

CSE565 Lab 1

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Before You Start:

*Please write a detailed lab report, with **screenshots**, to describe what you have **done** and what you have **observed**. You also need to provide an **explanation** to the observations that you noticed. Please also show the important **code snippets** followed by an explanation. Simply attaching a code without any explanation will **NOT** receive credits.*

After you finish, export this report as a **PDF** file and submit it on UBLearn.

Academic Integrity Statement:

I, Kishan Nagaraja, have read and understood the course academic integrity policy.

(Your report will not be graded without filling in your name in the above AI statement)

Task-1: Frequency Analysis

Since monosubstitution ciphers replace each letter in the alphabet with the same letter every time, it is easily possible to decrypt the cipher text by studying the pattern and occurrence of words in the respective language. In the English Language, the order of occurrence of letters from highest to lowest is given by the order “ETAOINSRHDLCMFYWGPVKXQJZ”. (reference: <https://pi.math.cornell.edu/~mec/2003-2004/cryptography/subs/frequencies.html>)

Below is the snapshot of the given cipher text:

```
ytn xqavhq yzhu xu qzupvd lmat qnncq vxzy hmrty vbynh ytmq ixur qyhurn  
vlvhpq yhme ytn gvrnnh bnniq imsn v uxuvrnuvhmvu yxx  
  
ytn vlvhpq hvan lvq gxxsnupnp gd ytn pncmqn xb tvhfnd lnmuqynmu vy myq xzyqny  
vup ytn veevhnu ymcnq ze givasrxlu eximymaq vhacavupd vaymfmcq vup  
v uvymxuvi axufnhqvyxu vq ghmn vup cvp vq v bfnh phnvc vgxy ltnytnh ytnhn  
xzrt yx gn v ehnqmpnu ymubhnd ytn qnvqxu pmpuy ozqy qnnc nkyhv ixur my lvq  
nkyhv ixur gnavzqn ytn xqavhq lnhn cxfnp yx ytn bmhqqy lnnsnup mu cvhat yx  
vfxmp axubimaymur lmyt ytn aixqmur anhncxud xb ytn lmuynh xidcemaq ytvsq  
ednxuratvur  
  
xun gmr jznqymxu qzhhxzupmur ytmq dnvhq vavpncd vlvhpq mq txl xh mb ytn  
anhncxud lmii vpphnnq cnyxx nqenamviid vbynh ytn rxipnu rixgnq lmat gnavcn  
v ozgimivuy axcmurxzy evhyd bxh ymcnq ze ytn cxfncnuy qenvhtnvpnp gd  
exlnhbzi txidlxp lxncu ltx tnienp hvmqn cmiimxuq xb pxivh yx bmrt ynkzvi  
tvhqqcnu ytvu v axzuyhd  
  
qmruvimir ytnmh qzeexhy rxipnu rixgnq vyyupnnq qlvytnp ytnqcnifnq mu givas  
qexhypn iveni emuq vup qxzupnp xbb vxzy qnkqy exlnh mcgvivuanq bhxc ytn hnp  
avheny vup ytn qyvrn xu ytn vmb n lvq aviinp xzy vxzy evd munjzmyd vbynh  
myq bxhcnuh vuatxh avyy qvpinh jzmy xuan qtn invhump ytvy qtn lvq cvsmur bv  
inqq ytvu v cvin axtxqy vup pzhmur ytn anhncxud uvyvimm exhycvu yxxs v gizuy  
vup qvymqbdmnr pmr vy ytn viicvin hxqynh xb uxcmuvynp pmhnayxq txl axzip  
ytvy gn yxeenp
```

1. Snapshot of one gram frequency extracted from the ciphertext:

```
[09/01/23] seed@VM:~/.../Files$ python3 freq.py  
-----  
1-gram (top 20):  
n: 488  
y: 373  
v: 348  
x: 291  
u: 280  
q: 276  
m: 264  
h: 235  
t: 183  
i: 166  
p: 156  
a: 116  
c: 104  
z: 95  
l: 90  
g: 83  
b: 83  
r: 82  
e: 76  
d: 59
```

Observation-1: We can observe that the letter ‘v’ is the most frequently appearing one-letter word in the cipher sentences that can correspond to the letter ‘a’ in the English language. The order of frequency of the letter ‘v’ in ciphertext also matches that of the letter ‘a’ in the English language. Therefore, we can conclude that the letter ‘v’ is substituted for the letter ‘a’ in the cipher.

Below is the snapshot of the top 20 bigrams and trigrams extracted from the ciphertext.

2-gram (top 20): yt: 115 tn: 89 mu: 74 nh: 58 vh: 57 hn: 57 vu: 56 nq: 53 xu: 52 up: 46 xh: 45 yn: 44 np: 44 vy: 44 nu: 42 qy: 39 vq: 33 vi: 32 gn: 32 av: 31	3-gram (top 20): ytn: 78 vup: 30 mur: 20 ynh: 18 xzy: 16 mxu: 14 gnq: 14 ytv: 13 nqy: 13 vii: 13 bxh: 13 lva: 12 nuy: 12 vyn: 12 uvy: 11 lmu: 11 nvh: 11 cmu: 11 tmq: 10 vhp: 10
---	--

The below snapshot shows the top bigram and trigram frequencies from the English language: (reference from Wikipedia)

Bigram frequency in the English language [edit]		
The frequency of the most common letter bigrams in a large English corpus is: ^[4]		
th 3.56%	of 1.17%	io 0.83%
he 3.07%	ed 1.17%	le 0.83%
in 2.43%	is 1.13%	ve 0.83%
er 2.05%	it 1.12%	co 0.79%
an 1.99%	al 1.09%	me 0.79%
re 1.85%	ar 1.07%	de 0.76%
on 1.76%	st 1.05%	hi 0.76%
at 1.49%	to 1.05%	ri 0.73%
en 1.45%	nt 1.04%	ro 0.73%
nd 1.35%	ng 0.95%	ic 0.70%
ti 1.34%	se 0.93%	ne 0.69%
es 1.34%	ha 0.93%	ea 0.69%
or 1.28%	as 0.87%	ra 0.69%
te 1.20%	ou 0.87%	ce 0.65%

Rank ^[1]	Trigram	Frequency ^[3] (Different source)
1	the	1.81%
2	and	0.73%
3	tha	0.33%
4	ent	0.42%
5	ing	0.72%
6	ion	0.42%
7	tio	0.31%
8	for	0.34%

Observation-2: We observe that the order of frequencies of bigrams ‘yt’ , ‘tn’ in ciphertext corresponds to bigrams ‘th’, ‘he’ in the English language. Moreover, the order of frequency of trigram ‘ytn’ in ciphertext corresponds to trigram ‘the’ in English. Therefore, from these observations, it can be concluded the letters ‘y’, ‘t’, and ‘n’ are replacements for letters ‘t’, ‘h’, and ‘e’ respectively.

Also, replace the letter ‘x’ in ciphertext with the letter ‘o’ since they have the same position in the order of frequency table and they are the top 4 most occurring letters in ciphertext and English language respectively.

Below is the snapshot of partial plain text after applying the substitutions from the above observations:

v	y	t	n	x
A	T	H	E	O

```
[09/02/23] seed@VM:~/.../Files$ tr 'vytnx' 'ATHE0' < ciphertext.txt > plain1.txt
[09/02/23] seed@VM:~/.../Files$ cat plain1.txt
THE OqaAhq Tzhu Ou qzupAd lHmaH qEEcq AgOzT hmrHT AbTEh THmq iOur qThAurE
AlAhpq Thme THE gArrEh bEEiq imsE A uOuArEuAhmAu T00

THE AlAhpq hAaE lAq g00sEupEp gd THE pEcmsgE Ob HAhfEd lEmuqTEmu AT mTq OzTqET
Aup THE AeeAhEuT mcei0qm0u Ob Hmq bmic aOceAud AT THE Eup Aup mT lAq qHAeEp gd
THE EcEhrEuaE Ob cET00 TmcEq ze giAasr0lu e0imTmaq AhcaAupd AaTmfmqc Aup
A uATm0uAi a0ufEhqATm0u Aq ghmEb Aup cAp Aq A bEfEh phEAc AgOzT lHETHEh THEhE
OzrHT TO gE A ehEqmpEuT lmubhEd THE qEAq0u pmpuT ozqT qEEc EkThA iOur mT lAq
EkThA iOur gEaAzqE THE OqaAhq lEhE c0fEp TO THE bmhqt lEEsEup mu cAhaH TO
Af0mp a0ubimaTm0u lmTH THE ai0qmur aEhEcoud Ob THE lmuTEh Oidcemaq THAusq
edEOuraHAur

OuE gmr jzEqTm0u qzhhozupmur THmq dEAhq AaApEcd AlAhpq mq H0l Oh mb THE
aEhEcoud lmii ApphEqq cET00 EqeEamAiid AbTEh THE r0ipEu ri0gEq lHmaH gEaAcE
A ozgmiAut a0cmur0zT eAhTd b0h TmcEq ze THE c0fEcEuT qeEAhHEApEp gd
e0lEhbzi H0iidl00p l0cEu lHO HEieEp hAmqE cmiim0uq Ob p0iiAhq TO bmrHT qEkzAi
HAhAqqcEut Ah0zup THE a0zuThd
```

From the above partial plaintext, we can deduce the below words in the ciphertext and make corresponding mappings based on knowledge of the usage of English words.

THEhE	THERE
OzrHT	OUGHT
gE	BE

Snapshot of second partial plaintext after making the substitutions based on the above words:

h	z	r	g
R	U	G	B

```
[09/02/23] seed@VM:~/.../Files$ tr 'hzrg' 'RUGB' < plain1.txt > plain2.txt
[09/02/23] seed@VM:~/.../Files$ cat plain2.txt
THE OqaARq TURu Ou qUupAd lHmaH qEEcq ABOUT RmGHT AbTER THmq i0uG qTRAuGE
ALARpq TRme THE BAGGER bEEiq imsE A uOuAGEuARmAu T00

THE ALARpq RAaE lAq B00sEupEp Bd THE pEcmqE Ob HARfEd lEmuqTEmu AT mTq OUTqET
Aup THE AeeAREuT mcei0qm0u Ob Hmq bmic a0ceAud AT THE Eup Aup mT lAq qHAeEp Bd
THE EcERGEuaE Ob cET00 TmcEq Ue BiAsGolu e0imTmaq ARcaAupd AaTmfmcq Aup
A uATm0uAi a0ufERqATm0u Aq BRmEb Aup cAp Aq A bEfER pREAc ABOUT lHETHER THERE
OUGHT TO BE A eREEqmpEuT lmubREd THE qEAq0u pmput oUqT qEEc EkTRA i0uG mT lAq
EkTRA i0uG BEaAUqE THE OqaARq lERE c0fEp T0 THE bmRqT lEEsEup mu cARaH T0
Afomp a0ubimaTmuG lmTH THE ai0qmuG aEREc0ud Ob THE lmuTER Oidcemaq THAusq
edEOuGaHAuG

OuE BmG jUEqTm0u qURROUupmuG THmq dEARq AaApEcd ALARpq mq H0l OR mb THE
aEREc0ud lmii AppREqq cET00 EqeEamAiid AbTER THE G0ipEu Gi0BEq lHmaH BEaAcE
A oUBmiAuT a0cmuG0UT eARTd b0R TmcEq Ue THE c0fEcEuT qeEARHEApEp Bd
e0lERbUi H0iidl00p l0cEu lHO HEieEp RAMqE cmiim0uq Ob p0iiARq T0 bmGHT qEkUAi
HARAqqcEuT AROUup THE aOUuTRd

qmGuAimuG THEmR qUeeORT G0ipEu Gi0BEq ATTEmqEEq qlATHEp THEEcqEifEq mu BiAs
qeORTEp iAeEi emuq Aup qOUupEp Obb ABOUT qEkmqT e0lER mcBAiAuaEq bR0c THE REp
aAReET Aup THE qTAGE Ou THE AmR E lAq aAiiEp OUT ABOUT eAd muEjUmTd AbTER
mTq b0RcER AuaHOR aATT qApiER jUmT OuaE qHE iEARuEp THAT qHE lAq cAsmuG bAR
iEqq THAu A cAiE a0H0qT Aup pURmuG THE aEREc0ud uATAimE e0RTcAu T00s A BiUuT
Aup qATmqbdmuG pmG AT THE AiicAiE R0qTER Ob u0cmuATEp pmREaT0Rq H0l aOUip
```

From the above partial plaintext, we can guess more words (boxed in red) and make appropriate substitutions for the letters in the ciphertext.

RmGHT	RIGHT
AbTER	AFTER
OUTqET	OUTSET
qURROUupmuG	SURROUNDING
AROUup	AROUND
qTAGE	STAGE
AuaHOR	ANCHOR
EkTRA	EXTRA

This leads to the following letter substitutions:

m	b	q	u	p	a	k
I	F	S	N	D	C	X

Snapshot of partial plaintext after making substitutions as above:

```
[09/02/23] seed@VM:~/.../Files$ tr 'mbqupak' 'IFSNDCX' < plain2.txt > plain3.txt
[09/02/23] seed@VM:~/.../Files$ cat plain3.txt
THE OSCARS TURN ON SUNDAd lHICH SEEcS ABOUT RIGHT AFTER THIS iONG STRANGE
ALARDS TRIe THE BAGGER FEEis iIsE A NONAGENARIAN TOO

THE ALARDS RACE LAS BOOsENDED Bd THE DEcISE OF HARfEd LEINSTEIN AT ITS OUTSET
AND THE AeeARENT IceiOSH OF HIS Fiic COceANd AT THE END AND IT LAS SHAeED Bd
THE EcERGENC OF cETOO TiCes Ue BiACsGOLN eOiTICS ARcCANDd ACTIfISc AND
A NATIONAl CONFERSATION AS BRIEF AND cAD AS A FEFeR DREAc ABOUT lHETHER THERE
OUGHT TO BE A eRESIDENT lINFREd THE SEASON DIDNT oUST SEEc EXTRA iONG IT LAS
EXTRA iONG BECAUSE THE OSCARS LERE cofED TO THE FIRST LEEsEND IN cARCH TO
AfOID CONFLICTING lITH THE CiOSING CEREcOND OF THE lINTER OidceICS THANsS
edEONGCHANG

ONE BIG jUESTION SURROUNDING THIS dEARS ACADEcd ALARDS IS HOL OR IF THE
CEREcOND lIii ADDRESS cETOO ESeCIAiid AFTER THE GOiDEN GiOBES lHICH BECAcE
A oUBIIANT COcINGOUT eARTd FOR TiCes Ue THE cOfEcENT SeEARHEADED Bd
eOLERFuI HOiidLOOD lOcEN lHO HEieED RAISE cIiIONS OF DOiiARS TO FIGHT SEXUAI
HARASScENT AROUND THE COUNTrd

SIGNAiNG THEIR SUeeORT GOiDEN GiOBES ATTENDEES SLATHED THEcSEifES IN BiACs
SeORTED iAeEi eINS AND SOUNDED OFF ABOUT SEXIST eOLER IcBAiANCES FROc THE RED
CAReET AND THE STAGE ON THE AIR E LAS CAiiED OUT ABOUT eAd INEjUITd AFTER
ITS FORcER ANCHOR CATT SADiER jUIT ONCE SHE iEARNED THAT SHE LAS cAsING FAR
iESS THAN A cAiE COHOST AND DURING THE CEREcOND NATAiIE eORTcAN TOO A BiUNT
AND SATISFding DIG AT THE AiicAiE ROSTER OF NOcINATED DIRECTORS HOL COUiD
```

The highlighted words give final clues about the remaining letters in the ciphertext.

SUNDAd	SUNDAY
lHICH	WHICH
AeeARENT	APPARENT
ACTIfISc	ACTIVISM
CiOSING	CLOSING
THANsS	THANKS
jUESTION	QUESTION
oUBILANT	JUBILANT
HARASScENT	HARASSMENT

Based on the above clue words, we can make substitutions for the remaining letters as below:

c	d	e	f	i	j	l	o	s	w
M	Y	P	V	L	Q	W	J	K	Z

On final substitution, we get the complete meaningful plaintext:

```
[09/02/23]seed@VM:~/.../Files$ tr 'cdefijlosw' 'MYPVLQWJKZ' < plain3.txt > plain4.txt
[09/02/23]seed@VM:~/.../Files$ cat plain4.txt
THE OSCARS TURN ON SUNDAY WHICH SEEMS ABOUT RIGHT AFTER THIS LONG STRANGE
AWARDS TRIP THE BAGGER FEELS LIKE A NONAGENARIAN TOO

THE AWARDS RACE WAS BOOKENDED BY THE DEMISE OF HARVEY WEINSTEIN AT ITS OUTSET
AND THE APPARENT IMPLOSION OF HIS FILM COMPANY AT THE END AND IT WAS SHAPED BY
THE EMERGENCE OF METOO TIMES UP BLACKGOWN POLITICS ARMCANDY ACTIVISM AND
A NATIONAL CONVERSATION AS BRIEF AND MAD AS A FEVER DREAM ABOUT WHETHER THERE
OUGHT TO BE A PRESIDENT WINFREY THE SEASON DIDNT JUST SEEM EXTRA LONG IT WAS
EXTRA LONG BECAUSE THE OSCARS WERE MOVED TO THE FIRST WEEKEND IN MARCH TO
AVOID CONFLICTING WITH THE CLOSING CEREMONY OF THE WINTER OLYMPICS THANKS
PYEONGCHANG

ONE BIG QUESTION SURROUNDING THIS YEARS ACADEMY AWARDS IS HOW OR IF THE
CEREMONY WILL ADDRESS METOO ESPECIALLY AFTER THE GOLDEN GLOBES WHICH BECAME
A JUBILANT COMINGOUT PARTY FOR TIMES UP THE MOVEMENT SPEARHEADED BY
POWERFUL HOLLYWOOD WOMEN WHO HELPED RAISE MILLIONS OF DOLLARS TO FIGHT SEXUAL
HARASSMENT AROUND THE COUNTRY

SIGNALING THEIR SUPPORT GOLDEN GLOBES ATTENDEES SWATHED THEMSELVES IN BLACK
SPORTED LAPEL PINS AND SOUNDED OFF ABOUT SEXIST POWER IMBALANCES FROM THE RED
CARPET AND THE STAGE ON THE AIR E WAS CALLED OUT ABOUT PAY INEQUITY AFTER
ITS FORMER ANCHOR CATT SADLER QUIT ONCE SHE LEARNED THAT SHE WAS MAKING FAR
LESS THAN A MALE COHOST AND DURING THE CEREMONY NATALIE PORTMAN TOOK A BLUNT
AND SATISFYING DIG AT THE ALLMALE ROSTER OF NOMINATED DIRECTORS HOW COULD
THAT BE TOPPED
```

The CipherKey used in the above encryption is
‘CFMYPVBRLQXWIEJDSGKHNAZOTU’ substituted for
‘ABCDEFGHIJKLMNOPQRSTUVWXYZ’.

Task 2: Encryption using Different Ciphers and Modes

Chosen Plain Text:

```
[09/02/23]seed@VM:~/.../Files$ cat plain4.txt
THE OSCARS TURN ON SUNDAY WHICH SEEMS ABOUT RIGHT AFTER THIS LONG STRANGE
AWARDS TRIP THE BAGGER FEELS LIKE A NONAGENARIAN TOO

THE AWARDS RACE WAS BOOKENDED BY THE DEMISE OF HARVEY WEINSTEIN AT ITS OUTSET
AND THE APPARENT IMPLOSION OF HIS FILM COMPANY AT THE END AND IT WAS SHAPED BY
THE EMERGENCE OF METOO TIMES UP BLACKGOWN POLITICS ARMCANDY ACTIVISM AND
A NATIONAL CONVERSATION AS BRIEF AND MAD AS A FEVER DREAM ABOUT WHETHER THERE
OUGHT TO BE A PRESIDENT WINFREY THE SEASON DIDNT JUST SEEM EXTRA LONG IT WAS
EXTRA LONG BECAUSE THE OSCARS WERE MOVED TO THE FIRST WEEKEND IN MARCH TO
AVOID CONFLICTING WITH THE CLOSING CEREMONY OF THE WINTER OLYMPICS THANKS
PYEONGCHANG

ONE BIG QUESTION SURROUNDING THIS YEARS ACADEMY AWARDS IS HOW OR IF THE
CEREMONY WILL ADDRESS METOO ESPECIALLY AFTER THE GOLDEN GLOBES WHICH BECAME
A JUBILANT COMINGOUT PARTY FOR TIMES UP THE MOVEMENT SPEARHEADED BY
POWERFUL HOLLYWOOD WOMEN WHO HELPED RAISE MILLIONS OF DOLLARS TO FIGHT SEXUAL
HARASSMENT AROUND THE COUNTRY

SIGNALING THEIR SUPPORT GOLDEN GLOBES ATTENDEES SWATHED THEMSELVES IN BLACK
SPORTED LAPEL PINS AND SOUNDED OFF ABOUT SEXIST POWER IMBALANCES FROM THE RED
CARPET AND THE STAGE ON THE AIR E WAS CALLED OUT ABOUT PAY INEQUITY AFTER
ITS FORMER ANCHOR CATT SADLER QUIT ONCE SHE LEARNED THAT SHE WAS MAKING FAR
LESS THAN A MALE COHOST AND DURING THE CEREMONY NATALIE PORTMAN TOOK A BLUNT
AND SATISFYING DIG AT THE ALLMALE ROSTER OF NOMINATED DIRECTORS HOW COULD
THAT BE TOPPED
```

1. Encryption using Advanced Encryption Standard AES-128 bit in CBC mode

- a. Key: 00112233445566778889aabcccddeeff
- b. IV: 0102030405060708
- c. Encryption:

```
[09/09/23]seed@VM:~/.../Files$ openssl enc -aes-128-cbc -e -in plain4.txt -out cipher_aes_128_cbc.bin \
> -K 00112233445566778889aabcccddeeff \
> -iv 0102030405060708
hex string is too short, padding with zero bytes to length
[09/09/23]seed@VM:~/.../Files$ hex

Command 'hex' not found, but can be installed with:

sudo apt install basez

[09/09/23]seed@VM:~/.../Files$ xxd cipher_aes_128_cbc.bin
00000000: e759 1f85 261e ee35 a9ec 730c 5efc 083e .Y..&..5..s.^..>
00000010: 9359 1122 f779 4d49 c9a4 8bdc 7caa d386 .Y.".yMI....|...
00000020: fe52 e7fe 9174 f8eb fef4 ca36 d109 6c4d .R...t.....6..lM
00000030: 6bdb 8 5cc6 579c bd17 ce94 b166 babc 76ab k.\.W.....f..v.
00000040: 6ac2 5161 89be 259d c852 5d43 dbde d352 j.Qa..%..R]C...R
00000050: 2215 64c1 dc81 d2c2 5a98 f4e5 0735 7033 ".d.....Z.....5p3
00000060: ae7c e125 f93c 9524 e6b3 1f62 0bcb 1cf3 .|.%.<.$....b....
00000070: a199 4b51 4cb5 c48a 8dcf 9840 036b 6d22 ..KQL.....@.km"
00000080: d28b 95fb eb72 1f3a 2b58 ed47 8e18 71e6 .....r.:+X.G..q.
00000090: 1889 1587 c1c4 8079 1c52 0084 c507 189b .....y.R.....
000000a0: c59e f599 b97c c744 78f3 d402 5568 aaab .....|.Dx...Uh..
000000b0: 3d7d 9507 6982 5b7d 8c03 a653 3671 7468 =}..i.[}...S6qth
000000c0: 7b5f 2371 0d6d bdb7 898f 110d 67aa cd91 { #q.m.....g...
000000d0: 6ac9 62b8 a344 f3dd 2e77 bd34 792f 3c31 j.b..D...w.4y/<1
000000e0: 8855 65f0 1a1a a579 1c7d e6f8 3828 2cdc .Ue....y}..8(,.
000000f0: dec7 abb9 3b38 1f5d 7508 3db1 c883 96ad ....;8.]u.=.....
00000100: c0b8 8419 8e01 806d 826d 11b7 5abe f195 .....m.m....Z...
00000110: b298 4bb2 696e 0e7a d0d5 bb5f e0aa 342d ..K.in.z.....4-
00000120: 26b4 41bf 20a8 2177 29d8 814a feb1 b095 &.A. .!w)..J....
```

d. Decryption

```
[09/09/23]seed@VM:~/.../Files$ openssl enc -aes-128-cbc -d -in cipher_aes_128_cbc.bin -out decrypted_aes_128_cbc.bin \
> -K 00112233445566778889aabcccddeeff \
> -iv 0102030405060708
hex string is too short, padding with zero bytes to length
[09/09/23]seed@VM:~/.../Files$ xxd decrypted_aes_128_cbc.bin
00000000: 5448 4520 4f53 4341 5253 2054 5552 4e20 THE OSCARS TURN
00000010: 204f 4e20 5355 4e44 4159 2057 4849 4348 ON SUNDAY WHICH
00000020: 2053 4545 4d53 2041 424f 5554 2052 4947 SEEMS ABOUT RIG
00000030: 4854 2041 4654 4552 2054 4849 5320 4c4f HT AFTER THIS LO
00000040: 4e47 2053 5452 414e 4745 0a41 5741 5244 NG STRANGE.AWARD
00000050: 5320 5452 4950 2054 4845 2042 4147 4745 S TRIP THE BAGGE
00000060: 5220 4645 454c 5320 4c49 4b45 2041 204e R FEELS LIKE A N
00000070: 4f4e 4147 454e 4152 4941 4e20 544f 4f0a ONAGENARIAN TOO.
00000080: 0a54 4845 2041 5741 5244 5320 5241 4345 .THE AWARDS RACE
00000090: 2057 4153 2042 4f4f 4b45 4e44 4544 2042 WAS BOOKENDED B
000000a0: 5920 5448 4520 4445 4d49 5345 204f 4620 Y THE DEMISE OF
000000b0: 4841 5256 4559 2057 4549 4e53 5445 494e HARVEY WEINSTEIN
000000c0: 2041 5420 4954 5320 4f55 5453 4554 0a41 AT ITS OUTSET.A
000000d0: 4e44 2054 4845 2041 5050 4152 454e 5420 ND THE APPARENT
000000e0: 494d 504c 4f53 494f 4e20 4f46 2048 4953 IMPLOSION OF HIS
000000f0: 2046 494c 4d20 434f 4d50 414e 5920 4154 FILM COMPANY AT
00000100: 2054 4845 2045 4e44 2041 4e44 2049 5420 THE END AND IT
00000110: 5741 5320 5348 4150 4544 2042 590a 5448 WAS SHAPED BY.TH
00000120: 4520 454d 4552 4745 4e43 4520 4f46 204d E EMERGENCE OF M
00000130: 4554 4f4f 2054 494d 4553 2055 5020 424c ET00 TIMES UP BL
00000140: 4143 4b47 4f57 4e20 504f 4c49 5449 4353 ACKGOWN POLITICS
00000150: 2041 524d 4341 4e44 5920 4143 5449 5649 ARMCANDY ACTIVI
00000160: 534d 2041 4e44 0a41 204e 4154 494f 4e41 SM AND.A NATIONA
00000170: 4c20 434f 4e56 4552 5341 5449 4f4e 2041 L CONVERSATION A
00000180: 5320 4252 4945 4620 414e 4420 4d41 4420 S BRIEF AND MAD
00000190: 4153 2041 2046 4556 4552 2044 5245 414d AS A FEVER DREAM
000001a0: 2041 424f 5554 2057 4845 5448 4552 2054 ABOUT WHETHER T
```

2. Encryption using Triple-DES (Data Encryption Standard) in CBC mode

- Key: 35539657483feadcb496408521acbfeee79bbcc3856444fc
- IV: 01020304 (uses only 8 bytes)
- Encryption:

```
[09/09/23] seed@VM:~/.../Files$ openssl enc -des-ede3-cbc -e -in plain4.txt -out cipher_3des_cbc.bin \
> -K 35539657483feadcb496408521acbfeee79bbcc3856444fc \
> -iv 01020304
hex string is too short, padding with zero bytes to length
[09/09/23] seed@VM:~/.../Files$ xxd cipher_3des_cbc.bin
00000000: 18b1 f221 7a6b 8806 9ac3 6de3 5ff9 7456 ...!zk....m._.tV
00000010: dc62 56e7 f5af 93f2 2c47 de73 760a 92b7 .bV.....,G.sv...
00000020: 2a8d 57f0 c52a 66d7 888b 9c4b 8259 8221 *.W..*f....K.Y.!.
00000030: 946d ded8 e7da 0a1d 7641 e1cd fa4f bd7e .m.....vA...0.~.
00000040: ea07 6702 1825 f119 1d42 6ec1 0502 9834 ..g.%...Bn....4
00000050: d5b0 53fa 8abf d481 85bf 1d30 d73c ff10 ..S.....0.<..
00000060: 2512 3fd8 a766 1a30 2366 b297 de08 4c6d %.?..f.0#f....Lm
00000070: 1bf9 2dda e66c 608e 2217 213e 61e3 ad0c ....l`.".!>a...
00000080: d01c aaeb 385d 61c3 3f0f 3cb7 bfc7 3d09 ..8]a.?.<...=.
00000090: b658 1d6c 2001 51a9 6dc0 b29d 10ae 367c .X.l .0.m.....6|
000000a0: 4e1d 478c 7f0c beda 76a2 54ce c02b a199 N.G.....v.T.+..
000000b0: 6933 0d7d 2bec 1620 4bf2 9aa2 4816 b313 i3.}+.. K...H...
000000c0: fbc7 6c56 484f 4f7e 3f35 a7d6 fd3b 1253 ..lVH00-?5...;S
000000d0: 3146 1e3b ba7f 6e8e 7c70 6e2e 2efb e70f 1F.;..n.|pn.....
000000e0: c596 2f89 104e 99ff 9004 eb75 8662 6479 ../.N.....u.bdy
000000f0: 1661 df10 13fe 2c4c 2a05 1b5e 8897 0e58 .a....,L*..^...X
00000100: c5ef 2add 5557 5f6b 8172 ef62 6ad8 580d ..*.UW_k.r.bj.X.
00000110: 8474 ce3e 5080 efc9 4e1d 9c95 f499 96da .t.>P...N.....
00000120: eaad 15aa f655 d4d6 2e12 1b42 8dd3 901b .....U....B....
00000130: 9f1d 0fa7 eae5 8077 0a97 9395 372a 9883 .....w....7*..
00000140: e8b2 562d fdc4 50c6 3c2f 5214 95c2 4fbf ..V...P.</R...0.
00000150: 74d8 7cce 971c 33cf a1f0 7a4f 3404 6e8f t.|...3...z04.n.
00000160: 3651 2eb1 3b1a bb58 8635 1d85 1f0e 8c3a 6Q...;..X.5....:
00000170: a6d6 d327 a7ba c355 d7b0 6fcf f035 7c75 ...'..U..o..5|u
00000180: e042 6573 bfbf 6365 19e8 50b6 92ec 3e7e .Bes..ce..P...>~
00000190: 9a84 9fe0 871e b93d 3c01 7758 aa5a ac7c .....=<.wX.Z.|_
000001a0: 10cf 4e12 def1 f677 1954 cc7d a5c8 10f3 ..N....w.T.}....
```

- Decryption

```
[09/09/23] seed@VM:~/.../Files$ openssl enc -des-ede3-cbc -d -in cipher_3des_cbc.bin -out decrypt_3des_cbc.bin \
> -K 35539657483feadcb496408521acbfeee79bbcc3856444fc \
> -iv 01020304
hex string is too short, padding with zero bytes to length
[09/09/23] seed@VM:~/.../Files$ xxd decrypt_3des_cbc.bin
00000000: 5448 4520 4f53 4341 5253 2054 5552 4e20 THE OSCARS TURN
00000010: 204f 4e20 5355 4e44 4159 2057 4849 4348 ON SUNDAY WHICH
00000020: 2053 4545 4d53 2041 424f 5554 2052 4947 SEEMS ABOUT RIG
00000030: 4854 2041 4654 4552 2054 4849 5320 4c4f HT AFTER THIS LO
00000040: 4e47 2053 5452 414e 4745 0a41 5741 5244 NG STRANGE AWARD
00000050: 5320 5452 4950 2054 4845 2042 4147 4745 S TRIP THE BAGGE
00000060: 5220 4645 454c 5320 4c49 4b45 2041 204e R FEELS LIKE A N
00000070: 4f4e 4147 454e 4152 4941 4e20 544f 4f0a ONAGENARIAN TOO.
00000080: 0a54 4845 2041 5741 5244 5320 5241 4345 .THE AWARDS RACE
00000090: 2057 4153 2042 4f4f 4b45 4e44 4544 2042 WAS BOOKENDED B
000000a0: 5920 5448 4520 4445 4d49 5345 204f 4620 Y THE DEMISE OF
000000b0: 4841 5256 4559 2057 4549 4e53 5445 494e HARVEY WEINSTEIN
000000c0: 2041 5420 4954 5320 4f55 5453 4554 0a41 AT ITS OUTSET.A
000000d0: 4e44 2054 4845 2041 5050 4152 454e 5420 ND THE APPARENT
000000e0: 494d 504c 4f53 494f 4e20 4f46 2048 4953 IMPLOSION OF HIS
000000f0: 2046 494c 4d20 434f 4d50 414e 5920 4154 FILM COMPANY AT
00000100: 2054 4845 2045 4e44 2041 4e44 2049 5420 THE END AND IT
00000110: 5741 5320 5348 4150 4544 2042 590a 5448 WAS SHAPED BY TH
00000120: 4520 454d 4552 4745 4e43 4520 4f46 204d E EMERGENCE OF M
00000130: 4554 4f4f 2054 494d 4553 2055 5020 424c ET00 TIMES UP BL
00000140: 4143 4b47 4f57 4e20 504f 4c49 5449 4353 ACKGOWN POLITICS
00000150: 2041 524d 4341 4e44 5920 4143 5449 5649 ARMCANDY ACTIVI
00000160: 534d 2041 4e44 0a41 204e 4154 494f 4e41 SM AND.A NATIONA
00000170: 4c20 434f 4e56 4552 5341 5449 4f4e 2041 L CONVERSATION A
00000180: 5320 4252 4945 4620 414e 4420 4d41 4420 S BRIEF AND MAD
00000190: 4153 2041 2046 4556 4552 2044 5245 414d AS A FEVER DREAM
```

3. Encryption using RC2 cipher in ECB (Electronic Code Block) mode

- a. Key: 365dcba26895168b7dca147895421768
- b. Iv: Not used by RC2 cipher
- c. Encryption

```
[09/09/23]seed@VM:~/.../Files$ openssl enc -rc2-ecb -e -in plain4.txt -out cipher_rc2_ecb.bin \
> -K 365dcba26895168b7dca147895421768
[09/09/23]seed@VM:~/.../Files$ xxd cipher_rc2_ecb.bin
00000000: e125 47f8 85ec ebce f5f7 2eb1 3828 8bdd %.G.....8(..%
00000010: 0399 60ed 0e43 2343 6121 2a7e 7d7b 7aa9 ...`..C#Ca!*~}{z.
00000020: 612e cfb9 a7f2 3233 8f27 3252 5fcc c2d2 a....23.'2R_...
00000030: 8a4e 0b89 f0b6 bb9c 5b11 08b8 d012 c5c9 .N.....[.....
00000040: a4b5 12bb 72dc c950 dae3 0c7b 2e10 b18b .....r..P....{....
00000050: db17 fe3c aa28 a67a 5212 fc4c 1615 d01f ...<.(.zR..L....
00000060: eac7 2b8b 960e 86aa 6c3d ce6d 6520 c71a ..+....l=.me ..
00000070: c856 ee7c 75d8 04a1 8065 bd1b da8b dbd1 .V.|u....e.....
00000080: 8e49 670d 995b af14 7843 54c8 2093 dd13 .Ig..[..xCT. ....
00000090: 4e8b 45bd bbe4 d90d 573d 9b9b 88f9 870f N.E....W=.....
000000a0: d143 6e9f 023d 35cb a0b8 2c53 8cf0 7e1d .Cn..=5...,S..~.
000000b0: accd 27f7 f03e 784a c657 21e8 d0d9 1e6a ...'..>xJ.W!....j
000000c0: b2e9 c167 2fcf c988 3dbc bfc3 326f 5b38 ...g,...=.2o[8
000000d0: c42b 8bc6 cd21 d29a 0dd0 432a cf2f 246e .+....!....C*./$n
000000e0: b514 1fde e350 a109 95e4 e2ae c2e5 adc1 .....P.....
000000f0: 157c 335c 35e8 ed9b 647a 979d 07f2 25d4 .|3\5...dz....%.
00000100: 07b6 5b77 bb87 270a d7fa 1b5e 770c 2135 ..[w..'....^w..!5
00000110: 02ac 1d46 e35b c54e 8587 2f6c c016 0d25 ...F.[.N../l...%
00000120: bed7 d9a3 a944 2473 9e34 5324 eb90 eb19 .....D$.4S$....
00000130: 268a 127c 0e52 a7c8 5547 588d 563a 61f4 &...|.R..UGX.V:a.
00000140: 1a06 4c5b 913c 7dfb 39c7 ab65 e5b4 aba7 ..L[.<}.9..e....
00000150: eedc c4e6 d712 243e 5886 9530 af45 504e .....$>X..0.EPN
00000160: e822 acb9 12cd 5702 8d8e 2e73 8a4e 1922 ."....W....s.N."
00000170: 5330 e650 3d6d 5062 882b 3066 04cc e3e3 S0.P=mPb.+0f....
```

d. Decryption

```
[09/09/23]seed@VM:~/.../Files$ openssl enc -rc2-ecb -d -in cipher_rc2_ecb.bin -out decrypt_rc2_ecb.bin \
> -K 365dcba26895168b7dca147895421768
[09/09/23]seed@VM:~/.../Files$ xxd decrypt_rc2_ecb.bin
00000000: 5448 4520 4f53 4341 5253 2054 5552 4e20 THE OSCARS TURN
00000010: 204f 4e20 5355 4e44 4159 2057 4849 4348 ON SUNDAY WHICH
00000020: 2053 4545 4d53 2041 424f 5554 2052 4947 SEEMS ABOUT RIG
00000030: 4854 2041 4654 4552 2054 4849 5320 4c4f HT AFTER THIS LO
00000040: 4e47 2053 5452 414e 4745 0a41 5741 5244 NG STRANGE.AWARD
00000050: 5320 5452 4950 2054 4845 2042 4147 4745 S TRIP THE BAGGE
00000060: 5220 4645 454c 5320 4c49 4b45 2041 204e R FEELS LIKE A N
00000070: 4f4e 4147 454e 4152 4941 4e20 544f 4f0a ONAGENARIAN TOO.
00000080: 0a54 4845 2041 5741 5244 5320 5241 4345 .THE AWARDS RACE
00000090: 2057 4153 2042 4f4f 4b45 4e44 4544 2042 WAS BOOKENDED B
000000a0: 5920 5448 4520 4445 4d49 5345 204f 4620 Y THE DEMISE OF
000000b0: 4841 5256 4559 2057 4549 4e53 5445 494e HARVEY WEINSTEIN
000000c0: 2041 5420 4954 5320 4f55 5453 4554 0a41 AT ITS OUTSET.A
000000d0: 4e44 2054 4845 2041 5050 4152 454e 5420 ND THE APPARENT
000000e0: 494d 504c 4f53 494f 4e20 4f46 2048 4953 IMPLOSION OF HIS
000000f0: 2046 494c 4d20 434f 4d50 414e 5920 4154 FILM COMPANY AT
00000100: 2054 4845 2045 4e44 2041 4e44 2049 5420 THE END AND IT
00000110: 5741 5320 5348 4150 4544 2042 590a 5448 WAS SHAPED BY.TH
00000120: 4520 454d 4552 4745 4e43 4520 4f46 204d E EMERGENCE OF M
00000130: 4554 4f4f 2054 494d 4553 2055 5020 424c ET00 TIMES UP BL
00000140: 4143 4b47 4f57 4e20 504f 4c49 5449 4353 ACKGOWN POLITICS
00000150: 2041 524d 4341 4e44 5920 4143 5449 5649 ARMCANDY ACTIVI
00000160: 534d 2041 4e44 0a41 204e 4154 494f 4e41 SM AND.A NATIONA
00000170: 4c20 434f 4e56 4552 5341 5449 4f4e 2041 L CONVERSATION A
00000180: 5320 4252 4945 4620 414e 4420 4d41 4420 S BRIEF AND MAD
00000190: 4153 2041 2046 4556 4552 2044 5245 414d AS A FEVER DREAM
000001a0: 2041 424f 5554 2057 4845 5448 4552 2054 ABOUT WHETHER T
```

Task-3: Encryption Mode – ECB vs. CBC

ECB (Electronic Code Block) and CBC (Cipher Block Chaining) are the two major modes of encryption. The main difference between ECB and CBC mode is that ECB encrypts each block independently whereas CBC encrypts each block with the previous block. Therefore, CBC mode is considered more secure than ECB. (Ref: <https://www.linkedin.com/pulse/what-ecb-cbc-encryption-edward-kiledjian/>)

Below is the screenshot of the original picture included in the Lab setup before encryption:



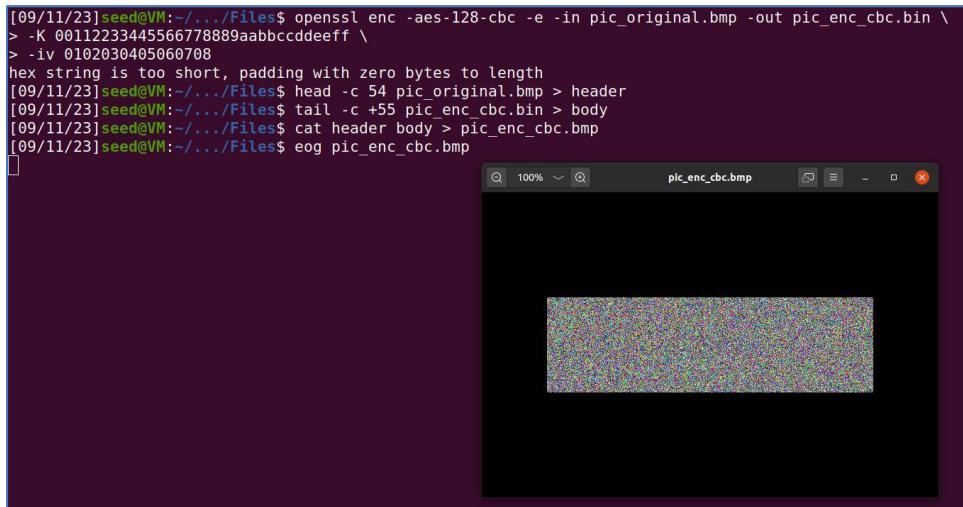
Now, I encrypt the above picture with AES-128 bit cipher in both ECB and CBC mode and replace headers in the encrypted content so that they can be loaded into the image viewing tool.

Key: 00112233445566778889aabbcdddeeff

- **CBC mode**

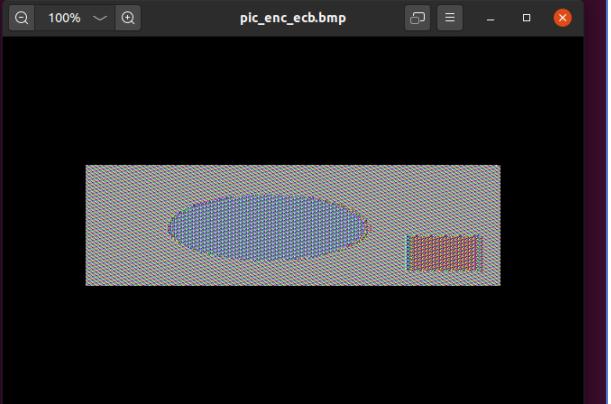
CBC mode also requires IV (Initialization Vector) for the purpose of encryption. Here, I take Iv to be “0102030405060708”.

Displaying the encrypted image content:



- **ECB Mode:**

```
[09/11/23] seed@VM:~/.../Files$ openssl enc -aes-128-ecb -e -in pic_original.bmp -out pic_enc_ecb.bin \
> -K 00112233445566778889aabbcdddeeff
[09/11/23] seed@VM:~/.../Files$ head -c 54 pic_original.bmp > header
[09/11/23] seed@VM:~/.../Files$ tail -c +55 pic_enc_ecb.bin > body
[09/11/23] seed@VM:~/.../Files$ cat header body > pic_enc_ecb.bmp
[09/11/23] seed@VM:~/.../Files$ eog pic_enc_ecb.bmp
```



Observation:

- From the above screenshots, we can observe that patterns from the original image are still visible in the encrypted content in the case of ECB mode while it is completely scrambled in the case of CBC mode.
- Since ECB mode encrypts each block in the plaintext independently using the same cipher key. So, if any block is repeated in the plaintext as we see in the icon above, we can see the same pattern in the encrypted content since the same key is used to encrypt each block. Therefore, in many cases, it is possible to break the ciphertext when encrypted in ECB mode by using **pattern recognition**.
- On the other side, the CBC mode uses the previous block for encrypting the present block. CBC encrypts each block with the same key but after encryption, it also applies XOR on the encrypted block with the previous block, and therefore, even if any block is repeated in the plaintext, the same pattern will not be repeated in the ciphertext because of the XOR operation involved.
- CBC also randomizes the plaintext by adding the Initial Vector (IV) supplied with the key at the beginning of the plaintext.
- Therefore, we can infer that CBC mode is much more secure than ECB mode.

Task-4: Padding

Padding may be required when the size of the plaintext is not a multiple of the block size for certain cipher modes.

For the demonstration purpose, I created three files of sizes 5 bytes, 10 bytes, and 16 bytes respectively.

```
[09/11/23] seed@VM:~/.../Padding$ echo -n "12345" > 5bytes.txt
[09/11/23] seed@VM:~/.../Padding$ echo -n "123456789a" > 10bytes.txt
[09/11/23] seed@VM:~/.../Padding$ echo -n "123456789abcdef0" > 16bytes.txt
[09/11/23] seed@VM:~/.../Padding$ █
```

I will use the AES-128-bit cipher for encryption having Key:

00112233445566778889aabcccddeeff and Initialization Vector: **0102030405060708**.

Encryption and Decryption of 5-byte file:

1. ECB

```
[09/11/23] seed@VM:~/.../Padding$ openssl enc -aes-128-ecb -e -in 5bytes.txt -out 5bytes_ecb.bin \
> -K 00112233445566778889aabcccddeeff
[09/12/23] seed@VM:~/.../Padding$ openssl enc -aes-128-ecb -d -nopad -in 5bytes_ecb.bin -out 5bytes_ecb_dec.bin \
> -K 00112233445566778889aabcccddeeff
[09/12/23] seed@VM:~/.../Padding$ hexdump -C 5bytes_ecb_dec.bin
00000000  31 32 33 34 35 0b |12345.....|
00000010
```

Observation: ECB added 11 bytes as padding.

2. CBC

Using IV: 0102030405060708

```
[09/12/23] seed@VM:~/.../Padding$ openssl enc -aes-128-cbc -e -in 5bytes.txt -out 5bytes_cbc.bin \
> -K 00112233445566778889aabcccddeeff \
> -iv 0102030405060708
hex string is too short, padding with zero bytes to length
[09/12/23] seed@VM:~/.../Padding$ openssl enc -aes-128-cbc -d -nopad -in 5bytes_cbc.bin -out 5bytes_cbc_dec.bin \
> -K 00112233445566778889aabcccddeeff \
> -iv 0102030405060708
hex string is too short, padding with zero bytes to length
[09/12/23] seed@VM:~/.../Padding$ hexdump -C 5bytes_cbc_dec.bin
00000000  31 32 33 34 35 0b |12345.....|
00000010
```

Observation: CBC added 11 bytes as padding

3. CFB

Using IV: 0102030405060708

```
[09/12/23] seed@VM:~/.../Padding$ openssl enc -aes-128-cfb -e -in 5bytes.txt -out 5bytes_cfb.bin \
> -K 00112233445566778889aabcccddeeff \
> -iv 0102030405060708
hex string is too short, padding with zero bytes to length
[09/12/23] seed@VM:~/.../Padding$ openssl enc -aes-128-cfb -d -nopad -in 5bytes_cfb.bin -out 5bytes_cfb_dec.bin \
> -K 00112233445566778889aabcccddeeff \
> -iv 0102030405060708
hex string is too short, padding with zero bytes to length
[09/12/23] seed@VM:~/.../Padding$ hexdump -C 5bytes_cfb_dec.bin
00000000  31 32 33 34 35 0b |12345|
00000005
```

Observation: No padding bytes are added

4. OFB

```
[09/12/23]seed@VM:~/.../Padding$ openssl enc -aes-128-ofb -e -in 5bytes.txt -out 5bytes_ofb.bin \
> -K 00112233445566778889aabbccddeeff \
> -iv 0102030405060708
hex string is too short, padding with zero bytes to length
[09/12/23]seed@VM:~/.../Padding$ openssl enc -aes-128-ofb -d -nopad -in 5bytes_ofb.bin -out 5bytes_ofb_dec.bin \
> -K 00112233445566778889aabbccddeeff \
> -iv 0102030405060708
hex string is too short, padding with zero bytes to length
[09/12/23]seed@VM:~/.../Padding$ hexdump -C 5bytes_ofb_dec.bin
00000000  31 32 33 34 35                           |12345|
00000005
```

Observation: No padding bytes are added

Encryption and Decryption of 10-byte file:

1. ECB

```
[09/12/23]seed@VM:~/.../Padding$ openssl enc -aes-128-ecb -e -in 10bytes.txt -out 10bytes_ecb.bin \
> -K 00112233445566778889aabbccddeeff
[09/12/23]seed@VM:~/.../Padding$ openssl enc -aes-128-ecb -d -nopad -in 10bytes_ecb.bin -out 10bytes_ecb_dec.bin \
> -K 00112233445566778889aabbccddeeff
[09/12/23]seed@VM:~/.../Padding$ hexdump -C 10bytes_ecb_dec.bin
00000000  31 32 33 34 35 36 37 38 39 61 06 06 06 06 06 06 |123456789a.....|
00000010
```

Observation: ECB added 6 bytes as padding

2. CBC

```
[09/12/23]seed@VM:~/.../Padding$ openssl enc -aes-128-cbc -e -in 10bytes.txt -out 10bytes_cbc.bin \
> -K 00112233445566778889aabbccddeeff \
> -iv 010203040506
hex string is too short, padding with zero bytes to length
[09/12/23]seed@VM:~/.../Padding$ openssl enc -aes-128-cbc -d -nopad -in 10bytes_cbc.bin -out 10bytes_cbc_dec.bin \
> -K 00112233445566778889aabbccddeeff \
> -iv 010203040506
hex string is too short, padding with zero bytes to length
[09/12/23]seed@VM:~/.../Padding$ hexdump -C 10bytes_cbc_dec.bin
00000000  31 32 33 34 35 36 37 38 39 61 06 06 06 06 06 06 |123456789a.....|
00000010
```

Observation: CBC added 6 bytes as padding

3. CFB

```
[09/12/23]seed@VM:~/.../Padding$ openssl enc -aes-128-cfb -e -in 10bytes.txt -out 10bytes_cfb.bin \
> -K 00112233445566778889aabbccddeeff \
> -iv 010203040506
hex string is too short, padding with zero bytes to length
[09/12/23]seed@VM:~/.../Padding$ openssl enc -aes-128-cfb -d -nopad -in 10bytes_cfb.bin -out 10bytes_cfb_dec.bin \
> -K 00112233445566778889aabbccddeeff \
> -iv 010203040506
hex string is too short, padding with zero bytes to length
[09/12/23]seed@VM:~/.../Padding$ hexdump -C 10bytes_cfb_dec.bin
00000000  31 32 33 34 35 36 37 38 39 61                           |123456789a|
00000001
```

Observation: No padding bytes are added by CFB

4. OFB

```
[09/12/23]seed@VM:~/.../Padding$ openssl enc -aes-128-ofb -e -in 10bytes.txt -out 10bytes_ofb.bin \
> -K 00112233445566778889aabcccddeeff \
> -iv 010203040506
hex string is too short, padding with zero bytes to length
[09/12/23]seed@VM:~/.../Padding$ openssl enc -aes-128-ofb -d -nopad -in 10bytes_ofb.bin -out 10bytes_ofb_dec.bin \
> -K 00112233445566778889aabcccddeeff \
> -iv 010203040506
hex string is too short, padding with zero bytes to length
[09/12/23]seed@VM:~/.../Padding$ hexdump -C 10bytes_ofb_dec.bin
00000000 31 32 33 34 35 36 37 38 39 61 |123456789a|
0000000a
```

Observation: No padding bytes are added by OFB

Encryption and Decryption of 16-byte file:

1. ECB

```
[09/12/23]seed@VM:~/.../Padding$ openssl enc -aes-128-ecb -e -in 16bytes.txt -out 16bytes_ecb.bin \
> -K 00112233445566778889aabcccddeeff
[09/12/23]seed@VM:~/.../Padding$ openssl enc -aes-128-ecb -d -nopad -in 16bytes_ecb.bin -out 16bytes_ecb_dec.bin \
> -K 00112233445566778889aabcccddeeff
[09/12/23]seed@VM:~/.../Padding$ hexdump -C 16bytes_ecb_dec.bin
00000000 31 32 33 34 35 36 37 38 39 61 62 63 64 65 66 30 |123456789abcdef0|
00000010 10 10 10 10 10 10 10 10 10 10 10 10 10 10 10 10 |.....|
00000020
```

Observation: ECB added 16 bytes of padding

2. CBC

```
[09/12/23]seed@VM:~/.../Padding$ openssl enc -aes-128-cbc -e -in 16bytes.txt -out 16bytes_cbc.bin \
> -K 00112233445566778889aabcccddeeff \
> -iv 010203040506
hex string is too short, padding with zero bytes to length
[09/12/23]seed@VM:~/.../Padding$ openssl enc -aes-128-cbc -d -nopad -in 16bytes_cbc.bin -out 16bytes_cbc_dec.bin \
> -K 00112233445566778889aabcccddeeff \
> -iv 010203040506
hex string is too short, padding with zero bytes to length
[09/12/23]seed@VM:~/.../Padding$ hexdump -C 16bytes_cbc_dec.bin
00000000 31 32 33 34 35 36 37 38 39 61 62 63 64 65 66 30 |123456789abcdef0|
00000010 10 10 10 10 10 10 10 10 10 10 10 10 10 10 10 10 |.....|
00000020
```

Observation: CBC added 16 bytes as padding

3. CFB

```
[09/12/23]seed@VM:~/.../Padding$ openssl enc -aes-128-cfb -e -in 16bytes.txt -out 16bytes_cfb.bin \
> -K 00112233445566778889aabcccddeeff \
> -iv 010203040506
hex string is too short, padding with zero bytes to length
[09/12/23]seed@VM:~/.../Padding$ openssl enc -aes-128-cfb -d -nopad -in 16bytes_cfb.bin -out 16bytes_cfb_dec.bin \
> -K 00112233445566778889aabcccddeeff \
> -iv 010203040506
hex string is too short, padding with zero bytes to length
[09/12/23]seed@VM:~/.../Padding$ hexdump -C 16bytes_cfb_dec.bin
00000000 31 32 33 34 35 36 37 38 39 61 62 63 64 65 66 30 |123456789abcdef0|
00000010
```

Observation: CFB has not added any padding bytes

4. OFB

```
[09/12/23] seed@VM:~/.../Padding$ openssl enc -aes-128-ofb -e -in 16bytes.txt -out 16bytes_ofb.bin \
> -K 00112233445566778889aabcccddeeff \
> -iv 010203040506
hex string is too short, padding with zero bytes to length
[09/12/23] seed@VM:~/.../Padding$ openssl enc -aes-128-ofb -d -nopad -in 16bytes_ofb.bin -out 16bytes_ofb_dec.bin \
> -K 00112233445566778889aabcccddeeff \
> -iv 010203040506
hex string is too short, padding with zero bytes to length
[09/12/23] seed@VM:~/.../Padding$ hexdump -C 16bytes_ofb_dec.bin
00000000  31 32 33 34 35 36 37 38  39 61 62 63 64 65 66 30  |123456789abcdef0|
00000010
```

Observation: OFB has not added any padding bytes

Overall Observation of Padding bytes added:

Cipher Mode	5 Bytes Input	10 Bytes Input	16 Bytes Input
ECB	11 bytes	6 bytes	16 bytes
CBC	11 bytes	6 bytes	16 bytes
CFB	No padding	No padding	No padding
OFB	No padding	No padding	No padding

Explanation:

- ECB and CBC modes encrypt data in fixed-size blocks. For AES cipher, the block size is always 128 bits or 16 bytes. If the size of the data is not a multiple of 16, then both ECB and CBC pad the data with the remainder number of bytes when the data size is divided by 16. The bytes used for padding will be the value of the remainder obtained.
- When the size of the data is a multiple of 16, ECB and CBC modes still add 16 bytes of padding to the plaintext. The reason is that when the blocks are full and no padding is applied, it becomes impossible to identify whether the trailing bytes are part of plaintext or padding. (For Eg: If the last 2 bytes of the plaintext contain 0x02, it is impossible to interpret if those 2 bytes belong to data or padding).
- CFB and OFB modes do not require padding because they XOR the final block of the plaintext with the first few bytes of the Keystream (generated by combining Key and Initialization Vector). These two modes produce ciphertext which is of the same size as the plaintext.

Task-5: Error Propagation – Corrupted Cipher Text

I am going to choose the below plaintext with 100 bytes size for the purpose of demonstration.

```
[09/14/23] seed@VM:~/.../Corruption$ cat plain.txt
THE OSCARS TURN ON SUNDAY WHICH SEEMS ABOUT RIGHT AFTER THIS LONG STRANGE
AWARDS TRIP THE BAGGER FEELS LIKE A NONAGENARIAN TOO

THE AWARDS RACE WAS BOOKENDED BY THE DEMISE OF HARVEY WEINSTEIN AT ITS OUTSET
AND THE APPARENT IMPLOSION OF HIS FILM COMPANY AT THE END AND IT WAS SHAPED BY
THE EMERGENCE OF METOO TIMES UP BLACKGOWN POLITICS ARMCANDY ACTIVISM AND
A NATIONAL CONVERSATION AS BRIEF AND MAD AS A FEVER DREAM ABOUT WHETHER THERE
OUGHT TO BE A PRESIDENT WINFREY THE SEASON DIDNT JUST SEEM EXTRA LONG IT WAS
EXTRA LONG BECAUSE THE OSCARS WERE MOVED TO THE FIRST WEEKEND IN MARCH TO
AVOID CONFLICTING WITH THE CLOSING CEREMONY OF THE WINTER OLYMPICS THANKS
PYEONGCHANG

ONE BIG QUESTION SURROUNDING THIS YEARS ACADEMY AWARDS IS HOW OR IF THE
CEREMONY WILL ADDRESS METOO ESPECIALLY AFTER THE GOLDEN GLOBES WHICH BECAME
A JUBILANT COMINGOUT PARTY FOR TIMES UP THE MOVEMENT SPEARHEADED BY
POWERFUL HOLLYWOOD WOMEN WHO HELPED RAISE MILLIONS OF DOLLARS TO FIGHT SEXUAL
HARASSMENT AROUND THE COUNTRY
```

I am encrypting the above content using 4 cipher modes ECB, CBC, CFB, and OFB as mentioned in the task, and changing the single bit of 0x55th byte in the ciphertext using bless hex editor.

I will use the AES-128-bit cipher for encryption having Key:

00112233445566778889aabcccddeeff and Initialization Vector: **0102030405060708**.

Encryption: (EDB, CBC, CFB, and OFB modes):

```
[09/14/23] seed@VM:~/.../Corruption$ openssl enc -aes-128-ecb -e -in plain.txt -out cipher_ecb.bin \
> -K 00112233445566778889aabcccddeeff
[09/14/23] seed@VM:~/.../Corruption$ openssl enc -aes-128-cbc -e -in plain.txt -out cipher_cbc.bin \
> -K 00112233445566778889aabcccddeeff \
> -iv 0102030405060708
hex string is too short, padding with zero bytes to length
[09/14/23] seed@VM:~/.../Corruption$ openssl enc -aes-128-cfb -e -in plain.txt -out cipher_cfb.bin \
> -K 00112233445566778889aabcccddeeff \
> -iv 0102030405060708
hex string is too short, padding with zero bytes to length
[09/14/23] seed@VM:~/.../Corruption$ openssl enc -aes-128-ofb -e -in plain.txt -out cipher_ofb.bin \
> -K 00112233445566778889aabcccddeeff \
> -iv 0102030405060708
hex string is too short, padding with zero bytes to length
```

Now, I will corrupt the single bit in the 0x55th byte in the ciphertext by changing the single bit of the byte and analyze the decrypted text for all 4 modes.

1. ECB Mode:

Before Corruption:

cipher_ecb.bin	
00000000	EC 3D 8D A4 F1 79 BA E7 01 29 88 0C 7E FB EB 98 D8 40 B0 7B 21 92 D9 23 38 C7 EC 88 C3 83 B0 8F C2 E2 A5 83 69 FE 6A B0 7F 77 8E .=...y...). .
0000002b	64 64 DB 7E CB 09 25 85 AE 08 D8 05 E3 27 9A AE 74 C7 8A CC AF C4 DB E1 A0 90 6A 5B 05 91 FB 08 A1 B2 79 D8 E3 7C EC 3F 6F 45 D5 dd..~..%....
00000056	47 04 C5 95 02 F3 FF 14 88 5C C7 FA 63 4A 67 A1 40 62 85 E9 92 F7 A8 B8 4A 1F E9 12 7F FD EC 7B 27 46 20 AB A8 79 2A DB 08 5C AC G.....\..
00000081	B0 71 14 7B E5 76 B0 B9 E8 F8 63 6E ED 68 D4 46 90 68 57 97 2B C4 24 90 D2 06 AB 5A 4B 92 CF 0E 09 AC 95 84 94 60 0E 83 B0 DC 1B .q.(v....cn
000000ac	24 94 BF AD 4C 02 B6 2A C6 8D 49 55 D0 73 FD 9F D7 FF D8 22 59 28 22 BA F4 E3 13 2D 05 E6 23 A9 45 CO 3C BB 63 78 23 70 30 F5 B7 \$...L..*..IU
000000d7	9F 83 45 93 DA 21 D2 E2 23 33 21 99 CE B5 00 91 78 17 6D 38 3E D9 27 FD 7B 7A OA C4 2B 56 6F 0D BF B5 B1 CC 95 57 C1 85 86 C3 AA ..E.!..#!..
00000102	94 4B 13 14 9F 29 83 98 4C A9 4B 9E D2 45 31 EB 82 38 A6 13 AB 4B 1B 8E AA 12 DD C6 26 52 4D FE 4B A2 8E 57 C6 A6 96 51 D2 40 3A .K...)..L.K.
0000012d	62 9A E3 64 F5 24 3F FF 65 37 43 D3 E2 68 B2 71 CE D9 8E A2 FD B3 C1 7B DD 93 D3 F2 42 B9 02 88 32 34 5E 3B FD 88 F0 06 b..d..~?..e7C
00000158	83 D9 CD 26 CF B2 EA 09 96 EB A7 37 19 44 F9 9E F7 10 E6 5A C6 0D 86 91 FA 2E A8 76 AF 98 70 0B 39 A8 BD BB DE 57 47 05 37 97 D6 ...&....7
00000183	73 44 FC 13 F3 D5 4B BC 09 12 45 03 FA BF E1 B9 71 76 41 A2 EC 17 7A 6B 6F 14 32 FD D0 99 7A 03 C7 4E EE C5 30 83 08 D0 42 16 sD....K..E..
000001ae	A6 9E 7F 19 60 87 40 8E FE 3A 75 53 86 2A F7 49 47 32 06 8D 7A 50 B2 F8 11 94 6B C3 50 30 20 DA C8 7F 08 85 3F D2 92 21 90 DF CD `..@..:us

The binary representation of the byte **D5** which is at 0x55th position is 11010101. By changing the last bit to 0, I get 11010100 which is **D4**.

Assumption before conducting the task:

- Since ECB encrypts each block of data independently with the same key, only the block containing the corrupted byte must be affected.

After Corruption:

cipher_ecb.bin	
00000000	EC 3D 8D A4 F1 79 BA E7 01 29 88 0C 7E FB EB 98 D8 40 B0 7B 21 92 D9 23 38 C7 EC 88 C3 83 B0 8F C2 E2 A5 83 69 FE 6A B0 7F 77 8E .=...y...). .
0000002b	64 64 DB 7E CB 09 25 85 AE 08 D8 05 E3 27 9A AE 74 C7 8A CC AF C4 DB E1 A0 90 6A 5B 05 91 FB 08 A1 B2 79 D8 E3 7C EC 3F 6F 45 D4 dd..~..%....
00000056	47 04 C5 95 02 F3 FF 14 88 5C C7 FA 63 4A 67 A1 40 62 85 E9 92 F7 A8 B8 4A 1F E9 12 7F FD EC 7B 27 46 20 AB A8 79 2A DB 08 5C AC G.....\..
00000081	B0 71 14 7B E5 76 B0 B9 E8 F8 63 6E ED 68 D4 46 90 68 57 97 2B C4 24 90 D2 06 AB 5A 4B 92 CF 0E 09 AC 95 84 94 60 0E 83 B0 DC 1B .q.(v....cn.h.i
000000ac	24 94 BF AD 4C 02 B6 2A C6 8D 49 55 D0 73 FD 9F D7 FF D8 22 59 28 22 BA F4 E3 13 2D 05 E6 23 A9 45 CO 3C BB 63 78 23 70 30 F5 B7 \$...L..*..IU.s.
000000d7	9F 83 45 93 DA 21 D2 E2 23 33 21 99 CE B5 00 91 78 17 6D 38 3E D9 27 FD 7B 7A OA C4 2B 56 6F 0D BF B5 B1 CC 95 57 C1 85 86 C3 AA ..E.!..#!..
00000102	94 4B 13 14 9F 29 83 98 4C A9 4B 9E D2 45 31 EB 82 38 A6 13 AB 4B 1B 8E AA 12 DD C6 26 52 4D FE 4B A2 8E 57 C6 A6 96 51 D2 40 3A .K...)..L.K..El
0000012d	62 9A E3 64 F5 24 3F FF 65 37 43 D3 E2 68 B2 71 CE D9 88 A2 FD B3 C1 7B DD 93 D3 F2 F2 44 B9 02 88 32 34 5E 3B FD 88 F0 06 b..d..~?..e7C..h
00000158	83 D9 CD 26 CF B2 EA 09 96 EB A7 37 19 44 F9 9E F7 10 E6 5A C6 0D 86 91 FA 2E A8 76 AF 98 70 0B 39 A8 BD BB DE 57 47 05 37 97 D6 ...&....7.D.
00000183	73 44 FC 13 F3 D5 4B BC 09 12 45 03 FA BF E1 B9 71 76 41 A2 EC 17 7A 6B 6F 14 32 FD D0 99 7A 03 C7 4E EE C5 30 83 08 D0 42 16 sD....K..E..
000001ae	A6 9E 7F 19 60 87 40 8E FE 3A 75 53 86 2A F7 49 47 32 06 8D 7A 50 B2 F8 11 94 6B C3 50 30 20 DA C8 7F 08 85 3F D2 92 21 90 DF CD `..@..:us.*.
000001d9	87 7C AE C1 94 41 98 83 49 3C 2B CE 5F 96 B1 E5 B6 C0 65 3B 71 59 9C 5D A5 9F A5 FD AD AD C5 E3 C7 B5 38 E6 FC C1 8D A5 35 92 A..I<+..._

Now, applying Decryption to the corrupted ciphertext (ECB):

[09/14/23] seed@VM:~/.../Corruption\$ openssl enc -aes-128-ecb -d -in cipher_ecb.bin -out dec_ecb.bin -K 00112233445566778889aabccdddeff
[09/14/23] seed@VM:~/.../Corruption\$ xxd dec_ecb.bin
00000000: 5448 4520 4f53 4341 5253 2054 5552 4e20 THE OSCARS TURN
00000010: 204f 4e20 5355 4e44 4159 2057 4849 4348 ON SUNDAY WHICH
00000020: 2053 4545 4d53 2041 424f 5554 2052 4947 SEEMS ABOUT RIG
00000030: 4854 2041 4654 4552 2054 4849 5320 4c4f HT AFTER THIS LO
00000040: 4e47 2053 5452 414e 4745 0a41 5741 5244 NG STRANGE.AWARD
00000050: 85ee 91a0 e863 984b 6113 1282 101d 66c8c.Ka.....f.
00000060: 5220 4645 454c 5320 4c49 4b45 2041 204e R FEELS LIKE A N
00000070: 4f4e 4147 454e 4152 4941 4e20 544f 4f0a ONAGENARIAN TOO.
00000080: 0a54 4845 2041 5741 5244 5320 5241 4345 .THE AWARDS RACE
00000090: 2057 4153 2042 4f4f 4b45 4e44 4544 2042 WAS BOOKENDED B

Observation:

- Here, only the highlighted block of 16 bytes from 0x50 to 0x5f is corrupted in the plaintext after decryption.

- The reason is that ECB mode encrypts each block of fixed size (16 bytes in case of AES) independently using the same encryption key.
- Since each block is independently encrypted, only the block that contains the corrupted byte is affected. In this case, since we corrupted 0x55th byte, only the block of data between 0x50 to 0x5f is affected. The rest of the data is unaffected by corruption.

How much information can be recovered?

- Since only the blocks that contain corrupted bytes in ciphertext are propagated to the plain text, much of the information will be unaffected.

2. CBC Mode

Before Corruption:

```
cipher_cbc.bin
E7 59 1F 85 26 1E EE 35 A9 EC 73 0C 5E FC 08 3E 93 59 11 22 F7 79 4D 49 C9 A4 8B DC 7C AA D3 86 FE 52 E7 FE 91 74 F8 EB FE F4 CA |.Y..&..5..s.^..>.Y.
3E 01 09 6C 4D 6B D8 5C C6 57 9C BD 17 CE 94 B1 66 BA BC 76 AB 6A C2 51 61 89 BE 25 9D C8 52 5D 43 DB DE D3 52 22 15 64 C1 DC |.80|6..1MK.\.W....f..
D2 C2 5A 98 F4 E5 07 35 70 33 AE 7C E1 25 F9 3C 95 24 E6 B3 1F 62 0B CB 1C F3 A1 99 4B 51 4C B5 C4 8A 8D CF 98 40 03 6B 6D 22 D2 |..Z...5p3.|.%<.$.
00000081 8B 95 FB EB 72 1F 3A 2B 58 ED 47 8E 18 71 E6 18 89 15 87 C1 C4 80 79 1C 52 00 84 C5 07 18 9B C5 9E F5 99 B9 7C C7 44 78 F3 D4 02 |...r.:X.G..q.|.
000000ac 55 68 AA AB 3D 7D 95 07 69 82 5B 7D 8C 03 A6 53 36 71 74 68 7B 5F 23 71 0D 6D BD B7 89 8F 11 0D 67 AA CD 91 6A C9 62 B8 A3 44 F3 |0h..=|.i.|...Sgqt
000000d7 DD 2E 77 BD 34 79 2F 3C 31 88 55 65 F0 1A 1A 59 79 1C 7D E6 F8 38 28 2C DC DE C7 AB B9 3B 38 1F 5D 75 08 3D B1 C8 83 96 AD CO B8 |..w.4y/<1.Ue...y.|
00000102 84 19 8E 01 80 6D 82 6D 11 B7 5A BE F1 95 B2 98 4B B2 69 6E 0E 7A 00 D5 BB 5E EO 34 2D 26 B4 41 BF 20 A8 21 77 29 D8 81 4A FE |....m.m.Z....R.i
B1 B0 95 A1 E7 66 90 FF 5B 8E E2 EC D2 EE B3 E9 6F C3 AF 45 EB F1 A2 F9 F6 F5 08 76 0F 4D C9 D1 72 9D 2A 63 4B A0 78 2B A9 46 EB |....f.[.....o.|
3F C3 9A EF 98 40 1F AA 35 A8 29 1D 7C 1F 2D 5C F7 CD 5d 42 9E 63 0B 6D 96 B0 3E EO 43 F3 B7 58 6C 97 20 B1 CA E7 6A A9 1E 87 54 |?...@..5.)|.-\..|
00000158 5F AC AA 0B 53 C2 0B OC 93 24 0F 58 25 60 96 58 A4 A6 EO D5 24 8B 5A 4F 1F CO 4A 84 28 89 58 94 49 9A A5 B2 38 77 96 85 33 59 EE |...S...$.X%^.X.|
000001ae 45 14 03 8E 84 B6 20 9B CE CF B6 44 B1 22 89 5C 4A 85 41 3B F9 F4 4A D8 EB 19 C7 8D CB AF F5 91 18 23 8C 84 E8 9A 46 43 55 91 D4 E.....D."|.J.A
000001d9 B1 4C F6 33 C1 F3 B8 09 B2 0F 15 FF 7F 6E C7 17 30 FF 37 96 DC 40 D5 2F 1A 5D C4 EF A6 C2 D6 BC 84 12 A9 0D 3C 02 61 DA 62 CF |.L.3.....n..0
B1 4C F6 33 C1 F3 B8 09 B2 0F 15 FF 7F 6E C7 17 30 FF 37 96 DC 40 D5 2F 1A 5D C4 EF A6 C2 D6 BC 84 12 A9 0D 3C 02 61 DA 62 CF |.L.3.....n..0
```

The binary representation of the byte **81** which is at 0x55th position is 10000001. By changing the last bit to 0, I get 10000000 which is **80**.

Assumption before conducting the task:

- In CBC, only the block containing the corrupted byte will be affected because the entire corrupted cipher block is decrypted with the same key. Also, the block that follows the cipher block containing the corrupted byte will be affected since it is XORed with the corrupted cipher block.

After Corruption:

```
cipher_cbc.bin
E7 59 1F 85 26 1E EE 35 A9 EC 73 0C 5E FC 08 3E 93 59 11 22 F7 79 4D 49 C9 A4 8B DC 7C AA D3 86 FE 52 E7 FE 91 74 F8 EB FE F4 CA |.Y..&..5..s.^..>.Y.
3E 01 09 6C 4D 6B D8 5C C6 57 9C BD 17 CE 94 B1 66 BA BC 76 AB 6A C2 51 61 89 BE 25 9D C8 52 5D 43 DB DE D3 52 22 15 64 C1 DC |.80|6..1MK.\.W....f..
D2 C2 5A 98 F4 E5 07 35 70 33 AE 7C E1 25 F9 3C 95 24 E6 B3 1F 62 0B CB 1C F3 A1 99 4B 51 4C B5 C4 8A 8D CF 98 40 03 6B 6D 22 D2 |..Z...5p3.|.%<.$.
00000081 8B 95 FB EB 72 1F 3A 2B 58 ED 47 8E 18 71 E6 18 89 15 87 C1 C4 80 79 1C 52 00 84 C5 07 18 9B C5 9E F5 99 B9 7C C7 44 78 F3 D4 02 |...r.:X.G..q.|.
000000ac 55 68 AA AB 3D 7D 95 07 69 82 5B 7D 8C 03 A6 53 36 71 74 68 7B 5F 23 71 0D 6D BD B7 89 8F 11 0D 67 AA CD 91 6A C9 62 B8 A3 44 F3 |0h..=|.i.|...Sgqt
000000d7 DD 2E 77 BD 34 79 2F 3C 31 88 55 65 F0 1A 1A 59 79 1C 7D E6 F8 38 28 2C DC DE C7 AB B9 3B 38 1F 5D 75 08 3D B1 C8 83 96 AD CO B8 |..w.4y/<1.Ue...y.|
00000102 84 19 8E 01 80 6D 82 6D 11 B7 5A BE F1 95 B2 98 4B B2 69 6E 0E 7A 00 D5 BB 5E EO 34 2D 26 B4 41 BF 20 A8 21 77 29 D8 81 4A FE |....m.m.Z....R.i
B1 B0 95 A1 E7 66 90 FF 5B 8E E2 EC D2 EE B3 E9 6F C3 AF 45 EB F1 A2 F9 F6 F5 08 76 0F 4D C9 D1 72 9D 2A 63 4B A0 78 2B A9 46 EB |....f.[.....o.|
3F C3 9A EF 98 40 1F AA 35 A8 29 1D 7C 1F 2D 5C F7 CD 5d 42 9E 63 0B 6D 96 B0 3E EO 43 F3 B7 58 6C 97 20 B1 CA E7 6A A9 1E 87 54 |?...@..5.)|.-\..|
00000158 5F AC AA 0B 53 C2 0B OC 93 24 0F 58 25 60 96 58 A4 A6 EO D5 24 8B 5A 4F 1F CO 4A 84 28 89 58 94 49 9A A5 B2 38 77 96 85 33 59 EE |...S...$.X%^.X.|
000001ae 45 14 03 8E 84 B6 20 9B CE CF B6 44 B1 22 89 5C 4A 85 41 3B F9 F4 4A D8 EB 19 C7 8D CB AF F5 91 18 23 8C 84 E8 9A 46 43 55 91 D4 E.....D."|.J.A
000001d9 B1 4C F6 33 C1 F3 B8 09 B2 0F 15 FF 7F 6E C7 17 30 FF 37 96 DC 40 D5 2F 1A 5D C4 EF A6 C2 D6 BC 84 12 A9 0D 3C 02 61 DA 62 CF |.L.3.....n..0
B1 4E A5 BF B7 8C 37 89 2C 30 53 6B 60 16 4A 31 9C 30 95 D8 0C 09 DF 92 8A 7A AF 1A FC DC 06 C6 89 BD EB 08 BD 58 8B |.N...7.,OSK*.J1.0.
```

Now, applying Decryption to the corrupted ciphertext (CBC):

```
[09/14/23]seed@VM:~/.../Corruption$ openssl enc -aes-128-cbc -d -in cipher_cbc.bin -out dec_cbc.bin -K 00112233445566778889aabcccddeeff -iv 0102030405060708
hex string is too short, padding with zero bytes to length
[09/14/23]seed@VM:~/.../Corruption$ xxd dec_cbc.bin
00000000: 5448 4520 4f53 4341 5253 2054 5552 4e20 THE OSCARS TURN
00000010: 204f e420 5355 4e44 4159 2057 4849 4348 ON SUNDAY WHICH
00000020: 2053 4545 4d53 2041 424f 5554 2052 4947 SEEKS ABOUT RIG
00000030: 4854 2041 4654 4552 2054 4849 5320 4c4f HT AFTER THIS LO
00000040: 4e47 2053 5452 414e 4745 0a41 5741 5244 NG STRANGE AWARD
00000050: 9f5b 05ce 3d30 b516 1a1f 5611 95d9 0cae [...]=0,...V.....
00000060: 5220 4645 454d 5320 4c49 4b45 2041 204e R FEEMS LIKE A N
00000070: 4f4e 4147 454e 4152 4941 4e20 544f 4f0a ONAGENARIAN TOO.
00000080: 0a54 4845 2041 5741 5244 5320 5241 4345 .THE AWARDS RACE
```

Observation:

- In CBC Mode, the entire 6th block between 0x51 to 0x60 containing the corrupted byte at location 0x55 is corrupted since the whole block in the cipher will be decrypted by the same key.
- In the 7th block between 0x61 to 0x70, only the byte corresponding to the corrupted byte in the previous block is corrupted (Highlighted in blue). The rest of the bytes in the decrypted block are unaffected because the corrupted cipher text is only XORed with the plain text block.

How much information can be recovered?

- Only the block that contains the corrupted byte and the byte corresponding to the corrupted byte in the next block will be affected. The rest of the data is unaffected.

3. CFB Mode:

Before Corruption:

```
cipher_cfb.bin x
00000000|D3 CB CA 05 8E 6C F9 90 FD CD 9F 49 20 42 12 4D 57 D3 EA 17 CA 4B 05 14 CC 46 07 7F 59 BA A0 FA 49 99 E6 62 D8 50 F2 55 28 50 C4|....l....I B.|.
0000002b|DB 65 AA 20 FC 83 B3 EA CE 38 A4 01 98 7D 0E CF 21 12 4E F1 D2 31 97 EE 80 D0 0B 94 38 55 A3 12 B7 E0 F6 3D 69 E1 41 38 8C B1 15|...e. ....8...|.
00000056|E8 C4 31 93 13 D3 AA 8E F2 1C 00 78 39 BA B9 6B 61 AF 56 E7 43 B9 D4 93 E7 65 DC 45 99 A0 DD 12 84 06 E0 B4 C2 F7 55 E4 81 D2 5A|..1.....x9...|.
00000081|68 72 96 41 45 5A 14 FA 75 A9 E9 1E 00 40 15 49 53 6F 40 CF E5 FD C1 01 22 56 C9 32 26 EA B5 C8 OB AD E0 6B 79 13 75 04 40 16 04 or.AEZ.u...@.
000000ac|69 49 46 11 E3 BF 8E EE 58 94 6F 2D F7 6E FA 41 9C FA E3 53 27 14 DD 81 47 9D AE CB 24 CA 29 CB DD D0 06 CE 52 38 80 EE E3 99 ii.F....X.o.-n
000000d7|7B D2 EC E1 72 44 DA C8 BE C6 01 63 4B BB F4 E0 37 00 60 74 28 84 FD E9 0F 33 99 DD 6B 33 7D 1C 60 D1 05 7D F1 F8 DF 3F 15 20 AA (...rD....ck..|
00000102|72 24 F1 95 1B 8D 21 0D AC 01 0E D6 96 29 D7 0B 4E 46 F7 4E 7C AD 9C 1D BC D4 48 2C 6D 8E 81 0F 3E D2 4F 52 A5 BA 66 C1 0F 35 C2 rs$....!....)
0000012d|27 32 9F 89 D2 9A 3A 58 E1 15 F3 C6 77 57 8D 2F F7 FE 6E 4C A8 D0 A3 73 C1 41 92 23 22 FB B7 C7 CF 39 A3 3D 98 CF 6B 90 B4 CU 14 '2....X....ww.
00000158|6D 11 7F 40 E9 21 4A 70 04 22 D7 BD 12 1B 71 26 C6 56 2C DB 1B B5 87 04 C1 AF 53 00 E8 41 OC FB 93 BA A6 2C 5F OC FB B8 AD 06 97 m..@.!p."....q
00000183|43 4A 86 51 60 DB 0F 80 74 A0 C7 64 84 FB 9E 96 3F C1 67 E9 90 02 73 26 46 03 58 1D 40 60 E8 F1 27 39 B2 2B 21 46 42 FE 02 21 53 CJ.Q'....d...
000001ae|38 6F DE C4 6A 22 54 6F 89 0A 0E 70 05 DB 63 41 32 EE F0 2A AD 1F 85 37 54 AA 15 7F 90 BA 74 60 F1 72 22 8F BD 81 AC 96 D4 1F AF 80..j"To...p..c.
000001d9|EF C5 C8 19 6D 01 BB 83 BB C9 28 DF 6F 6D A3 A9 19 D0 63 29 E1 78 79 57 F8 2C 55 44 68 82 11 32 CE 22 3C A6 49 03 92 91 OF 51 1E ...m....(.om.
```

The binary representation of the byte 15 which is at 0x55th position is 00010101. By changing the last bit to 0, I get 00010100 which is 14.

Assumption before completing the task:

- The CFB is similar to CBC. In CFB since the corrupted cipher block is decrypted with the same key, the entire plain text block will be corrupted. When the corrupted Cipher block is XORed with the decrypted Cipher block, only the byte corresponding to the corrupted byte in the previous block should be affected.

After Corruption:

```
cipher cfb.bin
00000000 D3 CE CA 05 8E 6C F9 90 FD CD 9F 49 20 42 12 4D 57 D3 EA 17 CA 4B 05 14 CC 46 07 7F 59 BA A0 FA 49 99 E6 62 D8 50 F2 55 28 50 C4|.....I.B.MW...
0000002b DB 65 AA 20 FC 83 B3 EA CE 38 A1 01 98 7D 0E CF 21 12 4F F1 D2 31 97 EE 80 D0 08 94 38 55 A3 12 B7 E0 F6 3D 69 E1 41 38 8C B1 [14].e. ....8...).!N..
00000056 EB C4 31 93 13 D3 AA 8B F2 1C 00 78 39 BA 98 6B 61 AF 50 E7 43 B9 D4 93 E7 65 DC 45 99 A0 DD 12 84 06 E0 B4 C2 FF 55 E1 81 D2 5A|..1.....x9.ka.V.C
00000081 6F 72 96 41 45 54 14 FA 75 A9 E1 1E 00 40 15 49 53 8F 40 CF E5 FD C1 01 22 56 C9 32 26 EA B5 C9 0B AD E0 6B 79 13 75 04 4C 16 04 or.AEZ..u....0.IS.0.
000000ac 69 49 CE 46 11 E3 BF 8F EE 58 94 6F 2D F7 6E FA 41 9C FA E3 53 27 14 DD 81 47 9D AE CB 24 CA 29 CB DD D4 06 CE 52 38 80 EE 3 99 11.F.....X.o..n.A..S
000000d7 7B D2 EC E1 72 44 DA C8 BE C6 01 63 4B BB F4 E0 37 00 60 74 28 84 FD E9 0F 33 99 DD 6B 33 7D 1C 60 D1 05 7D F1 F8 DF 3F 15 20 AA (...rD....cK..7..t(
00000102 72 24 F1 95 1B 8D 21 OD AC 01 08 D6 96 29 D7 0B 4E 46 F7 4E 7C AD 9C 1D BC D4 48 2C 6D 8B 81 0F 3E D2 4F 52 A5 B4 66 C1 0F 35 C2 r8.....l.....).NF.N
0000012d 27 32 9F 89 D2 9A 3A 58 E1 15 F3 C6 77 57 6D 2F F7 FE 6E 4C A8 D0 A3 73 C1 41 92 23 22 FB B7 C7 CF 39 A3 3D 98 CE 6B 90 B4 C0 14 |'2.....x1...wW./..NL.
00000158 6D 11 7E 40 E9 21 4A 70 04 22 D7 BD 12 1B 71 26 C6 56 2C D8 1B B5 87 04 C1 AF 53 00 E8 41 0C FB 93 BA A0 2C 5F 0C FB B0 AD 06 97 m..@.!Jp."....q&.V. ..
00000183 43 4A 86 51 60 0B DF 80 74 A0 C7 64 84 FB 9E 96 3F C1 67 E9 90 02 73 26 46 03 58 1D 40 60 E8 F1 27 39 B2 ZE 21 46 42 FE 02 21 53 CJ.Q ..t..d....?g..
000001ae 38 6F DE C4 6A 22 54 6F 89 OA 08 70 05 DE 63 41 32 EE F0 2A AD 1F 85 37 54 AA 15 7F 90 BA 74 60 F1 72 22 8F BD 81 AC 93 D4 1F AB 8o..j*T0..p..cA2..*.
000001d9 EF C5 C8 19 6D 01 BB 83 BB C9 28 DF 6F 6D A3 A9 19 D0 63 29 E1 78 79 57 F8 2C 55 44 68 82 11 32 CE 22 3C A6 49 03 92 91 OF 51 1E .....m....(om...c).
00000204 15 32 4C D2 D5 15 3E F8 AC 37 FE 91 C8 85 E2 C0 C7 1C AA B1 A9 05 FE 7B 0F 81 9E A7 75 37 D0 D2 47 BF EB BD C2 B3 CD E0 CF AA 92 .2L...>7....
```

Now, applying Decryption to the corrupted ciphertext (CFB):

```
[09/14/23]seed@VM:~/.../Corruption$ openssl enc -aes-128-cfb -d -in cipher_cfb.bin -out dec_cfb.bin -K 00112233445566778889aabccddeff -iv 0
102030405060708
hex string is too short, padding with zero bytes to length
[09/14/23]seed@VM:~/.../Corruption$ xxd dec_cfb.bin
00000000: 5448 4520 4f53 4341 5253 2054 5552 4e20 THE OSCARS TURN
00000010: 204f 4e20 5355 4e44 4159 2057 4849 4348 ON SUNDAY WHICH
00000020: 2053 4545 4d53 2041 424f 5554 2052 4947 SEEKS ABOUT RIG
00000030: 4854 2041 4654 4552 2054 4849 5320 4c4f HT AFTER THIS LO
00000040: 4e47 2053 5452 414e 4745 0a41 5741 5244 NG STRANGE.AWARD
00000050: 5320 5452 4951 2054 4845 2042 4147 4745 S TRIO THE BAGGE
00000060: a3a3 0383 358b 1671 c3a8 51ae a128 bd59 [.....5..a..0...Y
00000070: 4f4e 4147 454e 4152 4941 4e20 544f 4f0a ONAGENARIAN TOO.
00000080: 0a54 4845 2041 5741 5244 5320 5241 4345 .THE AWARDS RACE
00000090: 2057 4153 2042 4f4f 4b45 4e44 4544 2042 WAS BOOKENDED B
000000a0: 5920 5448 4520 4445 4d49 5345 204f 4620 Y THE DEMISE OF
```

Observation:

- When decrypted in CFB mode, we observe that the entire 6th block that contains the corrupted cipher text byte is not corrupted. Only the byte that was corrupted in the cipher text is corrupted in the plain text.
- However, we observe that the entire block that immediately follows the block (0x51 to 0x60) that contains the corrupted ciphertext byte is corrupted. The corruption has propagated to the next immediate block (0x61 to 0x70).
- This is because as the corrupted 6th block of the cipher text was decrypted using the key it will affect the corresponding plain text formed and when the same corrupted cipher block was XORed with the decrypted cipher block only the corresponding byte in the decrypted block would be corrupted.

How much information can be recovered?

- The block containing the corrupted byte will have the same byte corrupted in the plain text while the entire block that immediately follows the corrupted cipher block will be affected. The rest of the data is unaffected.

4. OFB Mode:

Before Corruption:

```
cipher_ofb.bin x
00000000 D3 CE CA 05 8E 6C F9 90 FD CD 9F 49 20 42 12 4D 11 1B D7 27 E8 0E 0C B7 B6 01 26 AB 1B 20 0B D7 F5 3E 79 24 28 36 7F 14 57 E9 3E .....l.....I.B.M....'.
0000002b D1 1D 6F EO EF 46 71 92 DB 23 C2 A0 74 1B 25 29 53 8C 5B 59 1B 33 35 FD 83 8C 35 04 B6 21 EB 4D A2 E4 80 99 0F 4F 78 34 AD 83 BD .....o..Fq..#..t.%$.^Y.
00000056 53 0A 61 B9 97 24 2A 5B 72 FD A7 7F 84 94 AC B0 30 3D 3B 4D A0 EA P3 D7 C7 6E 56 C7 15 9D 17 FF 7E EF B7 09 55 FA ED 75 68 80 UC S.a..$*[r.....0;M.
00000081 65 EF CB CF 9B 41 CD 53 C6 67 B6 OC CO AA B5 FA 18 D8 5D 7F 10 5C AD DD AF 01 CF EE 0E 10 AB A4 A2 AC 36 4A 50 5A 64 7E 1A FB EC e....A.S.g.....].
000000ac 10 F9 DD C8 E9 D7 D1 BE FB C7 98 C2 4D 09 6B CB 34 C3 07 DF 6E 44 94 A6 39 58 85 7D 6B 6C 39 13 FF 09 FB 6B 08 8F 88 4B 0C 32 E7 .....M.k.4..nL
000000d7 EA 08 87 75 84 2A D6 CE 82 74 D7 C8 8E 33 81 CC 61 C0 65 15 90 2A D4 CC 41 2F BB 03 ED 50 DB 09 1A A3 D3 1C 28 48 5A 99 FF 5D F8 .....u.*...t...3..a.e..
00000102 C4 02 EE 2D 30 2F BE 42 D5 F5 11 6D 03 67 E2 E7 2B 77 0E 04 3E 6D 75 D9 30 11 56 22 85 D8 DF BA C4 28 A3 C4 56 27 47 F4 DF 37 08 ...-0..B..m.g...>h
0000012d 0F C8 84 32 ED 75 9D 69 A4 13 2E C0 96 D5 37 B2 EC F1 A3 3C 29 EC 75 FA FE 19 07 EE 82 2C 83 2A C1 9C 98 D5 FC 8D 06 6F 72 EC 98 .....2.u.i.....7...<)
00000158 93 F3 81 78 8E F8 33 32 CO 5E B5 BC C6 C5 77 10 3F C6 1F 7F 1A B3 C5 FC D0 E1 OC F2 OC 6C 9D A3 29 C4 A2 F2 6E 96 A3 73 CD 1C 98 .....x..32.^.....w?...
00000183 90 A5 50 52 DD 07 9A 74 B2 41 6B 0E 96 A6 DF 9A 10 81 EE 5A 0D 05 76 A3 63 57 E2 AF 44 FF 28 AF 15 9F 0F 15 48 68 2E F0 7D 5B 03 ..PR..t.Ak.....Z.
000001ae 4B 1A F4 12 6E A6 AB 08 90 2B 79 51 E2 C0 15 D6 E9 8C 48 75 58 CC E7 14 C7 95 CE C1 29 73 AA 76 17 F8 F9 99 EO OF 1C 4B 4E 47 63 N...n...+yQ.....HuX
000001d9 74 E5 19 A3 E9 F1 1C A4 A6 03 94 12 FC 1C 8B C9 A3 6C 97 62 C2 AB 3C FA 96 47 47 F7 C9 46 51 72 AB 35 1A D3 27 00 FD 3A 67 E1 4D t.....l.b.
000001d9 74 E5 19 A3 E9 F1 1C A4 A6 03 94 12 FC 1C 8B C9 A3 6C 97 62 C2 AB 3C FA 96 47 47 F7 C9 46 51 72 AB 35 1A D3 27 00 FD 3A 67 E1 4D t.....l.b.
```

The binary representation of the byte **3C** which is at 0x55th position is 00111100. By changing the last bit to 1, I get 00111101 which is **3D**.

Assumption before completing the task:

- As only a single bit of the cipher block is corrupted, and therefore when the cipher block is XORed with the output block of the block cipher encryption, only a single bit of the corresponding plain text will be corrupted. The other cipher blocks will not be affected as they are not XORed with the corrupted cipher block.

After Corruption:

```
cipher_ofb.bin x
00000000 D3 CE CA 05 8E 6C F9 90 FD CD 9F 49 20 42 12 4D 11 1B D7 27 E8 0E 0C B7 B6 01 26 AB 1B 20 0B D7 F5 3E 79 24 28 36 7F 14 57 E9 3E .....l.....I.B.M....'.
0000002b D1 1D 6F EO EF 46 71 92 DB 23 C2 A0 74 1B 25 29 53 8C 5B 59 1B 33 35 FD 83 8C 35 04 B6 21 EB 4D A2 E4 80 99 0F 4F 78 34 AD 83 BD .....o..Fq..#..t.%$.^Y.
00000056 53 0A 61 B9 97 24 2A 5B 72 FD A7 7F 84 94 AC B0 30 3D 3B 4D A0 EA P3 D7 C7 6E 56 C7 15 9D 17 FF 7E EF B7 09 55 FA ED 75 68 80 UC S.a..$*[r.....0;M.
00000081 65 EF CB CF 9B 41 CD 53 C6 67 B6 OC CO AA B5 FA 18 D8 5D 7F 10 5C AD DD AF 01 CF EE 0E 10 AB A4 A2 AC 36 4A 50 5A 64 7E 1A FB EC e....A.S.g.....].
000000ac 10 F9 DD C8 E9 D7 D1 BE FB C7 98 C2 4D 09 6B CB 34 C3 07 DF 6E 44 94 A6 39 58 85 7D 6B 6C 39 13 FF 09 FB 6B 08 8F 88 4B 0C 32 E7 .....M.k.4..nL
000000d7 EA 08 87 75 84 2A D6 CE 82 74 D7 C8 8E 33 81 CC 61 C0 65 15 90 2A D4 CC 41 2F BB 03 ED 50 DB 09 1A A3 D3 1C 28 48 5A 99 FF 5D F8 .....u.*...t...3..a.e..
00000102 C4 02 EE 2D 30 2F BE 42 D5 F5 11 6D 03 67 E2 E7 2B 77 0E 04 3E 6D 75 D9 30 11 56 22 85 D8 DF BA C4 28 A3 C4 56 27 47 F4 DF 37 08 ...-0..B..m.g...>h
0000012d 0F C8 84 32 ED 75 9D 69 A4 13 2E C0 96 D5 37 B2 EC F1 A3 3C 29 EC 75 FA FE 19 07 EE 82 2C 83 2A C1 9C 98 D5 FC 8D 06 6F 72 EC 98 .....2.u.i.....7...<)
00000158 93 F3 81 78 8E F8 33 32 CO 5E B5 BC C6 C5 77 10 3F C6 1F 7F 1A B3 C5 FC D0 E1 OC F2 OC 6C 9D A3 29 C4 A2 F2 6E 96 A3 73 CD 1C 98 .....x..32.^.....w?...
00000183 90 A5 50 52 DD 07 9A 74 B2 41 6B 0E 96 A6 DF 9A 10 81 EE 5A 0D 05 76 A3 63 57 E2 AF 44 FF 28 AF 15 9F 0F 15 48 68 2E F0 7D 5B 03 ..PR..t.Ak.....Z.
000001ae 4B 1A F4 12 6E A6 AB 08 90 2B 79 51 E2 C0 15 D6 E9 8C 48 75 58 CC E7 14 C7 95 CE C1 29 73 AA 76 17 F8 F9 99 EO OF 1C 4B 4E 47 63 N...n...+yQ.....HuX
000001d9 74 E5 19 A3 E9 F1 1C A4 A6 03 94 12 FC 1C 8B C9 A3 6C 97 62 C2 AB 3C FA 96 47 47 F7 C9 46 51 72 AB 35 1A D3 27 00 FD 3A 67 E1 4D t.....l.b.
00000204 5F F4 3E 63 D6 37 19 81 E3 B4 F3 CD B8 4D 7A 5A C1 69 C6 7F 71 E7 D6 78 71 89 55 78 50 2C C9 70 D6 25 G9 60 7D 87 6F 64 F3 EC D6 ...>c.7.....Mzz.i..q.
```

Now, applying Decryption to the corrupted ciphertext (OFB):

```
[09/14/23] seed@VM:~/.../Corruption$ openssl enc -aes-128-ofb -d -in cipher_ofb.bin -out dec_ofb.bin \
> -K 00112233445566778889aabcccddeeff \
> -iv 0102030405060708
hex string is too short, padding with zero bytes to length
[09/14/23] seed@VM:~/.../Corruption$ xxd dec_ofb.bin
00000000: 5448 4520 4f53 4341 5253 2054 5552 4e20 THE OSCARS TURN
00000010: 204f 4e20 5355 4e44 4159 2057 4849 4348 ON SUNDAY WHICH
00000020: 2053 4545 4d53 2041 424f 5554 2052 4947 SEEKS ABOUT RIG
00000030: 4854 2041 4654 4552 2054 4849 5320 4c4f HT AFTER THIS LO
00000040: 4e47 2053 5452 414e 4745 0a41 5741 5244 NG STRANGE.AWARD
00000050: 5320 5452 4951 2054 4845 2042 4147 4745 S TRIQ THE BAGGE
00000060: 5220 4645 454c 5320 4c49 4b45 2041 204e R FEELS LIKE A N
00000070: 4f4e 4147 454e 4152 4941 4e20 544f 4f0a ONAGENARIAN TOO.
00000080: 0a54 4845 2041 5741 5244 5320 5241 4345 .THE AWARDS RACE
00000090: 2057 4153 2042 4f4f 4b45 4e44 4544 2042 WAS BOOKENDED B
```

Observation:

- Here, only the corrupted cipher text byte is affected. The rest of the bytes in the cipher block 0x51 to 0x60 is unaffected.

- This is because, only a single bit of the cipher block is corrupted, so therefore, when the corrupted cipher block is XORed with the output of the block encryption, only the corresponding byte will be affected.

How much information can be recovered?

- Much of the information remains intact. Only the bytes that are corrupted in the cipher text will be affected.

Task-6: Initial Vector (IV) and Common Mistakes

Task-6.1: IV Experiment

The Initialization Vector is mainly used to introduce randomness in the encryption process so that even if multiple messages containing identical plaintext encrypted with the same key generate different ciphertext.

IVs are generally sent openly along with the ciphertext. If the same IV is reused more than once, then an attacker with some background knowledge of the message may compare multiple ciphertexts encrypted using the same IV and gain important information about the message.

a) Using different IV

Consider the below plaintext:

```
[09/15/23] seed@VM:~/.../IVExperiment$ echo -n "THIS IS THE PLAINTEXT THAT IS BEING CONSIDERED FOR THE PURPOSE OF DEMONSTRATING THE CONSEQUENCES OF USING SAME IV" > plaintext
```

Now I encrypt the above plaintext in **AES-128-CBC mode** using the same key and two different Initialization vectors.

Key: 00112233445566778889aabbcdddeeff

IV-1: 0102030405060708090a

IV-2: 453fd127934701296cb4

Below is the screenshot of the encryption process generating 2 ciphers cipher1 and cipher2

```
[09/15/23]seed@VM:~/.../IVExperiment$ openssl enc -aes-128-cbc -e -in plaintext -out cipher1 \
> -K 00112233445566778889aabccddeeff \
> -iv 0102030405060708090a
hex string is too short, padding with zero bytes to length
[09/15/23]seed@VM:~/.../IVExperiment$ openssl enc -aes-128-cbc -e -in plaintext -out cipher2 \
> -K 00112233445566778889aabccddeeff \
> -iv 453fd127934701296cb4
hex string is too short, padding with zero bytes to length
[09/15/23]seed@VM:~/.../IVExperiment$ xxd cipher1
00000000: c621 e532 51c9 222d 6342 5a49 35f7 6bc7 ..!.20.-cBZ15.k.
00000010: e71b 5c3d 8895 a371 92d9 30a7 1afc a378 ..\=...q.0...x
00000020: ec7a 20b0 cdf0 a4ac df3b ade7 4c24 f162 .z .....;.L$.b
00000030: 0039 a8a2 6c68 ef4c 1e38 1261 dc22 e39a .9..lh.L.B.a."..
00000040: 5788 3875 1bcb 91a2 eef9 3019 7e70 f709 W.8u.....0.-p..
00000050: 5b45 b87f 6a92 caa8 23f9 632a 5297 b3cf [E..j...#.c*R...
00000060: 740b 5bf7 6950 9fd7 3f37 5cb5 5f05 6075 t.[.1P..?7\...u
00000070: 8084 da00 325d 055b 6e93 cc8c c75b fae3 ....2].[n....[..
[09/15/23]seed@VM:~/.../IVExperiment$ xxd cipher2
00000000: 6d6b c7e3 ac16 65d0 18a4 f327 99d0 6c36 mk....e...'l6
00000010: 15af 18bd dd91 9660 a26c d941 6188 ba9f .....`l.Aa...
00000020: 3664 6c5f e188 ae0 824d Geed a6cc 4afb 6d1.....Mn..J.
00000030: 489c cbda 56a7 cd88 02a2 9fc5 8cbe 1ad3 H...V.....
00000040: ece2 4b96 fbe2 d578 d260 2cb1 6e5f f4ea ..K....x.,.n...
00000050: acef 5486 dd5c 823d c55c 58a8 a294 1ae4 ..T....=.V. ....
00000060: 4929 0478 006e 1db9 3c9d 9b14 c074 7f6f I).x.n.<....t.o
00000070: 4e84 f8d5 faa2 02cf d81b 25c8 000e 5118 N.....%...Q.
```

Observation:

- There is little to observe here as the two different IVs generated completely different ciphertexts from the same plaintext using the same key.

b) Using Same IV

Consider the below plaintext:

```
[09/15/23]seed@VM:~/.../IVExperiment$ echo -n "THIS IS THE PLAINTEXT THAT IS BEING CONSIDERED FOR THE PURPOSE OF DEMONSTRATING THE CONSEQUENCES OF USING SAME IV" > plaintext
```

Now I encrypt the above plaintext twice in **AES-128-CBC mode** using the same key and same Initialization vectors.

Also, I slightly change the content of the plaintext message and encrypt using the same Key/IV pairs as above.

Key: 00112233445566778889aabccddeeff

IV: 0106d304057507bc090a

```
[09/15/23]seed@VM:~/.../IVExperiment$ openssl enc -aes-128-cbc -e -in plaintext -out cipher1 \
> -K 00112233445566778889aabccddeeff \
> -iv 0106d304057507bc090a
hex string is too short, padding with zero bytes to length
[09/15/23]seed@VM:~/.../IVExperiment$ openssl enc -aes-128-cbc -e -in plaintext -out cipher2 \
> -K 00112233445566778889aabccddeeff \
> -iv 0106d304057507bc090a
hex string is too short, padding with zero bytes to length
[09/15/23]seed@VM:~/.../IVExperiment$ xxd cipher1
00000000: e100 cdfd 8d78 d266 17fd a95e 6395 2caa ....x.f...^c...
00000010: 10b9 b23c 98ba 6f64 fc01 4494 cf95 88aa ...<..od..D....
00000020: 2f33 83c0 e26a 3d52 4f53 95e8 f63f 2fb4 /3...j=R0S...?/
00000030: 3b68 c717 b7ca e2d2 d5b9 f97f 50b5 3c4a ;h.....P.<]
00000040: 3b81 a021 f3c5 4ebd 55c8 f291 9c48 4a09 ;...!.N.U...H].
00000050: dc73 4dcl 369f 5675 0476 dc3a fdea 0a81 .sM.6.Vu.v. .....
00000060: a443 a984 d1a0 09c0 160f 4d3e 7bac 5fcf .C.....M>{.\.
00000070: 2caa e39e b2b5 d993 0a62 5d92 1caa d2b1 ,.....b].....
[09/15/23]seed@VM:~/.../IVExperiment$ xxd cipher2
00000000: e100 cdfc 8d78 d266 17fd a95e 6395 2caa ....x.f...^c...
00000010: 10b9 b23c 98ba 6f64 fc01 4494 cf95 88aa ...<..od..D....
00000020: 2f33 83c0 e26a 3d52 4f53 95e8 f63f 2fb4 /3...j=R0S...?/
00000030: 3b68 c717 b7ca e2d2 d5b9 f97f 50b5 3c4a ;h.....P.<]
00000040: 3b81 a021 f3c5 4ebd 55c8 f291 9c48 4a09 ;...!.N.U...H].
00000050: dc73 4dcl 369f 5675 0476 dc3a fdea 0a81 .sM.6.Vu.v. .....
00000060: a443 a984 d1a0 09c0 160f 4d3e 7bac 5fcf .C.....M>{.\.
00000070: 2caa e39e b2b5 d993 0a62 5d92 1caa d2b1 ,.....b].....
```

Observation:

- We can observe that both ciphers are identical.

Now, I slightly change the middle portions of the plaintext message and encrypt with the same Key/IV pair.

```
[09/15/23]seed@VM:~/.../IVExperiment$ cat plaintext
THIS IS THE PLAINTEXT THAT IS BEING CONSIDERED FOR THE PURPOSE OF DEMONSTRATING THE CONSEQUENCES OF USING SAME IV[09/1
periment$  
[09/15/23]seed@VM:~/.../IVExperiment$ cat plaintext2
THIS IS THE PLAINTEXT THAT IS CHANGING THE CONTENT HERE THE PURPOSE OF DEMONSTRATING THE CONSEQUENCES OF USING SAME IV
/IVExperiment$  
[09/15/23]seed@VM:~/.../IVExperiment$ openssl enc -aes-128-cbc -e -in plaintext2 -out cipher2 \
> -iv 0106d304057507bc090a
hex string is too short, padding with zero bytes to length
[09/15/23]seed@VM:~/.../IVExperiment$ xxd cipher1
00000000: e100 cdfc 8d78 d266 17fd a95e 6395 2caa .....x.f...^c.,.
00000010: 1069 b23c 98ba 6f64 fc1 4494 cf95 88aa ...<..od..D....
00000020: 2f33 83c0 e26a 3d52 4f53 95e8 f63f 2fb4 /3...j=R0S...?/.
00000030: 3b68 c717 b7ca e2d2 d5b9 f97f 50b5 3c4a ;h.....P.-J
00000040: 3b81 a021 f3c5 4ebd 55c8 f291 9c48 4a09 ;...N.U....HJ
00000050: dc73 4dc1 369f 5675 0476 dc3a fdea 0a81 .sM.6.Vu.v.v:...
00000060: a443 a984 dia0 09c0 160f 4d3e 7bac 5cf0 .C.....M>{.\.
00000070: 2caa e39e b2b5 d933 0a62 5d92 1caa d2b1 ,.....b]....
[09/15/23]seed@VM:~/.../IVExperiment$ xxd cipher2
00000000: e100 cdfc 8d78 d266 17fd a95e 6395 2caa .....x.f...^c.,.
00000010: e966 c7f4 8b49 d0b3 0838 5d5a 76be 202a ....I...8]Zv. *
00000020: 210a 9d53 e855 e5d3 cb9c df6e efa7 5def !..S.U.....n..].
00000030: 33f2 c4c0 360c f12b f12c b2d9 a31e ac65 3...6..+.....e
00000040: 963c f16b 3782 0d4a 11b8 dd85 0c8c b64b <.K7..J.....K
00000050: 516e 0174 076c 4b56 5c34 c7b6 4710 8389 Qn.t.lKV\4..G...
```

Observation:

- In this case, we can observe that both ciphers are identical up to the block where the plaintext message is not modified. Here, the first block in both ciphertexts is identical.
- The rest of the blocks of ciphertext starting from the block where the plaintext message is modified are changing despite the portions of plaintexts being identical because CBC mode applies XOR to the plaintext with the output of the previous block cipher encryption before applying encryption.

Why IV needs to be unique?

From the above observations, we can conclude that the identical plaintext messages encrypted with the same Key/IV pair produce identical ciphertexts. Therefore, using the same IVs can lead to potential security concerns such as:

- When multiple messages are encrypted with the same IV, the attacker can analyze the messages and with a certain background knowledge of the context of the message, the attacker can gain many valuable insights about the message.
- Additionally, using the same Key/IV pair can also lead to attackers launching **Replay Attacks** that work on vulnerable servers where the attacker can intercept, capture the message, and replay the message to the server. The attacker doesn't even need advanced skills to decrypt the message.

Task-6.2: Common Mistake: Use the same IV

Given Data:

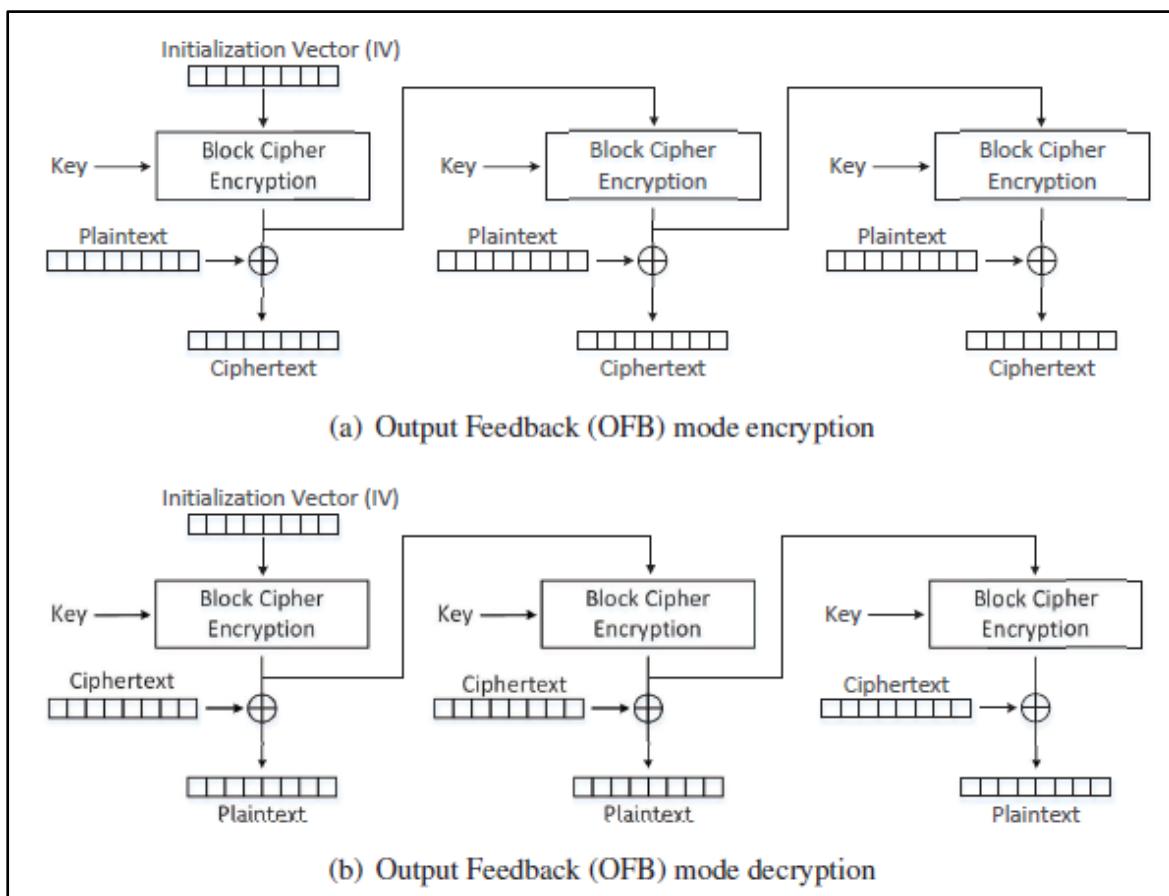
Plaintext(P1): This is a known message!

Ciphertext(C1): a469b1c502c1cab966965e50425438e1bb1b5f9037a4c159

Plaintext(P2): (unknown to you)

Ciphertext(C2): bf73bcd3509299d566c35b5d450337e1bb175f903fafc159

The encryption and Decryption process for the OFB mode is as shown below: (Copied from [Image link](#))



- The OFB mode applies XOR to the plaintext with the output of the block cipher encryption. The IV and Key go into the block cipher encryption. The same procedure is repeated for the other blocks of plaintexts as well.
- Therefore, it is very critical that a different IV is chosen for each block. If identical IVs are used, the equation for OFB literally turns out to be
 - $P_1 \text{ XOR } C_1 = P_2 \text{ XOR } C_2 = P_3 \text{ XOR } C_3 \dots$ and so on
 - Where, (P_1, C_1) and (P_2, C_2) , and (P_3, C_3) are corresponding plaintexts and ciphertexts respectively.

Using the weakness above, I designed the below code by slightly modifying the sample code as below:

```
[09/16/23]seed@VM:~/.../Files$ cat sample_code.py
#!/usr/bin/python3

# XOR two bytearrays
def xor(first, second):
    return bytearray(x^y for x,y in zip(first, second))

#Plaintext-1
P1 = "This is a known message!"
#Ciphertext-1
C1 = "a469b1c502c1cab966965e50425438e1bb1b5f9037a4c159"

#Ciphertext-2
C2 = "bf73bcd3509299d566c35b5d450337e1bb175f903fafc159"

#As we now have P1, C1 and C2, we can derive P2 (Decrypt ciphertext-2) by the relation P1 XOR C1 = P2 XOR C2
#P2 = (P1 XOR C1) XOR C2

# Convert ascii string to bytearray
D1 = bytes(P1, 'utf-8')

# Convert hex string to bytearray
D2 = bytearray.fromhex(C1)
D3 = bytearray.fromhex(C2)

#Compute P1 XOR C1 and convert to bytearray
r1 = bytearray(xor(D1, D2))
print("P1 XOR C1: ", r1)

#Compute "P2 = (P1 XOR C1) XOR C2 and convert to bytearray"
r2 = bytearray(xor(r1, D3))
print("Decrypted Plain Text: ", r2)[09/16/23]seed@VM:~/.../Files$
```

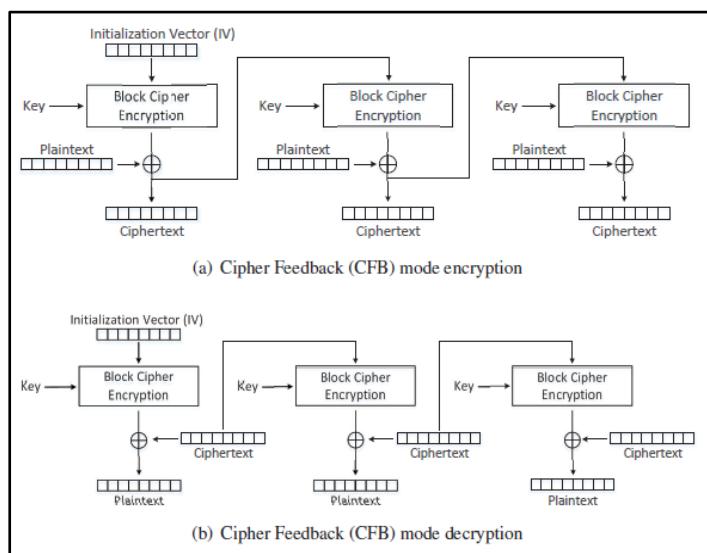
Now, running the above code to get the decrypted plaintext:

```
[09/16/23]seed@VM:~/.../Files$ python3 sample_code.py
P1 XOR C1: bytearray(b'\xf0\x01\xd8\xb6"\xa8\xb9\x99\x07\xb6>-#V\xc1\xd6~, \xe3V\xc3\x4x')
Decrypted Plain Text: bytearray(b'Order: Launch a missile!')
```

Decrypted Plaintext: Order: Launch a missile!

If we replace OFB in this experiment with CFB(CipherFeedback), how much of P2 can be revealed?

The operation of CFB is demonstrated in the below block diagram: (Copied from [Image Link](#))



From the above diagram, we observe that CFB supplies the XOR of the plaintext with the output of block cipher encryption as the input stream for the block cipher encryption for the next block of plaintext sequentially.

So, CFB is also as vulnerable as OFB. Since the plaintext and IV are known, it is possible for the attacker to get the input stream that is supplied to block encryption by applying XOR on the known plaintext and ciphertext.

Task-6.3: Common Mistake – Use a Predictable IV

One of the most important requirements in block cipher encryption is that the IV must be generated randomly.

If the attacker is able to predict the next IV beforehand, then he can launch the chosen plaintext attack and recover the original plaintext.

In AES_128_CBC mode, the encryption routine in the first block is given by:

$C_0 = \text{Enc} (P_0 \text{ XOR } IV)$ where,

C_0 is the Ciphertext,

P_0 is the plaintext, and

IV is the Initialization Vector

If the attacker can guess the IV used in the next message which I call as IV_{next} . Let P_{chosen} be the chosen plaintext block. Then, we have:

$C_{\text{next}} = \text{Enc} (P_{\text{chosen}} \text{ XOR } IV_{\text{next}})$

If the attacker makes a successful guess on the chosen plaintext in such a way that $C_0 = C_{\text{next}}$, then we have

$C_0 = C_{\text{next}}$

$\Rightarrow P_0 \text{ XOR } IV = P_{\text{chosen}} \text{ XOR } IV_{\text{next}}$

$\Rightarrow P_{\text{chosen}} = P_0 \text{ XOR } IV \text{ XOR } IV_{\text{next}}$

Running the encryption oracle,

```
[09/18/23] seed@VM:~/.../Crypto_Prog$ nc 10.9.0.80 3000
Bob's secret message is either "Yes" or "No", without quotations.
Bob's ciphertext: 5e0902995ff7f835d43d9f16c7985d7a
The IV used      : e3fd9e2832e8beead1819da5534ff6b2

Next IV          : c5a3df4532e8beead1819da5534ff6b2
Your plaintext : █
```

Here, we have

Initial IV = e3fd9e2832e8beead1819da5534ff6b2

Predicted IV = c5a3df4532e8beead1819da5534ff6b2

Ciphertext = 5e0902995ff7f835d43d9f16c7985d7a

Below is the snap of the Python code I used to get the chosen plaintext:

```
[09/18/23] seed@VM:~/.../Chosen_Plaintext$ cat choose_plaintext.py
#!/usr/bin/env python3

#Guessed Plaintext
guess = "Yes"

p1 = bytearray(guess, encoding='utf-8')

padding = 16 - len(p1)%16
p1.extend([padding]*padding)

Initial_IV = bytearray.fromhex("e3fd9e2832e8beead1819da5534ff6b2")

Next_IV = bytearray.fromhex("c5a3df4532e8beead1819da5534ff6b2")

P_chosen = bytearray(x^y^z for x,y,z in zip(p1,Initial_IV,Next_IV))

#Displaying the chosen plaintext
print("Plaintext to choose: ",P_chosen.hex(), "\n") [09/18/23] seed@Vi
```

Output:

Attaching the above plaintext into the encryption oracle:

As we can see from the screenshot the initial part of the ciphertext we obtained is matching with Bob's original ciphertext. Therefore, Bob's secret message must be "Yes".

Task-7: Programming Using Crypto Library

The below program reads the given plaintext, ciphertext, and initialization Vector and tries to find the key from the given list of English words saved in the file word_list.txt.

Given:

Plaintext: This is a top secret.

Ciphertext: 764aa26b55a4da654df6b19e4bce00f4 ed05e09346fb0e762583cb7da2ac93a2

IV: aabbccddeeff00998877665544332211

Code (written in Python):

```
#This is a program to find the key given plaintext, ciphertext and IV with
aes-128-cbc encryption

from binascii import unhexlify
from Crypto.Cipher import AES
from Crypto.Util.Padding import pad, unpad

PLAIN_TEXT = "This is a top secret."

CIPHER_TEXT =
"764aa26b55a4da654df6b19e4bce00f4 ed05e09346fb0e762583cb7da2ac93a2"

IV = "aabbccddeeff00998877665544332211"

print("Plaintext: ",PLAIN_TEXT)
print("Ciphertext: ",CIPHER_TEXT)
print("IV: ",IV)

#Reading Key List from the English word list
KEY_LIST = [line.strip() for line in open("word_list.txt","r")]

#Converting Hex String to Binary
CIPHER_TEXT = unhexlify(CIPHER_TEXT)
IV = unhexlify(IV)

#Encoding the plain text message
MESSAGE = pad(PLAIN_TEXT.encode(), AES.block_size)
```

```

key_found = 0

#Iterate through the keys list
for key in KEY_LIST:

    #If the size of key is more than 16 bytes, trim the key to include
    only 16 characters
    if len(key) > 16:
        corrected_key = key[0:16]

    #If size of key is less than 16, pad the rest of the size with
    character '#'
    else:
        corrected_key = key + "#"*(16 - len(key))

    #Encryption Process
    cipher = AES.new(corrected_key.encode(), AES.MODE_CBC, IV)
    ENCRYPTED_TEXT = cipher.encrypt(MESSAGE)

    #Check if the encrypted text matches with the given ciphertext
    if ENCRYPTED_TEXT == CIPHER_TEXT:
        print("Found the Key: ",key)
        key_found = 1
        break

if key_found == 0:
    print("Key Not found")

```

Output:

```
[09/18/23]seed@VM:~/.../Crypto_Prog$ python3 crypto_prog.py
Plaintext: This is a top secret.
Ciphertext: 764aa26b55a4da654df6b19e4bce00f4ed05e09346fb0e762583cb7da2ac93a2
IV: aabbccddeef00998877665544332211
Found the Key: Syracuse
```

Key Found: “Syracuse”

Explanation of the code procedure:

1. Declare variables and initialize values of the given plaintext message, ciphertext, and initialization vector to the respective variables.
2. Since the initialization vector is in Hex string, convert the string to binary code (I did this using the Python binascii library).
3. Add padding to the plaintext message after encoding corresponding to the AES block size using the pad() function available in the Python Crypto library.
4. Read the keys from the words.txt file supplied by the seed labs that contains the list of commonly used keys.
5. For each key read from the key file, check if the size of the key is 16 bytes. Add padding with the character '#' if the size is less than 16 or remove the excess characters from the key if the size exceeds 16.
6. Generate a new AES crypto cipher using the AES.new() function from the Python crypto library. The function takes the encoded key, cipher mode (AES_CBC), and the IV.
7. Apply encryption to the plaintext message and store the ciphertext obtained into a variable.
8. Compare the ciphertext obtained with the given ciphertext and output the key if they are matching.