Homework 2

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1.

- 1) Divide A[1..n] into three segments: A[1,1third], A[1third+1, 2third], A[2third+1,n].
- 2) Find three maximum subarrays which are entirely in A[1,1third]], A[1third+1, 2third], A[2third+1,n].
- 3) Find maximum subarray crossing boundarys:

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Find-Max-Cross-Subarray(A,low,1third,2third,high)
        Sum=0
        Left-sum= A[1third]
        For i=1third-1 downto low
                Sum=sum + A[i]
                If sum > left-sum
                        Left-sum=sum
                        Max-left=i
        Left-sum-2=-1000
        For i=2third downto 1third+1
                Sum= sum+ A[i]
                If sum>left-sum-2
                        Left-sum-2=sum
                        Max-left-2=i
        If left-sum-2 > left-sum
                Left-sum=left-sum-2
                Max-left=max-left-2
        Sum=0
        right-sum=-1000
        For j=1third+1 upto 2third
                Sum=sum + A[j]
                If sum>right-sum
                        Right-sum=sum
                        Max-right=i
                        If max-right=2third
                                For k=2third+1 upto high
                                        Sum = sum + A[k]
                                        If sum>right-sum
                                                Right-sum=sum
                                                Max-right=k
Return(max-left, max-right, left-sum + right-sum)
Total time: T(n)=3T(n/3)+\Theta(4n/3)
T(n) = \Theta(n \log n)
```

[Type here]

2.

a.

In each list, the worst case running time is $\Theta(k^2)$. There are n/k lists in total, so the worst case total running time is $\Theta(k^2 \frac{n}{k}) = \Theta(k)$

b.

The depth of merging from 1 element lists is $\log(n)$. We start from length k lists, so the depth is $\log(n)$ - $\log(k)$. And the running time in each level is still $\Theta(n)$. So total running time is $\Theta\left(n\log\left(\frac{n}{k}\right)\right)$

c.

The standard merge sort running time is $\Theta(nlogn)$. Let

 $\Theta(nlogn) = \Theta\left(nk + nlog\left(\frac{n}{k}\right)\right) = \Theta(nk + nlogn - nlogk)$. In the right term, nk already dominates nlogk, and as long as k grows no faster than logn, the equation stands. So k(n) = O(logn).

d.

Choose k = logn can minimize the running time.

3.

Input: two length n arrays A and B, each contains binary values of a and b

Output: a length n+1 array C, containing binary values of a+b

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 \begin{aligned} &\text{C=zeros}(1,\,\text{n+1})\\ &\text{For i=n downto 1}\\ &\text{C[i+1]=(A[i]+B[i]+c)mod2}\\ &\text{If A[i]+B[i]+c>= 2}\\ &\text{c=1}\\ &\text{else}\\ &\text{c=0}\\ &\text{C[1]=c} \end{aligned}
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[Type here]

4.

Bubble sort:

- [25, 13, 10, 16, 5, 8, 1]
- [25, 13, 10, 16, 5, 1, 8]
- [25, 13, 10, 16, 1, 5, 8]
- [25, 13, 10, 1, 16, 5, 8]
- [25, 13, 1, 10, 16, 5, 8]
- [25, 1, 13, 10, 16, 5, 8]
- [1, 25, 13, 10, 16, 5, 8]
- [1, 25, 13, 10, 5, 16, 8]
- [1, 25, 13, 5, 10, 16, 8]
- [1, 25, 5, 13, 10, 16, 8]
- [1, 5, 25, 13, 10, 16, 8]
- [1, 5, 25, 13, 10, 8, 16]
- [1, 5, 25, 13, 8, 10, 16]
- [1, 5, 25, 8, 13, 10, 16]
- [1, 5, 8, 25, 13, 10, 16]
- [1, 5, 8, 25, 13, 10, 16]
- [1, 5, 8, 25, 10, 13, 16]
- [1, 5, 8, 10, 25, 13, 16]
- [1, 5, 8, 10, 13, 25, 16]
- [1, 5, 8, 10, 13, 16, 25]

Insertion sort:

- [25, 13, 10, 16, 5, 8, 1]
- [13, 25, 10, 16, 5, 8, 1]
- [10, 13, 25, 16, 5, 8, 1]
- [10, 13, 16, 25, 5, 8, 1]
- [5, 10, 13, 16, 25, 8, 1]
- [5, 8, 10, 13, 16, 25, 1]
- [1, 5, 8, 10, 13, 16, 25]

5.

53	13	4	45	67	1	10	5
13, 53		4, 45		1, 67		5, 10	
4, 13, 45, 53				1, 5, 10, 67			
1, 4, 5, 10, 13, 45, 53, 67							