# Algorithms, Probability and Computing: Special Assignment 1

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Prof. Steger, Prof. Holenstein, Prof. Welzl

Kevin Klein

### Exercise 1

Avoiding points.

## Exercise 2

#### Exercise 3

(a)  $a_0 = 7, a_1 = 1 + 2 * a_0 = 15$  $\forall n \ge 2$ :

$$a_n = 1 + 2 * \sum_{i=1}^{n} a_{i-1} \tag{1}$$

$$a_{n-1} = 1 + 2 * \sum_{i=1}^{n} a_{i-1}$$
(2)

$$\Rightarrow a_n - a_{n-1} = 1 + 2 * \sum_{i=1}^n a_{i-1} - (1 + 2 * \sum_{i=1}^n a_{i-1}) = 2 * a_{n-1}$$
(3)

$$\Rightarrow a_n = 3 * a_{n-1} \tag{4}$$

$$=\dots$$
 (5)

$$=3^{n-1}*a_1\tag{6}$$

$$=3^{n-1}*15$$
 (7)

(8)

(b)  $L_n:=length\ of\ a\ central\ path\ in\ a\ BST\ of\ size\ n\ nodes$   $l_n:=\mathbb{E}(L_n)$   $l_0=0, l_1=1, l_2=1/2\cdot 2+1/2\cdot 1=3/2$   $\forall n\geq 3:$ 

$$l_n = \sum_{i=1}^{n} (\mathbb{E}[L_n|root = i] \cdot \Pr[root = i])$$
(9)

$$= \sum_{i=1}^{n} (\mathbb{E}[L_n|root = i] \cdot 1/n) \tag{10}$$

$$=1/n \cdot \sum_{i=1}^{n} (\mathbb{E}[L_n|root=i]) \tag{11}$$

$$= 1/n \cdot \sum_{i=1}^{n} (\mathbb{E}[L_{i-1}] + 1) \tag{12}$$

$$=1+1/n\cdot\sum_{i=1}^{n}(l_{i-1})$$
(13)

(14)

Furthermore,  $l_{n-1} = 1/n \cdot \sum_{i=1}^{n-1} (l_{i-1})$ 

$$n \cdot l_n - (n-1) \cdot l_{n-1} = n + \sum_{i=1}^n (l_{i-1}) - ((n-1) + \sum_{i=1}^{n-1} (l_{i-1})$$
(15)

$$= 1 + l_{n-1} \tag{16}$$

$$\Rightarrow n \cdot l_n = 1 + n \cdot l_{n-1} \tag{17}$$

$$\Rightarrow l_n = 1/n + l_{n-1} \tag{18}$$

$$= 1/n + 1/(n-1) + \dots + l_3 + l_2 \tag{19}$$

$$= 1/n + 1/(n-1) + \dots + l_3 + 3/2 \tag{20}$$

$$= 1/n + 1/(n-1) + \dots + l_3 + 1/2 + 1 \tag{21}$$

$$=H_n\tag{22}$$

(23)

We see that this also holds for n=2 and n=1 and therefore:  $\forall n \in \mathbb{N} - \{0\}$ .

(c) 
$$P_0^{(1)} = P_1^{(1)} = P_2^{(1)} = 0, P_3^{(1)} = 1/3, P_4^{(1)} = 1/4 \cdot 1/3 = 1/12$$
  
 $\forall n > 5$ .

$$P_n^{(1)} = \sum_{i=1}^n ((P_n^{(1)}|root = i) \cdot \Pr[root = i])$$
(24)

$$= \sum_{i=1}^{n} ((P_n^{(1)}|root = i) \cdot 1/n$$
(25)

$$= 1/n \cdot \sum_{i=1}^{n} ((P_n^{(1)}|root = i)$$
 (26)

$$=1/n \cdot \sum_{i=1}^{n} P_{i-1}^{(1)} \tag{27}$$

(28)

Furthermore,  $P_{n-1}^{(1)} = 1/(n-1) \cdot \sum_{i=1}^{n-1} ((P_{i-1}^{(1)}|root=i))$ 

$$\Rightarrow n \cdot P_n^{(1)} - (n-1) \cdot P_{n-1}^{(1)} = \sum_{i=1}^n P_{i-1}^{(1)} - (\sum_{i=1}^{n-1} P_{i-1}^{(1)})$$
 (29)

$$=P_{n-1}^{(1)} \tag{30}$$

$$\Rightarrow n \cdot P_n^{(1)} = n \cdot P_{n-1}^{(1)} \tag{31}$$

$$\Rightarrow P_n^{(1)} = P_{n-1}^{(1)} \tag{32}$$

$$= P_{n-2}^{(1)} = \dots = P_4^{(1)} \tag{33}$$

$$=1/12\tag{34}$$

We see that this also holds for n=4 and therefore:  $\forall n \in \mathbb{N}, n \geq 4$ .

(d)

## Exercise 4