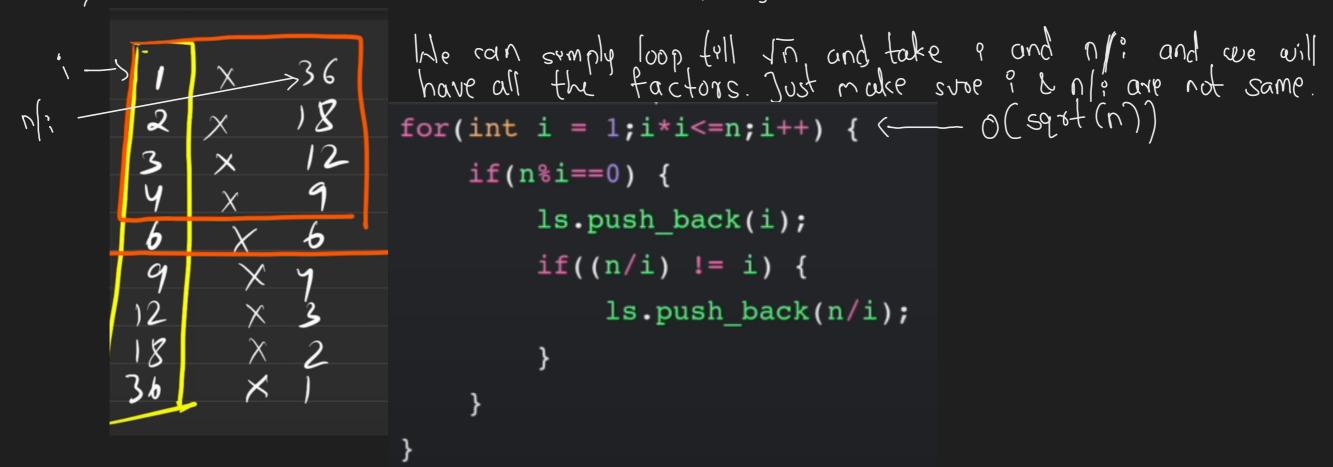
1) Digit extraction—To get last digit we do modulo 10)

Eig- 721 last digit 951. 721% 10=1 then we can divide number by 10 to get next digit.

TC-0(log to N) TC is log of N to the base 10 because each time we divide by 10.

- 2) Reverse number Almost semilar to above
- 3) Palendrome-Reverse et number is some ous original. Fond reverse l'compare with gaven number.
- 4) Flornstrong number Court of number of degits en number as power of each degit and their sum.

 Like 371= 32+73+13 = 371
- 5) Print all divisors = Divisor :s number which completely divide 1 (semainder zero)



6) Prome number - A number that is devesible by I and itself (I is not a prome number)

```
int n;
cin >> n;
int cnt = 0;
for(int i = 1; i*i<=n; i++) {
    if(n*i==0) {
        cnt++;
        if((n/i) != i) cnt++;
    }
}
if(cnt == 2) cout << "true";
else cout << "false";</pre>
We can find count of all factors of number and if
there count exceeds 2 then its not a prome number.

### There count exceeds 2 then its not a prome number.

#### There count exceeds 2 then its not a prome number.
```

7) acritical Here - areadest common Devisor OR Highest common Factors between two numbers is the largest common factor between the two which can divide both numbers

GCD (20,40) -> 20.

Flay two numbers always have a G.C.D of I at least.

A brute force approach would be to go from i-I to I-min(nino) & if both one divisible by P replace result with i.

We can also solve this problem using euclidean algorithm. The find gid (a-b, b) where a>b till one of them is zero and between the number that's not zero.

TC - O(log & (non (a,b)))

This rature keeps fluctualing depending on examples.

```
int gcd(int a,int b)
{
    while(a > 0 && b > 0) {
        if(a > b) a = a % b;
        else b = b % a;
    }
    if(a==0) return b;
    return a;
}
```

E) L(M - Lowest Common Multiple is the smallest value which can be divided by both no de no