

Fast Asynchronous Anti-TrustRank for Web Spam Detection

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Main Contributions

- Asynchronous Anti-TrustRank algorithms
 - Significantly reduce the number of arithmetic operations compared to the traditional synchronous Anti-TrustRank algorithm
 - Without degrading the performance in detecting Web spams
- Convergence of the asynchronous Anti-TrustRank algorithms
- Experiments on a real-world Web graph indexed by NAVER which is the most popular search engine in Korea.

Notation

- $G' = (\mathcal{V}, \mathcal{E}')$: a graph with reverse edges, i.e., if an edge $\{i, j\} \in \mathcal{E}$ then $\{j, i\} \in \mathcal{E}'$. Also, let \mathbf{A} denote the adjacency matrix of G' .
- $\mathbf{P} \equiv \mathbf{D}^{-1}\mathbf{A}$ (\mathbf{D} is the degree diagonal matrix)
- \mathcal{Q}_i : the set of incoming neighbors of node i on G'
- \mathcal{T}_i : the set of outgoing neighbors of node i on G'
- \mathbf{x} : a vector of the ATR scores, \mathbf{r} : a vector of the residuals
- \mathbf{e}_s : a vector with ones for the positions of the seed spam documents and zeros for other positions

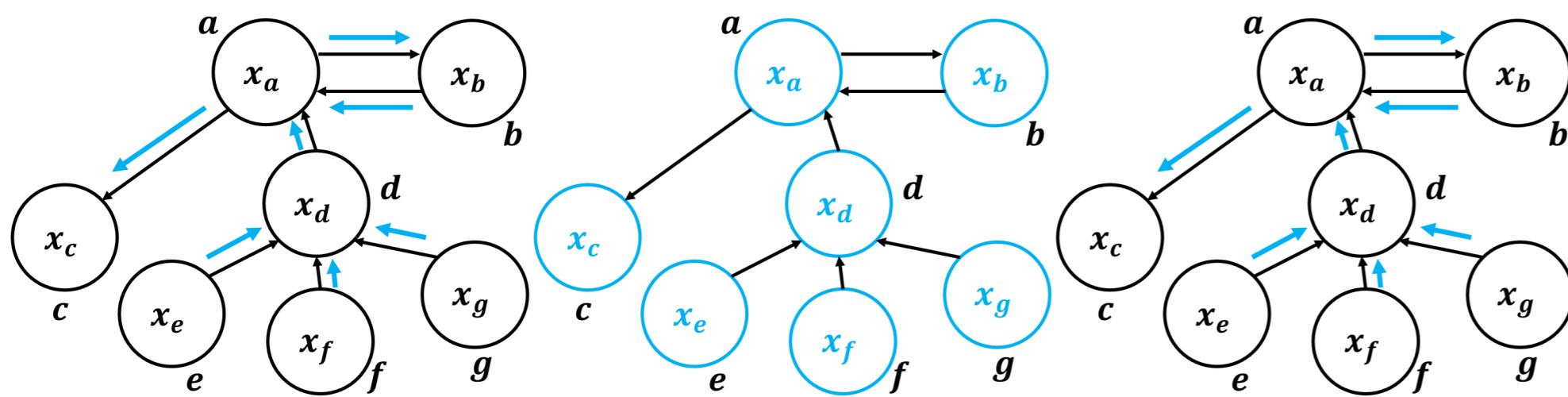
Anti-TrustRank*

- Spam pages are likely to be referred by other spam pages.
- Documents with high Anti-TrustRank (ATR) score \rightarrow spam pages
- From spam seeds, the ATR scores are propagated to incoming neighbors of the nodes so that the documents having links to the spam documents end up with having high ATR scores.

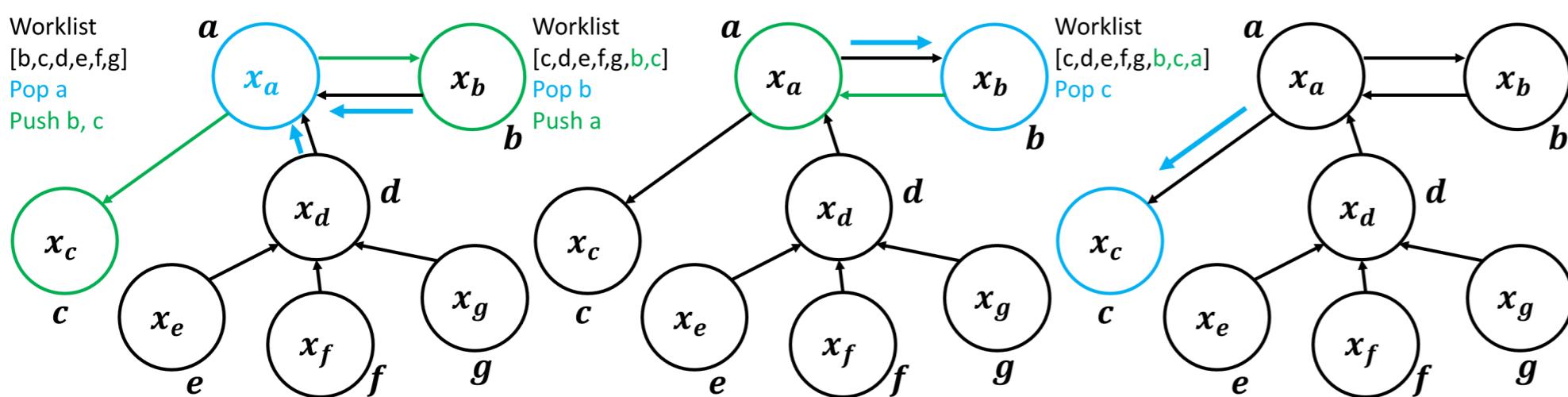
* V. Krishnan et al., Web spam detection with anti-trust rank. *AI/RWeb*, 2006.

Algorithms

- Synchronous Anti-TrustRank (SYNC ATR)
 - The scores are updated after all the nodes re-compute the scores.

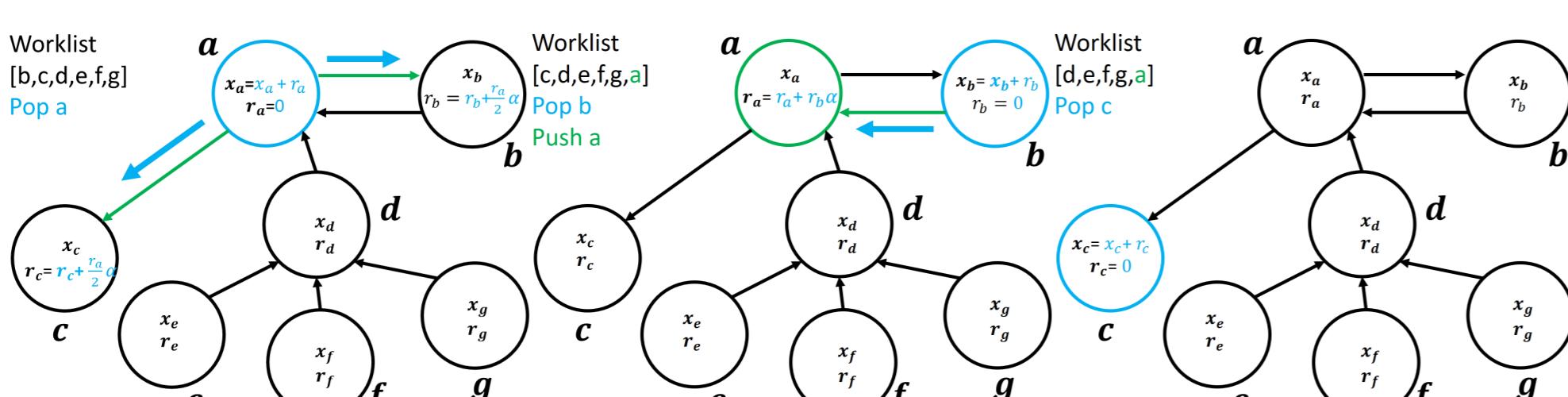


- Asynchronous Anti-TrustRank (ASYNC ATR)
 - worklist: a set of nodes whose ATR scores need to be updated.



- Residual-based Asynchronous Anti-TrustRank (RASYNC ATR)

- new ATR = current ATR + current residual (explicitly maintain the residual of each node)
- Filtering out unnecessary work in the worklist.



Pseudocodes

Algorithm: SYNC ATR

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Input:  $G' = (\mathcal{V}, \mathcal{E}')$ ,  $\mathcal{S}$ ,  $\alpha$ ,  $\epsilon$ 
Output: ATR vector  $\mathbf{x}$ 
1: Initialize  $\mathbf{x} = (1 - \alpha)\mathbf{e}_s$ 
2: while true do
3:   for  $i \in \mathcal{V}$  do
4:     if  $i \in \mathcal{S}$  then
5:        $x_i^{\text{new}} = \alpha \sum_{j \in \mathcal{Q}_i} \frac{x_j}{|\mathcal{T}_j|} + (1 - \alpha)$ 
6:     else
7:        $x_i^{\text{new}} = \alpha \sum_{j \in \mathcal{Q}_i} \frac{x_j}{|\mathcal{T}_j|}$ 
8:     end if
9:      $\delta_i = |x_i^{\text{new}} - x_i|$ 
10:   end for
11:    $\mathbf{x} = \mathbf{x}^{\text{new}}$ 
12:   if  $\|\delta\|_\infty < \epsilon$  then
13:     break;
14:   end if
15: end while
16:  $\mathbf{x} = \frac{\mathbf{x}}{\|\mathbf{x}\|_1}$ 

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Algorithm: ASYNC ATR

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Input:  $G' = (\mathcal{V}, \mathcal{E}')$ ,  $\mathcal{S}$ ,  $\alpha$ ,  $\epsilon$ 
Output: ATR vector  $\mathbf{x}$ 
1: Initialize  $\mathbf{x} = (1 - \alpha)\mathbf{e}_s$ 
2: for  $i \in \mathcal{V}$  do
3:    $wlist.push(i)$ 
4: end for
5: while ! $wlist.empty$  do
6:    $i = wlist.pop()$ 
7:   if  $i \in \mathcal{S}$  then
8:      $x_i^{\text{new}} = \alpha \sum_{j \in \mathcal{Q}_i} \frac{x_j}{|\mathcal{T}_j|} + (1 - \alpha)$ 
9:   else
10:     $x_i^{\text{new}} = \alpha \sum_{j \in \mathcal{Q}_i} \frac{x_j}{|\mathcal{T}_j|}$ 
11:   end if
12:   if  $|x_i^{\text{new}} - x_i| \geq \epsilon$  then
13:      $x_i = x_i^{\text{new}}$ 
14:     for  $j \in \mathcal{T}_i$  do
15:       if  $j$  is not in  $wlist$  then
16:          $wlist.push(j)$ 
17:       end if
18:     end for
19:   end if
20: end while
21:  $\mathbf{x} = \frac{\mathbf{x}}{\|\mathbf{x}\|_1}$ 

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Algorithm: RASYNC ATR

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Input:  $G' = (\mathcal{V}, \mathcal{E}')$ ,  $\mathcal{S}$ ,  $\alpha$ ,  $\epsilon$ 
Output: ATR vector  $\mathbf{x}$ 
1: Initialize  $\mathbf{x} = (1 - \alpha)\mathbf{e}_s$ 
2: Initialize  $\mathbf{r} = (1 - \alpha)\mathbf{P}^T \mathbf{e}_s$ 
3: for  $i \in \mathcal{V}$  do
4:    $wlist.push(i)$ 
5: end for
6: while ! $wlist.empty$  do
7:    $i = wlist.pop()$ 
8:    $x_i^{\text{new}} = x_i + r_i$ 
9:   for  $j \in \mathcal{T}_i$  do
10:     $r_j^{\text{old}} = r_j$ 
11:     $r_j = r_j + \frac{r_i \alpha}{|\mathcal{T}_j|}$ 
12:    if  $r_j \geq \epsilon$  and  $r_j^{\text{old}} < \epsilon$  then
13:       $wlist.push(j)$ 
14:    end if
15:  end for
16:   $r_i = 0$ 
17: end while
18:  $\mathbf{x} = \frac{\mathbf{x}}{\|\mathbf{x}\|_1}$ 

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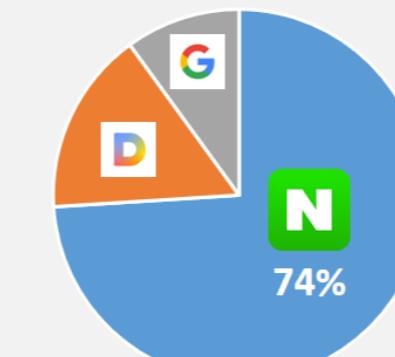
Experimental Results

Real-world Web graph from NAVER corporation

- 584,092 documents and 2,470,557 edges
- 437,386 (74.88%) normal docs and 45,641 (7.81%) spam docs
- 101,065 (17.30%) documents are unlabeled.

South Korea's No. 1 Search Portal

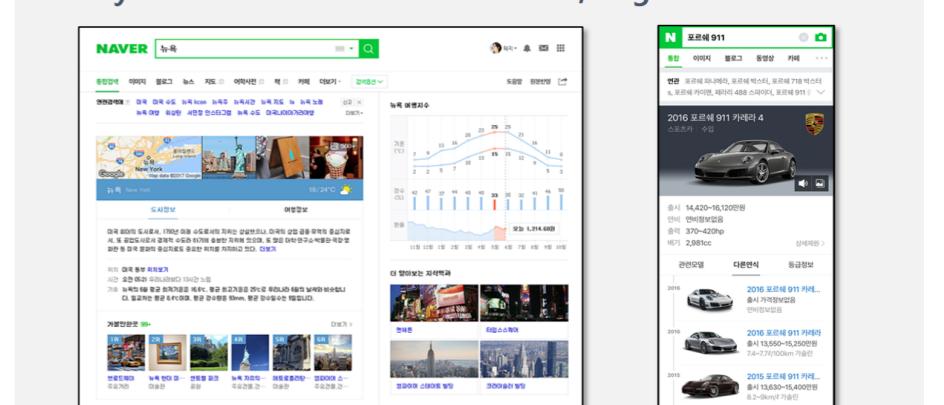
NAVER



(Nielsen KoreanClick Co., Ltd., 2016)

NAVER: Search Portal

- Most popular search portal in South Korea
- Search query share in Korea: 74.7% (As of June 2016)
- Daily visitors of mobile NAVER: 27M, Registered Users: 42M



- Most of the retrieved documents are correctly classified into spam.
- $|\mathcal{L}| = p|\mathcal{V}|$ where \mathcal{L} denotes the set of labeled documents
- Pick top $q|\mathcal{S}|$ documents where \mathcal{S} denotes the set of spam seeds

Table: Accuracy of the retrieved documents

| | | $q = 1$ | $q = 3$ | $q = 5$ |
|------------|----------------|--------------|------------------|------------------|
| $p = 0.01$ | spam docs | 1,367 (100%) | 4,099 (99.951%) | 6,833 (99.971%) |
| | normal docs | 0 (0%) | 0 (0%) | 0 (0%) |
| | unlabeled docs | 0 (0%) | 2 (0.049%) | 2 (0.029%) |
| $p = 0.02$ | spam docs | 3,083 (100%) | 9,113 (98.530%) | 15,279 (99.117%) |
| | normal docs | 0 (0%) | 107 (1.157%) | 107 (0.694%) |
| | unlabeled docs | 0 (0%) | 29 (0.314%) | 29 (0.188%) |
| $p = 0.03$ | spam docs | 3910 (100%) | 11,593 (98.832%) | 19,413 (99.299%) |
| | normal docs | 0 (0%) | 107 (0.912%) | 107 (0.547%) |
| | unlabeled docs | 0 (0%) | 30 (0.256%) | 30 (0.154%) |

- The asynchronous algorithms, ASYNC and RASYNC, make much fewer ATR updates than the synchronous algorithm, SYNC.
- RASYNC significantly reduces the number of arithmetic computations.

Table: No. of ATR updates and arithmetic operations

| p | ϵ | SYNC | | | ASYNC | | | RASYNC | | |
|------|------------|--------------------|--------------------|-----------|--------------------|--------------------|------------|--------------------|--------------------|------------|
| | | No. of ATR updates | No. of arithmetics | Sync | No. of ATR updates | No. of arithmetics | Sync | No. of ATR updates | No. of arithmetics | Sync |
| 0.01 | 10^{-8} | 2,336,368 | 24,442,660 | 20,361 | 2,336,368 | 24,424,970 | 20,361 | 2,336,368 | 24,424,970 | 20,361 |
| | 10^{-12} | 2,920,460 | 30,553,325 | 2,516,097 | 39,483 | 17,845,604 | 32,284,207 | 39,483 | 17,845,604 | 32,284,207 |
| 0.03 | 10^{-8} | 2,336,368 | 24,452,832 | 20,628 | 2,336,368 | 24,452,832 | 20,628 | 2,336,368 | 24,452,832 | 20,628 |
| | 10^{-12} | 2,920,460 | 30,566,040 | 2,703,630 | 39,804 | 17,845,604 | 32,284,207 | 39,804 | 17,845,604 | 32,284,207 |