# SQL Server: Why Physical Database Design Matters

Module 3: Data Types and Index Size

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# Introduction: Does Data Type Choice Impact Indexes?

- Again, you might be thinking: disk space is cheap who cares?...
- Once you understand the basic indexes structures, you'll see how profound the effect of your choice can be
- We'll discuss the impact of data type choice by index type:
  - Clustered index
  - Nonclustered index
  - Columnstore index
- Throughout this module, I'll be discussing/demonstrating:
  - Key considerations around index structures
  - What the physical structures look like and how to analyze them
  - How does SQL Server access data (based on index usage)
  - What is the effect on performance

### What Structures Exist for a Table?

#### Table structure is either:

- Unordered: the table is called a heap
- Ordered: through creation of a clustered index and the table is called a clustered table (not to be confused with other RDBMS' clustered table)

#### Indexes

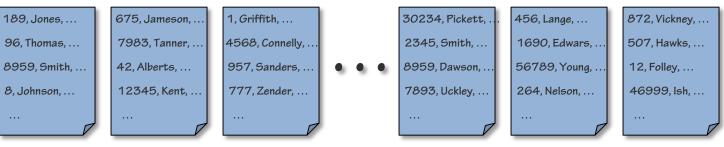
- Clustered: only one can exist per table as this defines the data's order
  - Not required, but <u>highly</u> recommended
  - The data is physically ordered at create/rebuild
  - The data is logically ordered through a doubly-linked list
  - Cluster key choice is CRITICAL!
- □ Nonclustered: not required, can have up to 999 of these (249 in 2000/2005)
  - DO NOT affect the base table's structure
  - Are affected by whether or not the table is clustered
- Hint: The nonclustered index dependency on the clustered index should impact your choice for the clustering key!

### What About Columnstore Indexes?

- Traditional clustered and nonclustered indexes are also known as "row-based" indexes
- Columnstore indexes are new in SQL Server 2012 and have a completely different internal structure as values for a single column are stored together
  - Data type has an impact but data distribution is even more interesting (in terms of possible column-level compression)
- These are beyond the scope of this course and have very specific/limited uses
  - SQL Server 2012: max of one nonclustered columnstore index per table and once created, the table is read-only
  - □ SOL Server 2014: columnstore indexes can be clustered and read/write
- Columnstore indexes will not be discussed here, check out Joe Sack's course SQL Server 2012: Nonclustered Columnstore Indexes

## **Table Structure: Heap**

- A table without a clustered index
- Records are not ordered and there is no doubly-linked list
- Accessed via allocation structures only so if no indexes exist then a full table scan is required for any SELECT query
- Imagine 80,000 records at 20 rows/per page = 4,000 pages
- Table scan costs at least 4,000 I/Os
  - Why "at least"?



File1, Page 497

File1, Page 498

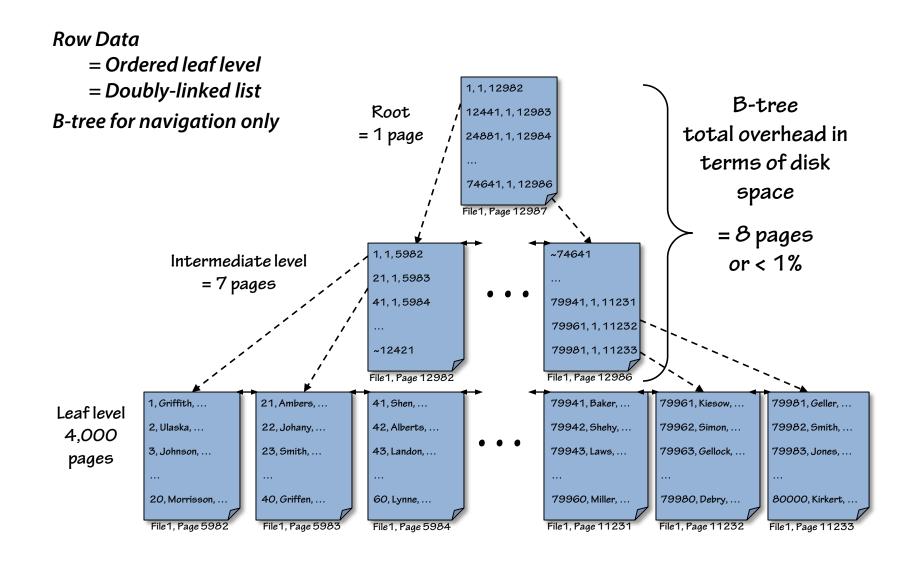
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File1, Page 5345

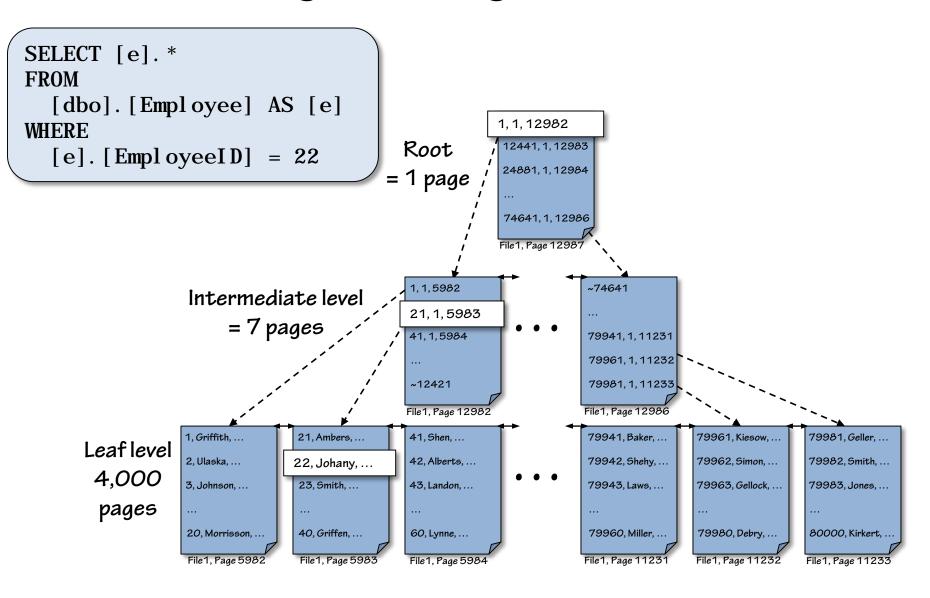
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4,000 pages of Employees in no specific order

### **Table Structure: Clustered Table**



## **Accessing Data Using a Clustered Index**



## **Nonclustered Indexes: The Book Analogy**

- Think of a book with indexes in the back
- The book has one form of logical ordering
- To lookup data, you use the indexes in the back
- Using a "Common Name" you look up that value in the index
- Once you find that index value then, you need to lookup the actual data based on its page in the book... i.e. a "bookmark" lookup
- The bookmark always depends on the book's content order

Index - Species Common Name

Index - Animal by Type, Name Bird, Mammal, Reptile, etc...

> Index - Animal by Country, Name

Index – Animals by Habitat, Name

Air, Land, Water

Index - Species Scientific Name

Index - Animal by Continent, Country, Name

## **Accessing Data Using a Nonclustered Index**

- Nonclustered indexes work a lot like an index in the back of a book:
  - Nonclustered index leaf level is the data defined by the index, in sorted order (exactly the same as a book)
  - Nonclustered index b-tree doesn't exist in a book but does in SQL Server.
    - Its use is just like the b-tree in a clustered index, i.e. navigational only
- Nonclustered indexes use the clustering key as the lookup value
  - Instead of a page number (or, physical locator which could change), SQL
     Server uses the clustering key as the lookup ID in a nonclustered index
    - Can be both good OR bad, depending on the key
- Remember: the nonclustered index dependency on the clustered index should impact your choice for the clustering key!
  - This is why...

## **Clustering Key Usage in Nonclustered Indexes**

Imagine the internals of a nonclustered index on SocialSecurityNumber on three different versions of the Employee table each with a different clustering key

SSN	Lookup	Uniquifier	ī	SSN	Lookup	Г	55N	Lookup
000-00-0184	Smith	0 (0 bytes)	Ш	000-00-0184	92CF41D7-17BF-49F7-		000-00-0184	31101
000-00-0236	Jones	1 (4 bytes)	Ш		B5C8-D3246C19B302 2F87EEBB-FBA1-4C06-		000-00-0236	22669
000-00-0395	Smith	1 (4 bytes)	Ш	000-00-0236	B7F1-BE63285B5935		000-00-0395	18705
000-00-0418	Jones	0 (0 bytes)	I	000-00-0395	2EF09CA4-6E48-47AA- A688-3D9FDEA220E0		:	:
The lookup value is non-unique (and wide as the column type of nvarchar (40)).  Also, what if there are multiple rows with the same lastname (Smiths/Jones/Anderson)?  CL: Lastname				The lookup value is a GUID = 16 bytes  CL: GUID		The lookup value is an int = 4 bytes  CL: EmployeeID		

Each table starts at 80,000 rows over 4,000 pages (due to the average row size of 400 bytes/row and therefore 20 rows/page). Then EACH/EVERY index must include the (entire) lookup value.

## **Clustering Key Widens Nonclustered Indexes**

- Imagine a real-world scenario
  - Table has 8 nonclustered indexes and 10 million rows
- What's the overhead required (and total space) for the bookmark lookups in the nonclustered indexes:
  - With a clustering key of an int (4 bytes)
  - With a clustering key of a GUID (16 bytes)
  - With a natural key (6 columns and ~64 bytes)
    - NOTE: This is just the overhead of the data type without factoring in nullable/non-unique.

Simple calculations for <u>overhead</u> in the LEAF level of the nonclustered indexes based on CL key columns defined					
CL Key Column(s)	Width of CL key (bytes)	МВ			
int	4	305.18			
datetime	8	610.35			
datetime, int	12	915.53			
guid	16	1,220.70			
composite	32	2,441.41			
composite	64	4,882.81			

### **Nonclustered Index Overhead**

- Table has 8 nonclustered indexes and 10 million rows
- What is the required disk space for placing the clustering key in each and every nonclustered index
- Add required overhead for nullability as well as whether the column is unique vs. non-unique
- Keys of: int / bigint / datetime, int /
   GUID are likely to be unique and non-nullable (marked with \*)

NOTE: Did not factor in additional overhead for composite keys and the number of variable-width columns they might have.

CL Key Column(s)	Bytes	MB
int*	4	305.18
int, nullable	7	534.06
int, non-unique (min)	4	305.18
int, non-unique (max)	12	915.53
int, non-unique (min), nullable	7	534.06
int, non-unique (max), nullable	15	1,144.41
bigint *	8	610.35
bigint, nullable	11	839.23
datetime, int *	12	915.53
datetime, int, nullable	15	1,144.41
guid *	16	1,220.70
guid, nullable	19	1,449.58
composite 32 bytes (comp32)*	32	2,441.41
comp32, nullable	35	2,670.29
comp32, non-unique (min)	32	2,441.41
comp32, non-unique (max)	40	3,051.76
comp32, non-unique (min), nullable	35	2,670.29
comp32, non-unique (max), nullable	43	3,280.64
composite 64 bytes (comp64)*	64	4,882.81
comp64, nullable	67	5,111.69
comp64, non-unique (min)	64	4,882.81
comp64, non-unique (max)	72	5,493.16
comp64, non-unique (min), nullable	67	5,111.69
comp64, non-unique (max), nullable	75	5,722.05
composite 128 bytes (comp128)*	128	9,765.63
comp128, nullable	131	9,994.51
comp128, non-unique (min)	128	9,765.63
comp128, non-unique (max)	136	10,375.98
comp128, non-unique (min), nullable	131	9,994.51
comp128, non-unique (max), nullable	139	10,604.86

## Is it Really That Much Space?

What about 100 million rows with 12 nonclustered indexes?

Simple calculations for <u>overhead</u> in the LEAF level of the nonclustered indexes based on CL key columns defined						
CL Key Column(s)	Width of CL key (bytes)	МВ				
int	4	4,577.64				
bigint	8	9,155.27				
datetime, int	12	13,732.91				
guid	16	18,310.55				
composite32, nullable	35	40,054.32				
composite64, nullable	67	76,675.42				
composite128, nullable	131	149,917.60				

- You're looking at <u>gigabytes</u> of storage, memory, backups
- Insert/update performance (logging)
- Maintenance requirements
- My point is that it really does add up
- It is something you need to strategize/analyze and DESIGN!

### **Clustered Index Criteria**

- How do you keep your clustering key as streamlined as possible?
  - Unique
    - Yes: No extra time/space overhead, data takes care of this criteria
    - NO: SQL Server must "uniquify" the rows on INSERT
  - Static
    - Yes: Reduces overhead
    - NO: Costly to maintain during updates to the key
  - Narrow
    - Yes: Keeps the nonclustered indexes narrow
    - NO: Unnecessarily wastes space
  - Non-nullable/fixed-width
    - Yes: Reduces overhead
    - NO: Adds overhead to ALL nonclustered indexes
  - Ever-increasing key value
    - Yes: Reduces index fragmentation
    - NO: Inserts/updates might cause significant index fragmentation

## **Choose a GOOD Clustering Key**

#### Identity column

 Adding this column and clustering on it can be extremely beneficial, even when you don't "use" this data

#### DateCol, identity

- Composite key defined in that order
  - Do not use date alone as that would need to be "uniquified"
- Great clustering key for partitioned tables
- Ideal where you have a lot of data-related queries (even if not partitioned)

#### GUID

- NO: if populated by client-side call to .NET client to generate the GUID
  - OK as the primary key but not as the clustering key
- NO: if populated by server-side NEWID() function
  - OK as the primary key but not as the clustering key
- Maybe: if populated by the server-side NEWSEQUENTIALID() function as it creates a more sequential pattern (and therefore less fragmentation)
  - But, this isn't really why you chose to use a GUID...

# Primary Key does NOT have to be the Clustering Key

- Primary key: relational integrity
- Clustering key: internal mechanism for looking up rows (bookmark lookup)
- SQL Server enforces uniqueness of a primary key through an index and defaults to clustered
  - 1 clustered index per table
  - 1 primary key per table
- If the primary key is a natural key then you probably want to enforce it with a nonclustered index
  - Might be very wide
  - Very expensive to duplicate in each and every nonclustered index
- If the table doesn't have a column (or small set of columns) that meets these criteria then consider adding a surrogate [identity] key and then create the clustered index on it

### **Scenario: What is the Real Cost?**

- AdventureWorksDW.dbo.FactInternetSales
- Clustered index (composite index of two seemingly narrow columns):
  - SalesOrderNumber type: nvarchar(20)
  - SalesOrderLineNumber type: tinyint
- Nonclustered indexes (all, single-column nonclustered):
  - IX\_FactIneternetSales\_ShipDateKey: ShipDateKey
  - IX\_FactInternetSales\_CurrencyKey: CurrencyKey
  - IX\_FactInternetSales\_CustomerKey: CustomerKey
  - IX\_FactInternetSales\_DueDateKey: DueDateKey
  - IX\_FactInternetSales\_OrderDateKey: OrderDateKey
  - IX\_FactInternetSales\_ProductKey: ProductKey
  - IX\_FactInternetSales\_PromotionKey: PromotionKey
- Everything seems narrow and somewhat optimal, until we talk about the data types of the clustering key columns

## Scenario: What's in That Key?

- What? What does the data look like?
  - SalesOrderNumber = 7 characters (SO12345) which is 14 bytes of data
    - □ EVERY row is SO + 5-digit number why?
    - Then, add variable-width overhead (each column has 2 bytes in the variable block as an offset)
    - When this is the first (or only) variable-width column then the addition of a variable block within the row adds 2 more bytes
  - $\Box$  Each of these 7 character "numbers" requires 18 bytes (14 + 2 + 2)
- If the clustering key requires it then EVERY nonclustered index requires it
- 7 nonclustered indexes:
  - Ironically, ALL columns (of all nonclustered indexes) are:
    - Non-nullable and fixed-width
  - They do not require a variable-block on their own, but when you add the clustering key

### **Scenario: What's the Total Cost?**

- What's the physical cost of this poorly defined column:
   SalesOrderNumber vs. an int (and ditching the type since there's only one type in all of the data)
- 14 bytes wasted per row, per index
- 7 nonclustered indexes x 14 = 98 bytes (completely wasted) per row
- What if the table were larger?
  - □ Imagine 10 million rows and 10 nonclustered indexes:
    - $\Box$  10,000,000 x 140 / 1024 / 1024 = **1.335 GB** of nonsense
  - Imagine 100 million rows and 10 nonclustered indexes:
    - $\Box$  100,000,000 x 140 / 1024 / 1024 = **13.35 GB** of nonsense
  - □ Imagine 1 billion rows and 12 nonclustered indexes:
    - $\Box$  1,000,000,000 x 154 / 1024 / 1024 = **143.42 GB** of nonsense

## **Summary: The Effect of Data Type Choice**

- Data type choice can have a profound affect on:
  - Column size
  - Row size
  - Index size
- Everything is less efficient/effective when design is poor
  - Waste disk space
  - Waste cache (and this is still very costly)
  - Larger backups, more time to backup
  - Logging is more costly with wider rows (DML is negatively affected)
  - Maintenance is more costly (and possibly required more)
  - Queries can be less efficient
    - Required to put data into memory than really needed
    - Returning more data than what's necessary