

Sheet 11

Exercise 29

Proof (Lower Bound): Suppose C_n has a matching of size k . $\Rightarrow n - 2k$ nodes are unmatched. These $n - 2k$ nodes fall into k „buckets“ between matched edges around C_n .

If a bucket contains more than one node, the matching is not maximal.

\Rightarrow If $n - 2k \geq k + 1 \Rightarrow$ matching is not maximal

\Rightarrow For a maximal matching it must hold: $n - 2k \leq k \Rightarrow k \geq \frac{n}{3}$.

\Rightarrow lower bound for the size of a maximal matching is $\lceil \frac{n}{3} \rceil$

□

Exercise 30

Let $G = (V, E)$ be a bipartite graph.

Define $(G' = (V', E'), s, t, c)$ as follows:

$$V' = V_1 \cup V_2 \cup \{s, t\}$$

$$E' = \{(s, i) | i \in V_1\} \cup \{(j, t) | v \in V_2\} \cup E$$

$$c(u, v) = \begin{cases} \infty, & \text{if } u \in V_1, v \in V_2 \\ 1, & \text{otherwise} \end{cases}$$

Proof: Let H be a vertex cover of G .

Define

$$Q = \{(s, i) | i \in H \cap V_1\} \cup \{(i, t) | i \in H \cap V_2\}$$

Since H is a vertex cover, each directed path from s to t in G' contains one edge in Q .

$\Rightarrow Q$ is a valid cut of size $|H|$.

Let (S, T) be a cut of capacity k in G' .

Since edges $e \in E \setminus E'$ have a capacity of ∞ , the cut (S, T) only contains edges from $E' \setminus E$. Define

$$H = \{i | (s, i) \in S \times T \cap E'\} \cup \{i | (i, t) \in S \times T \cap E'\}$$

Observe that each edge $(i, j) \in E \setminus E'$ defines a directed path $s \rightarrow i \rightarrow j \rightarrow t$ in G' .

$\Rightarrow (s, i) \in S \times T$ or $(j, t) \in S \times T$.

\Rightarrow either $i \in H$ or $j \in H$.

$\Rightarrow H$ is a vertex cover of size k in G .

By the max-flow-min-cut Theorem it now holds that the maximum number of independent edges in G equals the minimum number of nodes in a vertex cover of G . \square