



## IIUC Victory Day Programming Contest 2003

<b>Problem I</b>	<b>Right Words</b>
<b>Time Limit</b>	2 Seconds

A Context Free Grammar (CFG) consists of the followings: a set of *nonterminal* symbols **V**; a set of *terminal* symbols **T**; a special *nonterminal* symbol called the *root* and a set of *production rules*. If all the production rules are either of the form **A**  $\rightarrow$  **BC**, or **A**  $\rightarrow$  **a**, where **A**, **B**, **C** is a member of set **V** and **a** is a member of set **T** then we say that the grammar is in Chomsky Normal Form (CNF).

If we repeatedly apply the production rules over the *root* symbol we will finally end up with a string of *terminals*. Alternatively, we can start with a string of *terminals* and reduce it using given production rules.

For example the string **ab** can be obtained by the first CFG presented in the sample input in the following way:

```
S -> AB
AB -> aB ; because A -> a
aB -> ab ; because B -> b
```

But, we cannot obtain **a** from **S** by applying the production rules. The set of strings of *terminals* derivable from the *root* symbol of a CFG is called the Language of the CFG. In this problem you are required to determine whether a given string of *terminals* is in the Language of a CFG or not.

### Input

There will be several test cases in the input. Each test case describes a CFG in Chomsky Normal Form and will adhere to the following description. In the first line there will be the *root* symbol. It will always be an uppercase English letter. In line 2 the set **V** will be presented as a string of uppercase letters. Each character of the string will be identified as a member of **V**. The set **T** will be given as a string of printable characters (except # or any whitespace characters) in line 3. Each character of the string will be identified as a member of **T**. Then there will be several lines for each production rule. A production rule will be of the form **A**  $\rightarrow$  **BC** or of the form **A**  $\rightarrow$  **a**. Here **A**, **B**, **C** are from set **V** and **a** is from set **T**. A production rule of the form #  $\rightarrow$  # indicates the end of production rules. After that there will be several lines each containing a candidate string of printable characters. This string will not contain any character from **V** and there will be no more than 50 characters in it. The list of candidate strings will be terminated with a line containing a # in the first column.

### Output

For each candidate string **a** print "**a** is in L (G)" if it can derived from the given grammar otherwise print "**a** is not in L (G)". Output a blank line after each test case.



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