

Estimating Trending Topics on Twitter with Small Subsets of the Total Data

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1. INTRODUCTION

Since its founding in 2006, Twitter has grown into a corporation dominating a significant proportion of social media today. In recent years, it has become interesting to analyze tweets to determine interesting phenomena ranging from the viral spread of news to the detection of earthquakes [Burks2014]. The volume of existing tweets is prohibitively large for standard methods of analysis, and requires approaches that are able to either store large amounts of data in memory for fast access (including parallel approaches to data processing) or approaches that only require sublinear memory while retaining accuracy guarantees. In this paper, we investigate a problem utilizing the latter method.

1.1. Problem Statement

We would like to create an online algorithm that provides a real-time estimate of the frequencies of hashtags on Twitter time series data. Since we want to be able to run this algorithm in real time, we would like for it to be space-efficient: that is, it should not be required to store a large number of hashtag counts at any given time. Ideally, the algorithm's space complexity should be sublinear in the number of hashtags seen. Therefore, our estimate should not require knowledge of the exact frequency history of hashtags in tweets.

More specifically, we will attempt to approximate the k most popular Twitter hashtags in the time intervals of length y in order to tell what is trending during that time period. The idea is that estimating the top k hashtags in some time interval provides a model for the topics that are trending on Twitter in that time period.

1.2. Previous Work

Estimating frequencies with high probability and low space complexity has been an interest in algorithms development for a fair amount of time. The Count-Min sketch, developed in the early 2000s, provided an approximate data structure able to maintain counts of keys seen in an online setting in space sublinear to the number of keys, with the size of the structure instead dependent on parameters tunable for different accuracy and guarantee settings [Cormode2005].

More recently in 2012, Matuskevych et al. developed the Hokusai structure, which gives frequency counts for keys over a range of time [Matuskevych2012]. The rough idea behind this approach is that in the distant past, we should only care about heavy-hitters, i.e. hashtags with high frequencies in order to estimate the likelihood that the hashtag is trending again. The goal of the time-aggregated Hokusai system is to store older data at decreased precision since the older data also has decreased value in the computation. The time-aggregated Hokusai system works by storing aggregate data in Count-Min sketches each with a 2^i day resolution for the past 2^i days. Each of these Count-Min sketches computes d hash functions from $\{0, 1\}^*$ to $\{0, 1, \dots, m - 1\}$.

Analyzing what is trending on Twitter is a more recent research area. Twitter itself must give its users hashtags which it deems trending, to both spread the hashtag and to provide an awareness of what is going

on in the world to its users. Knowledge of what is trending and how that distribution looks over time is also useful for the business model, which relies on ads to a certain extent. Twitter itself stores all of its frequency data, and looks at recent spikes while conditioning on all of the past frequency data in order to determine estimates of trending likelihood. In 2012, MIT Professor Devavrat Shah built a machine-learning algorithm to predict which topics will trend on Twitter an hour and a half before they do, with 95% accuracy, though the datasets they tested their algorithm on were small[MITNews2012].

As a preliminary starting point for our approach to approximating the past frequencies, we will utilize concepts from the Hokusai paper to generalize the Count-Min Sketch scheme to time-series data [Matusevych2012].

1.3. Our Approaches

Our approach was to explore whether or not adding historical data would improve the accuracy of trending hashtag detection. To that end, we implemented two algorithms: one called the naive algorithm, the other called the history-sensitive algorithm. We wanted to test whether an algorithm that uses the history (in condensed form) to help understand what is trending at the current time could outperform an algorithm with a very short term memory.

The reason why the history is important is as follows: Some hashtags are constantly present. For example, in our experiments, the hashtags #porn, #gameinsight, #teamfollowback had nearly always had a high frequency in the current time unit. This fact is not surprising given the nature of the first hashtag. The second hashtag, #gameinsight, is the name of popular video game maker that people tweet when they win achievements in the game. Finally, #teamfollowback is a method Twitter users apply to gather Twitter followers quickly. The rule is if you follow that tag, you follow people who post that tag. Since these sorts of tags are consistently present, they should not be considered trending: they are not new. Therefore, we need to filter these out with the history. The necessity of maintaining approximate counts for a longer history depends on exactly how constant these consistently-appearing tweets actually are.

We had two different algorithms for finding out which hashtags were trending in real time, both space efficient. The naive algorithm does not take into account the history and simply looks at a y -unit before the current y -unit interval to find the current trending hashtags. To determine whether a hashtag was trending, we tested

$$\frac{\text{freq}(\text{hashtag}, \text{present } y\text{-unit})}{\text{freq}(\text{hashtag}, \text{past } y\text{-unit})} > c \quad (1)$$

where c is a threshold we picked to determine whether the hashtag was considered trending. We believe that in order to accurately determine trending hashtags, we will need to use more of the past in order to eliminate the hashtags which are ever-present, and therefore not trending.

The second approach differs in that we do not store all the data of the past, but instead use several Count-Min Sketches to approximate the past frequencies. We call this approach the history-sensitive algorithm. The idea is to remember the past with less accuracy to save space, while weighting the frequency counts of a hashtag in the past appropriately. We used the rule of thumb that the more often a hashtag shows up in the past, the less likely it is trending. However, we also discount the importance of the past by weighting the times long ago less and less as time moves forward.

2. DESCRIBING THE HISTORY-SENSITIVE ALGORITHM AND DATA STRUCTURE

We separate the problem of finding trending topics on Twitter into two parts. First, we need to maintain a data structure that efficiently stores data about all occurrences of every hashtag seen in the past. We also maintain a separate data structure that allows us to quickly gather information about the most recent hashtags seen.

We want the former data structure to be very space efficient since it must store data about a very large dataset. For this structure, space efficiency is more important than accuracy since small deviations in such a

large dataset should not be significant because the deviations in past data should not greatly affect what is currently trending.

For the latter data structure, accuracy is more important than space efficiency since the structure contains data which more closely relates to which topics are currently trending and the size of the dataset is much smaller.

2.1. Data Structures

2.1.1 History Data Structure

To store data about all occurrences of every hashtag seen in the past, we use a modified version of the time-aggregated Hokusai system [Matusevych2012], which is an extension of the Count-Min sketch. We previously described this data structure in [Previous Work](#). To the Hokusai structure we add another Count-Min sketch that combines the information from the Hokusai Count-Min sketches. We call this external Count-Min sketch the Kernel, since it acts as a weighting function on the CM sketches in the Hokusai structure. Its role is to depreciate the value of older data. Denoting the CM sketches of the Hokusai structure as a vector $\mathbf{M} = \{\bar{M}, M^0, M^1, \dots, M^{\lfloor \log(T) \rfloor}\}$, where \bar{M} is the Count-Min sketch storing the data of a unit time step (of size z), M^j is the Count-Min sketch storing the data for the sketch with a 2^j time resolution, and T is an upper bound on the number of z -units stored by the data structure. Then, the kernel function is given by

$$k(\mathbf{M}) = \bar{M} + \sum_{j=0}^{\lfloor \log(T) \rfloor} \frac{M^j}{2^j} \quad (2)$$

For each hashtag, the kernel sketch stores a positive value between 0 and 2 (though typically ≤ 1) that approximates how often the hashtag showed up in the past. In [Correctness](#) we will show that this aggregate Count-Min sketch weights the hashtags that occurred $i > 0$ days ago with weight approximately $\frac{1}{i}$.

We will refer to the combined data structure of Hokusai and the Kernel as the History.

2.1.2 Current-Window Data Structure

Our data structure for holding hashtags seen in the last y -unit consists of three components: a max Fibonacci heap, a queue, and a hash table. We refer to these components collectively as the Current-Window. The keys for the hash table are the hashtags, and the values stored in the table are (frequency, pointer to corresponding node in heap) pairs. The hash table exists so that we can keep track of exact frequency counts for each y -unit. The queue contains (hashtag, timestamp) pairs, each of which is inserted upon seeing a hashtag in the input stream. The purpose of the queue is to determine when a hashtag count is too old to be maintained in the present y -unit window. Finally, we have a Fibonacci heap to prioritize the hashtags based on how trending they currently are. We provide a heap priority function, which defines what trending means for the whole data structure. For a specific hashtag, we define the priority function to be

$$\text{priority}(\text{hashtag}) = \frac{\text{freq}(\text{hashtag}, \text{present } y\text{-unit})}{\text{History}[\text{hashtag}]} \quad (3)$$

The idea is that again, we want to be inversely proportional to the history, here corrected by more knowledge and clever approximation schemes to save space.

2.2. Algorithm Pseudocode

2.2.1 Updating History Data Structure

Algorithm 1 describes the necessary steps to maintain the History data structure as new input is provided. We perform time aggregation to fill the Kernel, K . This algorithm is only called when a z -unit has passed since the last time the algorithm was called.

Algorithm 1 Update History

```
1: for all  $i$  do
2:   Initialize Count-Min sketch  $M^i = 0$ 
3: Initialize  $t = 0$ 
4: Initialize Count-Min sketches  $\overline{M} = 0$  and  $K = 0$ 
5: while data arrives do
6:   Aggregate data into sketch  $\overline{M}$  for current  $z$ -unit while also adding this data to  $K$ 
7:    $t \leftarrow t + 1$  (increment counter)
8:    $K \leftarrow K - \overline{M}$ 
9:   for  $j = 0$  to  $\text{argmax } \{l \text{ where } t \bmod 2^l = 0\}$  do
10:     $K \leftarrow K + 2^{-j}(\overline{M} - M^j)$ 
11:     $T \leftarrow \overline{M}$  (back up temporary storage)
12:     $\overline{M} \leftarrow \overline{M} + M^j$  (increment cumulative sum)
13:     $M^j \leftarrow T$  (new value for  $M^j$ )
14:    $\overline{M} \leftarrow 0$  (reset aggregator)
```

2.2.2 Updating heap and hash tables

We run Algorithm 2 in the main loop, every time a single piece of data arrives. For each new piece of data, we also check whether the data we are storing has grown outdated – that is, older than a y -unit before the most current piece of data arrived.

2.2.3 Finding the trending hashtags

We use Algorithm 3 to calculate the top k hashtags by querying the heap. The heap is the structure responsible for maintaining the priority of the various hashtags, and provides an efficient way of implementing a changing priority queue.

3. ANALYZING THE HISTORY-SENSITIVE ALGORITHM

Symbol	Meaning
d	Number of hash tables for each Count-Min sketch
m	Size of each hash table in each Count-Min sketch
s	Number of distinct hashtags in the Current-Window
z	Time resolution of the unit-sized Count-Min sketches
T	Upper bound on the number of z -units of data stored in History
x	Total number of hashtags in the Current-Window
y	Time interval of data contained in the Current-Window

Table 1: Variables referenced in this section

3.1. Correctness

This algorithm finds the k hashtags that have the maximum value of (frequency in last y -unit) / (value in History data structure).

Algorithm 2 Update Current-Window

```
1: while data arrives do
2:   if the current z-unit is different than that of the last hashtag seen then
3:     Do all the end-of-z aggregation for the History structure as detailed in Algorithm 1.
4:   for all elements in the Fibonacci heap do
5:     look up the hashtag corresponding to this node in the hash table
6:     Update the key of the node to:
        
$$\frac{\text{frequency in last } y - \text{length time period found at the table entry}}{\text{new value in History data structure}}$$

7:   if queue is not empty then
8:     Peek at end of queue.
9:     if the timestamp +  $y$  is before the current time then
10:      Look up the hashtag in the hash table and decrement the stored frequency.
11:      if the frequency is now 0 then
12:        Delete the node in the heap pointed to by this entry in the table.
13:        Delete this entry in the hash table.
14:      else
15:        Update the key of this node pointed to by this entry in the table to the proper
        value given the new frequency.
16:   if hashtag is in hash table then
17:     Increment the frequency stored at that entry.
18:     Update the key of the node in the Fibonacci heap.
19:   else
20:     Insert a new node into the Fibonacci heap with the appropriate key and value.
21:     Insert the hashtag into the hash table with a pointer to this node.
```

Algorithm 3 Top k trending hashtags

```
1: Perform  $k$  delete-max operations on the Fibonacci heap, storing each of the nodes deleted in a
   list  $L$ .
2: for all nodes  $N$  in  $L$  do
3:   Announce that the hashtag associated with  $N$  is trending.
4:   Insert  $N$  into the Fibonacci heap.
```

Claim 1. *The value for hashtag x in the aggregate Count-Min sketch of the History data structure is within a factor of 4 of $\overline{M}(x) + \sum_{i=1}^T \frac{1}{i} * (\text{value for } x \text{ in a Count-Min sketch using the same hash functions for all hashtags occurring } i \text{ days ago})$.*

Proof. First, we use Theorem 4 in [Matusevych2012] which states that “At t , the sketch M^j contains statistics for the period $[t - \delta, t - \delta - 2^j]$ where $\delta = t \bmod 2^j$.”

Let b be the location in the aggregate Count-Min sketch containing the value returned when x is queried.

Let h be any instance of any hashtag that appeared on day p such that seeing h incremented counters in position b .

Case 1: $p = t$

Then the statistics for h are recorded in \overline{M} and are not in any M^j .

Case 2: $2^i > t - p \geq 2^{i-1}$ for some $i > 0$

By Theorem 4, for all $j \leq i - 2$, M^j does not contain statistics for p since $p \leq t - 2^{i-1} \leq t - \delta - 2^{i-2}$.

Therefore, the increments that occurred in the Count-Min sketch for hashtags occurring i days ago contribute at most $\sum_{j=i-1}^T 2^{-j} < 2^{2-i}$ to the value in position b .

Let q be the largest j such that $t - \delta - 2^j \geq p$

Then $p \leq t - \delta - 2^q$. Let $\lambda = t \bmod 2^{q+1}$. Then $\lambda = \delta$ or $\lambda = \delta + 2^q$.

Since q is the largest j such that $t - \delta - 2^j \geq p$, $p > t - \lambda - 2^{q+1}$.

Also, $p \leq t - \delta - 2^q \leq t - \lambda$, so M^{q+1} contains statistics about h .

For all $j \geq i$, $t - \delta - 2^j < t - 2^j < p$, so $q \leq i - 1$.

Thus, incrementing the counter in M^{q+1} contributed at least 2^{-i} to the value in position b .

Thus, the contributions to the sum are within a factor of 4 of $\frac{1}{t-p}$.

Therefore, summing over all hashtags that increment counters in position b gives $\overline{M}(x)$ for all hashtags that occurred on day t , and within a factor of 4 of $\sum_{i=1}^T \frac{1}{i} * (\text{value for } x \text{ in a Count-Min sketch using the same hash functions for all hashtags occurring } i > 0 \text{ days ago})$. \square

This value is approximately (frequency in last y -unit) / (freq today + $\sum_{i=1}^T \frac{1}{i} * (\text{frequency of hashtag } i \text{ z-units ago})$). This seems to be a desirable function to maximize since it finds hashtags that are common in the last y -unit that have been comparatively infrequent in the past. This function is good since it especially emphasizes hashtags that are different than those seen in the past few days. This ensures that the same items do not stay trending for too long.

3.2. Runtime Analysis

Processing the insertion of a hashtag takes the following time. It takes amortized $O(d)$ time to update the History. It takes expected $O(1)$ time to check if in hash table.

If it is, it requires $O(1)$ time to increment the frequency, $O(d)$ time to compute the new key, and $O(1)$ amortized time to update the key since it will be nondecreasing.

Otherwise, it requires $O(d)$ time to compute the key value, $O(1)$ time to insert the new node in the heap, and $O(1)$ time to insert into the hash table.

Thus, our algorithm takes $O(d)$ amortized time + expected $O(1)$ time to process a new hashtag.

Processing the removal of a hashtag from the Current-Window takes the following time. It takes $O(1)$ time to verify that the queue is not empty. It takes $O(1)$ time to look at the end of the queue and verify that the timestamp + y is before the current time. It takes $O(1)$ time in expectation to look up this hashtag and decrement its frequency. Then, it takes $O(1)$ time to check if the frequency is 0.

If so, it takes $O(\log(s))$ amortized time to delete the node in the heap and $O(1)$ time to delete the entry in the hash table.

Otherwise, it takes $O(d)$ amortized time to compute the new key for the hash table and $O(\log(s))$ amortized time to update the heap given this key.

Case	Amortized Time
Processing the insertion of a new hashtag	$O(d)$
Processing the removal of a hashtag from the Current-Window	$O(d + \log(s))$
Updating History and Current-Window at the end of a z-unit	$O(md + ds + s \log(s))$
Querying for the top k trending items	$O(k \log(s))$

Table 2: Time analysis summary

Thus, our algorithm requires $O(\log(s))$ amortized time + expected $O(1)$ time + $O(d)$ time to remove a hashtag from the Current-Window.

The end-of-day updates to the History and the resulting updates to the heap take the following time. By Lemma 5 of the Hokusai paper, the amortized time required to do all end-of-day aggregation is $O(md)$. Then, for each of the s nodes in the heap, it takes $O(d)$ time to compute each updated key and $O(\log(s))$ amortized time to update the heap given the new key.

Thus, it takes $O(md + s \log(s))$ amortized time + $O(ds)$ time to do all necessary updates at the end of the day.

Querying the data structure for the top k trending items takes $O(k \log(s))$ amortized time for delete-max operations, $O(k)$ time to announce that these items are trending, and $O(k)$ amortized time to reinsert these nodes into the heap.

Thus, it takes $O(k \log(s))$ amortized time to determine what's trending.

3.3. Spatial Analysis

The History requires $O(md)$ space for each Count-Min sketch, so it requires a total of $O(md \log(T))$ space.

Each node in the heap requires a constant amount of space, so the heap requires $O(s)$ space.

The hash table always contains at most s entries with each entry requiring a constant amount of space. Also, in order to maintain an expected $O(1)$ lookup time, the hash table needs to have $O(s)$ bins. Thus, the hash table requires $O(s)$ space.

The queue requires an entry for every hashtag still in the current window, so it requires $O(x)$ space.

Thus, everything requires $O(md \log(T) + x)$ space since $s < x$.

4. DESIGN CHOICES

4.1. Choosing the parameters y and z

For y , the unit time period we attempted to find the trending hashtags for, we chose a 3-hour interval. For z , the unit time period for the History data structure, we chose day-length intervals. These values made sense because the testing data we used occurred over a 30-day interval, and we wanted to avoid too fine a temporal resolution so that we could run our code on the data in reasonable time.

4.2. Parameters of the History Data Structure

pick some parameters here $n = 6$, $m = 3500$ (size of hash tables), $d = 20$ (number of tables) n : we do it for a month, so $\sim 2^{(6-1)}$ days is definitely enough $m = \lceil e/\epsilon \rceil$, $d = \lceil \ln(1/\delta) \rceil$ error in query is in factor ϵ with probability $1 - \delta$ for $d = 20$, we have with high probability $(1 - 1/e^{20})$ we have error factor of $e/3500$ of the total number of hashtags total number of hashtags is ~ 15 million thus error bounds are in range ~ 12000 on a given count this error provides a 'good' amount of smoothing

4.3. Choosing the Flavor of Heap

We chose a Fibonacci heap in the hopes that there would be a larger number of **decrease-key** operations, which has a better theoretical bound in Fibonacci heaps.

5. TESTING PROCEDURE

We downloaded the [September 2014 archive of Twitter data](#) to serve as test data for our algorithms [Twitter2014]. We then parsed the data into (timestamp, hashtag) format for all thirty days worth of data. Using these files as pseudo-real-time input, we used the last week of the data as a benchmark for comparison between the naive and the history-sensitive algorithms. We divided the data into three-hour chunks, and queried for the top ten hashtags in each of the three-hour time blocks in the last week of the data using both algorithms. Note that the history-sensitive algorithm was required to parse through the first three weeks as well in order to build its history, whereas the naive algorithm was only required to pay attention to six-hour chunks at a time, and thus never needed to look a bit beyond the scope of that week.

We used the Jaccard similarity to arrive at a similarity score between the two methods. We also investigated the hashtags that both algorithms came up with for a given 3-hour interval, and verified that they were in fact trending at that time using news and Twitter media.

6. RESULTS

6.1. Performance Measurements

We use the space complexities of the data structures and the sizes of the respective inputs to determine approximately how much space is used by each approach, and compare them.

6.1.1 Naive Algorithm

Space measurements (running on one week's worth of hashtags) (estimate 6 hours worth of hashtags)

Time (running over one week): ~ 15 minutes

6.1.2 History-Sensitive Algorithm

(we had to build a history over the first three weeks of the month in order to determine the History data structure's values for the week we ran the algorithm over)

Space measurements (running on one week's worth of hashtags):

Time (running over one week): 2 – 3 hours

6.2. Are the Algorithm Outputs Actually Trending?

6.2.1 The Intersection Distribution

First, we see how many hashtags were marked as trending by both the naive and history-sensitive algorithms in the same time slots. For each 3-hour interval, we calculate the Jaccard Similarity, a value in the range $[0, 1]$, given by

$$J(A, B) = \frac{|A \cap B|}{|A \cup B|} \quad (4)$$

where A and B are the sets of top 10 hashtags for a given 3-hour interval for each algorithm. We then calculate the mean and standard deviation of the Jaccard similarity over the 56 different 3-hour intervals of the last week of September 2014.

This value gives us a metric for understanding how similarly the two algorithms performed.

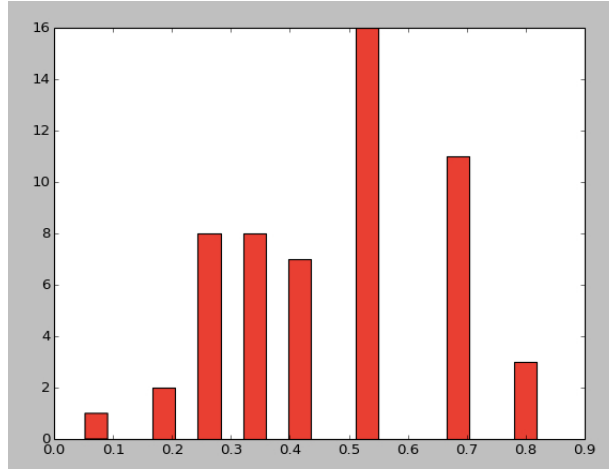


Figure 1: *Distribution of Jaccard Similarities*

We plot it below in the histogram shown in [Figure 1](#). The x -axis is the Jaccard similarity, and the y -axis is the count. The mean was 0.47, and the standard Deviation was 0.18. From the plot, we can see that the distribution of the Jaccard similarities is roughly a centered normal distribution around 0.5. A Jaccard similarity of 0.5 means about half the hashtags are the same in each 3-hour time interval between the two algorithms.

6.2.2 Checking the Algorithm-Defined Trends Against Real Life

In [Figure 2](#), we have superimposed pictures from news articles that the tweets describe on a time axis, as the hashtag referencing the article starts trending. The Jaccard similarities for each of these three-hour intervals were all above 0.5, so both algorithms found the hashtags that we graph in the figure. We checked that the news was released (via Yahoo! News, Twitter, Google News, ESPN, and others) at around the time our algorithms say the hashtag relating to the news begins to trend. A fair number of the trending hashtags related to sports and TV shows.

As we can see in the figure, our algorithm detected the start of the TV show "How to Get Away With Murder", which was released on September 25, 2014. Another TV-related hashtag was *#dailyshowgonetoo far*, when Jon Stewart called One Direction member Zayn Malik a terrorist (on the 26th) and caused a huge outrage. Sports events that occur real-time on TV also have the trending profile. There were three in this section of six three-hour intervals. First, Derek Jeter hit a walk-off homerun in the final time at bat on the Yankees, and naturally on Twitter, fans were excited. This event occurred on the 25th. The Ryder Cup is a major golf tournament that began on the 26th and lasted through the 28th, and its beginning launched a wave of tweets. The last sports-related hashtag was the beginning of the ToyotaEMR car race that took place in India on September 26th. Finally, *#blowthewhistle* related to a kids cancer campaign that had lasted throughout September, which perhaps got a boost as September ended to raise some last minute funds.

7. DISCUSSION

7.1. Comparison of Naive and History-Sensitive Algorithms

observed that history-sensitive did about the same in terms of getting trending tweets (based on Jaccard), but naive uses much less space.

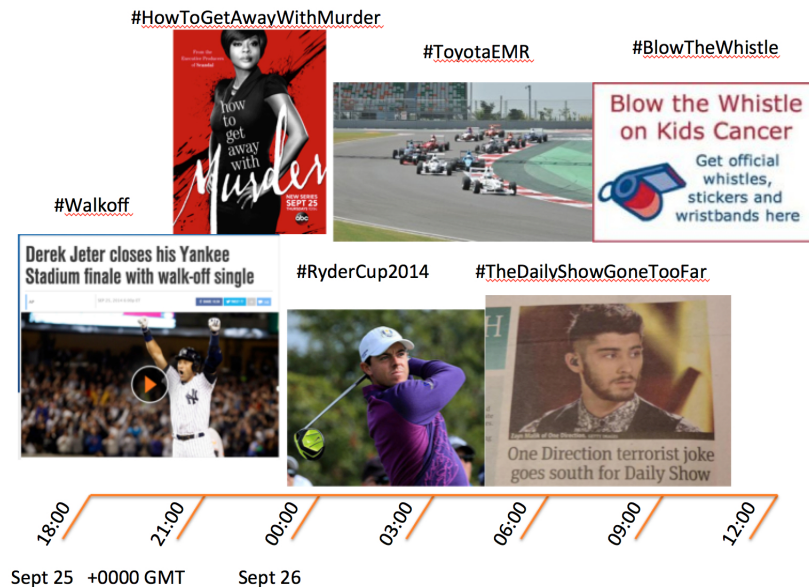


Figure 2: News Visualization of Trending Tweets

7.2. Applicability to Real-Time Data

do well to get month's worth of data in a few hours (2-3)

implementation wouldnt run in python, do it in C or something

for real data from firehose (all of it), we would use better infrastructure dedicated to processing all the data would use parallel setup (this scheme is adaptable to parallel setup)

seems reasonable if we have proper data storage setup.

8. FUTURE WORK

With respect to the algorithms specifically, in the future we would want to adapt our approach to the history-sensitive algorithm and play with the parameters a bit more. We would also like to try different kernel functions to see how they performed. Another approach to utilizing the history information would be instead of weighting by inverse exponentially weighted frequency in the past for a given hashtag, attempt to calculate some score of periodicity for the hashtag. If the hashtag has low periodicity over a given time scale, then even if it appears multiple times in the past (because the past is long), our algorithm will avoid discarding it as non-trending.

Then, we would like to try adapting our algorithms to real-time data: that is, hook up our algorithms to a significant portion of the Twitter firehose, after setting up the necessary hardware to process that amount of data. In order to aid with the processing, we would also ideally use a distributed system and parallelize our code for better performance.

We also know that an estimation the top k hashtags in a given time interval could provide a topic model for tweets. It could be interesting to use the top k hashtags to build a time-sensitive topic model and analyze quantities including the speed of topic change, the variety of topics, and so on.

Even though not perhaps irrelevant to the specific task of identifying trending hashtags, the approach of the history-sensitive algorithm may be useful in other domains. Future work could also investigate the performance of the history-sensitive algorithm on other data streams of interest, for instance, Wikipedia edits – our goal would be to estimate which Wikipedia topics are being edited the most at any given time interval of some length. Another interesting application of this algorithm to test would be the real-time estimation

of the hottest selling stocks on Wall Street to see if either of these settings results in a significant improvement over the naive algorithm.

9. APPENDIX I: CODE AND VISUALIZATIONS

We provide a link to [all our code](#), as well as to an online hosting of the [frequency visualization](#).

In our code, we made use of two libraries written by other people, which we modified. One of these was a Count-Min sketch implementation in Python, and another was a Fibonacci heap implementation, also in Python. These program files are in our codebase, with appropriate citations at the top of each document. Furthermore, we used multiple standard Python libraries in our code.

10. APPENDIX II: TOP- k HASHTAGS IN 6 INTERESTING TIME CHUNKS

To see all 56 3-hour intervals for both algorithms, see the results folder of our [Github](#). We display the top 10 hashtags for each 3-hour interval for both the naive algorithm and the history-sensitive algorithm, along with their heap priorities. Note that the heap priorities span a reasonable range from 0 to 1, and are not clustered around 1 or 0, implying the algorithms are actually differentiating the hashtags as trending and not trending.

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