

# KTH Challenge 2013

April 21, 2013

# Jury

KTH  
Challenge  
2013

Solutions

Further  
Information

- Lukáš Poláček (KTH, Spotify), head of jury
- Per Austrin (KTH)
- Oskar Werkelin Ahlin (Spotify)
- Ulf Lundström (KTH)
- Marc Vinyals (KTH)
- Erik Aas (KTH)
- Emma Enström (KTH)
- Andreas Lundblad (KTH)

# B – Peragrams

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Solutions

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Information

- Only one letter can have odd number of occurrences in a palindrome.
- We need to remove  $o - 1$  letters, where  $o$  is the number of letters with odd number of occurrences.
- Don't print  $-1$ !

Problem author: Oskar Werkelin Ahlin

Statistics: 89 submissions, 51 correct, first at 0:04:10.

S	A	T	O	R
A	R	E	P	O
T	E	N	E	T
O	P	E	R	A
R	O	T	A	S

Photo by Ross Beresford

# F – Bank Queue

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Solutions

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Information

- Create  $T$  time slots for  $T$  minutes.
- Put each person into their time slot.
- Process times slots from  $T - 1$  to 0 and add all the people to the set of candidates.
- At each time slot pick the person with the most money which hasn't been picked yet.
- Need fast data structure to get  $O(N \log N)$  time.



Problem author: Lukáš Poláček

Statistics: 93 submissions, 33 correct, first at 0:09:07.

# A – Car Game

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- Keep track of the first word for each possible licence plate (there are only 17 576 of them).
- For each word, list all license plates that fit.
  - Go through the word keeping a list of letters you have seen.
  - Use this to also keep a list of ordered pairs of letters.
  - Each such pair combined with a new letter gives a possible license plate.

Problem author: Ulf Lundström

Statistics: 135 submissions, 29 correct, first at 0:26:34.

# I – Flag Quiz

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Solutions

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Information

- Calculate  $d(i, j)$ , the distance between answer  $i$  and  $j$ .
- For each answer calculate the incongruousity – the maximum distance to other answers.
- Print all answers with the smallest maximum distance (incongruousity).



Problem authors: Ulf Lundström and Emma Enström  
Statistics: 57 submissions, 23 correct, first at 1:10:43.

# C – Vacuum Tubes

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Information

- Sort tubes:  $l_1 \leq \dots \leq l_L$ .
- For each tube  $i$ , find an index  $p_i$  such that  $l_i + l_{p_i} \leq L_1$  and  $p_i$  is as large as possible.
- Find similar index  $q_i$  for  $L_2$ .
- Try all  $i, j \in \{1, \dots, N\}$ .
- Try pairing  $i$  with  $p_i, p_i - 1, p_i - 2, p_i - 3$  and  $j$  with  $q_i, q_i - 1, q_i - 2, q_i - 3$ .  
Make sure we don't use a tube twice.
- Also possible in  $O(N \log N)$ .



Problem author: Ulf Lundström

Statistics: 47 submissions, 13 correct, first at 0:37:14.

# D – Chicken Joggers

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- Traverse the tree by depth-first-search, keep track of the distance from the root.
- If we can't visit any more intersections after visiting intersection  $u$  (leaf), we can decide if we need to put a lamp here.
- If  $u$  is not a leaf, we can decide whether we need a lamp by looking at already processed sons and edges going away from the root.



Photo by David Spencer

Problem author: Oskar Werkelin Ahlin

Statistics: 39 submissions, 7 correct, first at 0:44:26.

# H – Free Cell

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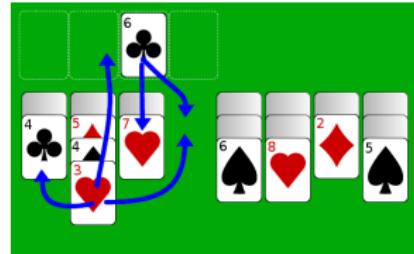
Solutions

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- We can move twice as many cards using  $M + 1$  empty stacks than using only  $M$ .
- We can move  $N + 1$  cards using 0 empty stacks and  $N$  free cells.
- Hence we can move at most  $(N + 1) \cdot 2^M$  cards.
- Slower solutions also worked.

Problem author: Andreas Lundblad

Statistics: 18 submissions, ?? correct, first at 1:42:57.



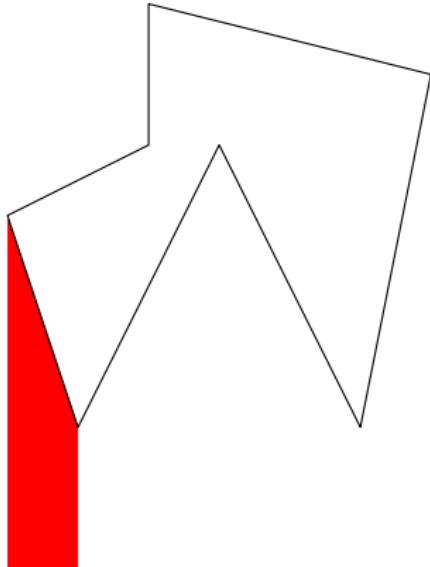
# Forest (1/2)

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Trapezoid method:



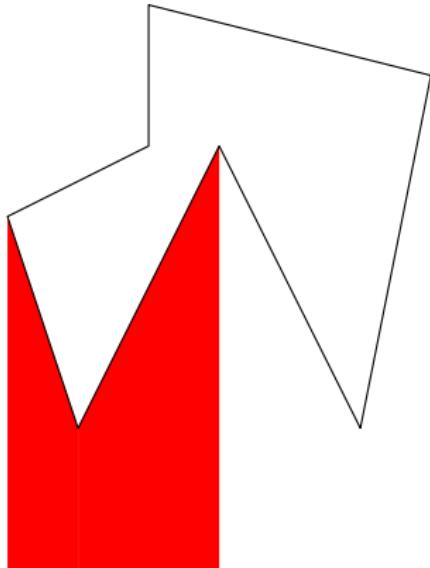
# Forest (1/2)

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Information

Trapezoid method:



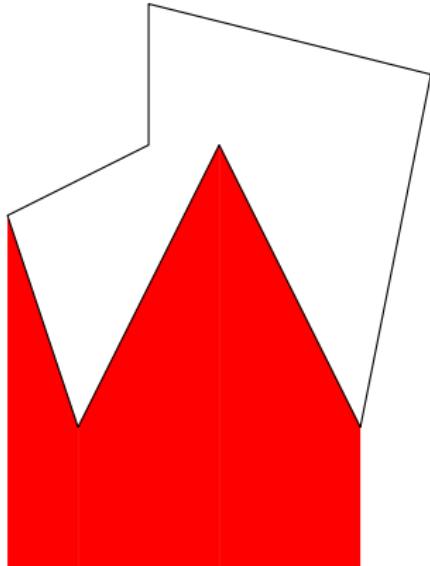
# Forest (1/2)

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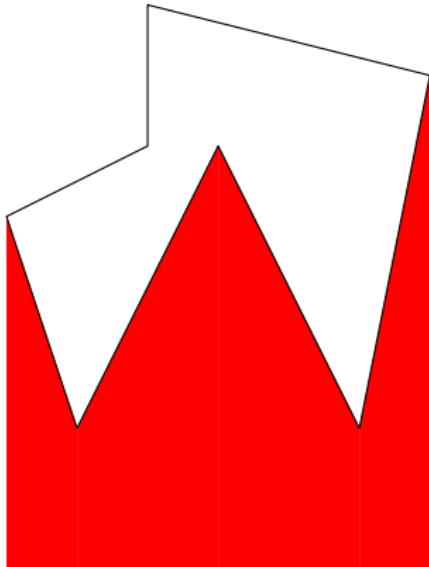
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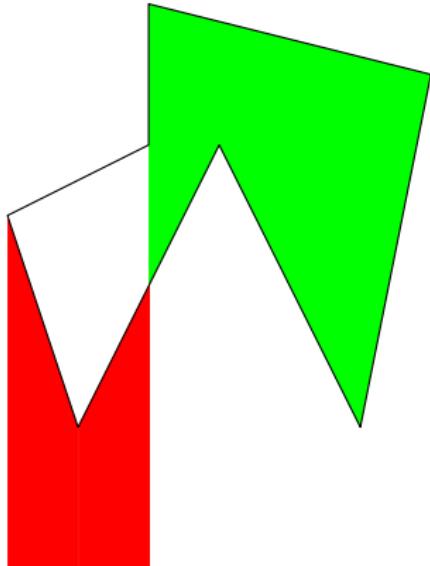
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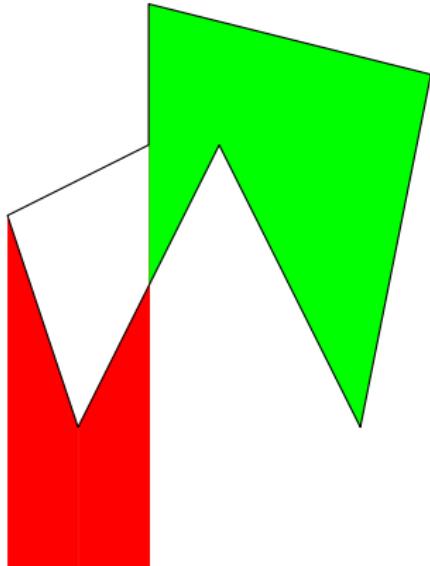
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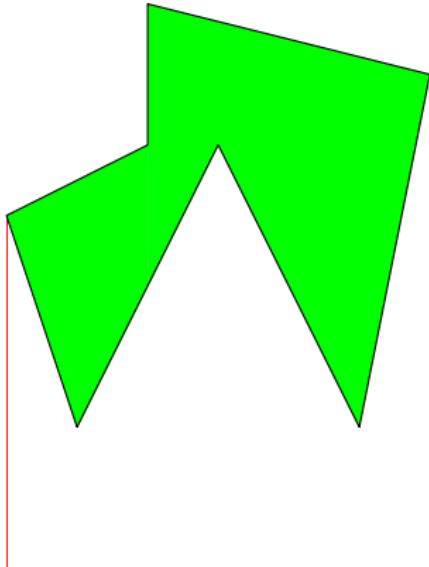
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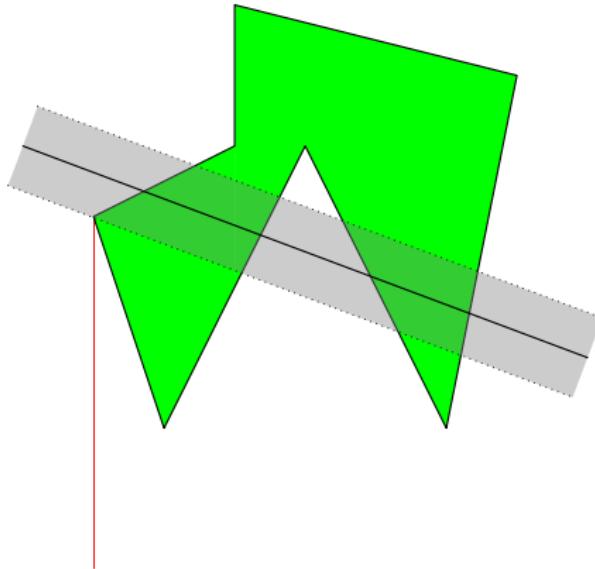
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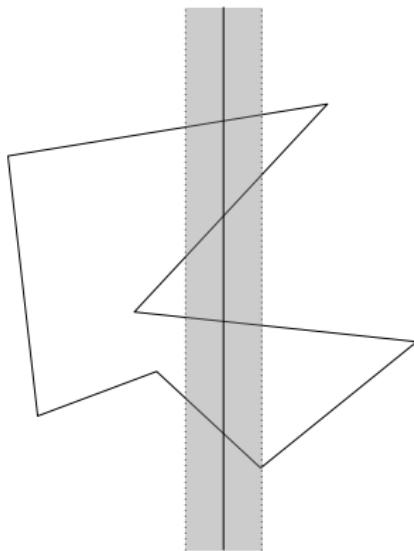
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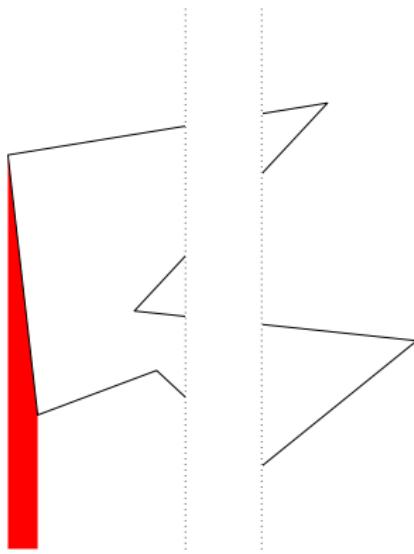
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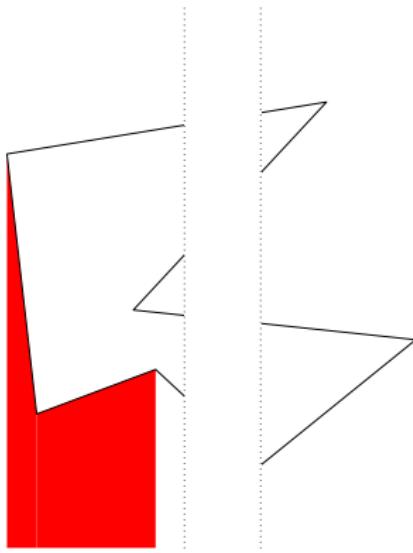
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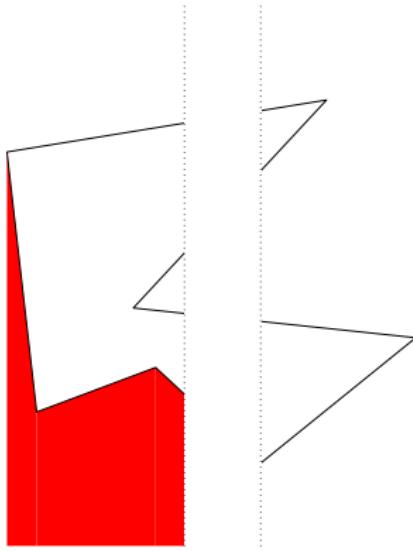
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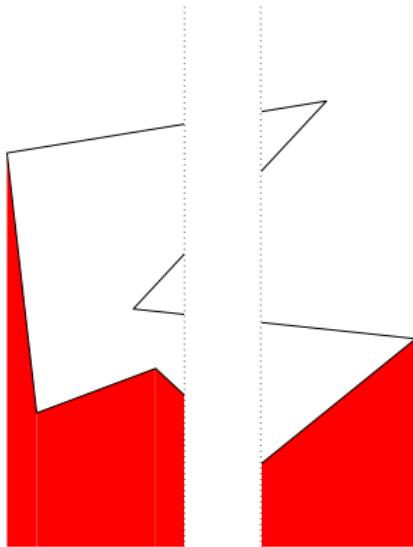
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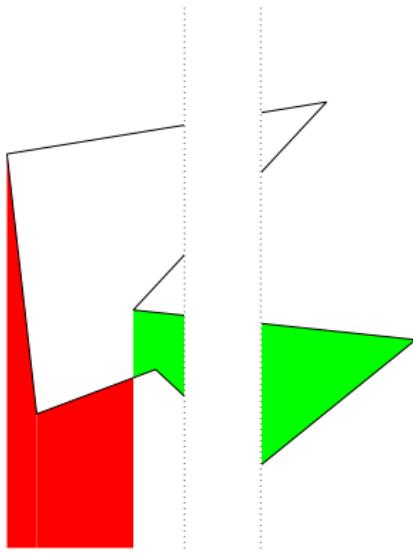
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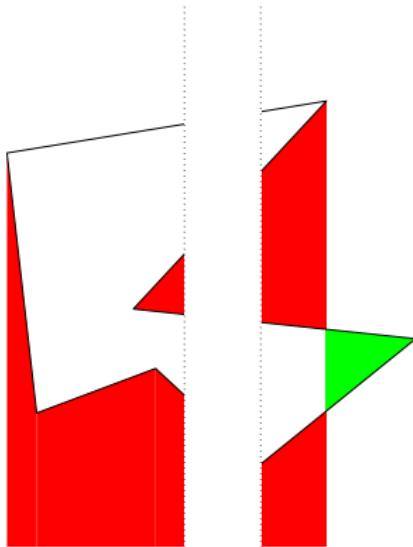
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Trapezoid method:

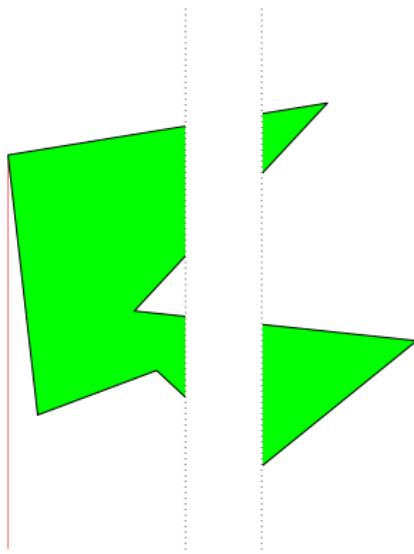


# Forest (1/2)

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Trapezoid method:



# Forest (2/2)

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Further  
Information

- Rotation is easy using complex numbers.

Problem author: Lukáš Poláček

Statistics: 5 submissions, ?? correct, first at ??.

# E – Hogwarts (1/2)

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Further  
Information

- Use black color for an existing edge, white color for missing edge.
- It's possible to rotate colors on a cycle  $u, v, w$  or a path of length 3  $u, v, w, x$ .
- Process edges in lexicographic order  $(0, 1), \dots, (0, N - 1), (1, 2), \dots, (N - 2, N - 1)$ :
  - Try to fix color of  $(i, j)$  by rotating a cycle  $i, j, k$ , such that  $i < j < k$ .
  - Otherwise try a path  $i, j, k, l$ , such that  $i < k < l$ .
  - We never change an edge that was already processed.



Photo by erinjudge

# E – Hogwarts (2/2)

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Solutions

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- In the end we might not be able to fix the last node – repeat the same process backwards.
- If this didn't succeed, randomly change labels from  $i$  to  $(i + c) \bmod N$  and try again.
- This works for big graphs, for small graphs use brute force.

Problem author: Erik Aas

Statistics: 17 submissions, ?? correct, first at ??.

# This was fun! When is the next contest?

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Information

- We train every two weeks at KTH, check [www.csc.kth.se/contest](http://www.csc.kth.se/contest).
- Next training on Wednesday at 17:15 in Orange.
- Nordic Championships in October, North-western Europe qualifier in November.
- Plenty of other online competitions every week.
- Subscribe to our calendar and RSS feed.

# Boot camp June 7 – June 9

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- 3 days  
on Möja in the archipelago.
- Lectures,  
trainings and fun activities.
- By invitation only.
- Also camp for Swedish IOI  
team and Linköping University.



Photo by The U.S. Army

# Guide To Programming Contests

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Solutions

Further  
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- <http://contest-wiki.csc.kth.se/>
- Written by Lukáš.
- Chapters “How to get better?” and “Team strategy” almost complete. More to come.
- The first training program for programming contests.
- Well received in the contest community.