#### Traversing graphs

Most basic algorithms on graphs will be applications of graph traversal.

- Printing or validating each edge/vertex.
- Copying a graph or converting between representations.
- Counting the number of edges/vertices.
- Identifying connected components.
- Finding paths between two vertices, or cycles.



#### **Efficiency and correctness**

- **Efficiency:** Don't loop or visit vertices repeatedly.
- Correctness: Don't miss any vertex.

We need to mark vertices as we traverse the graph.

- 1. Undiscovered: the initial state, before we've seen it.
- 2. Discovered: we've seen the vertex but not all of its incident edges.
- 3. Finished: all incident edges have been visited.



#### Order of exploration

The order in which we explore vertices depends on the container used for storing *discovered* but not *finished* vertices.

There are two types of containers used:

- Queue: leads to so called breadth-first search.
- Stack: leads to so called depth-first search.



#### **Breadth-first search**

- Properties
- The algorithm
- Example



#### Breadth-first search properties

- Input: A directed/undirected graph with source vertex s.
- Output: The shortest distance from s.
- ullet Only visits vertices reachable from s.
- Running time:  $\Theta(V+E)$

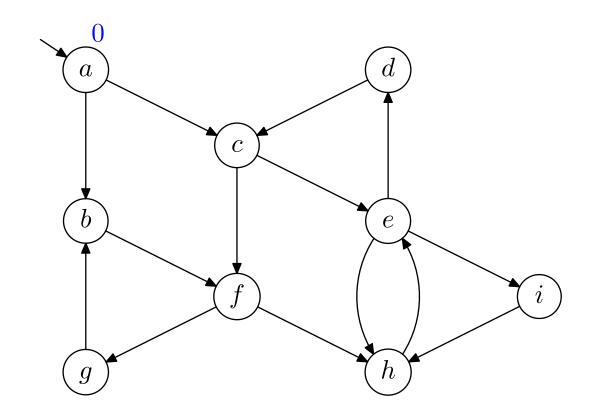
Linear time with respect to adjacency list.



#### Breadth-first search algorithm

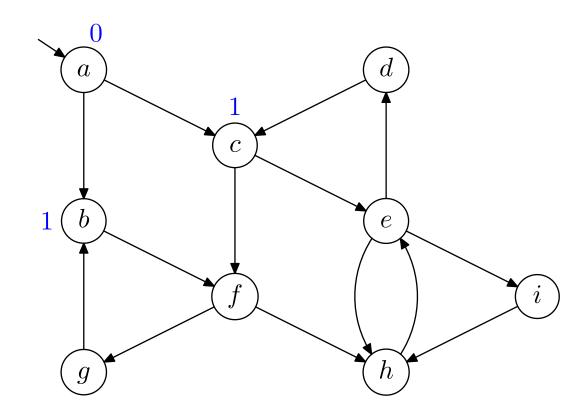
BFS(s)for each  $u \in V - \{s\}$ do  $d[u] \leftarrow \infty$  $d|s| \leftarrow 0$  $Q \leftarrow \{s\}$ while  $Q \neq \emptyset$ **do** remove u from Qfor each  $v \in Adj[u]$ do if  $d[v] = \infty$ then  $d[v] \leftarrow d[u] + 1$ put v onto Q





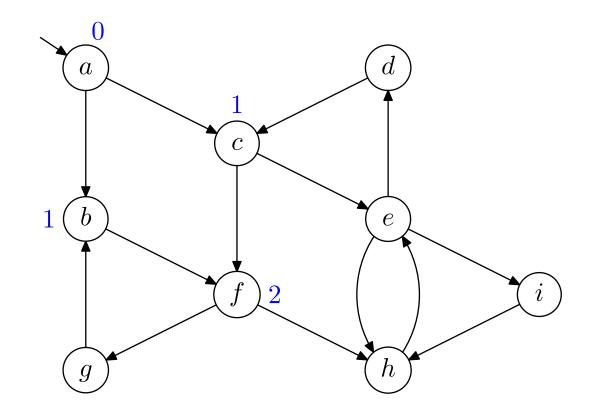
$$Q = \{a\}$$





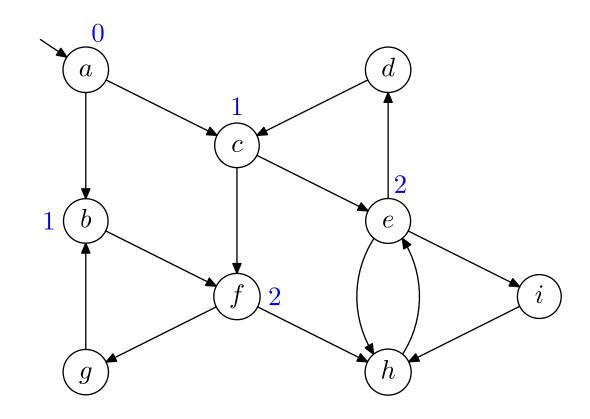
$$Q = \{ \not a, b, c \}$$





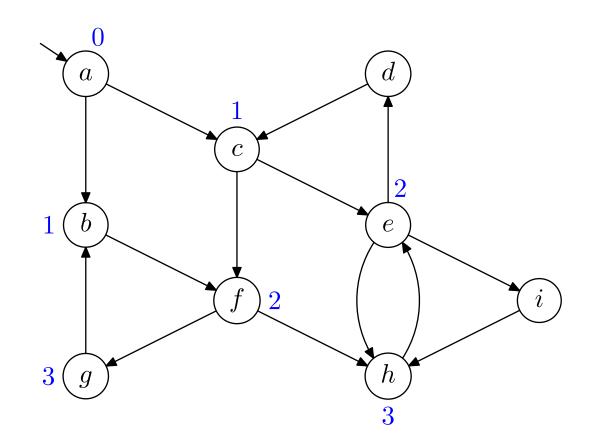
$$Q = \{ \not a, \not b, c, f \}$$





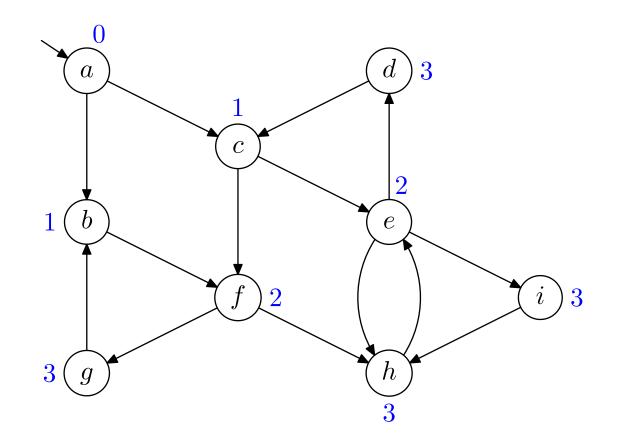
$$Q = \{ \not \Delta, \not b, \not c, f, e \}$$





$$Q = \{ \not a, \not b, \not c, \not f, e, g, h \}$$





$$Q = \{ \not a, \not b, \not c, \not f, \not e, g, h, d, i \}$$



#### Depth-first search

- Properties
- The algorithm
- Classification of vertices
- Classification of edges
- Example



#### Depth-first search properties

• **Input:** A directed/undirected graph.

#### • Output:

- ▷ Vertices are time-stamped: discover and finish.
- $\triangleright$  Edges  $\in \{ tree, back, forward or cross edge <math>\}$ .
- Visits all vertices.
- Running time:  $\Theta(V+E)$

Linear time with respect to adjacency list.



#### Depth-first search: algoritmen

```
DFS
```

```
\begin{aligned} & \textbf{for } \mathsf{each} \ u \in V \\ & \textbf{do} \ color[u] \leftarrow \mathsf{WHITE} \\ & time \leftarrow 0 \\ & \textbf{for } \mathsf{each} \ u \in V \\ & \textbf{do } \mathsf{if} \ color[u] = \mathsf{WHITE} \\ & \textbf{then } \mathsf{DFS-VISIT}(u) \end{aligned}
```

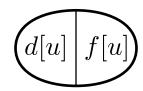


$$\begin{aligned} \operatorname{DFS-Visit}(u) \\ \operatorname{color}[u] &\leftarrow \operatorname{GRAY} & \rhd u \text{ has just been discovered} \\ \operatorname{time} &\leftarrow \operatorname{time} + 1 \\ d[u] &\leftarrow \operatorname{time} \\ \text{for each } v \in \operatorname{Adj}[u] \\ \text{do if } \operatorname{color}[v] &= \operatorname{WHITE} \\ \text{then } \operatorname{DFS-Visit}(v) \\ \operatorname{color}[u] &\leftarrow \operatorname{BLACK} & \rhd \text{ finished with } u \\ f[u] &\leftarrow \operatorname{time} \leftarrow \operatorname{time} + 1 \end{aligned}$$

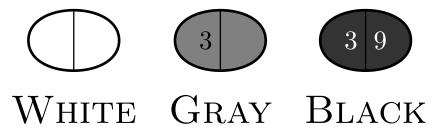


#### Depth-first search: vertex classification

Time-stamp vertices when they are discovered/finished:



Then vertex colors are equivalent to the following cases:



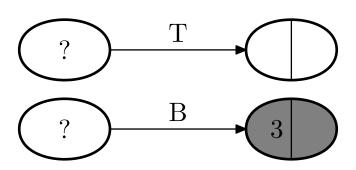


#### Depth-first search: edge classification

Edges are classified according to the following cases:

Tree edge

Back edge



$$d[u] < d[v]$$
: Forward edge (F)

$$d[u] > d[v]$$
: Cross edge (C)

