This chapter looks at some of the more advanced features of MapReduce, including counters and sorting and joining datasets.

### **Counters**

There are often things that you would like to know about the data you are analyzing but that are peripheral to the analysis you are performing. For example, if you were counting invalid records and discovered that the proportion of invalid records in the whole dataset was very high, you might be prompted to check why so many records were being marked as invalid—perhaps there is a bug in the part of the program that detects invalid records? Or if the data was of poor quality and genuinely did have very many invalid records, after discovering this, you might decide to increase the size of the dataset so that the number of good records was large enough for meaningful analysis.

Counters are a useful channel for gathering statistics about the job: for quality control or for application-level statistics. They are also useful for problem diagnosis. If you are tempted to put a log message into your map or reduce task, it is often better to see whether you can use a counter instead to record that a particular condition occurred. In addition to counter values being much easier to retrieve than log output for large distributed jobs, you get a record of the number of times that condition occurred, which is more work to obtain from a set of logfiles.

#### **Built-in Counters**

Hadoop maintains some built-in counters for every job, and these report various metrics. For example, there are counters for the number of bytes and records processed, which allow you to confirm that the expected amount of input was consumed and the expected amount of output was produced.

Counters are divided into groups, and there are several groups for the built-in counters, listed in <u>Table 9-1</u>.

Table 9-1. Built-in counter groups

Group	Name/Enum	Reference
MapReduce task counters	org.apache.hadoop.mapreduce.TaskCounter	<u>Table 9-2</u>
Filesystem counters	org.apache.hadoop.mapreduc e.FileSystemCounter	Table 9-3
FileInputForma t counters	<pre>org.apache.hadoop.mapreduc e.lib.input.FileInputFormat Counter</pre>	<b>Table 9-4</b>
FileOutputForm at counters	<pre>org.apache.hadoop.mapreduc e.lib.output.FileOutputForm atCounter</pre>	<u>Table 9-5</u>
Job counters	org.apache.hadoop.mapreduce.JobCounter	<u>Table 9-6</u>

Each group either contains *task counters* (which are updated as a task progresses) or *job counters* (which are updated as a job progresses). We look at both types in the following sections.

#### Task counters

Task counters gather information about tasks over the course of their execution, and the results are aggregated over all the tasks in a job. The MAP\_INPUT\_RECORDS counter, for example, counts the input records read by each map task and aggregates over all map tasks in a job, so that the final figure is the total number of input records for the whole job.

Task counters are maintained by each task attempt, and periodically sent to the application master so they can be globally aggregated. (This is described in **Progress and Status Updates**.) Task counters are sent in full every time, rather than sending the counts since the last transmission, since this guards against errors due to lost messages. Furthermore, during a job run, counters may go down if a task fails.

Counter values are definitive only once a job has successfully completed. However, some counters provide useful diagnostic information as a task is progressing, and it can be useful to monitor them with the web UI. For example, PHYSICAL\_MEMORY\_BYTES, VIRTUAL\_MEMORY\_BYTES, and COMMITTED\_HEAP\_BYTES provide an indication of how memory usage varies over the course of a particular task attempt.

The built-in task counters include those in the MapReduce task counters group (**Table 9-2**) and those in the file-related counters groups (**Tables 9-3**, **9-4**, and **9-5**).

Counter	Description
Map input records ( MAP_IN PUT_RECORDS )	The number of input records consumed by all the maps in the job. Incremented every time a record is read from a RecordReader and passed to the map's map() method by the framework.
Split raw bytes (SPLIT_RAW_BYT ES)	The number of bytes of input-split objects read by maps. These objects represent the split metadata (that is, the offset and length within a file) rather than the split data itself, so the total size should be small.
Map output records ( MAP_OU TPUT_RECORDS )	The number of map output records produced by all the maps in the job. Incremented every time the collect() method is called on a map's OutputCollector.
Map output bytes (MAP_OUTP UT_BYTES)	The number of bytes of uncompressed output produced by all the maps in the job.  Incremented every time the collect() method is called on a map's OutputCollecto r.
Map output materialized bytes (MAP_OUTP UT_MATERIALIZE D_BYTES)	The number of bytes of map output actually written to disk. If map output compression is enabled, this is reflected in the counter value.
Combine input records ( COMBIN E_INPUT_RECORD S )	The number of input records consumed by all the combiners (if any) in the job. Incremented every time a value is read from the combiner's iterator over values. Note that this

count is the number of values consumed by

the combiner, not the number of distinct key

Counter	Description
	groups (which would not be a useful metric, since there is not necessarily one group per key for a combiner; see <u>Combiner Functions</u> , and also <u>Shuffle and Sort</u> ).
Combine output records ( COMBIN E_OUTPUT_RECOR DS )	The number of output records produced by all the combiners (if any) in the job.  Incremented every time the collect() method is called on a combiner's OutputColl ector.
Reduce input groups ( REDUCE_ INPUT_GROUPS )	The number of distinct key groups consumed by all the reducers in the job. Incremented every time the reducer's reduce() method is called by the framework.
Reduce input records ( REDUCE _INPUT_RECORD S )	The number of input records consumed by all the reducers in the job. Incremented every time a value is read from the reducer's iterator over values. If reducers consume all of their inputs, this count should be the same as the count for map output records.
Reduce output records ( REDUCE _OUTPUT_RECORD S )	The number of reduce output records produced by all the maps in the job.  Incremented every time the collect() method is called on a reducer's OutputCollector.
Reduce shuffle bytes ( REDUCE_S HUFFLE_BYTES )	The number of bytes of map output copied by the shuffle to reducers.
Spilled records (SPILLED_RECOR DS)	The number of records spilled to disk in all map and reduce tasks in the job.

Counter	Description
CPU milliseconds ( CP U_MILLISECOND S )	The cumulative CPU time for a task in milliseconds, as reported by /proc/cpuinfo.
Physical memory bytes (PHYSICAL_MEMO RY_BYTES)	The physical memory being used by a task in bytes, as reported by /proc/meminfo.
Virtual memory bytes (VIRTUAL_ MEMORY_BYTES)	The virtual memory being used by a task in bytes, as reported by /proc/meminfo.
Committed heap bytes ( COMMITTE D_HEAP_BYTES )	The total amount of memory available in the JVM in bytes, as reported by Runtime.getRun time().totalMemory().
GC time milliseconds (GC _TIME_MILLIS)	The elapsed time for garbage collection in tasks in milliseconds, as reported by Garbage CollectorMXBean.getCollectionTime().
Shuffled maps ( SHUFFLED_MAP S )	The number of map output files transferred to reducers by the shuffle (see <b>Shuffle and Sort</b> ).
Failed shuffle (FAILED_SHUFFLE)	The number of map output copy failures during the shuffle.
Merged map outputs ( MERGED _MAP_OUTPUTS )	The number of map outputs that have been merged on the reduce side of the shuffle.

Table 9-3. Built-in filesystem task counters

Counter	Description
Filesystem bytes read (BYTES_READ)	The number of bytes read by the filesystem by map and reduce tasks. There is a counter for each filesystem, and Filesystem may be Local, HDFS, S3, etc.
Filesystem bytes written (BYTES_WRITTEN)	The number of bytes written by the filesystem by map and reduce tasks.
Filesystem read ops ( READ_OPS )	The number of read operations (e.g., open, file status) by the filesystem by map and reduce tasks.
Filesystem large read ops ( LARGE_ READ_OPS )	The number of large read operations (e.g., list directory for a large directory) by the filesystem by map and reduce tasks.
Filesystem write ops (WRITE_OPS)	The number of write operations (e.g., create, append) by the filesystem by map and reduce tasks.

Table 9-4. Built-in FileInputFormat task counters

Counter	Description
Bytes read ( BYTES	The number of bytes read by map tasks via
_READ )	the FileInputFormat.

Table 9-5. Built-in FileOutputFormat task counters

Counter	Description	
Bytes written (B	The number of bytes written by map tasks	
YTES_WRITTEN)	(for map-only jobs) or reduce tasks via the Fi	
	leOutputFormat.	

# Job counters

Job counters (**Table 9-6**) are maintained by the application master, so they don't need to be sent across the network, unlike all other counters, including user-defined ones. They measure job-level statistics, not values that change while a task is running. For example, TOTAL\_LAUNCHED\_MAPS counts the number of map tasks that were launched over the course of a job (including tasks that failed).

Counter	Description
Launched map tasks (TOTAL_LA UNCHED_MAPS)	The number of map tasks that were launched. Includes tasks that were started speculatively (see <b>Speculative Execution</b> ).
Launched reduce tasks ( T OTAL_LAUNCHED _REDUCES )	The number of reduce tasks that were launched. Includes tasks that were started speculatively.
Launched uber tasks ( TOTAL_LA UNCHED_UBERTA SKS )	The number of uber tasks (see <u>Anatomy of a</u> <u>MapReduce Job Run</u> ) that were launched.
Maps in uber tasks ( NUM_UBER _SUBMAPS )	The number of maps in uber tasks.
Reduces in uber tasks ( NUM_UBER _SUBREDUCES )	The number of reduces in uber tasks.
Failed map tasks ( NUM_FAIL ED_MAPS )	The number of map tasks that failed. See <u>Task</u> <u>Failure</u> for potential causes.
Failed reduce tasks (NUM_FAIL ED_REDUCES)	The number of reduce tasks that failed.
Failed uber tasks (NUM_FAIL ED_UBERTASKS)	The number of uber tasks that failed.

Counter	Description
Killed map tasks ( NUM_KILL ED_MAPS )	The number of map tasks that were killed. See  Task Failure for potential causes.
Killed reduce tasks ( NUM_KILL ED_REDUCES )	The number of reduce tasks that were killed.
Data-local map tasks ( DATA_LOC AL_MAPS )	The number of map tasks that ran on the same node as their input data.
Rack-local map tasks ( RACK_LOC AL_MAPS )	The number of map tasks that ran on a node in the same rack as their input data, but were not data-local.
Other local map tasks (OTHER_LO CAL_MAPS)	The number of map tasks that ran on a node in a different rack to their input data. Interrack bandwidth is scarce, and Hadoop tries to place map tasks close to their input data, so this count should be low. See <u>Figure 2-2</u> .
Total time in map tasks (MIL LIS_MAPS)	The total time taken running map tasks, in milliseconds. Includes tasks that were started speculatively. See also corresponding counters for measuring core and memory usage ( VCORE S_MILLIS_MAPS and MB_MILLIS_MAPS ).
Total time in reduce tasks ( M ILLIS_REDUCE S )	The total time taken running reduce tasks, in milliseconds. Includes tasks that were started speculatively. See also corresponding counters for measuring core and memory usage ( VCORE S_MILLIS_REDUCES and MB_MILLIS_REDUCES ).

## **User-Defined Java Counters**

MapReduce allows user code to define a set of counters, which are then incremented as desired in the mapper or reducer. Counters are defined by a Java enum, which serves to group related counters. A job may define an arbitrary number of enums, each with an arbitrary number of fields. The name of the enum is the group name, and the enum's fields are the counter names. Counters are global: the MapReduce framework aggregates them across all maps and reduces to produce a grand total at the end of the job.

We created some counters in <u>Chapter 6</u> for counting malformed records in the weather dataset. The program in <u>Example 9-1</u> extends that example to count the number of missing records and the distribution of temperature quality codes.

Example 9-1. Application to run the maximum temperature job, including counting missing and malformed fields and quality codes

```
public class MaxTemperatureWithCounters extends Configured implements Tool {
  enum Temperature {
   MISSING,
   MALFORMED
  static class MaxTemperatureMapperWithCounters
      extends Mapper<LongWritable, Text, Text, IntWritable> {
    private NcdcRecordParser parser = new NcdcRecordParser();
    @Override
    protected void map(LongWritable key, Text value, Context context)
        throws IOException, InterruptedException {
      parser.parse(value);
      if (parser.isValidTemperature()) {
        int airTemperature = parser.getAirTemperature();
        context.write(new Text(parser.getYear()),
            new IntWritable(airTemperature));
      } else if (parser.isMalformedTemperature()) {
        System.err.println("Ignoring possibly corrupt input: " + value);
        context.getCounter(Temperature.MALFORMED).increment(1);
```

```
} else if (parser.isMissingTemperature()) {
        context.getCounter(Temperature.MISSING).increment(1);
     }
     // dynamic counter
     context.getCounter("TemperatureQuality", parser.getQuality()).increment(1)
   }
  }
 @Override
 public int run(String[] args) throws Exception {
    Job job = JobBuilder.parseInputAndOutput(this, getConf(), args);
   if (job == null) {
     return -1;
   }
    job.setOutputKeyClass(Text.class);
    job.setOutputValueClass(IntWritable.class);
    job.setMapperClass(MaxTemperatureMapperWithCounters.class);
    job.setCombinerClass(MaxTemperatureReducer.class);
    job.setReducerClass(MaxTemperatureReducer.class);
   return job.waitForCompletion(true) ? 0 : 1;
  }
 public static void main(String[] args) throws Exception {
    int exitCode = ToolRunner.run(new MaxTemperatureWithCounters(), args);
   System.exit(exitCode);
 }
}
```

The best way to see what this program does is to run it over the complete dataset:

% hadoop jar hadoop-examples.jar MaxTemperatureWithCounters \
input/ncdc/all output-counters

When the job has successfully completed, it prints out the counters at the end (this is done by the job client). Here are the ones we are interested in:

```
Air Temperature Records
Malformed=3
```

```
Missing=66136856
TemperatureQuality
0=1
1=973422173
2=1246032
4=10764500
5=158291879
6=40066
9=66136858
```

Notice that the counters for temperature have been made more readable by using a resource bundle named after the enum (using an underscore as a separator for nested classes)—in this case MaxTemperatureWithCounters\_Temperature.properties, which contains the display name mappings.

### **Dynamic counters**

The code makes use of a dynamic counter—one that isn't defined by a Java enum. Because a Java enum's fields are defined at compile time, you can't create new counters on the fly using enums. Here we want to count the distribution of temperature quality codes, and though the format specification defines the values that the temperature quality code *can* take, it is more convenient to use a dynamic counter to emit the values that it *actually* takes. The method we use on the Context object takes a group and counter name using String names:

```
public Counter getCounter(String groupName, String counterName)
```

The two ways of creating and accessing counters—using enums and using strings—are actually equivalent because Hadoop turns enums into strings to send counters over RPC. Enums are slightly easier to work with, provide type safety, and are suitable for most jobs. For the odd occasion when you need to create counters dynamically, you can use the String interface.

#### **Retrieving counters**

In addition to using the web UI and the command line (using mapred job -counter), you can retrieve counter values using the Java API. You can do

this while the job is running, although it is more usual to get counters at the end of a job run, when they are stable. **Example 9-2** shows a program that calculates the proportion of records that have missing temperature fields.

Example 9-2. Application to calculate the proportion of records with missing temperature fields

```
import org.apache.hadoop.conf.Configured;
import org.apache.hadoop.mapreduce.*;
import org.apache.hadoop.util.*;
public class MissingTemperatureFields extends Configured implements Tool {
  @Override
  public int run(String[] args) throws Exception {
    if (args.length != 1) {
      JobBuilder.printUsage(this, "<job ID>");
      return -1;
    }
    String jobID = args[0];
    Cluster cluster = new Cluster(getConf());
    Job job = cluster.getJob(JobID.forName(jobID));
    if (job == null) {
      System.err.printf("No job with ID %s found.\n", jobID);
      return -1;
    }
    if (!job.isComplete()) {
      System.err.printf("Job %s is not complete.\n", jobID);
      return -1;
    }
    Counters counters = job.getCounters();
    long missing = counters.findCounter(
        MaxTemperatureWithCounters.Temperature.MISSING).getValue();
    long total = counters.findCounter(TaskCounter.MAP_INPUT_RECORDS).getValue();
    System.out.printf("Records with missing temperature fields: %.2f%%\n",
        100.0 * missing / total);
    return 0;
  public static void main(String[] args) throws Exception {
    int exitCode = ToolRunner.run(new MissingTemperatureFields(), args);
    System.exit(exitCode);
```

```
}
```

First we retrieve a Job object from a Cluster by calling the getJob() method with the job ID. We check whether there is actually a job with the given ID by checking if it is null. There may not be, either because the ID was incorrectly specified or because the job is no longer in the job history.

After confirming that the job has completed, we call the Job's getCounters() method, which returns a Counters object encapsulating all the counters for the job. The Counters class provides various methods for finding the names and values of counters. We use the findCounter() method, which takes an enum to find the number of records that had a missing temperature field and also the total number of records processed (from a built-in counter).

Finally, we print the proportion of records that had a missing temperature field. Here's what we get for the whole weather dataset:

```
% hadoop jar hadoop-examples.jar MissingTemperatureFields job_1410450250506_0007
Records with missing temperature fields: 5.47%
```

# **User-Defined Streaming Counters**

A Streaming MapReduce program can increment counters by sending a specially formatted line to the standard error stream, which is co-opted as a control channel in this case. The line must have the following format:

```
reporter:counter:group,counter,amount
```

This snippet in Python shows how to increment the "Missing" counter in the "Temperature" group by 1:

```
sys.stderr.write("reporter:counter:Temperature,Missing,1\n")
```

In a similar way, a status message may be sent with a line formatted like this:

```
reporter:status:message
```

# **Sorting**

The ability to sort data is at the heart of MapReduce. Even if your application isn't concerned with sorting per se, it may be able to use the sorting stage that MapReduce provides to organize its data. In this section, we examine different ways of sorting datasets and how you can control the sort order in MapReduce. Sorting Avro data is covered separately, in **Sorting Using Avro MapReduce**.

## **Preparation**

We are going to sort the weather dataset by temperature. Storing temperatures as Text objects doesn't work for sorting purposes, because signed integers don't sort lexicographically. [61] Instead, we are going to store the data using sequence files whose IntWritable keys represent the temperatures (and sort correctly) and whose Text values are the lines of data.

The MapReduce job in Example 9-3 is a map-only job that also filters the input to remove records that don't have a valid temperature reading. Each map creates a single block-compressed sequence file as output. It is invoked with the following command:

% hadoop jar hadoop-examples.jar SortDataPreprocessor input/ncdc/all \
input/ncdc/all-seq

Example 9-3. A MapReduce program for transforming the weather data into SequenceFile format

```
context.write(new IntWritable(parser.getAirTemperature()), value);
     }
   }
  }
 @Override
  public int run(String[] args) throws Exception {
    Job job = JobBuilder.parseInputAndOutput(this, getConf(), args);
    if (job == null) {
     return -1;
    }
    job.setMapperClass(CleanerMapper.class);
    job.setOutputKeyClass(IntWritable.class);
    job.setOutputValueClass(Text.class);
    job.setNumReduceTasks(0);
    job.setOutputFormatClass(SequenceFileOutputFormat.class);
    SequenceFileOutputFormat.setCompressOutput(job, true);
    SequenceFileOutputFormat.setOutputCompressorClass(job, GzipCodec.class);
    SequenceFileOutputFormat.setOutputCompressionType(job,
        CompressionType.BLOCK);
    return job.waitForCompletion(true) ? 0 : 1;
 public static void main(String[] args) throws Exception {
    int exitCode = ToolRunner.run(new SortDataPreprocessor(), args);
    System.exit(exitCode);
 }
}
```

#### **Partial Sort**

In <u>The Default MapReduce Job</u>, we saw that, by default, MapReduce will sort input records by their keys. <u>Example 9-4</u> is a variation for sorting sequence files with IntWritable keys.

Example 9-4. A MapReduce program for sorting a SequenceFile with IntWritable keys using the default HashPartitioner

```
public class SortByTemperatureUsingHashPartitioner extends Configured
  implements Tool {
    @Override
    public int run(String[] args) throws Exception {
        Job job = JobBuilder.parseInputAndOutput(this, getConf(), args);
}
```

```
return -1;
    }
    job.setInputFormatClass(SequenceFileInputFormat.class);
    job.setOutputKeyClass(IntWritable.class);
    job.setOutputFormatClass(SequenceFileOutputFormat.class);
    SequenceFileOutputFormat.setCompressOutput(job, true);
    SequenceFileOutputFormat.setOutputCompressorClass(job, GzipCodec.class);
    SequenceFileOutputFormat.setOutputCompressionType(job,
        CompressionType.BLOCK);
   return job.waitForCompletion(true) ? 0 : 1;
  }
  public static void main(String[] args) throws Exception {
    int exitCode = ToolRunner.run(new SortByTemperatureUsingHashPartitioner(),
        args);
   System.exit(exitCode);
 }
}
```

#### **CONTROLLING SORT ORDER**

The sort order for keys is controlled by a RawComparator, which is found as follows:

if (job == null) {

- 1. If the property mapreduce.job.output.key.comparator.class is set, either explicitly or by calling setSortComparatorClass() on Job, then an instance of that class is used. (In the old API, the equivalent method is setOutputKeyComparatorClass() on JobConf.)
- 2. Otherwise, keys must be a subclass of WritableComparable, and the registered comparator for the key class is used.
- 3. If there is no registered comparator, then a RawComparator is used. The RawComparator deserializes the byte streams being compared into objects and delegates to the Writable Comparable's compareTo() method.

These rules reinforce the importance of registering optimized versions of RawComparator's for your own custom Writable classes (which is covered in <a href="Implementing a RawComparator for speed">Implementing a RawComparator for speed</a>), and also show that it's straightforward to override the sort order by setting your own comparator (we do this in <a href="Secondary Sort">Secondary Sort</a>).

% hadoop jar hadoop-examples.jar SortByTemperatureUsingHashPartitioner \
-D mapreduce.job.reduces=30 input/ncdc/all-seq output-hashsort

This command produces 30 output files, each of which is sorted. However, there is no easy way to combine the files (by concatenation, for example, in the case of plain-text files) to produce a globally sorted file.

For many applications, this doesn't matter. For example, having a partially sorted set of files is fine when you want to do lookups by key. The SortByTemperatureToMapFile and LookupRecordsByTemperature classes in this book's example code explore this idea. By using a map file instead of a sequence file, it's possible to first find the relevant partition that a key belongs in (using the partitioner), then to do an efficient lookup of the record within the map file partition.

### **Total Sort**

How can you produce a globally sorted file using Hadoop? The naive answer is to use a single partition. But this is incredibly inefficient for large files, because one machine has to process all of the output, so you are throwing away the benefits of the parallel architecture that MapReduce provides.

Instead, it is possible to produce a set of sorted files that, if concatenated, would form a globally sorted file. The secret to doing this is to use a partitioner that respects the total order of the output. For example, if we had four partitions, we could put keys for temperatures less than  $-10^{\circ}$ C in the first partition, those between  $-10^{\circ}$ C and  $0^{\circ}$ C in the second, those between  $0^{\circ}$ C and  $10^{\circ}$ C in the fourth.

Although this approach works, you have to choose your partition sizes carefully to ensure that they are fairly even, so job times aren't dominated by a single reducer. For the partitioning scheme just described, the relative sizes of the partitions are as follows:

Temperature range	< -	[-10°C,	[0°C,	>=
	10°C	0°C)	10°C)	10°C
	4.407	100/	4.707	<b>300</b> /
Proportion of	11%	13%	17%	59%
records				

These partitions are not very even. To construct more even partitions, we need to have a better understanding of the temperature distribution for the whole dataset. It's fairly easy to write a MapReduce job to count the number of records that fall into a collection of temperature buckets. For example, <u>Figure 9-1</u> shows the distribution for buckets of size 1°C, where each point on the plot corresponds to one bucket.

Although we could use this information to construct a very even set of partitions, the fact that we needed to run a job that used the entire dataset to construct them is not ideal. It's possible to get a fairly even set of partitions by *sampling* the key space. The idea behind sampling is that you look at a small subset of the keys to approximate the key distribution, which is then used to construct partitions. Luckily, we don't have to write the code to do this ourselves, as Hadoop comes with a selection of samplers.

The InputSampler class defines a nested Sampler interface whose implementations return a sample of keys given an InputFormat and Job:

```
public interface Sampler<K, V> {
   K[] getSample(InputFormat<K, V> inf, Job job)
        throws IOException, InterruptedException;
}
```

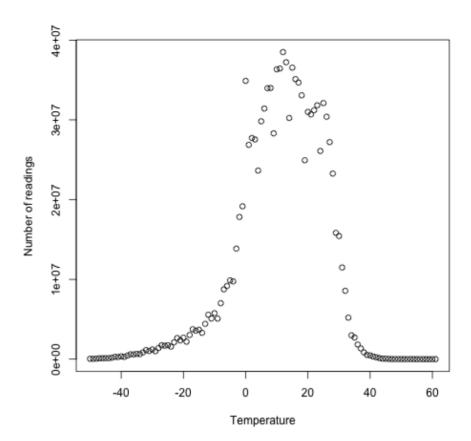


Figure 9-1. Temperature distribution for the weather dataset

This interface usually is not called directly by clients. Instead, the writePartitionFile() static method on InputSampler is used, which creates a sequence file to store the keys that define the partitions:

```
public static <K, V> void writePartitionFile(Job job, Sampler<K, V> sampler)
    throws IOException, ClassNotFoundException, InterruptedException
```

The sequence file is used by TotalOrderPartitioner to create partitions for the sort job. **Example 9-5** puts it all together.

Example 9-5. A MapReduce program for sorting a SequenceFile with IntWritable keys using the TotalOrderPartitioner to globally sort the data

```
}
    job.setInputFormatClass(SequenceFileInputFormat.class);
    job.setOutputKeyClass(IntWritable.class);
    job.setOutputFormatClass(SequenceFileOutputFormat.class);
    SequenceFileOutputFormat.setCompressOutput(job, true);
    SequenceFileOutputFormat.setOutputCompressorClass(job, GzipCodec.class);
    SequenceFileOutputFormat.setOutputCompressionType(job,
        CompressionType.BLOCK);
    job.setPartitionerClass(TotalOrderPartitioner.class);
    InputSampler.Sampler<IntWritable, Text> sampler =
        new InputSampler.RandomSampler<IntWritable, Text>(0.1, 10000, 10);
    InputSampler.writePartitionFile(job, sampler);
   // Add to DistributedCache
    Configuration conf = job.getConfiguration();
    String partitionFile = TotalOrderPartitioner.getPartitionFile(conf);
   URI partitionUri = new URI(partitionFile);
    job.addCacheFile(partitionUri);
   return job.waitForCompletion(true) ? 0 : 1;
  }
 public static void main(String[] args) throws Exception {
    int exitCode = ToolRunner.run(
        new SortByTemperatureUsingTotalOrderPartitioner(), args);
   System.exit(exitCode);
 }
}
```

We use a RandomSampler, which chooses keys with a uniform probability—here, 0.1. There are also parameters for the maximum number of samples to take and the maximum number of splits to sample (here, 10,000 and 10, respectively; these settings are the defaults when InputSampler is run as an application), and the sampler stops when the first of these limits is met. Samplers run on the client, making it important to limit the number of splits that are downloaded so the sampler runs quickly. In practice, the time taken to run the sampler is a small fraction of the overall job time.

The InputSampler writes a partition file that we need to share with the tasks running on the cluster by adding it to the distributed cache (see **Distributed Cache**).

On one run, the sampler chose –5.6°C, 13.9°C, and 22.0°C as partition boundaries (for four partitions), which translates into more even partition sizes than the earlier choice:

Temperature	< -	[-5.6°C,	[13.9°C,	>=
range	5.6°C	13.9°C)	22.0°C)	22.0°C
<b>Proportion of</b>	29%	24%	23%	24%
records				

Your input data determines the best sampler to use. For example, SplitSampler, which samples only the first n records in a split, is not so good for sorted data, 64 because it doesn't select keys from throughout the split.

On the other hand, IntervalSampler chooses keys at regular intervals through the split and makes a better choice for sorted data.

RandomSampler is a good general-purpose sampler. If none of these suits your application (and remember that the point of sampling is to produce partitions that are approximately equal in size), you can write your own implementation of the Sampler interface.

One of the nice properties of InputSampler and

TotalOrderPartitioner is that you are free to choose the number of partitions—that is, the number of reducers. However,

TotalOrderPartitioner will work only if the partition boundaries are distinct. One problem with choosing a high number is that you may get collisions if you have a small key space.

Here's how we run it:

% hadoop jar hadoop-examples.jar SortByTemperatureUsingTotalOrderPartitioner \
-D mapreduce.job.reduces=30 input/ncdc/all-seq output-totalsort

The program produces 30 output partitions, each of which is internally sorted; in addition, for these partitions, all the keys in partition i are less than the keys in partition i + 1.

# **Secondary Sort**

The MapReduce framework sorts the records by key before they reach the reducers. For any particular key, however, the values are not sorted. The order in which the values appear is not even stable from one run to the next, because they come from different map tasks, which may finish at different times from run to run. Generally speaking, most MapReduce programs are written so as not to depend on the order in which the values appear to the reduce function. However, it is possible to impose an order on the values by sorting and grouping the keys in a particular way.

To illustrate the idea, consider the MapReduce program for calculating the maximum temperature for each year. If we arranged for the values (temperatures) to be sorted in descending order, we wouldn't have to iterate through them to find the maximum; instead, we could take the first for each year and ignore the rest. (This approach isn't the most efficient way to solve this particular problem, but it illustrates how secondary sort works in general.)

To achieve this, we change our keys to be composite: a combination of year and temperature. We want the sort order for keys to be by year (ascending) and then by temperature (descending):

```
1900 35°C
1900 34°C
1900 34°C
...
1901 36°C
1901 35°C
```

If all we did was change the key, this wouldn't help, because then records for the same year would have different keys and therefore would not (in general) go to the same reducer. For example, (1900, 35°C) and (1900, 34°C) could go to different reducers. By setting a partitioner to partition by the year part of the key, we can guarantee that records for the same year go to the same reducer. This still isn't enough to achieve our goal, however. A partitioner ensures only that one reducer receives all the

records for a year; it doesn't change the fact that the reducer groups by key within the partition:

	Partition	Group
1900 359 1900 349		-
1900 34		
 1901 36	ν <b>(  </b>	1
1901 35°	·C	- 1

The final piece of the puzzle is the setting to control the grouping. If we group values in the reducer by the year part of the key, we will see all the records for the same year in one reduce group. And because they are sorted by temperature in descending order, the first is the maximum temperature:

	Partition		Group	
1900 35	PC			
1900 34	C			
1900 34	°C			
1901 369	C			
1901 359	°C			

To summarize, there is a recipe here to get the effect of sorting by value:

- Make the key a composite of the natural key and the natural value.
- The sort comparator should order by the composite key (i.e., the natural key and natural value).

• The partitioner and grouping comparator for the composite key should consider only the natural key for partitioning and grouping.

Java code

Putting this all together results in the code in **Example 9-6**. This program uses the plain-text input again.

Example 9-6. Application to find the maximum temperature by sorting temperatures in the key

```
public class MaxTemperatureUsingSecondarySort
    extends Configured implements Tool {
  static class MaxTemperatureMapper
      extends Mapper<LongWritable, Text, IntPair, NullWritable> {
    private NcdcRecordParser parser = new NcdcRecordParser();
    @Override
    protected void map(LongWritable key, Text value,
        Context context) throws IOException, InterruptedException {
      parser.parse(value);
      if (parser.isValidTemperature()) {
        context.write(new IntPair(parser.getYearInt(),
            parser.getAirTemperature()), NullWritable.get());
      }
    }
  }
  static class MaxTemperatureReducer
      extends Reducer<IntPair, NullWritable, IntPair, NullWritable> {
    @Override
    protected void reduce(IntPair key, Iterable<NullWritable> values,
        Context context) throws IOException, InterruptedException {
      context.write(key, NullWritable.get());
    }
  }
  public static class FirstPartitioner
      extends Partitioner<IntPair, NullWritable> {
```

```
@Override
  public int getPartition(IntPair key, NullWritable value, int numPartitions)
    // multiply by 127 to perform some mixing
    return Math.abs(key.getFirst() * 127) % numPartitions;
 }
}
public static class KeyComparator extends WritableComparator {
  protected KeyComparator() {
    super(IntPair.class, true);
 }
 @Override
  public int compare(WritableComparable w1, WritableComparable w2) {
    IntPair ip1 = (IntPair) w1;
    IntPair ip2 = (IntPair) w2;
    int cmp = IntPair.compare(ip1.getFirst(), ip2.getFirst());
    if (cmp != 0) {
     return cmp;
    }
    return -IntPair.compare(ip1.getSecond(), ip2.getSecond()); //reverse
 }
}
public static class GroupComparator extends WritableComparator {
 protected GroupComparator() {
    super(IntPair.class, true);
 }
 @Override
  public int compare(WritableComparable w1, WritableComparable w2) {
    IntPair ip1 = (IntPair) w1;
    IntPair ip2 = (IntPair) w2;
    return IntPair.compare(ip1.getFirst(), ip2.getFirst());
 }
}
@Override
public int run(String[] args) throws Exception {
 Job job = JobBuilder.parseInputAndOutput(this, getConf(), args);
 if (job == null) {
    return -1;
  }
  job.setMapperClass(MaxTemperatureMapper.class);
  job.setPartitionerClass(FirstPartitioner.class);
  job.setSortComparatorClass(KeyComparator.class);
```

```
job.setGroupingComparatorClass(GroupComparator.class);
job.setReducerClass(MaxTemperatureReducer.class);
job.setOutputKeyClass(IntPair.class);
job.setOutputValueClass(NullWritable.class);

return job.waitForCompletion(true) ? 0 : 1;
}

public static void main(String[] args) throws Exception {
  int exitCode = ToolRunner.run(new MaxTemperatureUsingSecondarySort(), args);
  System.exit(exitCode);
}
```

In the mapper, we create a key representing the year and temperature, using an IntPair Writable implementation. (IntPair is like the TextPair class we developed in Implementing a Custom Writable.) We don't need to carry any information in the value, because we can get the first (maximum) temperature in the reducer from the key, so we use a NullWritable. The reducer emits the first key, which, due to the secondary sorting, is an IntPair for the year and its maximum temperature. IntPair's toString() method creates a tab-separated string, so the output is a set of tab-separated year-temperature pairs.

#### NOTE

Many applications need to access all the sorted values, not just the first value as we have provided here. To do this, you need to populate the value fields since in the reducer you can retrieve only the first key. This necessitates some unavoidable duplication of information between key and value.

We set the partitioner to partition by the first field of the key (the year) using a custom partitioner called FirstPartitioner. To sort keys by year (ascending) and temperature (descending), we use a custom sort comparator, using setSortComparatorClass(), that extracts the fields and performs the appropriate comparisons. Similarly, to group keys by year, we set a custom comparator, using setGroupingComparatorClass(), to extract the first field of the key for comparison. [65]

Running this program gives the maximum temperatures for each year:

```
% hadoop jar hadoop-examples.jar MaxTemperatureUsingSecondarySort \
  input/ncdc/all output-secondarysort
% hadoop fs -cat output-secondarysort/part-* | sort | head
1901
        317
1902
        244
1903
        289
1904
       256
1905
       283
1906
        294
1907
       283
1908
       289
1909
        278
        294
1910
```

### Streaming

To do a secondary sort in Streaming, we can take advantage of a couple of library classes that Hadoop provides. Here's the driver that we can use to do a secondary sort:

```
% hadoop jar $HADOOP_HOME/share/hadoop/tools/lib/hadoop-streaming-*.jar \
    -D stream.num.map.output.key.fields=2 \
    -D mapreduce.partition.keypartitioner.options=-k1,1 \
    -D mapreduce.job.output.key.comparator.class=\
org.apache.hadoop.mapred.lib.KeyFieldBasedComparator \
    -D mapreduce.partition.keycomparator.options="-k1n -k2nr" \
    -files secondary_sort_map.py,secondary_sort_reduce.py \
    -input input/ncdc/all \
    -output output-secondarysort-streaming \
    -mapper ch09-mr-features/src/main/python/secondary_sort_map.py \
    -partitioner org.apache.hadoop.mapred.lib.KeyFieldBasedPartitioner \
    -reducer ch09-mr-features/src/main/python/secondary_sort_reduce.py
```

Our map function (**Example 9-7**) emits records with year and temperature fields. We want to treat the combination of both of these fields as the key, so we set stream.num.map.output.key.fields to 2. This means that values will be empty, just like in the Java case.

Example 9-7. Map function for secondary sort in Python

```
#!/usr/bin/env python
```

```
import sys

for line in sys.stdin:
  val = line.strip()
  (year, temp, q) = (val[15:19], int(val[87:92]), val[92:93])
  if temp == 9999:
    sys.stderr.write("reporter:counter:Temperature,Missing,1\n")
  elif re.match("[01459]", q):
    print "%s\t%s" % (year, temp)
```

However, we don't want to partition by the entire key, so we use KeyFieldBasedPartitioner, which allows us to partition by a part of the key. The specification mapreduce.partition.keypartitioner.options configures the partitioner. The value -k1,1 instructs the partitioner to use only the first field of the key, where fields are assumed to be separated by a string defined by the mapreduce.map.output .key.field.separator property (a tab charac-

mapreduce.map.output .key.field.separator property (a tab character by default).

Next, we want a comparator that sorts the year field in ascending order and the temperature field in descending order, so that the reduce function can simply return the first record in each group. Hadoop provides KeyFieldBasedComparator, which is ideal for this purpose. The comparison order is defined by a specification that is like the one used for GNU sort. It is set using the mapreduce.partition.keycomparator.options property. The value -k1n -k2nr used in this example means "sort by the first field in numerical order, then by the second field in reverse numerical order." Like its partitioner cousin, KeyFieldBasedPartitioner, it uses the map output key separator to split a key into fields.

In the Java version, we had to set the grouping comparator; however, in Streaming, groups are not demarcated in any way, so in the reduce function we have to detect the group boundaries ourselves by looking for when the year changes (**Example 9-8**).

Example 9-8. Reduce function for secondary sort in Python

```
#!/usr/bin/env python
```

```
last_group = None
for line in sys.stdin:
  val = line.strip()
  (year, temp) = val.split("\t")
  group = year
  if last_group != group:
    print val
    last_group = group
```

When we run the Streaming program, we get the same output as the Java version.

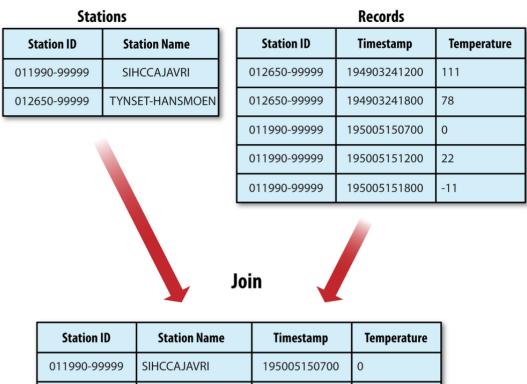
Finally, note that KeyFieldBasedPartitioner and KeyFieldBasedComparator are not confined to use in Streaming programs; they are applicable to Java MapReduce programs, too.

# **Joins**

MapReduce can perform joins between large datasets, but writing the code to do joins from scratch is fairly involved. Rather than writing MapReduce programs, you might consider using a higher-level framework such as Pig, Hive, Cascading, Cruc, or Spark, in which join operations are a core part of the implementation.

Let's briefly consider the problem we are trying to solve. We have two datasets—for example, the weather stations database and the weather records—and we want to reconcile the two. Let's say we want to see each station's history, with the station's metadata inlined in each output row. This is illustrated in <u>Figure 9-2</u>.

How we implement the join depends on how large the datasets are and how they are partitioned. If one dataset is large (the weather records) but the other one is small enough to be distributed to each node in the cluster (as the station metadata is), the join can be effected by a MapReduce job that brings the records for each station together (a partial sort on station ID, for example). The mapper or reducer uses the smaller dataset to look up the station metadata for a station ID, so it can be written out with each record. See **Side Data Distribution** for a discussion of this approach, where we focus on the mechanics of distributing the data to nodes in the cluster.



011990-99999 SIHCCAJAVRI 195005151200 22 011990-99999 SIHCCAJAVRI 195005151800 -11 012650-99999 TYNSET-HANSMOEN 194903241200 111 012650-99999 TYNSET-HANSMOEN 194903241800 78

Figure 9-2. Inner join of two datasets

If the join is performed by the mapper it is called a *map-side join*, whereas if it is performed by the reducer it is called a *reduce-side join*.

If both datasets are too large for either to be copied to each node in the cluster, we can still join them using MapReduce with a map-side or reduce-side join, depending on how the data is structured. One common example of this case is a user database and a log of some user activity (such as access logs). For a popular service, it is not feasible to distribute the user database (or the logs) to all the MapReduce nodes.

# **Map-Side Joins**

A map-side join between large inputs works by performing the join before the data reaches the map function. For this to work, though, the inputs to each map must be partitioned and sorted in a particular way. Each input dataset must be divided into the same number of partitions, and it must be sorted by the same key (the join key) in each source. All the

records for a particular key must reside in the same partition. This may sound like a strict requirement (and it is), but it actually fits the description of the output of a MapReduce job.

A map-side join can be used to join the outputs of several jobs that had the same number of reducers, the same keys, and output files that are not splittable (by virtue of being smaller than an HDFS block or being gzip compressed, for example). In the context of the weather example, if we ran a partial sort on the stations file by station ID, and another identical sort on the records, again by station ID and with the same number of reducers, then the two outputs would satisfy the conditions for running a map-side join.

You use a CompositeInputFormat from the org.apache.hadoop.mapreduce.join package to run a map-side join. The input sources and join type (inner or outer) for Composite Input Format are configured through a join expression that is written according to a simple grammar. The package documentation has details and examples.

The org.apache.hadoop.examples.Join example is a general-purpose command-line program for running a map-side join, since it allows you to run a MapReduce job for any specified mapper and reducer over multiple inputs that are joined with a given join operation.

# **Reduce-Side Joins**

A reduce-side join is more general than a map-side join, in that the input datasets don't have to be structured in any particular way, but it is less efficient because both datasets have to go through the MapReduce shuffle. The basic idea is that the mapper tags each record with its source and uses the join key as the map output key, so that the records with the same key are brought together in the reducer. We use several ingredients to make this work in practice:

## Multiple inputs

The input sources for the datasets generally have different formats, so it is very convenient to use the MultipleInputs class (see MultipleInputs) to separate the logic for parsing and tagging each source.

### Secondary sort

As described, the reducer will see the records from both sources that have the same key, but they are not guaranteed to be in any particular order. However, to perform the join, it is important to have the data from one source before that from the other. For the weather data join, the station record must be the first of the values seen for each key, so the reducer can fill in the weather records with the station name and emit them straightaway. Of course, it would be possible to receive the records in any order if we buffered them in memory, but this should be avoided because the number of records in any group may be very large and exceed the amount of memory available to the reducer.

We saw in <u>Secondary Sort</u> how to impose an order on the values for each key that the reducers see, so we use this technique here. To tag each record, we use TextPair (discussed in <u>Chapter 5</u>) for the keys (to store the station ID) and the tag. The only requirement for the tag values is that they sort in such a way that the station records come before the weather records. This can be achieved by tagging station records as <code>0</code> and weather records as <code>1</code>. The mapper classes to do this are shown in Examples <code>9-9</code> and <code>9-10</code>.

Example 9-9. Mapper for tagging station records for a reduce-side join

Example 9-10. Mapper for tagging weather records for a reduce-side join

```
public class JoinRecordMapper
    extends Mapper<LongWritable, Text, TextPair, Text> {
    private NcdcRecordParser parser = new NcdcRecordParser();

@Override
    protected void map(LongWritable key, Text value, Context context)
        throws IOException, InterruptedException {
        parser.parse(value);
        context.write(new TextPair(parser.getStationId(), "1"), value);
    }
}
```

The reducer knows that it will receive the station record first, so it extracts its name from the value and writes it out as a part of every output record (**Example 9-11**).

Example 9-11. Reducer for joining tagged station records with tagged weather records

```
public class JoinReducer extends Reducer<TextPair, Text, Text, Text> {
    @Override
    protected void reduce(TextPair key, Iterable<Text> values, Context context)
        throws IOException, InterruptedException {
        Iterator<Text> iter = values.iterator();
        Text stationName = new Text(iter.next());
        while (iter.hasNext()) {
            Text record = iter.next();
            Text outValue = new Text(stationName.toString() + "\t" + record.toString()
            context.write(key.getFirst(), outValue);
        }
    }
}
```

The code assumes that every station ID in the weather records has exactly one matching record in the station dataset. If this were not the case, we would need to generalize the code to put the tag into the value objects, by using another <code>TextPair</code>. The <code>reduce()</code> method would then be able to tell which entries were station names and detect (and handle) missing or duplicate entries before processing the weather records.

#### WARNING

Because objects in the reducer's values iterator are reused (for efficiency purposes), it is vital that the code makes a copy of the first Text object from the values iterator:

```
Text stationName = new Text(iter.next());
```

If the copy is not made, the stationName reference will refer to the value just read when it is turned into a string, which is a bug.

Tying the job together is the driver class, shown in **Example 9-12**. The essential point here is that we partition and group on the first part of the key, the station ID, which we do with a custom Partitioner (KeyPartitioner) and a custom group comparator, First Comparator (from TextPair).

Example 9-12. Application to join weather records with station names

```
public class JoinRecordWithStationName extends Configured implements Tool {
  public static class KeyPartitioner extends Partitioner<TextPair, Text> {
    @Override
    public int getPartition(TextPair key, Text value, int numPartitions) {
      return (key.getFirst().hashCode() & Integer.MAX VALUE) % numPartitions;
    }
  }
 @Override
  public int run(String[] args) throws Exception {
    if (args.length != 3) {
      JobBuilder.printUsage(this, "<ncdc input> <station input> <output>");
      return -1;
    }
    Job job = new Job(getConf(), "Join weather records with station names");
    job.setJarByClass(getClass());
    Path ncdcInputPath = new Path(args[0]);
    Path stationInputPath = new Path(args[1]);
    Path outputPath = new Path(args[2]);
```

```
MultipleInputs.addInputPath(job, ncdcInputPath,
        TextInputFormat.class, JoinRecordMapper.class);
   MultipleInputs.addInputPath(job, stationInputPath,
        TextInputFormat.class, JoinStationMapper.class);
    FileOutputFormat.setOutputPath(job, outputPath);
    job.setPartitionerClass(KeyPartitioner.class);
    job.setGroupingComparatorClass(TextPair.FirstComparator.class);
    job.setMapOutputKeyClass(TextPair.class);
    job.setReducerClass(JoinReducer.class);
    job.setOutputKeyClass(Text.class);
   return job.waitForCompletion(true) ? 0 : 1;
  }
 public static void main(String[] args) throws Exception {
    int exitCode = ToolRunner.run(new JoinRecordWithStationName(), args);
   System.exit(exitCode);
 }
}
```

Running the program on the sample data yields the following output:

```
011990-99999SIHCCAJAVRI0067011990999991950051507004...011990-99999SIHCCAJAVRI0043011990999991950051512004...011990-99999SIHCCAJAVRI0043011990999991950051518004...012650-99999TYNSET-HANSMOEN0043012650999991949032412004...012650-99999TYNSET-HANSMOEN0043012650999991949032418004...
```

## **Side Data Distribution**

Side data can be defined as extra read-only data needed by a job to process the main dataset. The challenge is to make side data available to all the map or reduce tasks (which are spread across the cluster) in a convenient and efficient fashion.

## Using the Job Configuration

You can set arbitrary key-value pairs in the job configuration using the various setter methods on Configuration (or JobConf in the old

MapReduce API). This is very useful when you need to pass a small piece of metadata to your tasks.

In the task, you can retrieve the data from the configuration returned by Context's getConfiguration() method. (In the old API, it's a little more involved: override the configure() method in the Mapper or Reducer and use a getter method on the JobConf object passed in to retrieve the data. It's very common to store the data in an instance field so it can be used in the map() or reduce() method.)

Usually a primitive type is sufficient to encode your metadata, but for arbitrary objects you can either handle the serialization yourself (if you have an existing mechanism for turning objects to strings and back) or use Hadoop's Stringifier class. The Default String ifier uses Hadoop's serialization framework to serialize objects (see Serialization).

You shouldn't use this mechanism for transferring more than a few kilobytes of data, because it can put pressure on the memory usage in MapReduce components. The job configuration is always read by the client, the application master, and the task JVM, and each time the configuration is read, all of its entries are read into memory, even if they are not used.

## **Distributed Cache**

Rather than serializing side data in the job configuration, it is preferable to distribute datasets using Hadoop's distributed cache mechanism. This provides a service for copying files and archives to the task nodes in time for the tasks to use them when they run. To save network bandwidth, files are normally copied to any particular node once per job.

#### **Usage**

For tools that use GenericOptionsParser (this includes many of the programs in this book; see GenericOptionsParser, Tool, and ToolRunner), you can specify the files to be distributed as a comma-separated list of URIs as the argument to the -files option. Files can be on the local filesystem, on HDFS, or on another Hadoop-readable filesystem (such as S3). If no scheme is supplied, then the files are assumed to be local. (This is true even when the default filesystem is not the local filesystem.)

You can also copy archive files (JAR files, ZIP files, tar files, and gzipped tar files) to your tasks using the -archives option; these are unarchived on the task node. The -libjars option will add JAR files to the classpath of the mapper and reducer tasks. This is useful if you haven't bundled library JAR files in your job JAR file.

Let's see how to use the distributed cache to share a metadata file for station names. The command we will run is:

```
% hadoop jar hadoop-examples.jar \
MaxTemperatureByStationNameUsingDistributedCacheFile \
-files input/ncdc/metadata/stations-fixed-width.txt input/ncdc/all output
```

This command will copy the local file *stations-fixed-width.txt* (no scheme is supplied, so the path is automatically interpreted as a local file) to the task nodes, so we can use it to look up station names. The listing for MaxTemperatureByStationNameUsing DistributedCacheFile appears in **Example 9-13**.

Example 9-13. Application to find the maximum temperature by station, showing station names from a lookup table passed as a distributed cache file

```
static class MaxTemperatureReducerWithStationLookup
    extends Reducer<Text, IntWritable, Text, IntWritable> {
  private NcdcStationMetadata metadata;
  @Override
  protected void setup(Context context)
      throws IOException, InterruptedException {
    metadata = new NcdcStationMetadata();
    metadata.initialize(new File("stations-fixed-width.txt"));
  }
  @Override
  protected void reduce(Text key, Iterable<IntWritable> values,
      Context context) throws IOException, InterruptedException {
    String stationName = metadata.getStationName(key.toString());
    int maxValue = Integer.MIN VALUE;
    for (IntWritable value : values) {
      maxValue = Math.max(maxValue, value.get());
    }
    context.write(new Text(stationName), new IntWritable(maxValue));
}
@Override
public int run(String[] args) throws Exception {
  Job job = JobBuilder.parseInputAndOutput(this, getConf(), args);
  if (job == null) {
    return -1;
  }
  job.setOutputKeyClass(Text.class);
  job.setOutputValueClass(IntWritable.class);
  job.setMapperClass(StationTemperatureMapper.class);
  job.setCombinerClass(MaxTemperatureReducer.class);
  job.setReducerClass(MaxTemperatureReducerWithStationLookup.class);
  return job.waitForCompletion(true) ? 0 : 1;
}
public static void main(String[] args) throws Exception {
  int exitCode = ToolRunner.run(
```

```
new MaxTemperatureByStationNameUsingDistributedCacheFile(), args);
System.exit(exitCode);
}
```

The program finds the maximum temperature by weather station, so the mapper (StationTemperatureMapper) simply emits (station ID, temperature) pairs. For the combiner, we reuse MaxTemperatureReducer (from Chapters 2 and 6) to pick the maximum temperature for any given group of map outputs on the map side. The reducer

(MaxTemperatureReducerWithStationLookup) is different from the combiner, since in addition to finding the maximum temperature, it uses the cache file to look up the station name.

We use the reducer's setup() method to retrieve the cache file using its original name, relative to the working directory of the task.

#### NOTE

You can use the distributed cache for copying files that do not fit in memory. Hadoop map files are very useful in this regard, since they serve as an on-disk lookup format (see <a href="MapFile">MapFile</a>). Because map files are collections of files with a defined directory structure, you should put them into an archive format (JAR, ZIP, tar, or gzipped tar) and add them to the cache using the -archives option.

Here's a snippet of the output, showing some maximum temperatures for a few weather stations:

PEATS RIDGE WARATAH	372
STRATHALBYN RACECOU	410
SHEOAKS AWS	399
WANGARATTA AERO	409
MOOGARA	334
MACKAY AERO	331

#### How it works

When you launch a job, Hadoop copies the files specified by the -files, -archives, and -libjars options to the distributed filesystem (normally HDFS). Then, before a task is run, the node manager copies the files from the distributed filesystem to a local disk—the cache—so the task can

access the files. The files are said to be *localized* at this point. From the task's point of view, the files are just there, symbolically linked from the task's working directory. In addition, files specified by -libjars are added to the task's classpath before it is launched.

The node manager also maintains a reference count for the number of tasks using each file in the cache. Before the task has run, the file's reference count is incremented by 1; then, after the task has run, the count is decreased by 1. Only when the file is not being used (when the count reaches zero) is it eligible for deletion. Files are deleted to make room for a new file when the node's cache exceeds a certain size—10 GB by default—using a least-recently used policy. The cache size may be changed by setting the configuration property

```
yarn.nodemanager.localizer.cache.target-size-mb.
```

Although this design doesn't guarantee that subsequent tasks from the same job running on the same node will find the file they need in the cache, it is very likely that they will: tasks from a job are usually scheduled to run at around the same time, so there isn't the opportunity for enough other jobs to run to cause the original task's file to be deleted from the cache.

## The distributed cache API

```
public void addCacheFile(URI uri)
public void addCacheArchive(URI uri)
public void setCacheFiles(URI[] files)
public void setCacheArchives(URI[] archives)
public void addFileToClassPath(Path file)
public void addArchiveToClassPath(Path archive)
```

Recall that there are two types of objects that can be placed in the cache: files and archives. Files are left intact on the task node, whereas archives are unarchived on the task node. For each type of object, there are three methods: an addCache XXXX () method to add the file or archive to the

distributed cache, a setCache XXXX s() method to set the entire list of files or archives to be added to the cache in a single call (replacing those set in any previous calls), and an add XXXX ToClassPath() method to add the file or archive to the MapReduce task's classpath. Table 9-7 compares these API methods to the GenericOptionsParser options described in Table 6-1.

Job API method	GenericOptionsParser equivalent	Description
addCache File(URI uri) set CacheFil es(URI[] files)	-files file1,file 2,	Add files to the distributed cache to be copied to the task node.
addCache Archive (URI ur i) setCa cheArchi ves(URI [] file s)	-archives archive1,a rchive2,	Add archives to the distributed cache to be copied to the task node and unarchived there.
addFileT oClassPa th(Path file)	-libjars jar1,jar 2,	Add files to the distributed cache to be added to the MapReduce task's classpath. The files are not unarchived, so this is a useful way to add JAR files to the classpath.
addArchi veToClas sPath(Pa th archi ve)	None	Add archives to the distributed cache to be unarchived and added to the MapReduce task's classpath. This can be useful when you want to add a directory of files to the classpath, since

you can create an

## Job API method

# GenericOptionsParser equivalent

## Description

archive containing the files. Alternatively, you could create a JAR file and use addFileToClas sPath(), which works equally well.

#### NOTE

The URIs referenced in the add or set methods must be files in a shared filesystem that exist when the job is run. On the other hand, the filenames specified as a GenericOptionsParser option (e.g., -files) may refer to local files, in which case they get copied to the default shared filesystem (normally HDFS) on your behalf.

This is the key difference between using the Java API directly and using GenericOptionsParser: the Java API does not copy the file specified in the add or set method to the shared filesystem, whereas the GenericOptionsParser does.

Retrieving distributed cache files from the task works in the same way as before: you access the localized file directly by name, as we did in **Example 9-13**. This works because MapReduce will always create a symbolic link from the task's working directory to every file or archive added to the distributed cache.  $\overline{^{[67]}}$  Archives are unarchived so you can access the files in them using the nested path.

# **MapReduce Library Classes**

Hadoop comes with a library of mappers and reducers for commonly used functions. They are listed with brief descriptions in <u>Table 9-8</u>. For further information on how to use them, consult their Java documentation.

#### **Classes**

## **Description**

ChainMapper, ChainReducer

Run a chain of mappers in a single mapper and a reducer followed by a chain of mappers in a single reducer, respectively.

(Symbolically, M+RM\*, where M is a mapper and R is a reducer.) This can substantially reduce the amount of disk I/O incurred compared to running multiple MapReduce jobs.

FieldSelectionMapReduc e (old API): FieldSelecti onMapper and FieldSelec tionReducer (new API) A mapper and reducer that can select fields (like the Unix cut command) from the input keys and values and emit them as output keys and values.

IntSumReducer, LongSumR
educer

Reducers that sum integer values to produce a total for every key.

InverseMapper

A mapper that swaps keys and values.

MultithreadedMapRunner (old API), MultithreadedM apper (new API)

A mapper (or map runner in the old API) that runs mappers concurrently in separate threads. Useful for mappers that are not CPU-bound.

TokenCounterMapper

A mapper that tokenizes the input value into words (using Java's Str ingTokenizer) and emits each word along with a count of 1.

Classes	Description
RegexMapper	A mapper that finds matches of a regular expression in the input value and emits the matches along with a count of 1.
Streaming applications—is to a	round for this problem—particularly in text-based dd an offset to eliminate all negative numbers and observe are the same number of characters.
<sup>62]</sup> See <mark>Sorting and merging Se</mark> he sort program example that	equenceFiles for how to do the same thing using comes with Hadoop.
	g ( <u>Sorting Data</u> ), Hive ( <u>Sorting and Aggregating</u> ), can sort with a single command.
or at least partially sorted. For $\epsilon$	mmon for some of the input to already be sorted, example, the weather dataset is ordered by time, tases, making the RandomSampler a safer choice.
	comparators as shown are not optimized; see ator for speed for the steps we would need to take

org. apache .hadoop .file cache.DistributedCache.

<sup>[67]</sup> In Hadoop 1, localized files were not always symlinked, so it was sometimes necessary to retrieve localized file paths using methods on <code>JobContext</code>. This limitation was removed in Hadoop 2.