



COMPUTER ORGANIZATION



CODE TRANSLATION I

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CASE #1: AN APP

- A single class
Hence, no instantiation of objects and no linking
- No attributes
Hence no heap and no data segment
- A single method, `main`
Hence no stack needed for method invocation
- Few integer variables
Hence no spilling —everything fits in registers

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I/O

- Each O/S comes with a host of syscalls. Also known as interrupts
- SPIM uses `$v0` for the service number.
- See the SPIM Syscall Sheet in the course resources.

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THE MIPS REGISTERS

2	\$0, \$ra	Reserved by hardware
10	\$t0 - \$t9	Callee-saved (temporary)
8	\$s0 - \$s7	Caller-saved (global)
4	\$a0 - \$a3	Method parameters
2	\$v0 - \$v1	Method returns
2	\$sp, \$fp	stack/frame pointers
1	\$gp	Reserved for global data
1	\$at	Reserved by assembler
2	\$k0 - \$k1	Reserved by O/S

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EXAMPLE

Java main method:

- Read x and y , both int
- Compute and output: $z + t * r + 1$, where

```
z = x + y
t = max(x, y)
r = x - y - 10
```

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GENERAL PATTERNS

- **The u suffix:**
Stands for un-trapped for add / sub
Stands for unsigned for mult, div, and slt
- **Handling Immediates**
5-bit, zero-extended for shifts
16-bit, zero-extended for logical and lui
26-bit, sign-extended for jump
16-bit, sign-extended for all the rest
- **Large Immediates**
The instruction lui

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THE ARITHMETIC FAMILY

- **add/sub**

Why three operands? Why addi but not subi?

- **slt**

Why do we need it? Why sltu?

- **mult**

Why two operands? When to ignore HI?

- **div**

Why two operands?

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THE LOGICAL BITWISE FAMILY

- **and, or, and xor**

Is there an immediate version?

- **nor**

What about not?

- **sll and srl**

Is there a variable version?

- **sra**

Why do we need it?

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THE JUMP FAMILY

- **j and jr**

What are their algorithms? Do we need both?

- **jal and jalr**

What are their algorithms? Do we need both?

This family is used for unconditional branching to skip the else fragment of if statements and to implement method invocation / return.

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THE BRANCH FAMILY

- **beq and bne**

What are their algorithms?

- **bltz, bltz, bltz**

Signed implied here.

Used for conditional branching to implement if statements and loops.

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THE LOAD / STORE FAMILY

- **lw and sw**

What are their algorithms?

- **lb and sb**

What does the u suffix do to lb?

Used to transfer data to/from DRAM to implement read/write to .data, to the heap, and to the stack.

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CASE #2: A UTILITY CLASS

- **A single class plus its client**

Hence, two linked classes but no objects

- **Static Attributes only**

Hence .data is needed but no heap

- **Can have several static methods**

Hence stack is needed

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STATIC ATTRIBUTES IN .data

- To allocate static int $x = 5$ in .data:
`x: .word* 5`
- To transfer the value of x to register r :
`lb/h/w $r, x($0)`
- To transfer the value of register r to x :
`sb/h/w $r, x($0)`

Can use byte/half/word; ascii/ascii; or space for declaration.

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```
.data
x: .byte 65
y: .word 123
z: .byte -30
u: .half 120
v: .ascii "York"
s: .ascii "York"
t: .float 2.45
```

```
.text
lb $t0, x($0)
lw $t0, y($0)
lb $t0, z($0)
lbu $t0, z($0)
lhu $t0, u($0)
lbu $t0, v($0)
```

Example #1

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HR/14

```
.data
y: .word 123, 150, 22
z: .byte -30, 12
p: .word y
```

```
.text
addi $t1, $0, 8
lw $t0, y($t1) # index-like
addi $t1, $0, 1
lb $t0, z($t1) # index-like
la $t1, y
lw $t0, 0($t1) # pointer-like
lw $t0, 4($t1) # pointer-like
lw $t1, p($0)
addi $t1, $t1, 8
lw $t0, 0($t1) # pointer-like
```

Example #2

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```

.data
x: .byte 65
y: .word 123, 150, 22
z: .byte -30, 12
u: .half 120
v: .ascii "York"
s: .asciiz "York"
p: .word y

.text
la    $t1, x
addi  $t1, $t1, 8
lw     $t0, 0($t1)    # alignment!

lw     $t1, p($0)
addi  $t1, $t1, 16    # endianness
lw     $t0, 0($t1)

```

Example #3

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STACK USAGE

- To push the content of register *r* on the stack:

```
sw    $r, 0($sp)
addi  $sp, $sp, -4
```
- To pop the word at the top of the stack into *r*:

```
addi  $sp, $sp, 4
lw     $r, 0($sp)
```
- Every method must preserve *\$sp*, *\$ra*, and *\$s*? plus any other register it needs after a call.

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EXAMPLE

	O/S		main		method
16	...	100	main:	600	method:
20	jal main	104	...	604	...
24	...	108	jal method	608	...
28	...	112	...	612	...
32	...	116	...	616	...
36	...	120	jr \$ra	620	jr \$ra

PC	20	100	108	600	608	620	112	120	112
\$ra	24	24	112	112	112	112	112	112	112

CASE #3: ANY CLASS

We need to be able to accommodate

- **Non-Static Attributes**
Storage allocated on the heap
- **Multi-Class Applications**
Multiple classes loaded and linked

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HEAP USAGE

To allocate four bytes on the heap:

```

.text
main: ...

    addi    $a0, $0, 4
    addi    $v0, $0, 9
    syscall
    sw      $s0, 0($v0)    # heap store
    ...
    lw      $s0, 0($v0)    # heap load

```

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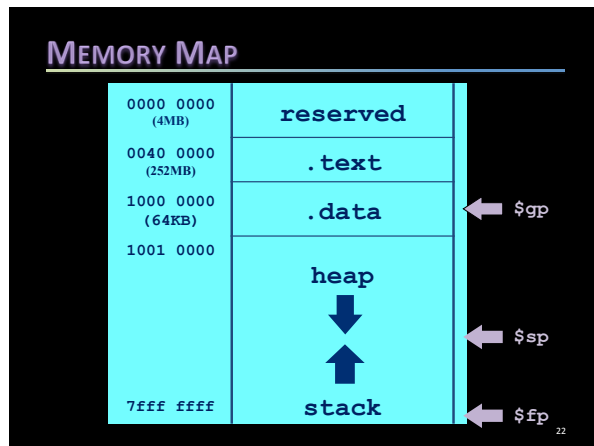
OPP AND THE HEAP

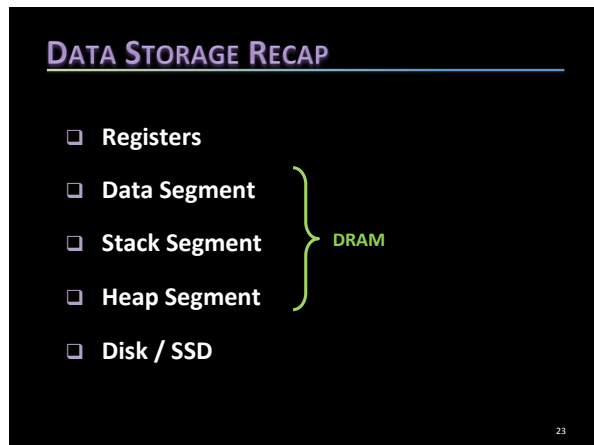
1. Class A instantiates class B via: `B b = new B(...);`
2. Class A invokes a method `m` in B via: `b.m();`

Class B constructor:

- Determines #of bytes needed to hold the state;
i.e. the sum of the sizes of all attributes.
- Requests a block of that many bytes on the heap.
- Let `x` = beginning address of the returned block.
- Store all attribute values beginning at `x`.
- Return `x`

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EXERCISES

For each data storage option, determine:

- Its latency (time to retrieve one byte) in seconds
- Its typical size (in bytes)
- The lifetime of its content
- What is it used for in Java / C
- How to allocate storage in it
- How to transfer data to/from it

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