

# Unit 2: Logic and Proofs

Originals slides by Dr. Baek and Dr. Still, adapted by J. Stelovsky

Based on slides Dr. M. P. Frank and Dr. J.L. Gross

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# Propositional Logic

What is a proposition?

- ▶ It is raining
- ▶  $1 + 1 = 2$
- ▶ every odd number is a prime
- ▶  $2^{67} - 1$  is a prime
- ▶  $(n + 1)(n - 1) = (n^2 - 1)$  for any integer  $n$

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What is common between these statements?

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A **proposition** is a statement that is either true or false (but not both).

**Definition:** A *proposition* (denoted  $p, q, r, \dots$ ) is simply:

- a *statement* (i.e., a declarative sentence)
  - *with some definite meaning*,  
(not vague or ambiguous)
- having a *truth value* that's either *true* (**T**) or *false* (**F**)
  - it is **never** both, neither, or somewhere “in between!”
    - However, you might not *know* the actual truth value,
    - and, the truth value might *depend* on the situation or context.
- Later, we will study *probability theory*, in which we assign *degrees of certainty* (“between” **T** and **F**) to propositions.
  - But for now: think True/False only! (or in terms of **1** and **0**)

# Examples of Propositions

- It is raining. (In a given situation)
- Beijing is the capital of China. (T)
- $2 + 2 = 5$ . (F)
- $1 + 2 = 3$ . (T)
- A fact-based declaration is a proposition, even if no one knows whether it is true
  - 11213 is prime.
  - There exists an odd perfect number.

- Give an example of a statement that is not a proposition.

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- $x + 1 = 8$



- Give an example of a statement that is not a proposition.

- $x + 1 = 8$

- propositions are statements that are either true or false.
- Just as we use variables  $x; y; : : :$  for numbers, we will use variables  $p; q; : : :$  for propositions.

“if it is raining, it will be wet” :  $p \rightarrow q$

# Examples of Non-Propositions

The following are **NOT** propositions:

- Who's there? (interrogative, question)
- Just do it! (imperative, command)
- La la la la la. (meaningless interjection)
- Yeah, I sorta dunno, whatever... (vague)
- $1 + 2$  (expression with a non-true/false value)
- $x + 2 = 5$  (declaration about semantic tokens of non-constant value)

# Truth Tables

- An *operator* or *connective* combines one or more *operand* expressions into a larger expression. (e.g., “+” in numeric expressions.)
- **Unary** operators take *one* operand (e.g.,  $-3$ );  
**Binary** operators take *two* operands (e.g.  $3 \times 4$ ).
- **Propositional** or **Boolean operators** operate on propositions (or their truth values) instead of on numbers.
- The **Boolean domain** is the set  $\{T, F\}$ . Either of its elements is called a **Boolean value**.  
An  $n$ -tuple  $(p_1, \dots, p_n)$  of Boolean values is called a **Boolean  $n$ -tuple**.
- An  $n$ -operand truth table is a table that assigns a Boolean value to the set of all Boolean  $n$ -tuples.

# Some Popular Boolean Operators

Formal Name	Nickname	Arity	Symbol
Negation operator	NOT	Unary	$\neg$
Conjunction operator	AND	Binary	$\wedge$
Disjunction operator	OR	Binary	$\vee$
Exclusive-OR operator	XOR	Binary	$\oplus$
Implication operator	IMPLIES	Binary	$\rightarrow$
Biconditional operator	IFF	Binary	$\leftrightarrow$

# The Negation Operator

- The unary ***negation operator*** “ $\neg$ ” (*NOT*) transforms a proposition into its logical *negation*.
- *E.g.* If  $p$  = “I have brown hair.”  
then  $\neg p$  = “It is not the case that I have brown hair” or “I do **not** have brown hair.”
- The *truth table* for NOT:

$p$	$\neg p$
T	F
F	T

Operand  
column

Result  
column

# The Conjunction Operator

- The binary ***conjunction operator*** “ $\wedge$ ” (***AND***) combines two propositions to form their logical ***conjunction***.
- *E.g.* If  $p$  = “I will have salad for lunch.” and  $q$  = “I will have steak for dinner.”  
then,  $p \wedge q$  = “I will have salad for lunch **and** I will have steak for dinner.”

# Conjunction Truth Table

Operand columns

$p$	$q$	$p \wedge q$
T	T	T
T	F	F
F	T	F
F	F	F

- Note that a conjunction  $p_1 \wedge p_2 \wedge \dots \wedge p_n$  of  $n$  propositions will have  $2^n$  rows in its truth table

# The Disjunction Operator

- The binary ***disjunction operator*** “ $\vee$ ” (*OR*) combines two propositions to form their logical *disjunction*.

- *E.g.* If  $p$  = “My car has a bad engine.” and  $q$  = “My car has a bad carburetor.”

then,  $p \vee q$  = “My car has a bad engine, **or** my car has a bad carburetor.”

Meaning is like “and/or” in informal English.



# Disjunction Truth Table

$p$	$q$	$p \vee q$
T	T	T
T	F	<b>T</b>
F	T	<b>T</b>
F	F	F

Note difference from AND

- Note that  $p \vee q$  means that  $p$  is true, or  $q$  is true, **or both** are true!
- So, this operation is also called ***inclusive or***, because it **includes** the possibility that both  $p$  and  $q$  are true.

# The Exclusive-Or Operator

- The binary **exclusive-or operator** “ $\oplus$ ” (*XOR*) combines two propositions to form their logical “exclusive or”
- *E.g.* If  $p$  = “I will earn an A in this course.” and  $q$  = “I will drop this course.”, then
$$p \oplus q = \text{“I will either earn an A in this course, or I will drop it (but not both!)”}$$

# Exclusive-Or Truth Table

$p$	$q$	$p \oplus q$
T	T	<b>F</b>
T	F	T
F	T	T
F	F	F

Note difference  
from OR.

- Note that  $p \oplus q$  means that  $p$  is true, or  $q$  is true, but **not both**!
- This operation is called **exclusive or**, because it **excludes** the possibility that both  $p$  and  $q$  are true.

# Natural Language is ambiguous

- Note that the English “or” can be ambiguous regarding the “both” case!

- “Pat is a singer or Pat is a writer.” -  $\vee$

- “Pat is a man or Pat is a woman.” -  $\oplus$

$p$	$q$	$p$ "or" $q$
T	T	?
T	F	T
F	T	T
F	F	F

- Need context to disambiguate the meaning!
- For this class, assume “or” means inclusive ( $\vee$ ).

# The Implication Operator

- The conditional statement (aka ***implication***)  
 $p \rightarrow q$  states that  $p$  implies  $q$ .
- */e.*, If  $p$  is true, then  $q$  is true; but if  $p$  is not true, then  $q$  could be either true or false.
- *E.g.*, let  $p$  = “You study hard.”  
 $q$  = “You will get a good grade.”  
 $p \rightarrow q$  = “If you study hard, then you will get a good grade.” (else, it could go either way)
  - $p$ : *hypothesis* or *antecedent* or *premise*
  - $q$ : *conclusion* or *consequence*

# Implication Truth Table

$p$	$q$	$p \rightarrow q$
T	T	T
T	F	<b>F</b>
F	T	T
F	F	T

The only False case!

- $p \rightarrow q$  is **false** only when  $p$  is true but  $q$  is **not** true.
- $p \rightarrow q$  does **not** require that  $p$  or  $q$  are ever true!
- E.g. “ $(1=0) \rightarrow$  pigs can fly” is TRUE!

# Examples of Implications

- “If this lecture ever ends, then the sun will rise tomorrow.” *True or False?* ( $T \rightarrow T$ )
- “If  $1+1=6$ , then Obama is president.” *True or False?* ( $F \rightarrow T$ )
- “If the moon is made of green cheese, then I am richer than Bill Gates.” *True or False?* ( $F \rightarrow F$ )
- “If Tuesday is a day of the week, then I am a penguin.” *True or False?* ( $T \rightarrow F$ )

# English Meaning Phrases of $p \rightarrow q$

“if  $p$ , then  $q$ ”

“if  $p$ ,  $q$ ”

“ $p$  is sufficient for  $q$ ”

“ $q$  if  $p$ ”

“ $q$  when  $p$ ”

“a necessary condition for  $p$  is  $q$ ”

“ $q$  unless  $\neg p$ ”

“ $p$  implies  $q$ ”

“ $p$  only if  $q$ ”

“a sufficient condition for  $q$  is  $p$ ”

“ $q$  whenever  $p$ ”

“ $q$  is necessary for  $p$ ”

“ $q$  follows from  $p$ ”



# Understanding the Implication Operation

- An obligation or an contract
- Example:
  - *“If I am elected, then I will lower taxes.”*
  - *“If you get 100% on the final, then you will get an A”.*
- Home exercise
  - Equivalence of *“if  $p$ , then  $q$ ”* and *“ $p$ , only if  $q$ ”*.
  - Equivalence of *“if  $p$ , then  $q$ ”* and *“ $q$  unless  $\neg p$ ”*.

# Exercise

- 22.** Write each of these statements in the form “if  $p$ , then  $q$ ” in English. [*Hint:* Refer to the list of common ways to express conditional statements provided in this section.]
- a)** It is necessary to wash the boss’s car to get promoted.
  - b)** Winds from the south imply a spring thaw.
  - c)** A sufficient condition for the warranty to be good is that you bought the computer less than a year ago.
  - d)** Willy gets caught whenever he cheats.
  - e)** You can access the website only if you pay a subscription fee.
  - f)** Getting elected follows from knowing the right people.
  - g)** Carol gets seasick whenever she is on a boat.

# Converse, Inverse and Contrapositive

- Some terminology, for an implication  $p \rightarrow q$ :
- Its **converse** is:  $q \rightarrow p$ .
- Its **inverse** is:  $\neg p \rightarrow \neg q$ .
- Its **contrapositive**:  $\neg q \rightarrow \neg p$ .

$p$	$q$	$p \rightarrow q$	$q \rightarrow p$	$\neg p \rightarrow \neg q$	$\neg q \rightarrow \neg p$
T	T	T	T	T	T
T	F	F	T	T	F
F	T	T	F	F	T
F	F	T	T	T	T

- One of these three has the *same meaning* (same truth table) as  $p \rightarrow q$ . Can you figure out which?

**Contrapositive**

# Examples

- $p$ : Today is Easter  
 $q$ : Tomorrow is Monday
- $p \rightarrow q$  :  
If today is Easter then tomorrow is Monday.
- **Converse**:  $q \rightarrow p$   
If tomorrow is Monday then today is Easter.
- **Inverse**:  $\neg p \rightarrow \neg q$   
If today is not Easter then tomorrow is not Monday.
- **Contrapositive**:  $\neg q \rightarrow \neg p$   
If tomorrow is not Monday then today is not Easter.

# The Biconditional Operator

- The **biconditional** statement  $p \leftrightarrow q$  states that  $p$  **if and only if** (iff)  $q$ .
- $p =$  "It is below freezing."  
 $q =$  "It is snowing."  
 $p \leftrightarrow q =$  "It is below freezing if and only if it is snowing."  
or  
 $=$  "That it is below freezing is necessary and sufficient for it to be snowing"

# Biconditional Truth Table

$p$	$q$	$p \leftrightarrow q$
T	T	T
T	F	F
F	T	F
F	F	T

- $p$  is necessary and sufficient for  $q$
- If  $p$  then  $q$ , and conversely
- $p$  iff  $q$

- $p \leftrightarrow q$  is equivalent to  $(p \rightarrow q) \wedge (q \rightarrow p)$ .
- $p \leftrightarrow q$  means that  $p$  and  $q$  have the **same** truth value.
- $p \leftrightarrow q$  does **not** imply that  $p$  and  $q$  are true.
- Note this truth table is the exact **opposite** of  $\oplus$ 's!  
Thus,  $p \leftrightarrow q$  means  $\neg(p \oplus q)$ .

# Combining Propositions

- Conjunction:  $p \wedge q$ , (read  $p$  and  $q$ ), “discrete math is a required course **and** I am a computer science major”.
- Disjunction: ,  $p \vee q$ , (read  $p$  or  $q$ ), “discrete math is a required course **or** I am a computer science major”.
- Exclusive or:  $p \oplus q$ , “discrete math is a required course **or** I am a computer science major **but not both**”.
- Implication:  $p \rightarrow q$ , “**if** discrete math is a required course **then** I am a computer science major”.
- Biconditional:  $p \leftrightarrow q$ , “discrete math is a required course **if and only if** I am a computer science major”.

# Boolean Operations Summary

- We have seen 1 unary operator and 5 binary operators. What are they? Their truth tables are below.

$p$	$q$	$\neg p$	$p \wedge q$	$p \vee q$	$p \oplus q$	$p \rightarrow q$	$p \leftrightarrow q$
T	T	F	T	T	F	T	T
T	F	F	F	T	T	F	F
F	T	T	F	T	T	T	F
F	F	T	F	F	F	T	T

- For an implication  $p \rightarrow q$
- Its **converse** is:  $q \rightarrow p$
- Its **inverse** is:  $\neg p \rightarrow \neg q$
- Its **contrapositive**:  $\neg q \rightarrow \neg p$



# Compound Propositions

- A **propositional variable** is a variable such as  $p, q, r$  (possibly subscripted, e.g.  $p_j$ ) over the Boolean domain.
- An **atomic proposition** is either Boolean constant or a propositional variable: e.g.  $T, F, p$
- A **compound proposition** is derived from atomic propositions by application of propositional operators: e.g.  $\neg p, p \vee q, (p \vee \neg q) \rightarrow q$
- Precedence of logical operators:  $\neg, \wedge, \vee, \rightarrow, \leftrightarrow$
- Precedence also can be indicated by parentheses.
  - e.g.  $\neg p \wedge q$  means  $(\neg p) \wedge q$ , not  $\neg(p \wedge q)$

# An Exercise

- Any compound proposition can be evaluated by a truth table
- $(p \vee \neg q) \rightarrow q$

$p$	$q$	$\neg q$	$p \vee \neg q$	$(p \vee \neg q) \rightarrow q$
T	T	F	T	T
T	F	T	T	F
F	T	F	F	T
F	F	T	T	F

# Translating English Statements

- Let  $p$  = “It rained last night”,  
 $q$  = “The sprinklers came on last night,”  
 $r$  = “The lawn was wet this morning.”

Translate each of the following into English:

$\neg p$  = “It didn’t rain last night.”

$r \wedge \neg p$  = “The lawn was wet this morning,  
and it didn’t rain last night.”

$\neg r \vee p \vee q$  = “The lawn wasn’t wet this  
morning, or it rained last night, or  
the sprinklers came on last night.”

# Exercise

- Find the converse of the following statement.

*“Raining tomorrow is a sufficient condition for my not going to town.”*

# Another Example

- Find the converse of the following statement.
  - “Raining tomorrow is a sufficient condition for my not going to town.”
- **Step 1:** Assign propositional variables to component propositions.
  - $p$ : It will rain tomorrow
  - $q$ : I will not go to town
- **Step 2:** Symbolize the assertion:  $p \rightarrow q$
- **Step 3:** Symbolize the converse:  $q \rightarrow p$
- **Step 4:** Convert the symbols back into words.
  - “If I don’t go to town then it will rain tomorrow” or
  - “Raining tomorrow is a *necessary condition* for my not going to town.”

# Logic and Bit Operations

- A **bit** is a **b**inary (base 2) dig**it**: 0 or 1.
- Bits may be used to represent truth values.
  - By convention:  
0 represents “False”; 1 represents “True”.
- A **bit string of length  $n$**  is an ordered sequence of  $n \geq 0$  bits.
- By convention, bit strings are (sometimes) written left to right:
  - e.g. the “first” bit of the bit string “1001101010” is 1.
  - What is the length of the above bit string?

# Bitwise Operations

- Boolean operations can be extended to operate on bit strings as well as single bits.

- Example:

01 1011 0110

11 0001 1101

11 1011 1111 Bit-wise OR

01 0001 0100 Bit-wise AND

10 1010 1011 Bit-wise XOR

# Propositional Equivalences

- A **tautology** is a compound proposition that is **true** *no matter what* the truth values of its atomic propositions are!
  - e.g.  $p \vee \neg p$  (“Today the sun will shine or today the sun will not shine.”) [What is its truth table?]
- A **contradiction** is a compound proposition that is **false** no matter what!
  - e.g.  $p \wedge \neg p$  (“Today is Wednesday and today is not Wednesday.”) [Truth table?]
- A **contingency** is a compound proposition that is neither a tautology nor a contradiction.
  - e.g.  $(p \vee q) \rightarrow \neg r$



# Logical Equivalence

- Compound proposition  $p$  is **logically equivalent** to compound proposition  $q$ , written  $p \equiv q$  or  $p \Leftrightarrow q$ , **iff** the compound proposition  $p \leftrightarrow q$  is a tautology.
- Compound propositions  $p$  and  $q$  are logically equivalent to each other **iff**  $p$  and  $q$  contain the same truth values as each other in all corresponding rows of their truth tables.

# Proving Equivalence Via Truth Tables

- Prove that  $\neg(p \wedge q) \equiv \neg p \vee \neg q$ . (De Morgan's law)

$p$	$q$	$p \wedge q$	$\neg p$	$\neg q$	$\neg p \vee \neg q$	$\neg(p \wedge q)$
T	T	T	F	F	F	F
T	F	F	F	T	T	T
F	T	F	T	F	T	T
F	F	F	T	T	T	T

- Show that [Check out the solution in the textbook!](#)
  - $\neg(p \vee q) \equiv \neg p \wedge \neg q$  (De Morgan's law)
  - $p \rightarrow q \equiv \neg p \vee q$
  - $p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$  (distributive law)

# Equivalence Laws

- These are similar to the arithmetic identities you may have learned in algebra, but for propositional equivalences instead.
- They provide a pattern or template that can be used to match part of a much more complicated proposition and to find an equivalence for it and possibly simplify it.

# Equivalence Laws

- *Identity:*  $p \wedge \mathbf{T} \equiv p$      $p \vee \mathbf{F} \equiv p$
- *Domination:*  $p \vee \mathbf{T} \equiv \mathbf{T}$      $p \wedge \mathbf{F} \equiv \mathbf{F}$
- *Idempotent:*  $p \vee p \equiv p$      $p \wedge p \equiv p$
- *Double negation:*  $\neg\neg p \equiv p$
- *Commutative:*  $p \vee q \equiv q \vee p$      $p \wedge q \equiv q \wedge p$
- *Associative:*  $(p \vee q) \vee r \equiv p \vee (q \vee r)$   
 $(p \wedge q) \wedge r \equiv p \wedge (q \wedge r)$

# More Equivalence Laws

- *Distributive:*  
$$p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$$
$$p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$$
- *De Morgan's:*  
$$\neg(p \wedge q) \equiv \neg p \vee \neg q$$
$$\neg(p \vee q) \equiv \neg p \wedge \neg q$$
- *Absorption*  
$$p \vee (p \wedge q) \equiv p \qquad p \wedge (p \vee q) \equiv p$$
- *Trivial tautology/contradiction:*  
$$p \vee \neg p \equiv \mathbf{T} \qquad p \wedge \neg p \equiv \mathbf{F}$$

# Defining Operators via Equivalences

Using equivalences, we can *define* operators in terms of other operators.

- Exclusive or:  $p \oplus q \equiv (p \wedge \neg q) \vee (\neg p \wedge q)$   
 $p \oplus q \equiv (p \vee q) \wedge \neg(p \wedge q)$
- Implies:  $p \rightarrow q \equiv \neg p \vee q$
- Biconditional:  $p \leftrightarrow q \equiv (p \rightarrow q) \wedge (q \rightarrow p)$   
 $p \leftrightarrow q \equiv \neg(p \oplus q)$

This way we can “normalize” propositions

# An Example Problem

- Show that  $\neg(p \rightarrow q)$  and  $p \wedge \neg q$  are logically equivalent.

$$\neg(p \rightarrow q) \quad [\text{Expand definition of } \rightarrow]$$

$$\equiv \neg(\neg p \vee q) \quad [\text{DeMorgan's Law}]$$

$$\equiv \neg(\neg p) \wedge \neg q \quad [\text{Double Negation}]$$

$$\equiv p \wedge \neg q$$

# Another Example

- Check using a symbolic derivation whether

$$(p \wedge \neg q) \rightarrow (p \oplus r) \equiv \neg p \vee q \vee \neg r$$



# Solution ...

$$\begin{aligned}(p \wedge \neg q) &\rightarrow (p \oplus r) \quad [\text{Expand definition of } \rightarrow] \\ &\equiv \neg(p \wedge \neg q) \vee (p \oplus r) \quad [\text{Expand definition of } \oplus] \\ &\equiv \neg(p \wedge \neg q) \vee ((p \vee r) \wedge \neg(p \wedge r)) \\ &\hspace{15em} [\text{DeMorgan's Law}] \\ &\equiv (\neg p \vee q) \vee ((p \vee r) \wedge \neg(p \wedge r)) \\ &\hspace{15em} \textit{cont.}\end{aligned}$$

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$$\begin{aligned}
& (\neg p \vee q) \vee ((p \vee r) \wedge \neg(p \wedge r)) \text{ [\vee Commutative]} \\
& \equiv (q \vee \neg p) \vee ((p \vee r) \wedge \neg(p \wedge r)) \text{ [\vee Associative]} \\
& \equiv q \vee ( \neg p \vee ((p \vee r) \wedge \neg(p \wedge r)) ) \text{ [Distribute } \vee \text{ over } \wedge] \\
& \equiv q \vee ( ( \neg p \vee (p \vee r) ) \wedge ( \neg p \vee \neg(p \wedge r) ) ) \text{ [\vee Assoc.]} \\
& \equiv q \vee ( ( \neg p \vee p ) \vee r ) \wedge ( \neg p \vee \neg(p \wedge r) ) ) \text{ [Trivial taut.]} \\
& \equiv q \vee ( (\mathbf{T} \vee r) \wedge ( \neg p \vee \neg(p \wedge r) ) ) \text{ [Domination]} \\
& \equiv q \vee ( \mathbf{T} \wedge ( \neg p \vee \neg(p \wedge r) ) ) \text{ [Identity]} \\
& \equiv q \vee ( \neg p \vee \neg(p \wedge r) )
\end{aligned}$$

$$\begin{aligned}
& q \vee (\neg p \vee \neg(p \wedge r)) && [\text{DeMorgan's Law}] \\
& \equiv q \vee (\neg p \vee (\neg p \vee \neg r)) && [\vee \text{ Associative}] \\
& \equiv q \vee ((\neg p \vee \neg p) \vee \neg r) && [\text{Idempotent}] \\
& \equiv q \vee (\neg p \vee \neg r) && [\text{Associative}] \\
& \equiv (q \vee \neg p) \vee \neg r && [\vee \text{ Commutative}] \\
& \equiv \neg p \vee q \vee \neg r \quad \blacksquare
\end{aligned}$$

# Review: Propositional Logic

- Atomic propositions:  $p, q, r, \dots$
- Boolean operators:  $\neg \wedge \vee \oplus \rightarrow \leftrightarrow$
- Compound propositions:  $(p \wedge \neg q) \vee r$
- Equivalences:  $p \wedge \neg q \leftrightarrow \equiv \neg(p \rightarrow q)$
- Proving equivalences using:
  - Truth tables
  - Symbolic derivations (series of logical equivalences)  $p \equiv q \equiv r \equiv \dots$