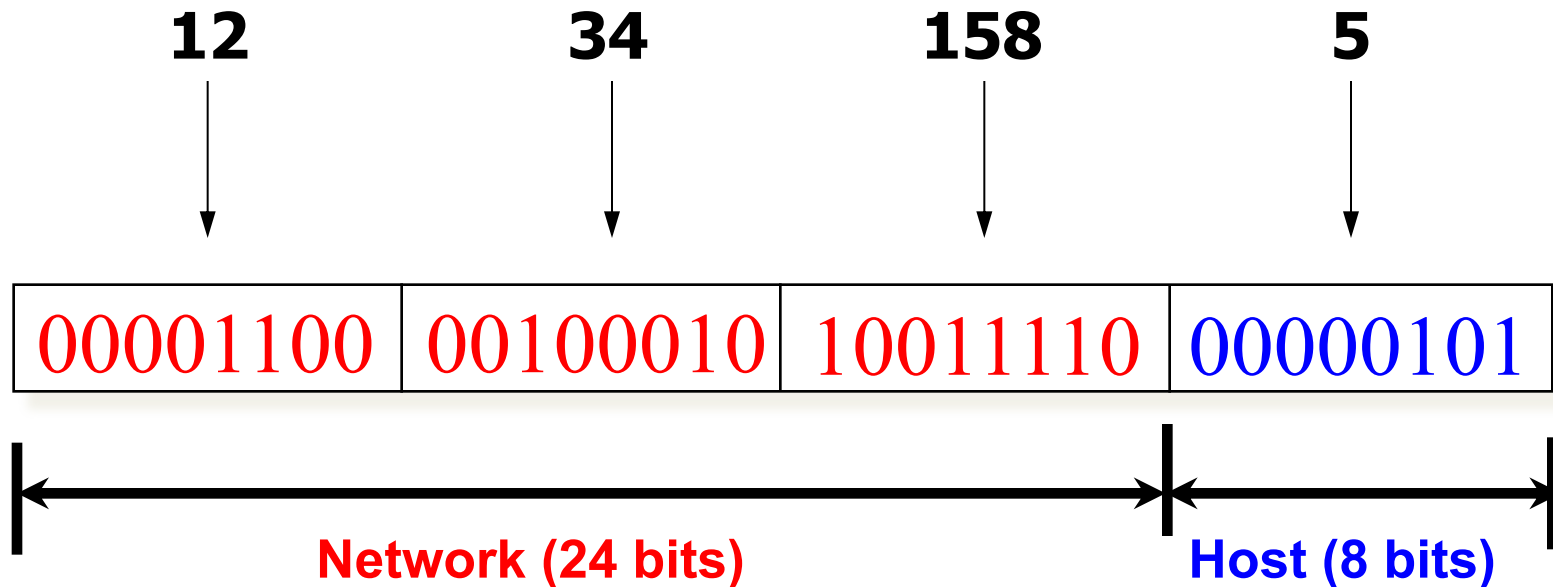
- Deliver to the post office in the zip code
- Assign to mailman covering the building
- Drop letter into mailbox for building/room
- Give letter to the appropriate person



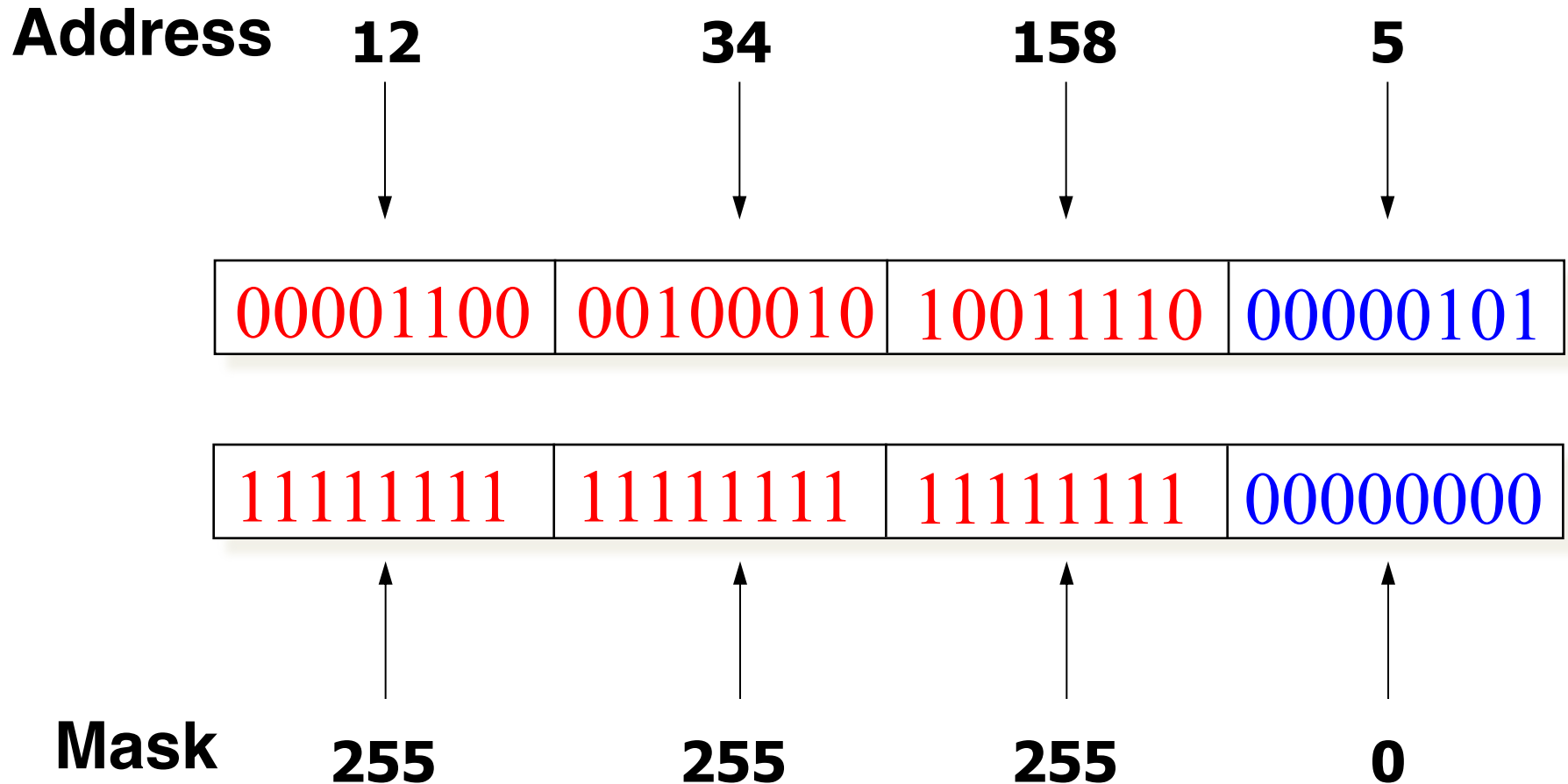


# Hierarchical Addressing: IP Prefixes

- Network and host portions (left and right)
- 12.34.158.0/24 is a 24-bit **prefix** with  $2^8$  addresses



# IP Address and 24-bit Subnet Mask



# Network Address and Broadcast Address

- IP address: 10.10.1.97/23
- Network address – IP & Netmask

IP: 00001010.00001010.00000001.01100001

& Netmask: 11111111.11111111.11111110.00000000

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Network Addr: 00001010.00001010.00000000.00000000

Network address: 10.10.0.0

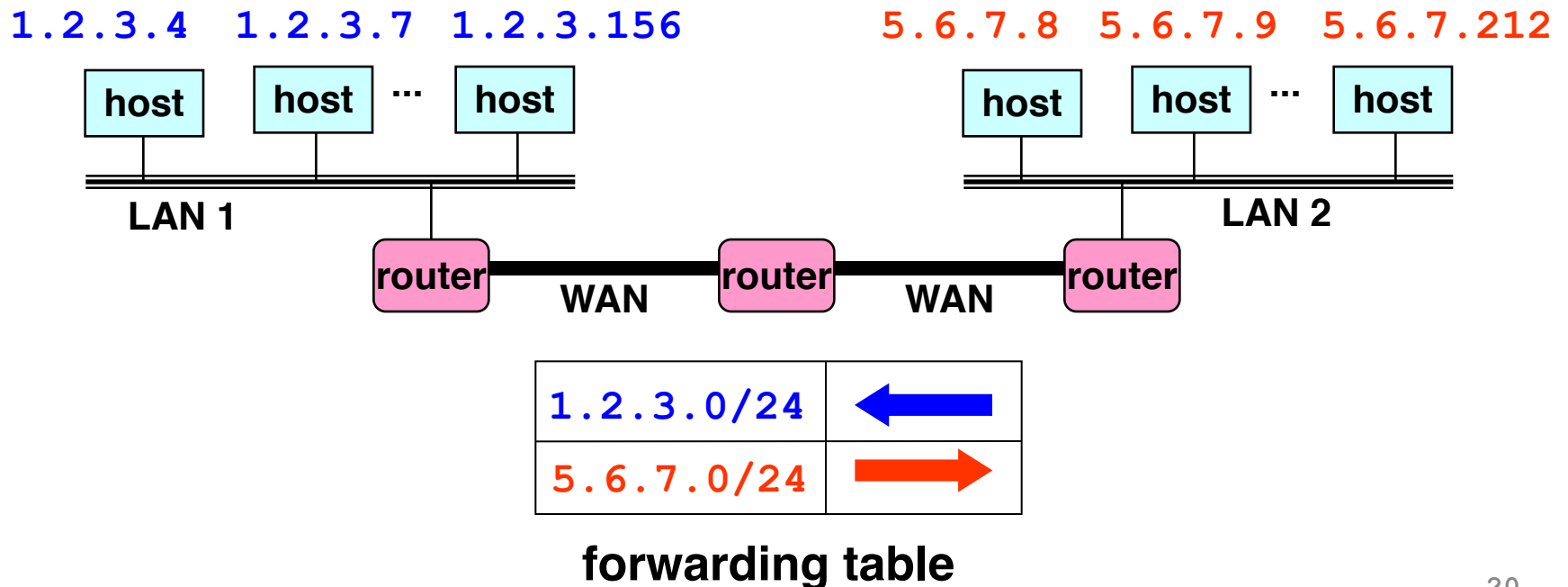
- Broadcast address – IP | ~Netmask

Broadcast address: 10.10.1.255

Please check the example at <https://www.countryipblocks.net/identifying-the-network-and-broadcast-address-of-a-subnet>

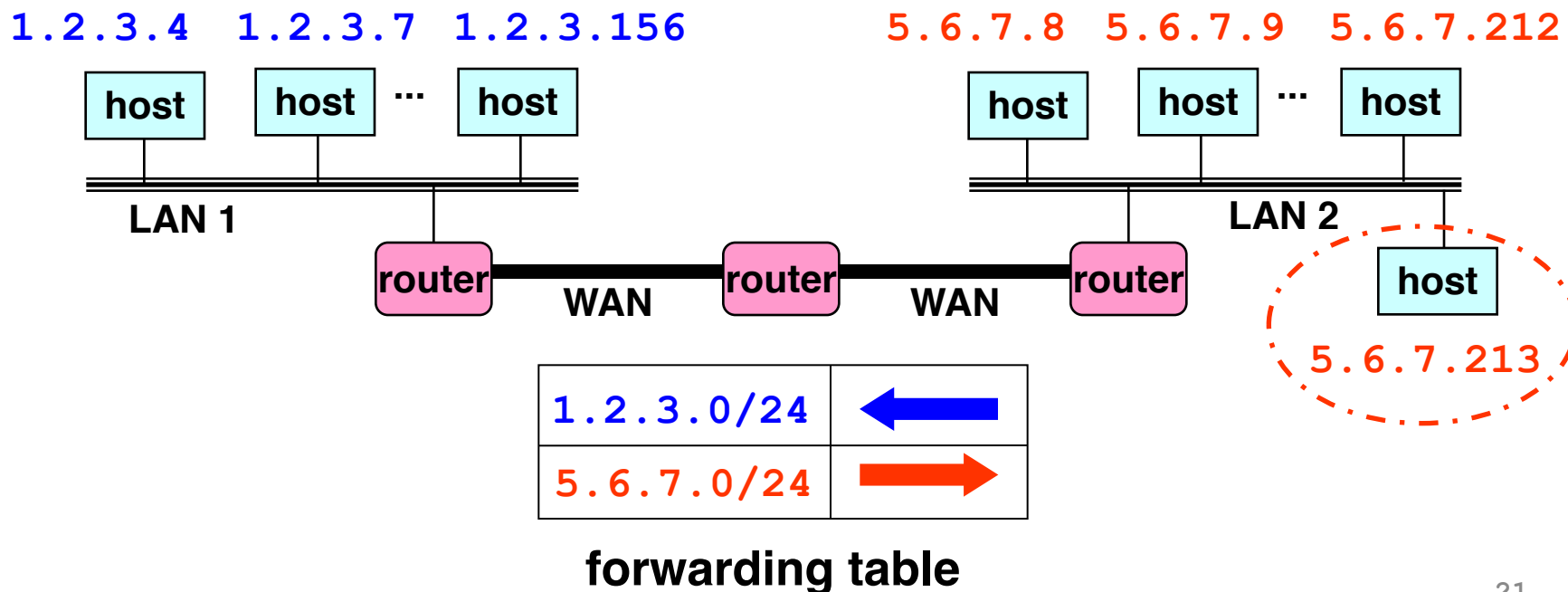
# Scalability Improved

- Number related hosts from a common subnet
  - 1.2.3.0/24 on the left LAN
  - 5.6.7.0/24 on the right LAN



# Easy to Add New Hosts

- No need to update the routers
  - E.g., adding a new host 5.6.7.213 on the right
  - Doesn't require adding a new forwarding-table entry



# History of IP Address Allocation

# Classful Addressing

- In the olden days, only fixed allocation sizes
  - Class A: 0\*
    - Very large /8 blocks
  - Class B: 10\*
    - Large /16 blocks
  - Class C: 110\*
    - Small /24 blocks (e.g., AT&T Labs has 192.20.225.0/24)
  - Class D: 1110\* for multicast groups
  - Class E: 11110\* reserved for future use
- This is why folks use dotted-quad notation!

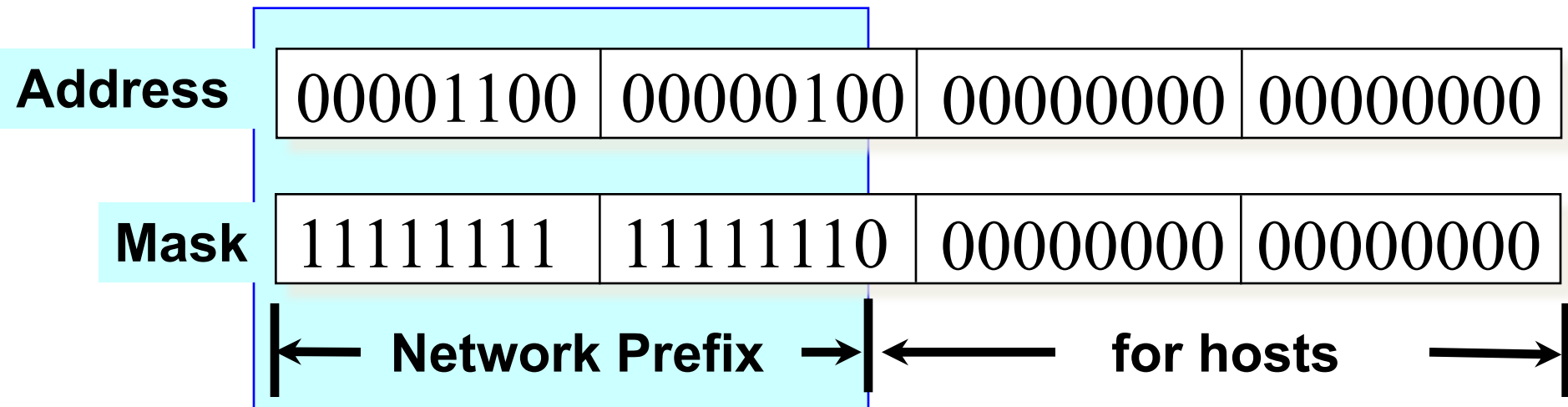
# Classless Inter-Domain Routing (CIDR)

- Use two 32-bit numbers to represent network:

Network number = IP address + Mask

**IP Address : 12.4.0.0**

**IP Mask: 255.254.0.0**

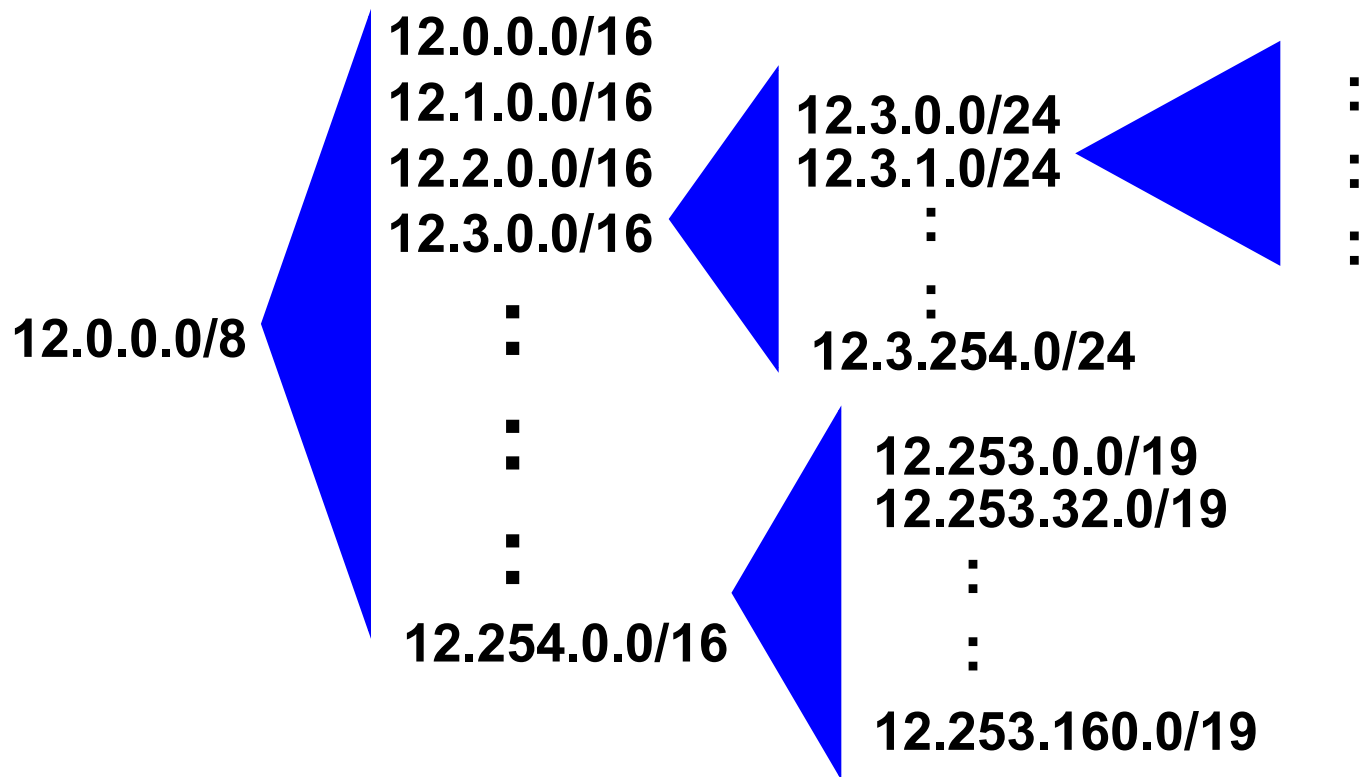


**Written as 12.4.0.0/15**



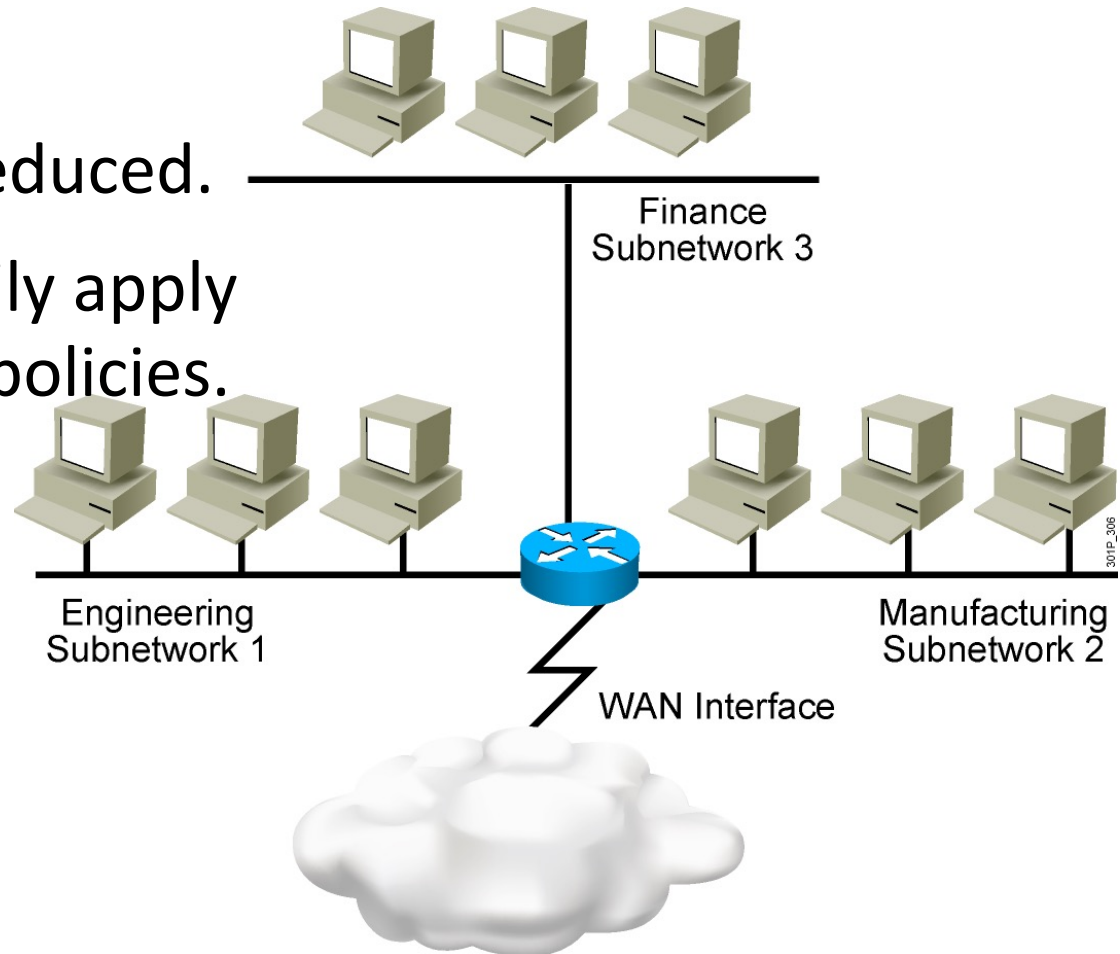
# Hierarchical Address Allocation

- **Hierarchy is key to scalability**
  - Address allocated in contiguous chunks (prefixes)
  - Today, the Internet has about 400,000 prefixes

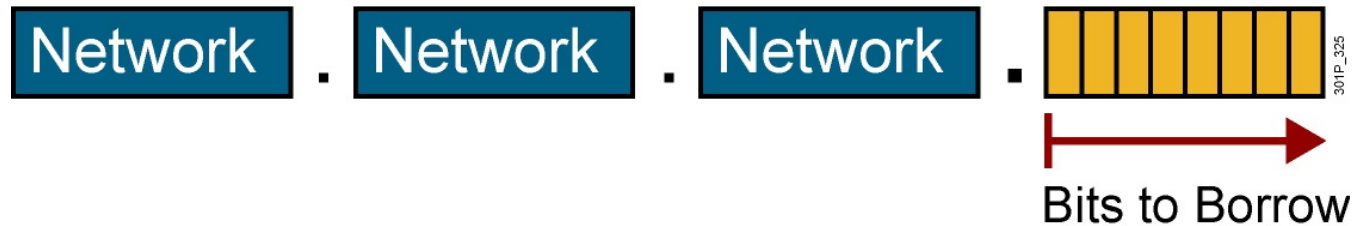


# Subnetworks (subnets)

- Smaller networks are easier to manage.
- Overall traffic is reduced.
- You can more easily apply network security policies.

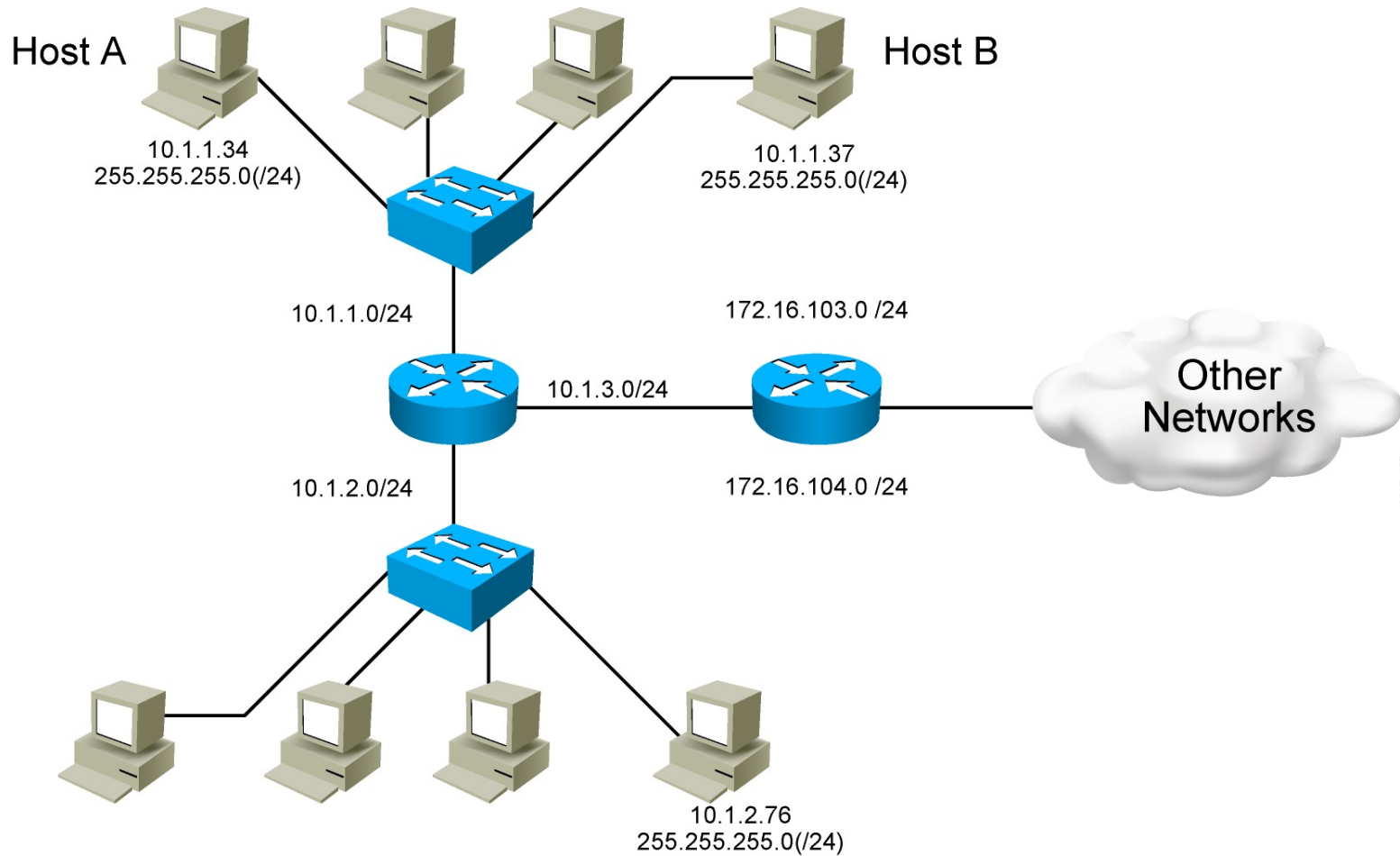


# Possible Subnets for a Class C Network

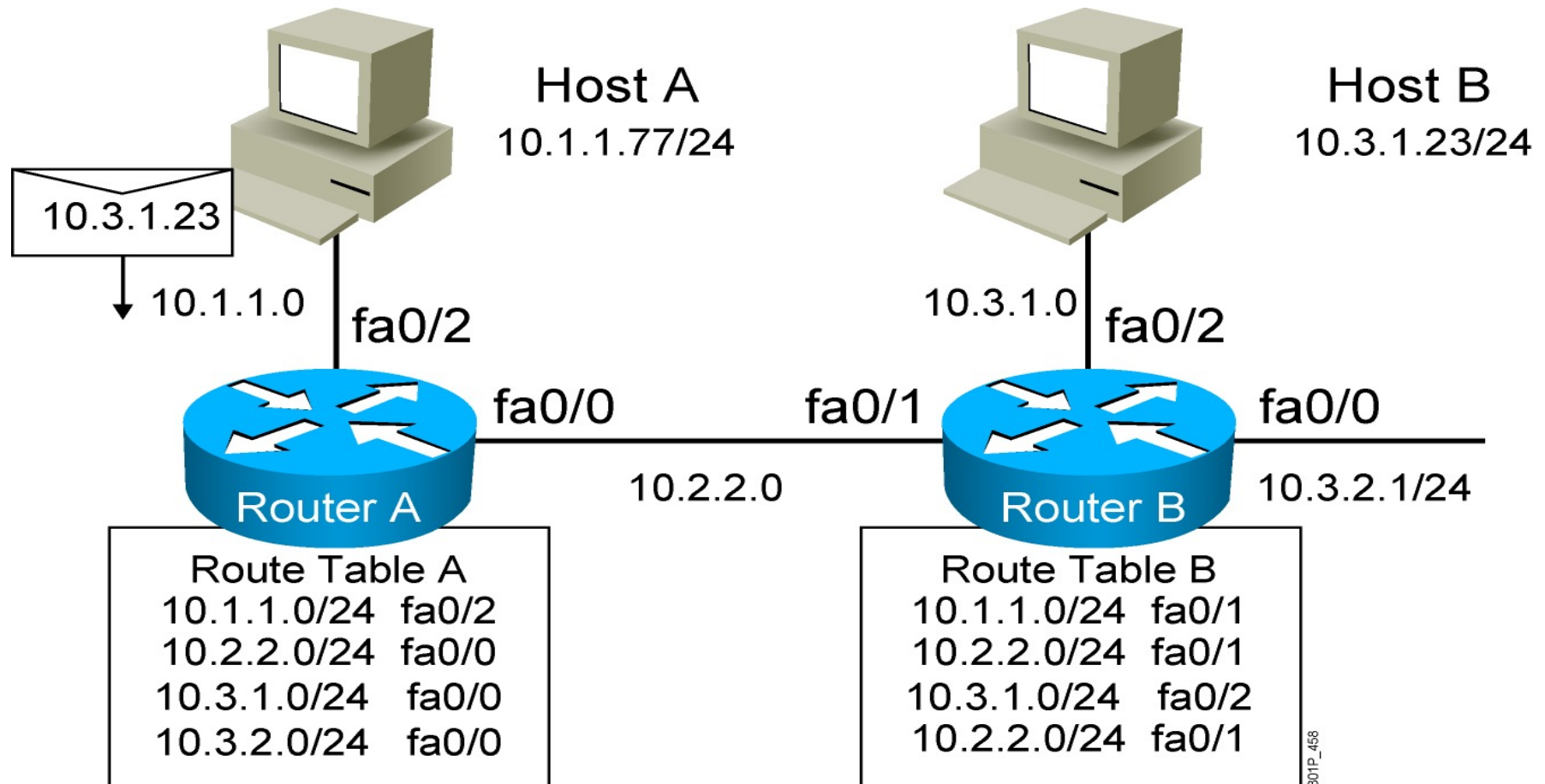


Number of Bits Borrowed (s)	Number of Subnets Possible ( $2^s$ )	Number of Bits Remaining in Host ID ( $8 - s = h$ )	Number of Hosts Possible Per Subnet ( $2^h - 2$ )
1	2	7	126
2	4	6	62
3	8	5	30
4	16	4	14
5	32	3	6
6	64	2	2
7	128	1	2

# End System Subnet Mask Operation



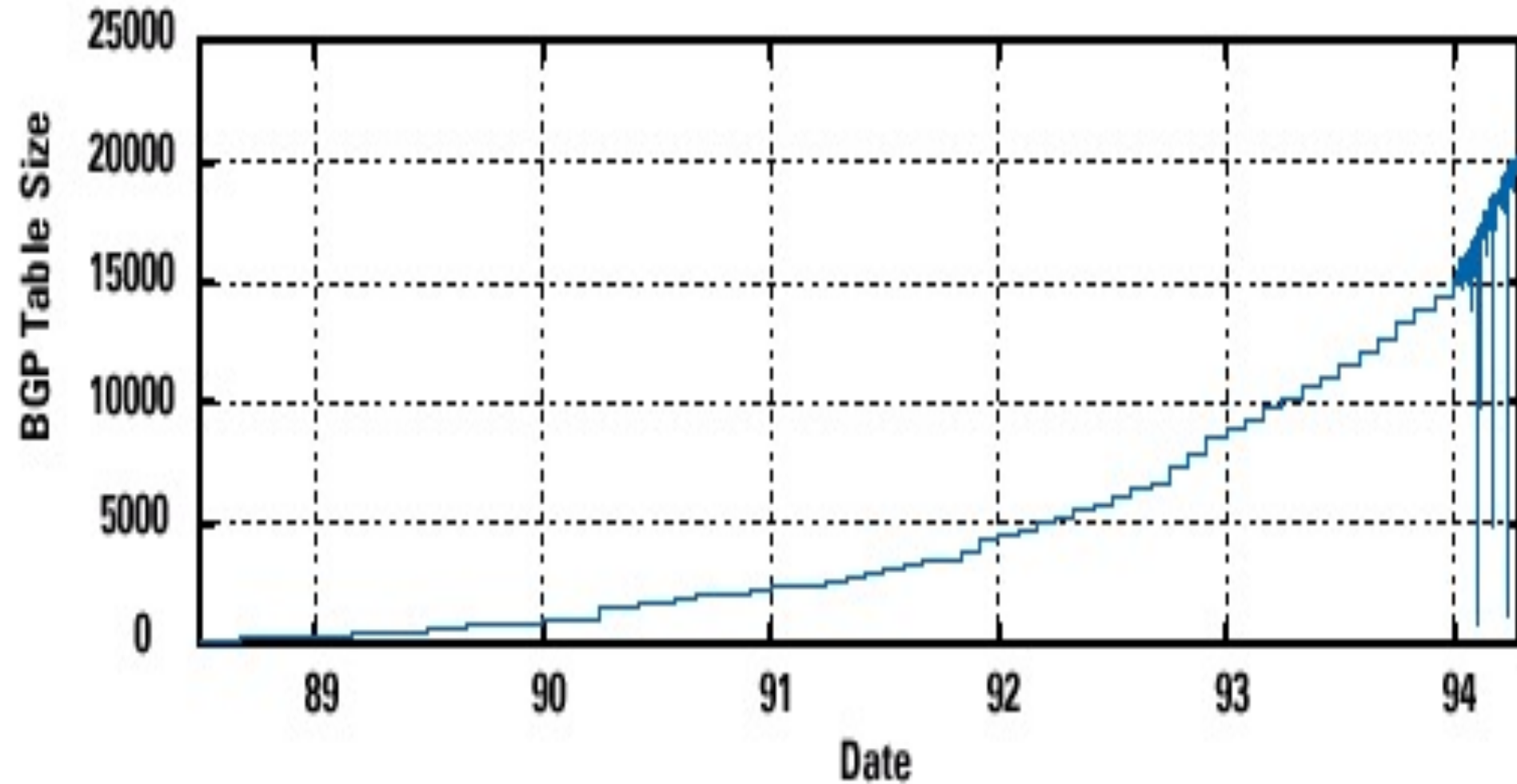
# How Routers Use Subnet Masks



# Obtaining a Block of Addresses

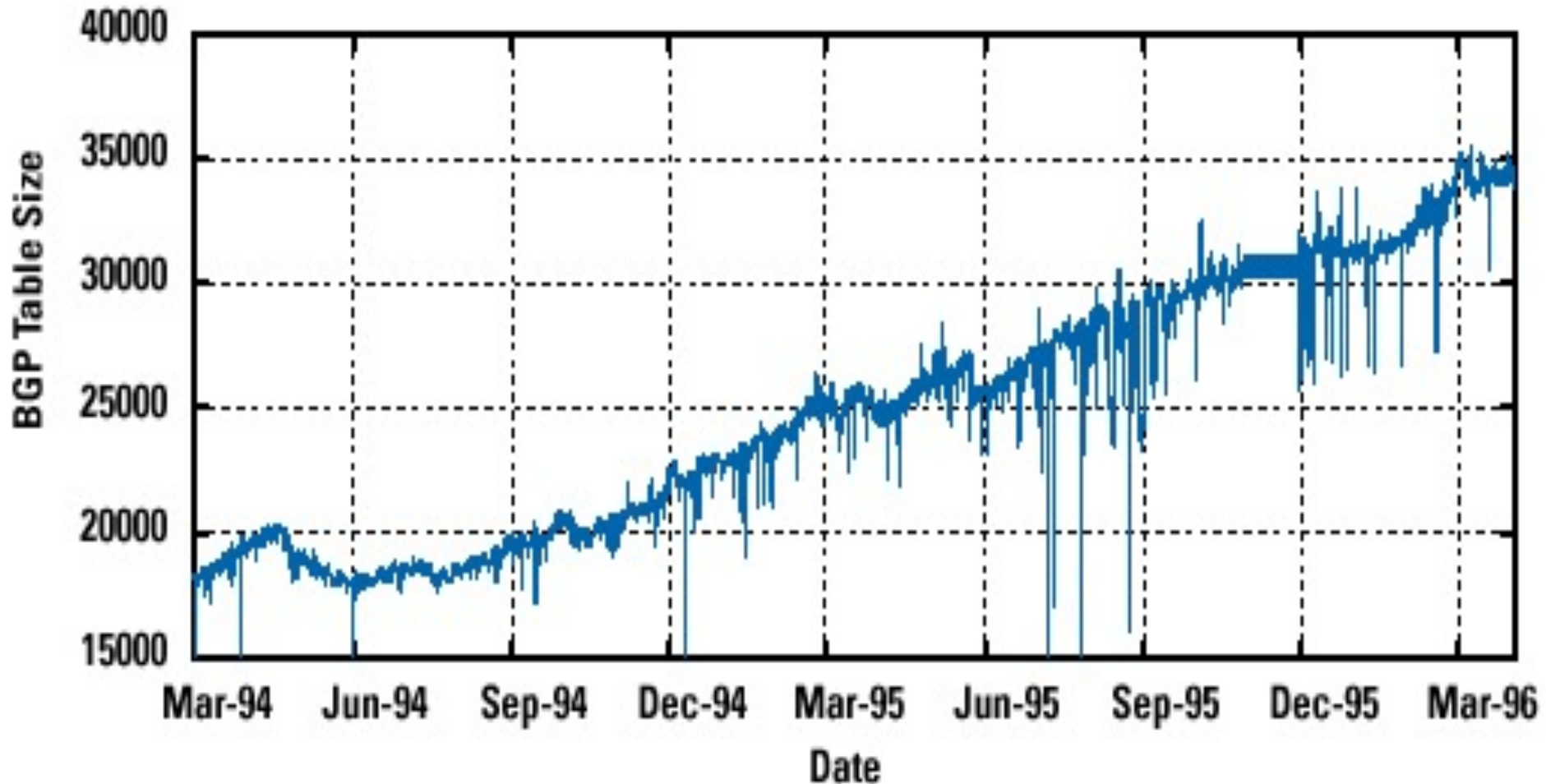
- Internet Corporation for Assigned Names and Numbers (ICANN)
  - Allocates large blocks to Regional Internet Registries
- Regional Internet Registries (RIRs)
  - E.g., ARIN (American Registry for Internet Numbers)
  - Allocates to ISPs and large institutions
- Internet Service Providers (ISPs)
  - Allocate address blocks to their customers
  - Who may, in turn, allocate to their customers...

# Pre-CIDR (1988-1994): Steep Growth



Growth faster than improvements in equipment capability

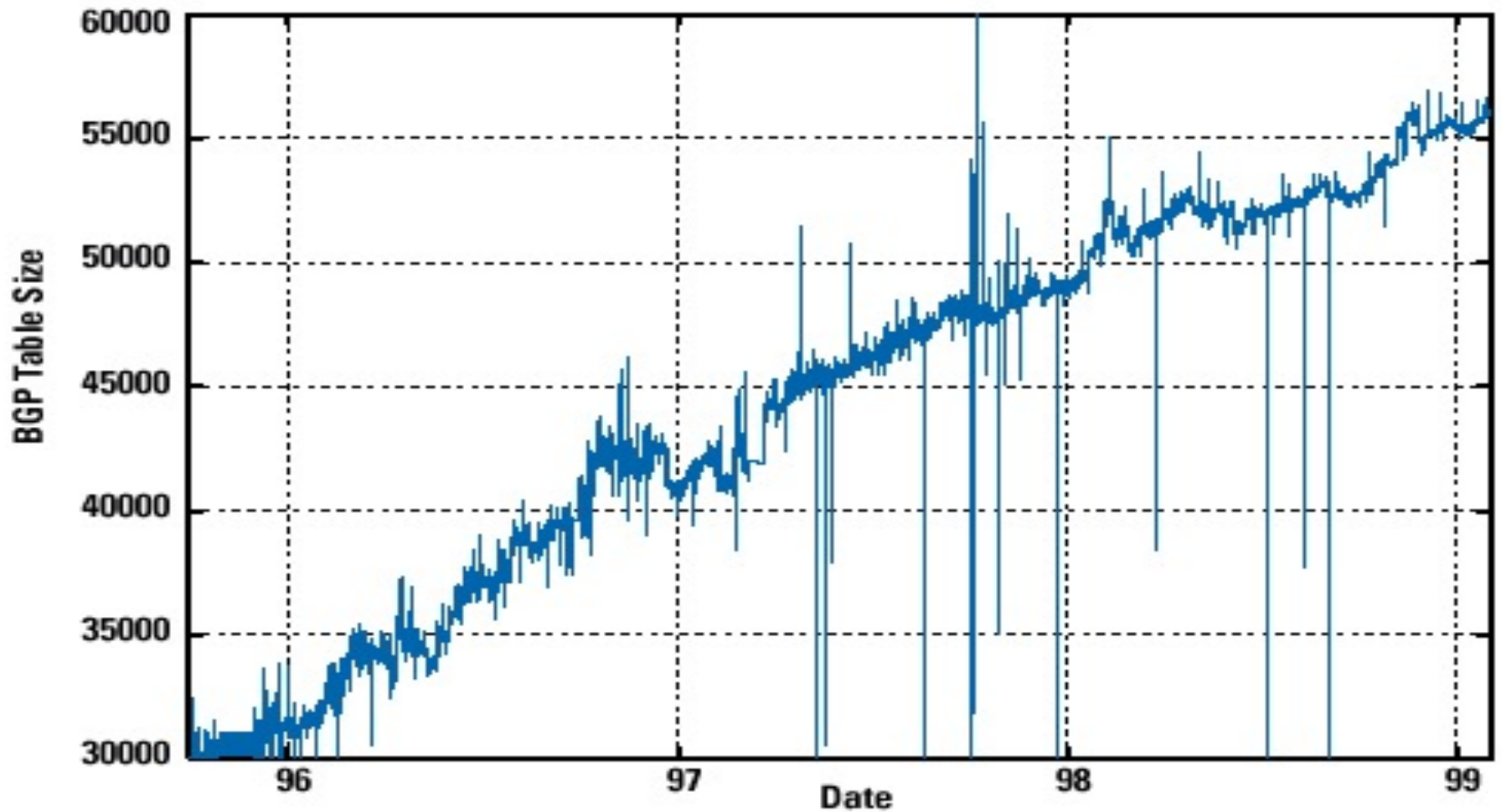
# CIDR (1994-1996): Much Flatter



Efforts to aggregate

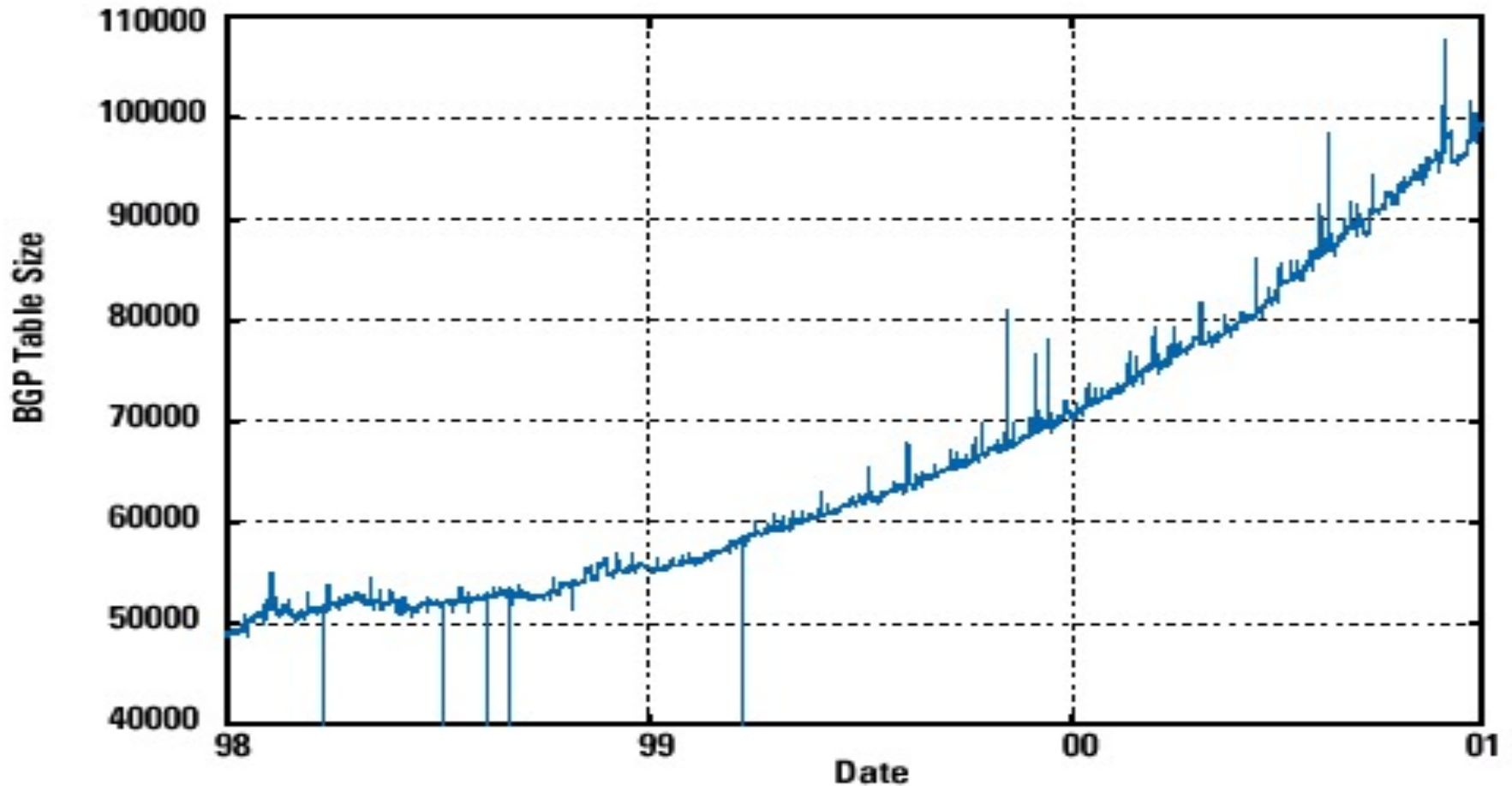


# CIDR Growth (1996-1998): Roughly Linear



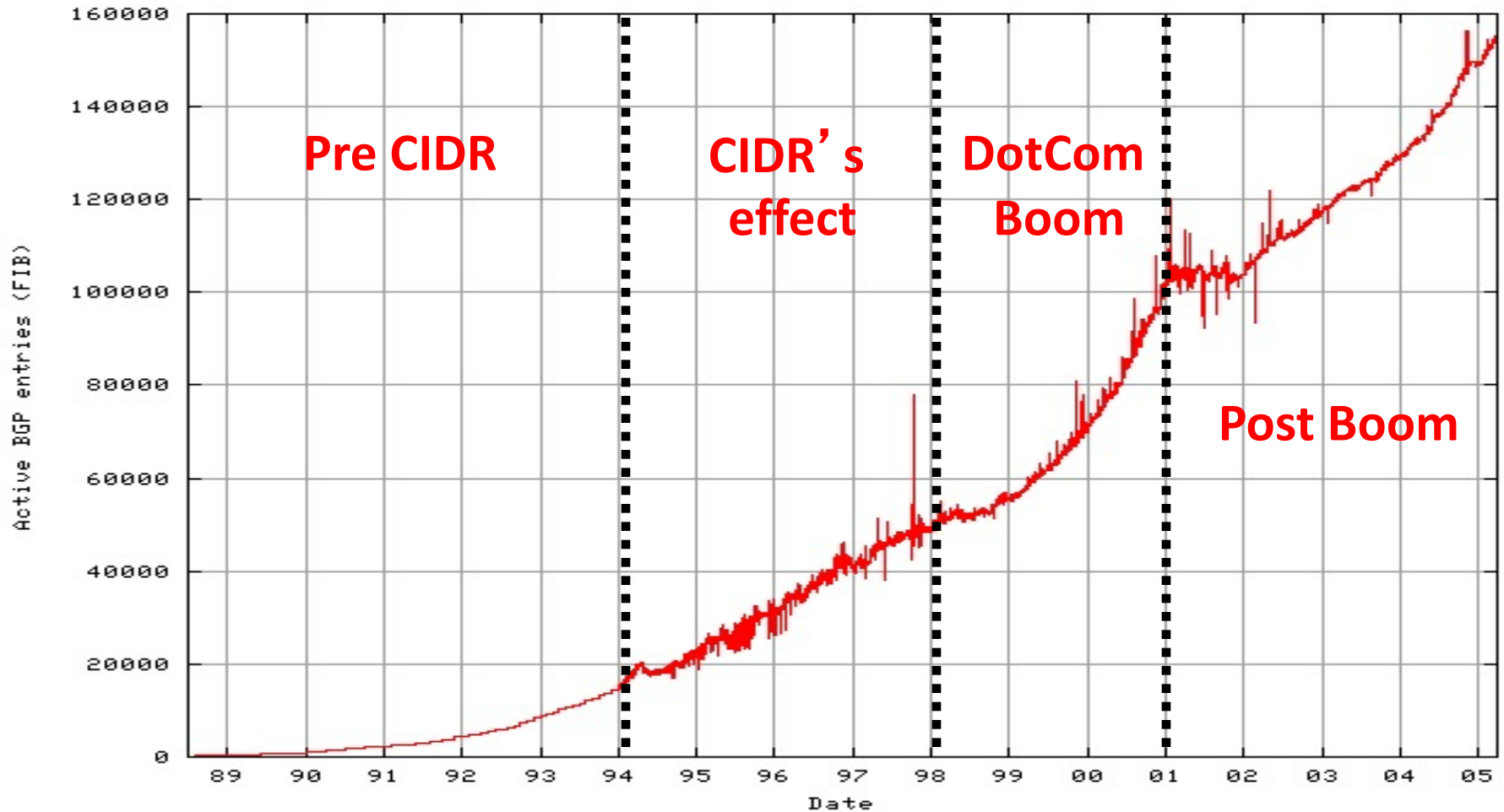
Good use of aggregation, and peer pressure!

# DotCom Boom (1998-2001): Steep Growth



Internet boom and increased multi-homing

# Long Term Growth (1989-2005)



Today we are up to ~400,000 prefixes