我的算法模板

动态规划部分

LCS 最长公共子序列

这里犯懒了,直接贴的是两个字符串的公共子序列的长度,其实道理是一样的。注意, 这里两个序列的存储都是从0开始的,而dp数组的含义是前i个和前j个,也就是说,这里 dp数组比序列往后错一个。

```
#include <iostream>
#include <cstring>
#include <algorithm>
#include <string>
using namespace std;
#define Max 505
unsigned long dp[Max][Max];
unsigned long max_sub_len(string a, string b){
    memset(dp,0, sizeof(dp));
    unsigned long len=a.size();
    for(int i=1;i<=len;i++)</pre>
        for(int j=1;j<=len;j++){</pre>
            if(a[i-1]==b[j-1])
                 dp[i][j]=dp[i-1][j-1]+1;
            else
                 dp[i][j]=max(dp[i-1][j],dp[i][j-1]);
    return dp[len][len];
}
```

LCS的进阶版本: Edit-Distance

```
int dp[1005][1005];
class Solution {
  public:
    int minDistance(string word1, string word2) {
        int len1=(int)word1.length();
        int len2=(int)word2.length();
        word1 = "0" + word1;
        word2 = "0" + word2;
        //memset(dp,0,sizeof(dp));
        dp[0][0]=0;
        for(int i=1;i<=len1;i++)</pre>
```

```
dp[i][0]=i;
                                                               for(int j=1;j<=len2;j++)</pre>
                                                                                              dp[0][j]=j;
                                                               for(int i=1;i<=len1;i++){</pre>
                                                                                              for(int j=1;j<=len2;j++){</pre>
                                                                                                                              if(word1[i]==word2[j])
                                                                                                                                                            dp[i][j] = dp[i-1][j-1];
                                                                                                                              else{
                                                                                                                                                           dp[i][j] = dp[i-1][j] < dp[i][j-1]?dp[i-1][j]+1:dp[i][j-1]
1]+1;
                                                                                                                                                            dp[i][j] = dp[i][j] < dp[i-1][j-1] + 1?dp[i][j] : dp[i-1][j-1] = dp[i][j] = dp[i-1][j-1] = dp[i][j] = dp[i][
1]+1;
                                                                                                                            }
                                                                                              }
                                                               return dp[len1][len2];
                               }
};
```

LIS 最长上升子序列

```
const int MAXX=100000+5;
const int INF=INT MAX;
int a[MAXX], dp[MAXX]; // a数组为数据, dp[i]表示长度为i+1的LIS结尾元素的最小值
int main()
{
   int n;
   while(cin>>n)
       for(int i=0; i<n; i++)</pre>
           cin>>a[i];
           dp[i]=INF; // 初始化为无限大
       int pos=0; // 记录dp当前最后一位的下标
       dp[0]=a[0]; // dp[0]值显然为a[0]
       for(int i=1; i<n; i++)</pre>
       {
           if(a[i]>dp[pos]) // 若a[i]大于dp数组最大值,则直接添加
               dp[++pos] = a[i];
                  // 否则找到dp中第一个大于等于a[i]的位置,用a[i]替换之。
              dp[lower_bound(dp,dp+pos+1,a[i])-dp]=a[i]; // 二分查找
       cout<<pos+1<<endl;
   return 0;
}
```

LIS讲阶:间隔为k的最长上升子序列

```
const int Max = 6000005;
int dp[Max]; //dp[i]存放的是长度为i的间隔为k的递增序列最后一个元素可能的最小值
vector<int> nums;
int location[Max]; //location[i]存放的是如果把当前的元素插入到最终的递增序列中应该插入
的位置
int LIS_k(int k, vector<int>& nums){
   int n = (int)nums.size();
   for(int i = 1; i \le n; i++) dp[i] = INT32_MAX;
   dp[0] = 0;
    int ans = 0;
    for(int i = 1; i <= n; i++){
       location[i] = int(lower_bound(dp+1,dp+ans+1,nums[i-1]) - dp) ;
       ans = max(location[i], ans);
       int j = i-k;
       if(j>0 && dp[location[j]] > nums[j-1]) dp[location[j]] = nums[j-1];
    }
   return ans;
}
```

n个字符中插入m个乘号,求最大值

```
#include <iostream>
#include <cstring>
#include <algorithm>
#include <string>
#include <cstdio>
#include <climits>
#include <cstdlib>
using namespace std;
#define Max 20
unsigned long long nums[Max][Max];
unsigned long long dp[Max][Max];
//主要使用的字符串,其实输入的是tmp,father是经过简单处理之后的
string father;
unsigned long long my_atoull(int left,int right){
    unsigned long long ans=0;
    for(int i=left;i<=right;i++){</pre>
        ans *= 10;
        ans+=(father[i]-'0');
    }
    return ans;
}
int main(){
```

```
freopen("in.txt", "r", stdin);
    int m,n;
    while(cin>>m){
        string tmp;
        cin>>tmp;
        father = "0"+tmp;
        n=(int)tmp.length();
        //nums[i][j]表示从第i个字符到第j个字符组成的整数的大小
        for(int i=1;i<=n;i++)</pre>
             for(int j=i;j<=n;j++)</pre>
                 nums[i][j]=my_atoull(i,j);
        //初始化dp数组
        for(int j=1; j<=n; j++){</pre>
             dp[0][j]=nums[1][j];
        for(int i=1;i<=m;i++){</pre>
             for(int j=1;j<=i;j++)</pre>
                 dp[i][j]=0;
        }
        //打表开始
        for(int i=1;i<=m;i++){</pre>
             for(int j=1;j<=n;j++){</pre>
                 unsigned long long now=0;
                 unsigned long long now_tmp;
                 for(int k=i; k<=j-1; k++){
                      now_tmp = dp[i-1][k]*nums[k+1][j];
                      if(now_tmp>now)
                          now=now_tmp;
                 }
                 dp[i][j]=now;
             }
        }
        cout<<dp[m][n]<<endl;</pre>
    }
    return 0;
}
```

背包(全)

```
#include <iostream>
#include <cstring>
#include <algorithm>
using namespace std;
const int Max_volume=100005;
```

```
int dp[Max_volume];
int V; //背包的最大容量
inline void ZeroOnePack(int value, int volume){
    for(int v=V;v>=volume;v--)
        if(v>=volume)
            dp[v]=dp[v]>dp[v-volume]+value?dp[v]:dp[v-volume]+value;
}
inline void CompletePack(int value, int volume){
    for(int v=volume; v<=V; v++)</pre>
            dp[v]=dp[v]>dp[v-volume]+value?dp[v]:dp[v-volume]+value;
}
inline void MultiplePack(int value,int volume, int number){
    if(volume*number>=V){
        CompletePack(value, volume);
        return;
    }
    int k=1;
    while(k<number){</pre>
        ZeroOnePack(value*k, volume*k);
        number-=k;
        k<<=1;
    if(number)
        ZeroOnePack(value*number, volume*number);
}
```

背包问题中,如果要求全部装满,那么在初始化的时候,除了dp[0] = 0之外,其余均初始化为 $-\infty$,当然,如果题目要求的是价值最小的方案,应该初始化为 $+\infty$ 。

股票系列

限制整个过程最多交易k次

```
int dp[2][1000005];
int maxProfit(int k, vector<int> &prices) {
    int len = (int) prices.size();
    if (len <= 0)
        return 0;
    k = min(k, len);
    if (k>len/2){
        int ans = 0;
        for (int i=1; i<len; ++i){
            ans += max(prices[i] - prices[i-1],0);
        }
        return ans;
    }
}</pre>
```

```
memset(dp, 0, sizeof(dp));
int cur = 0;
for (int ii = 1; ii <= k; ii++) {
    cur = cur ^ 1;
    int min_prices = prices[0];
    for (int i = 1; i < len; i++) {
        min_prices = min(min_prices , prices[i] - dp[cur][i-1]);
        dp[cur^1][i] = max(dp[cur^1][i - 1] , prices[i] - min_prices);
    }
}
return dp[cur^1][len-1];
}</pre>
```

限制一次交易后应该至少休息一次,解决方案是那个有趣的状态转移图

```
int buy[1000005];
int rest[1000005];
int sell[1000005];
int maxProfit(vector<int>& prices) {
    int len = (int) prices.size();
    if(len<2)
        return 0;
    rest[0]=0;
    buy[0]=0-prices[0];
    sell[0]=INT_MIN;
    for(int i=1;i<len;i++){</pre>
        rest[i] = max(rest[i-1], sell[i-1]);
        buy[i] = max(rest[i-1]-prices[i], buy[i-1]);
        sell[i] = buy[i-1]+prices[i];
    return max(rest[len-1], sell[len-1]);
}
```

概率DP

伯努利分布

就是抛硬币一类的,每次实验成功的概率都是二分之一,求做k次实验,成功i次的概率,也就是求 $C_k^i(\frac{1}{2})^i(\frac{1}{2})^{(k-i)}$,在以下代码中对应的就是Bernoulli_number[k][i]的值。

```
vector<double> Bernoulli_number[Max];

void compute_Bernoulli(int k){
   if(!Bernoulli_number[k].empty())
        return ;
   double res = 1.0;
   for(int i=1;i<=k;i++)
        res*=2;
   res = 1.0/res;
   Bernoulli_number[k].push_back(1.0*res);</pre>
```

```
int middle = k/2:
    double up = k, down = 1 , gap = DBL_MAX/60;
    for(int i=1;i<=middle;i++){</pre>
        Bernoulli_number[k].push_back(res*(up/down));
        if(up>gap){
            up/=down;
            up*=(k-i);
            down = i+1;
        } else {
            up*=(k-i);
            down *= (i+1);
        }
    }
    for(int i = middle +1; i \le k; i++)
        Bernoulli_number[k].push_back(Bernoulli_number[k][k-i]);
}
```

几何分布

如果做成一件事的概率是p,那么做成n次这件事的总次数的期望是 $\frac{1}{n}$ 次。

1021实验

贪心, 简单题。

题目已经说明有且只有一种方法表示所求数,简单列举几项可以发现只由前i个砝码会可以表示 $[1,\Sigma Wi]$ 的所有数的。先找到最大需要的砝码Wi,问题变成了表示(n-Wi),可递归,可循环。见看考代码一。

简单讲讲为什么, 贪心在哪里。令W[6]={1,3,9,27,81,243}表示砝码重量, Sum[6]= {1,4,13,40,121,364}表示前i个砝码总和。

可以发现W[i+1]=Sum[i]*2+1(i \in [0,4]),这表明若Sum[i-1]<n \le Sum[i],能用的最大砝码只能是W[i]。如果最大使用W[i+1],就算减去所有更小的(即-Sum[i])结果依然大于W[i];如果最大使用W[i-1],就算加上所有更小的(即+Sum[i-2])变成Sum[i-1],亦不能表示n。所以判断最大使用的一定是W[i],而且是输出正的W[i],废话!

接下来用n-W[i]后,问题就变成了表示(n-Wi),负数并不是问题,变成正数,输出取负可解决。

```
#include <cstdio>
const int Weight[6] = {1,3,9,27,81,243};
const int Sum[6] = {1,4,13,40,121,364};

int main()
{
   int n;
   while(~scanf("%d", &n))
   {
     int flag = 0;
     int i = 0;
}
```

```
while(n > 0)
        {
            for (i = 0; i < 6; ++i) {
                if(n <= Sum[i])</pre>
                    break;
            }
            if(flag == 0)//第一个数
                n -= Weight[i];
                printf("%d", Weight[i]);
            }
            if(flag == 1) {
                n -= Weight[i];
                printf("+%d", Weight[i]);
            if(flag == -1) {
                 n = Weight[i] - n;
                printf("-%d", Weight[i]);
            }
            if(n > 0) flag = 1;
            else if(n < 0)
                flag = -1;
                n = -n;
            }
        printf("\n");
    }
}
```

图论

DFS-邻接表实现

```
#include <cstdio>
#include <iostream>
#include <queue>
#include <functional>
#include <cstring>
#include <string>
#include <queue>
#include <queue>
#include <algorithm>
using namespace std;
const int Max = 1005;
vector<int> g[Max];
```

```
bool used[Max];
int dist[Max];
//图的编号依旧是从1到n, 很正常
void DFS(vector<int> g[], int start, int dist[], int n, int end){
    for(int i = 1; i \le n; i++) dist[i] = -1;
    for(int i = 1; i \le n; i++) used[i] = false;
    queue<int> que;
    while(!que.empty()) que.pop();
    que.push(start);
    used[start] = true;
    dist[start] = 0;
    while(!que.empty()){
        int u = que.front();
        que.pop();
        for(auto item : g[u]){
            if(!used[item]){
                used[item]=true;
                que.push(item);
                dist[item] = dist[u] + 1;
            if(item == end) return;
        }
    }
}
```

二分图-匈牙利算法-邻接表实现

需要注意的是,这里的uN表示二分图左边的顶点个数,跟右边图定点个数以及编号方式 随意。

```
#include <cstdio>
#include <iostream>
#include <cstring>
#include <algorithm>
#include <queue>
#include <vector>
using namespace std;
const int MaxN = 1005;
vector<int> g[MaxN];
int linker[MaxN];
bool used[MaxN];
int uN;
bool dfs(int u){
    for(auto v:g[u]){
        if(!used[v]){
            used[v] = true;
            if(linker[v] == -1 || dfs(linker[v])){
                linker[v] = u;
                return true;
```

```
}
}
return false;

int hungry(){
   int res = 0;
   memset(linker, -1, sizeof(linker));
   for(int u = 0;u<uN;u++){
       memset(used, false, sizeof(used));
       if(dfs(u)) res++;
   }
   return res;
}</pre>
```

按字典序输出二分图左边已经配对的顶点,相关信息存储在数组ress中,按照以下代码输出的将是最小字典序,如若输出最大字典序的话,只需要将hungry中的第一个循环反向即可。

```
#define _CRT_SECURE_NO_WARNINGS
#include <cstdio>
#include <iostream>
#include <cstring>
#include <cmath>
#include <algorithm>
#include <vector>
#include <string>
#include <stack>
#include <cstdlib>
using namespace std;
const int mod = 100007;
const int MAXN = 1005;
typedef long long 11;
const int inf = 0x3f3f3f3f;
int uN, vN;
int Graph[MAXN][MAXN];
int linker[MAXN];
bool used[MAXN];
int ress[MAXN];
bool dfs(int u) {
    for (int v = 1; v \le 1000; v++) {
        if (Graph[u][v] && !used[v]) {
            used[v] = true;
            if (linker[v] == -1 || dfs(linker[v])) {
                linker[v] = u;
                ress[u] = 1;
                return true;
            }
        }
```

```
return false;
}
int hungary() {
   int res = 0;
    memset(linker, -1, sizeof(linker));
    for (int u = 1; u \le 1000; u++) {
        memset(used, false, sizeof(used));
        if (dfs(u)) res++;
    for (int i = 1; i <= uN; i++) {
        int cnt = 0;
        if (ress[i] == 1) {
            cnt++;
            if (cnt == res) printf("%d\n", i); //这里就是保证结尾没有多于空格, 其实
没什么实际意义,还费时间
            else printf("%d ", i);
        }
   return res;
}
int main() {
    freopen("Text.txt", "r", stdin);
    while (~scanf("%d%d", &uN, &vN)) {
        memset(Graph, 0, sizeof(Graph));
        for (int i = 1; i <= vN; i++) {
            int num;
            scanf("%d", &num);
            for (int j = 1; j \le num; j++) {
                int tmp;
                scanf("%d", &tmp);
                Graph[i][tmp] = 1;
            }
        printf("%d\n", hungary());
   return 0;
}
```

最大流-EK算法-邻接表实现

```
#include <iostream>
#include <cstdio>
#include <cstring>
#include <vector>
#include <algorithm>
#include <queue>
#include <map>
#include <unordered_map>
using namespace std;
```

```
#define Max 10005
#define INF 0x7ffffff
int flow[Max], father[Max], vertex, E;
vector<pair<int,int>> g[Max];
inline void change_element_add(int a, int b, int c){
    int i=0;
    int len = (int) g[a].size();
    for(;i<len;i++)
        if(g[a][i].first == b)
            break;
    if(i<len)</pre>
        g[a][i].second += c;
    else
        g[a].emplace_back(b,c);
}
inline void change_element_sub(int a,int b,int c){
    int i=0;
    int len = (int) g[a].size();
    for(;i<len;i++)
        if(g[a][i].first == b){
            g[a][i].second -= c;
            break;
        }
}
inline int BFS(int s,int t){
    queue<int> q;
    while(!q.empty())
        q.pop();
    memset(father, -1, sizeof(int)*(vertex+5));
    flow[s] = INF;
    q.push(s);
    while(!q.empty()){
        int v = q.front();
        q.pop();
        for(auto item:g[v]){
            int i=item.first;
            if(father[i]==-1 && item.second>0){
                 flow[i] = min(flow[v], item.second);
                 father[i] = v;
                 if(i==t)
                     return flow[t];
                q.push(i);
            }
        }
    if(father[t] == -1)
        return -1;
    else
```

```
return flow[t];
}
int EK(int s,int t){
    int ans = 0;
    int increase= BFS(s,t), k=t, last;
    while(increase!=-1){
        while(k!=s){
            last = father[k];
            change_element_sub(last,k,increase);
            change_element_add(k,last,increase);
            k=last;
        }
        ans += increase;
        k=t;
        increase = BFS(s,t);
    return ans;
}
```

最小生成树-Kruskal算法-邻接表实现

```
#include <iostream>
#include <cstring>
#include <vector>
#include <cstdio>
#include <climits>
#include <cfloat>
#include <queue>
#include <functional>
#include <algorithm>
using namespace std;
struct edge{
public:
    int u;
    int v;
    int cost;
    edge(int u,int v,int cost):u(u),v(v),cost(cost){}
    bool operator < (const edge & b) const {</pre>
        return cost<b.cost;</pre>
};
vector<edge> g;
const int Max = 10005;
int Find[Max];
int find(int x){
    if(Find[x] == -1)
```

```
return x;
    else
        return Find[x] = find(Find[x]);
}
//返回最小生成树的最小费用
int Kruskal(int n){
    memset(Find, -1, sizeof(Find));
    sort(g.begin(),g.end());
    int cnt = 0;
    int ans = 0;
    for(auto item : g){
        int u = item.u;
        int v = item.v;
        int cost = item.cost;
        int t1 = find(u);
        int t2 = find(v);
        if(t1!=t2){
            ans += cost;
            Find[t1] = t2;
            cnt ++;
        if(cnt == n-1)
            break;
    if(cnt<n-1) return -1;
    else return ans;
}
```

并查集重要函数

```
int Find[Max];

//寻找根节点编号
int find(int x){
   if(Find[x] == -1)
       return x;
   else
      return Find[x] = find(Find[x]);
}
```

单源最短路-Dijkstra算法-邻接表实现

```
#include <cstdio>
#include <iostream>
#include <queue>
#include <functional>
#include <cstring>
```

```
#include <string>
#include <queue>
#include <algorithm>
using namespace std;
const int Max = 10005;
const int INF = 0x3f3f3f3f;
struct edge{
    int to;
    int cost;
    edge(int first,int to):to(first),cost(to){}
    edge(){}
    bool operator < (const edge & b) const {</pre>
        return cost>b.cost;
    }
};
vector<edge> g[Max]; //这个图本身
int dist[Max]; //最终每个点对应的距离
int Path[Max]; //记录路径
int vis[Max]; //记录是不是已经判断过了
//节点编号从1到n
void Dijkstra(int n, int start){
    memset(Path, -1, sizeof(Path));
    memset(vis, false, sizeof(vis));
    for(int i=1;i<=n;i++) dist[i] = INF;</pre>
    priority_queue<edge> que;
    while(!que.empty()) que.pop();
    dist[start] = 0;
    que.push(edge(start,0));
    edge tmp;
    while(!que.empty()){
        tmp = que.top();
        que.pop();
        int u = tmp.to;
        if(vis[u]) continue;
        vis[u] = true;
        for(auto item:g[u]){
            int v = item.to;
            int cost = item.cost;
            if(!vis[v] && dist[v]>dist[u]+cost){
                dist[v] = dist[u] + cost;
                que.push(edge(v,dist[v]));
                Path[v] = u;
            }
        }
    }
}
```

计算几何

凸包-Jarvis步进法

```
#include<cstdio>
#include<vector>
#include<cmath>
#include<algorithm>
using namespace std;
//精度判断
const double eps = 1e-10;
double dcmp(double x) {
    if(fabs(x) < eps) return 0;
    else return x < 0 ? -1 : 1;
}
struct Point {
    double x, y;
    Point(double x=0, double y=0):x(x),y(y) {}
Point operator - (const Point& A, const Point& B) {
    return Point(A.x-B.x, A.y-B.y);
double Cross(const Point& A, const Point& B) {
    return A.x*B.y - A.y*B.x;
double Dot(const Point& A, const Point& B) {
    return A.x*B.x + A.y*B.y;
bool operator < (const Point& p1, const Point& p2) {</pre>
    return p1.x < p2.x \mid | (p1.x == p2.x && p1.y < p2.y);
bool operator == (const Point& p1, const Point& p2) {
    return p1.x == p2.x && p1.y == p2.y;
}
//点集凸包, Jarvis步进法,注意,是逆时针方向,并且首尾不相连
vector<Point> ConvexHull(vector<Point> p) {
    //预处理,删除重复点
    sort(p.begin(), p.end());
    p.erase(unique(p.begin(), p.end()), p.end());
    int n = p.size();
    int m = 0;
    vector<Point> ch(n+1);
    for(int i = 0; i < n; i++) {
        while(m > 1 \&\& Cross(ch[m-1]-ch[m-2], p[i]-ch[m-2]) <= 0) m--;
```

```
ch[m++] = p[i];
}
int k = m;
for(int i = n-2; i >= 0; i--) {
    while(m > k && Cross(ch[m-1]-ch[m-2], p[i]-ch[m-2]) <= 0) m--;
    ch[m++] = p[i];
}
if(n > 1) m--;
ch.resize(m);
return ch;
}
```

凸包-最后一次练习赛最后一题-包含判断两个凸包是否相交

```
#include<cstdio>
#include<vector>
#include<cmath>
#include<algorithm>
using namespace std;
//精度判断
const double eps = 1e-10;
double dcmp(double x) {
    if(fabs(x) < eps) return 0;
    else return x < 0? -1 : 1;
}
struct Point {
    double x, y;
    Point(double x=0, double y=0):x(x),y(y) {}
};
Point operator - (const Point& A, const Point& B) {
    return Point(A.x-B.x, A.y-B.y);
}
double Cross(const Point& A, const Point& B) {
    return A.x*B.y - A.y*B.x;
}
double Dot(const Point& A, const Point& B) {
    return A.x*B.x + A.y*B.y;
bool operator < (const Point& p1, const Point& p2) {</pre>
    return p1.x < p2.x \mid | (p1.x == p2.x \&\& p1.y < p2.y);
bool operator == (const Point& p1, const Point& p2) {
    return p1.x == p2.x \&\& p1.y == p2.y;
//判断两条线段是否相离
bool SegmentProperIntersection(const Point& a1, const Point& a2, const Point&
b1, const Point& b2) {
    double c1 = Cross(a2-a1, b1-a1), c2 = Cross(a2-a1, b2-a1),
```

```
c3 = Cross(b2-b1, a1-b1), c4=Cross(b2-b1, a2-b1);
    return dcmp(c1)*dcmp(c2)<0 \&\& dcmp(c3)*dcmp(c4)<0;
}
bool OnSegment(const Point& p, const Point& a1, const Point& a2) {
    return dcmp(Cross(a1-p, a2-p)) == 0 && dcmp(Dot(a1-p, a2-p)) < 0;
}
//点集凸包, Jarvis步进法
vector<Point> ConvexHull(vector<Point> p) {
    //预处理, 删除重复点
    sort(p.begin(), p.end());
    p.erase(unique(p.begin(), p.end()), p.end());
    int n = p.size();
    int m = 0;
    vector<Point> ch(n+1);
    for(int i = 0; i < n; i++) {
        while(m > 1 && Cross(ch[m-1]-ch[m-2], p[i]-ch[m-2]) <= 0) m--;
        ch[m++] = p[i];
    }
    int k = m;
    for(int i = n-2; i \ge 0; i--) {
        while(m > k \&\& Cross(ch[m-1]-ch[m-2], p[i]-ch[m-2]) <= 0) m--;
        ch[m++] = p[i];
    }
    if(n > 1) m--;
    ch.resize(m);
    return ch;
}
//判断点与凸多边形是否相离
int IsPointInPolygon(const Point& p, const vector<Point>& poly) {
    int wn = 0;
    int n = poly.size();
    for(int i=0; i<n; ++i) {
        const Point& p1 = poly[i];
        const Point& p2 = poly[(i+1)%n];
        if(p1 == p || p2 == p || OnSegment(p, p1, p2)) return -1;//在边界上
        int k = dcmp(Cross(p2-p1, p-p1));
        int d1 = dcmp(p1.y - p.y);
        int d2 = dcmp(p2.y - p.y);
        if(k > 0 \&\& d1 \le 0 \&\& d2 > 0) wn++;
        if(k < 0 \&\& d2 <= 0 \&\& d1 > 0) wn--;
    if(wn != 0) return 1;//内部
    return 0;//外部
}
bool ConvexPolygonDisjoint(const vector<Point> ch1, const vector<Point> ch2) {
    int c1 = ch1.size();
    int c2 = ch2.size();
```

```
for(int i=0; i<c1; ++i)
        if(IsPointInPolygon(ch1[i], ch2) != 0) return false;
    for(int i=0; i<c2; ++i)
        if(IsPointInPolygon(ch2[i], ch1) != 0) return false;
    for(int i=0; i<c1; ++i)</pre>
        for(int j=0; j<c2; ++j)</pre>
            if(SegmentProperIntersection(ch1[i], ch1[(i+1)%c1], ch2[j],
ch2[(j+1)%c2])) return false;
    return true;
}
int main()
    freopen("in.txt", "r", stdin);
    int n, m;
    while(scanf("%d %d", &n, &m) == 2 && n > 0 && m > 0)
        vector<Point> P1, P2;
        double x, y;
        for(int i = 0; i < n; i++) {
            scanf("%lf %lf", &x, &y);
            P1.push_back(Point(x, y));
        for(int i = 0; i < m; i++) {
            scanf("%lf %lf", &x, &y);
            P2.push_back(Point(x, y));
        }
        if(ConvexPolygonDisjoint(ConvexHull(P1), ConvexHull(P2)))
            printf("YES\n");
        else
            printf("NO\n");
    return 0;
}
```

数论

FFT

高精度超长整数乘法

```
#include <cstdio>
#include <vector>
#include <cmath>
#include <algorithm>
#include <iostream>
#include <string>
#include <cstring>
#include <complex>
```

```
using namespace std;
const int Max = 200005;
const double PI = acos(-1.0);
void change(complex<double> y[], int len) {
    int i, j, k;
    for (i = 1, j = len / 2; i < len - 1; i++) {
        if (i < j) swap(y[i], y[j]);
        k = len / 2;
        while (j >= k) {
            j -= k;
            k /= 2;
        }
        if (j < k) j += k;
    }
}
void FFT(complex<double> y[], int len, int on) {
    change(y, len);
    for (int h = 2; h \le len; h \le 1) {
        complex<double> wn(cos(-on * 2 * PI / h), sin(-on * 2 * PI / h));
        for (int j = 0; j < len; <math>j += h) {
            complex<double> w(1, 0);
            for (int k = j; k < j + h / 2; k++) {
                complex<double> u = y[k];
                complex<double> t = w * y[k + h / 2];
                y[k] = u + t;
                y[k + h / 2] = u - t;
                w = w * wn;
            }
        }
    }
    if (on == -1) {
        for (int i = 0; i < len; i++)
            y[i] = complex<double>(y[i].real() / len, y[i].imag());
    }
}
complex<double> x1[Max], x2[Max];
int sum[Max];
int main() {
    freopen("in.txt", "r", stdin);
    string a, b;
    while (cin >> a >> b) {
        int len1 = (int) a.size();
        int len2 = (int) b.size();
        int len = 1;
        while (len < len1 * 2 || len < len2 * 2) len <<= 1;
```

```
for (int i = 0; i < len1; i++) x1[i] = complex < double > (a[len1 - 1 - i])
- 'O', O);
        for (int i = len1; i < len; i++) x1[i] = complex < double > (0, 0);
        for (int i = 0; i < len2; i++) x2[i] = complex < double > (b[len2 - 1 - i])
- 'O', O);
        for (int i = len2; i < len; i++) x2[i] = complex < double > (0, 0);
        FFT(x1, len, 1);
        FFT(x2, len, 1);
        for (int i = 0; i < len; i++) x1[i] = x1[i] * x2[i];
        FFT(x1, len, -1);
        for (int i = 0; i < len; i++) sum[i] = (int) lround(x1[i].real());
        for (int i = 0; i < len; i++) {
            sum[i + 1] += sum[i] / 10;
            sum[i] %= 10;
        len = len1 + len2 - 1;
        while (sum[len] \le 0 \&\& len > 0) len--;
        for (int i = len; i \ge 0; i--) printf("%c", sum[i] + '0');
        cout << endl;
    }
    return 0;
}
```

其他

OJ必备

```
#ifndef ONLINE_JUDGE
    freopen("in", "r", stdin);
    ios_base::sync_with_stdio(false);
    cin.tie(nullptr);
    cout.tie(nullptr);
#endif
```

输入输出控制

```
#include <iostream>
#include <iomanip>
using namespace std;
void main()
{

int x=1000;
double y=1.23456789;
cout<<"默认x值:"<<x<<endl;
cout<<"十进制:"<<dec<<x<<endl;
hex 会一直作用。
cout<<"八进制:"<<oct<<x<<endl;
```

```
cout<<"十六进制:"<<hex<<x<<endl;
  cout<<"十六进制(大写字母):"<<hex<<uppercase<<x<<endl<
  cout<<"默认y值(左对齐且有效数字位数为6):"<<y<end1;
  cout<<"宽度为10并右对齐:"<<setw(10)<<right<<y<<endl;
  cout<<"宽度为8:"<<setw(8)<<y<<end1;
  cout<<"宽度为4:"<<setw(4)<<y<<end1;
  cout<<"用*号填充空位(10位宽度):"<<setfill('*')<<setw(10)<<y<<endl; //setw(int
i) 只对紧随的数据显示有影响。控制多个数据要多个setw()method.
  cout << "保留3位小数" << fixed << setprecision(3) << y << endl;
  cout<<"设精度为3输出y(不包括小数点):"<<setprecision(3)<<y<<end1;
 //setprecision(int i) 会一直作用
  cout<<"设精度为8输出y(不包括小数点):"<<setprecision(8)<<y<<end1;
  cout<<"显示正负号:"<<showpos<<y<endl;
  cout<<"用科学计数法表示y:"<<scientific<<y<<endl;
  cout<<"用科学计数法表示y(控制E前数据的小数点后位数):"<<scientific
      <<setprecision(3)<<y<<endl;
}
```

VS2013产品密钥-所有版本

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