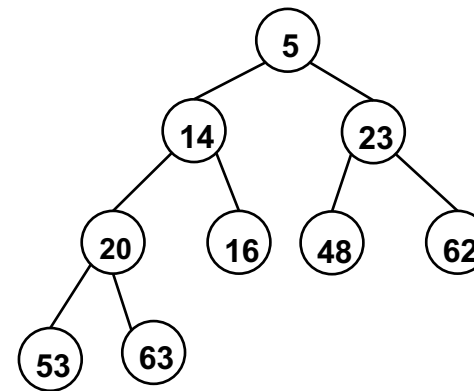


DS: **Binary Heap**

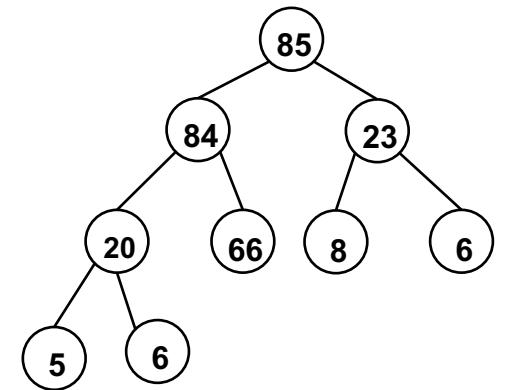
Liwei

What is Binary Heap

- Binary Heap is a special Binary Tree
 - Heap property
 - Value of any given node must be \leq value of its children (Min-Heap)
 - Value of any given node must be \geq value of its children (Max-Heap)
 - Complete Tree
 - All levels are completely filled except possibly the last level and the last level has all keys as left as possible
 - This makes Binary Heap ideal candidate for Array Implementation



Min-Heap



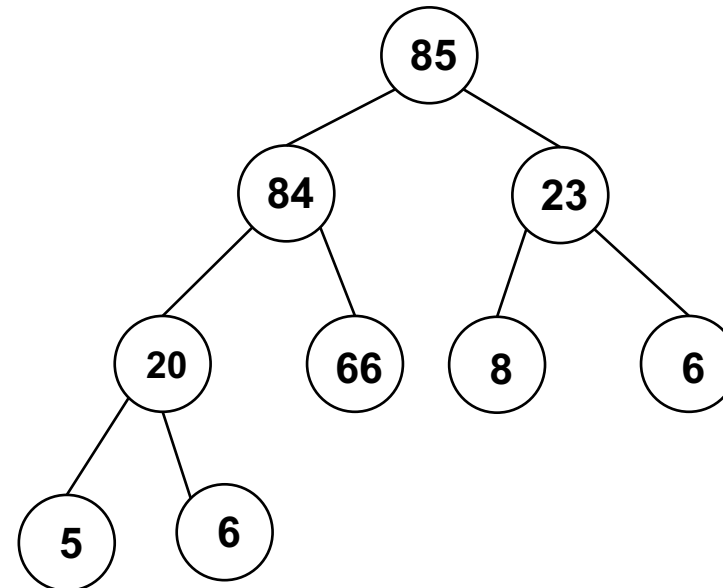
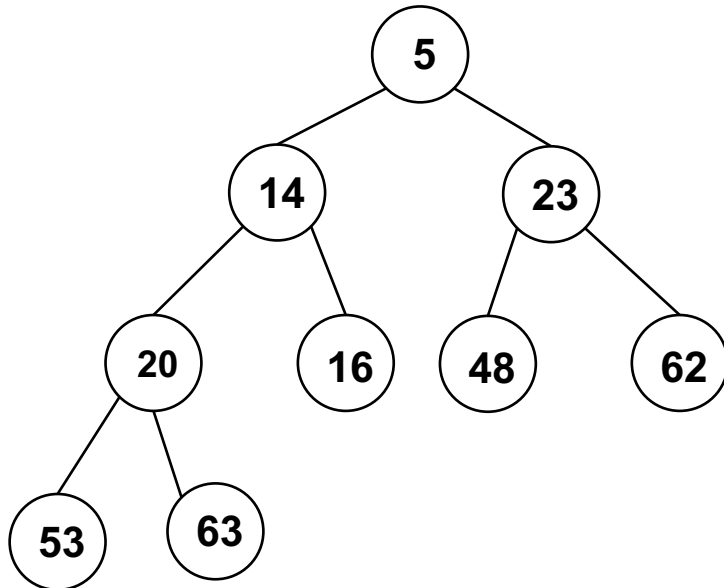
Max-Heap

Why should we learn Binary Heap

- There are cases when we want to find min/max number among set of numbers in less than $O(\log n)$ time. Also, we want to make sure that inserting additional numbers dose not take more than $O(\log n)$ time.
- Possible solutions:
 - Store the numbers in sorted array
 - Issue here is that once we insert/delete a new number, our array needs to be adjusted again to keep it sorted which will take $O(n)$ time
 - Store the numbers in linked list in sorted manner
 - Search takes $O(n)$
- Practical use
 - Prime's Algorithm, Heap Sort, Priority Queue

Types of Binary Heap

- Min-Heap: if the value of each node is less than or equal to value of both of its children
- Max-Heap: if the value of each node is more than or equal to value of both of its children



Common operations

- createHeap – creates a blank array to be used for storing heap
- peekTopOfHeap – returns min/max from Heap
- extractMin / extractMax – extracts Min/Max from Heap.
- sizeOfHeap – returns the size of the Heap
- insertValueInHeap – inserts a value in Heap
- deleteHeap – Deletes the entire Heap

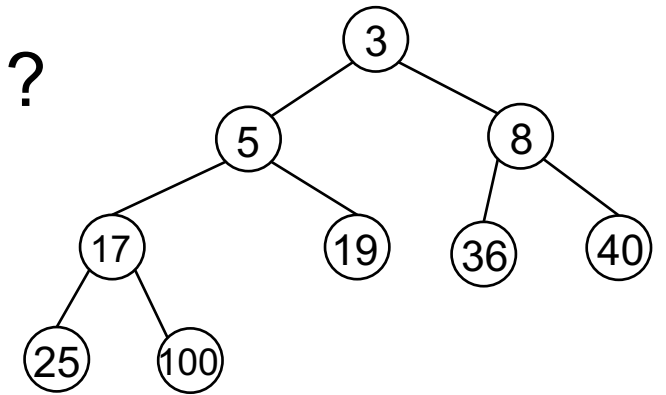
We can extract only this node.

Implementation Options

- Array based implementation
- Reference / Pointer based implementation

Binary Heap – Array Representation

- How does Binary Heap looks like at logical level?



- How does Binary Heap looks when implemented via Array

cell #	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17
value	X	3	5	8	17	19	36	40	25	100								

Left Child – $\text{cell}[2x]$

Right Child – $\text{cell}[2x + 1]$

Creation of Heap

`createHeap(size)`

Create a blank array of “size+1”

Initialize `sizeOfHeap` with 0

cell #	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17
value	X																	

Time complexity: $O(1)$

Peak of Heap

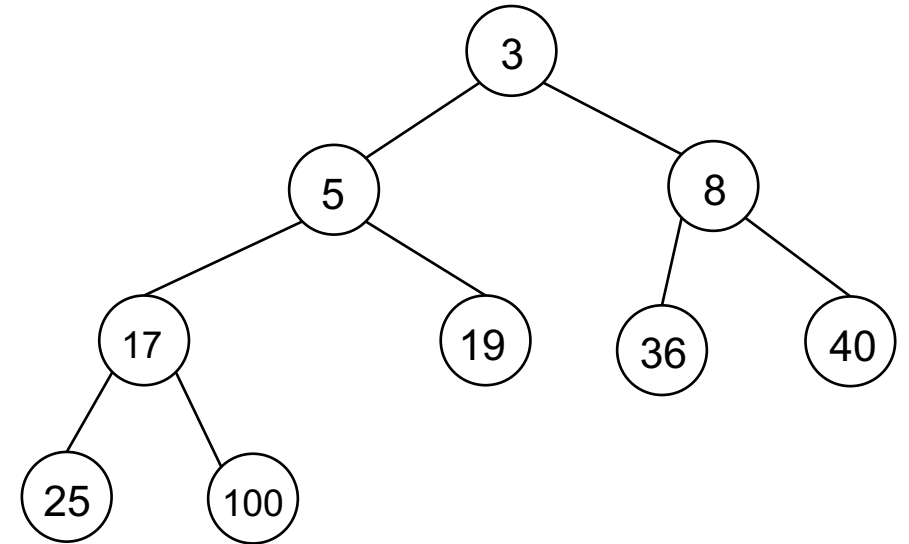
peakTopOfHeap()

If tree does not exist

return error message

else

1st cell of the array

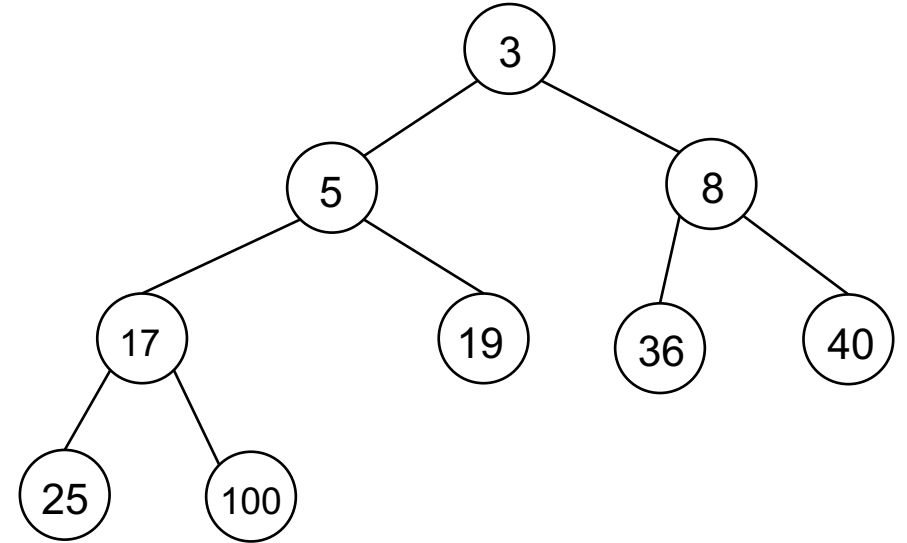


cell #	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17
value	X	3	5	8	17	19	36	40	25	100								

Time complexity: $O(1)$

Size of Heap

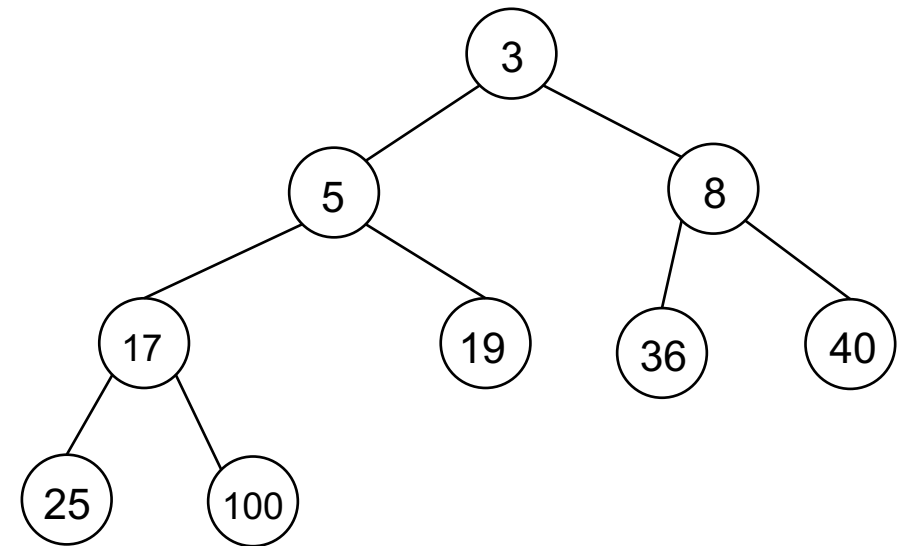
sizeofHeap()
return sizeofHeap



cell #	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17
value	X	3	5	8	17	19	36	40	25	100								

Time complexity: $O(1)$

Insertion in Heap



cell #	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17
value	X	3	5	8	17	19	36	40	25	100								

Insertion in Heap

insertValueInHeap (value)

If tree does not exist

return error message

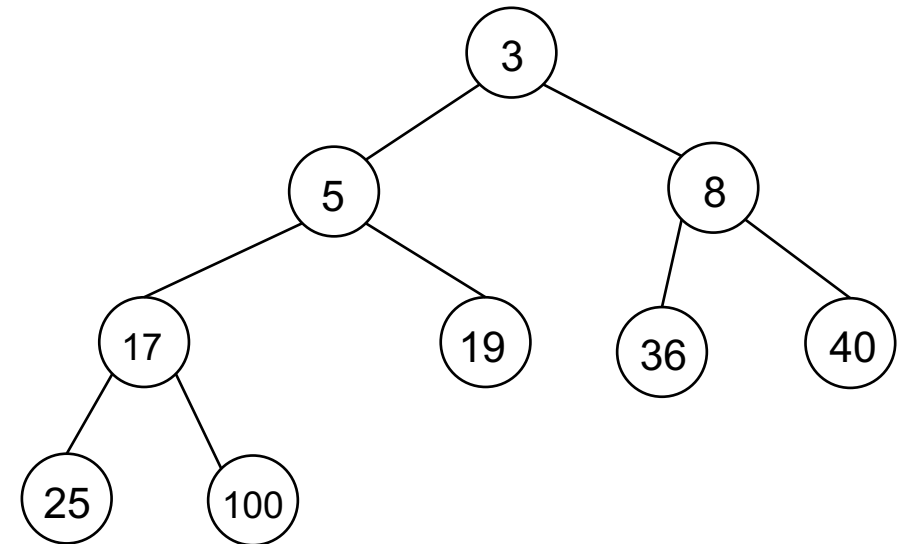
else

insert value in first unused cell of array

sizeOfHeap ++

heapifyBottomToTop(sizeOfHeap)

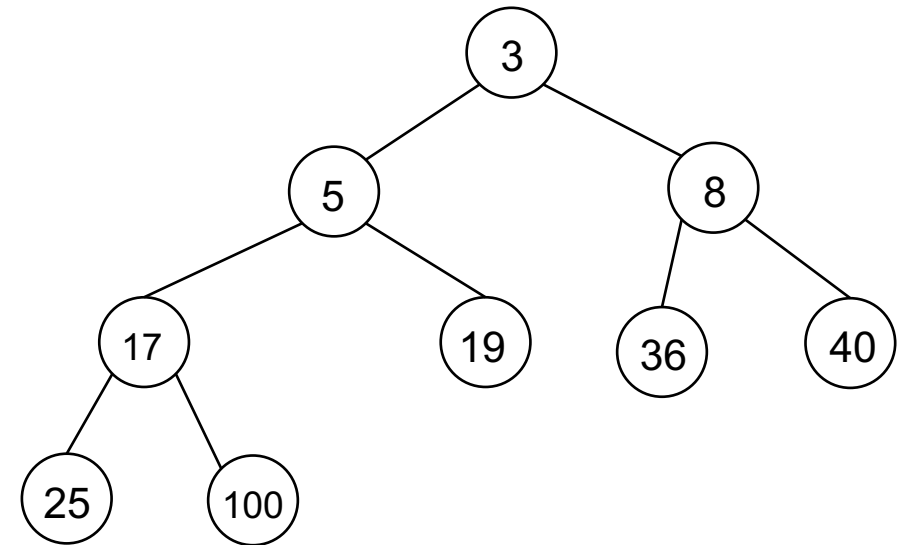
$O(\log n)$



cell #	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17
value	X	3	5	8	17	19	36	40	25	100								

Time complexity: $O(\log n)$

ExtractMin from Heap



cell #	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17
value	X	3	5	8	17	19	36	40	25	100								

ExtractMin from Heap

extractMin ()

 If tree does not exist

 return error message

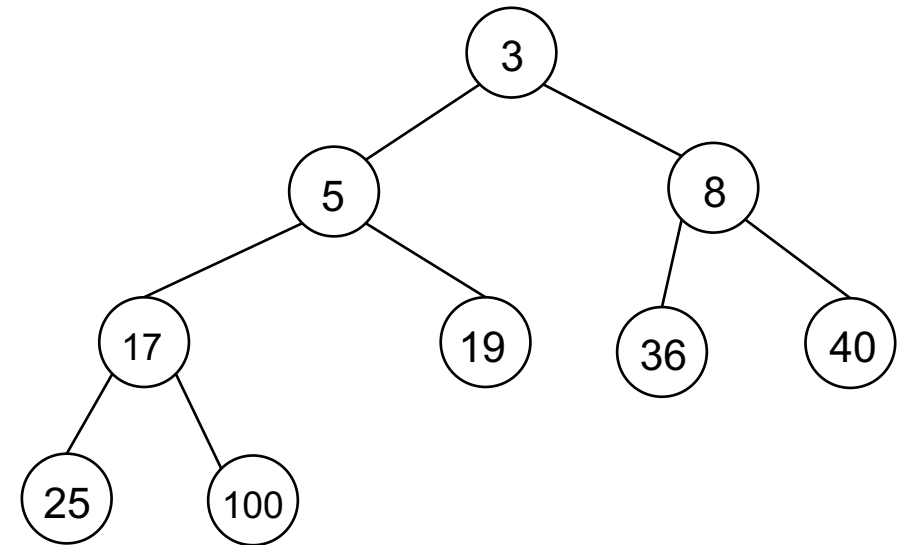
 else

 extract 1st cell of array

 promote last element to first

 sizeOfHeap --

 heapifyTopToBottom (1) $O(\log n)$



cell #	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17
value	X	3	5	8	17	19	36	40	25	100								

Time complexity: $O(\log n)$

Delete Heap

`deleteHeap ()`

Free the array from memory

$O(1)$

Time complexity: $O(n \log n)$

Time and Space Complexity (Binary Heap)

	Time Complexity
CreateHeap	$O(1)$
PeekTopOfHeap	$O(1)$
sizeOfHeap	$O(1)$
insertValueInHeap	$O(\log n)$
extractMin / extractMax	$O(\log n)$
deleteHeap	$O(1)$

Binary Heap

Reference-based Implementation

Common operations

- createHeap – creates a blank array to be used for storing heap
- peekTopOfHeap – returns min/max from Heap
- extractMin / extractMax – extracts Min/Max from Heap.
- sizeOfHeap – returns the size of the Heap
- insertValueInHeap – inserts a value in Heap
- deleteHeap – Deletes the entire Heap

ExtractMin from Heap

extractMin ()

 If tree does not exist

 return error message

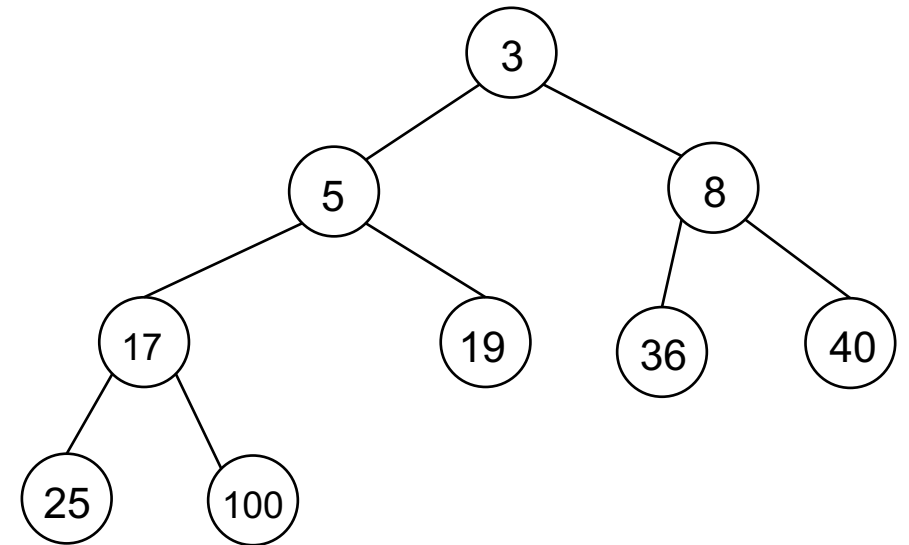
 else

 extract 1st cell of array

 promote last element to first

 sizeOfHeap --

 heapifyTopToBottom (1) → O(n)



cell #	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17
value	X	3	5	8	17	19	36	40	25	100								