# Module 1 .NET runtime

# .NET runtime - CIL, JIT, assembler

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- specified in "ECMA-335. Common Intermediate Language (CIL)"

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- This specification name is **Common Language Infrastructure**...
- specified in "ECMA-335. Common Intermediate Language (CIL)" do not confuse with "ECMA-334. C# Language Specification."
- obviously, above that we have the whole *ecosystem*:
  - o compilers C# (Roslyn), F#, ...
  - libraries Base Class Library, ...
  - frameworks ASP.NET, Entity Framework, ...
  - tools MSBuild, Visual Studio, NuGet, ...

Common Language Infrastructure:

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And... CLI-compatible programs compiles to Common Intermediate Language (CIL) (we will return to that)

# Common Type System (CTS) vs Common Language Specification (CLS)...

- CLS is a subset of CTS!
- our code may by CLS-compliant or not... Eg. public field starting with underscore is not:

```
[assembly: CLSCompliant(true)]

□ namespace ConsoleApp

{

Oreferences
public class SomeType
{

public int Field = 4;
}

② (field) int SomeType._Field

CS3008: Identifier '_Field' is not CLS-compliant
Show potential fixes (Ctrl+.)
```

- usually we don't care 99% is written in C# anyway
- uint is also not-CLS compliant and so on, and so forth

# What is .NET runtime aka Virtual Execution System (VES)?

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- piece of code executing CIL code...
  - knowing about Common Type System and Metadata
  - we may ignore CLS:)
- ... from assembly/module
  - o a physical container for code and metadata

#### ECMA-335, II.6:

- "an **assembly** is a set of one or more files deployed as a unit, (...) contains **manifest** defining:
  - version, name, culture, ...
  - which other files, if any, belong to the assembly

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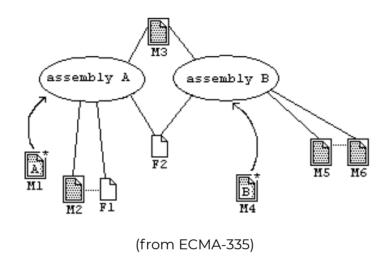
- "an **assembly** is a set of one or more files deployed as a unit, (...) contains **manifest** defining:
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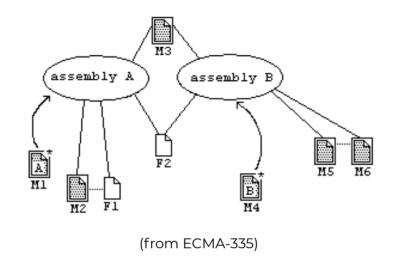
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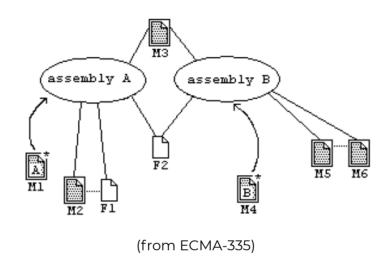
#### **II.6.2**:

• "assembly is specified as a module that contains a manifest in the metadata"

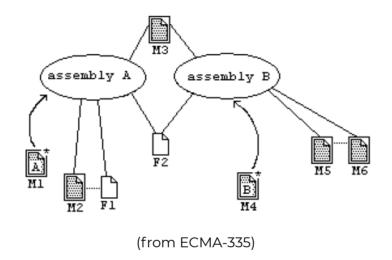




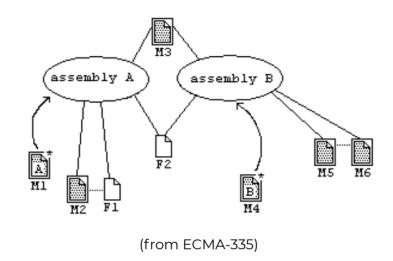
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- Mx module files
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- "The Visual Studio IDE for C# and Visual Basic can only be used to create single-file assemblies."
  - o so, just a regular one-to-one correspondence between assembly and module

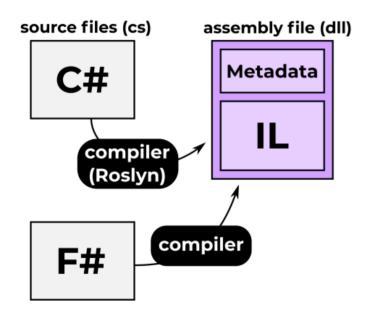
For example, compiling a ".NET Console App", produces single module file/assembly, with the manifest as:

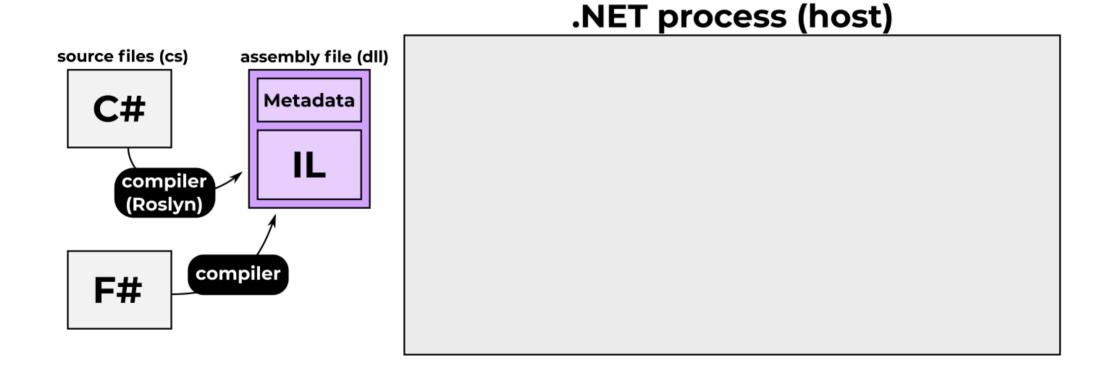
```
// Metadata version: v4.0.30319
.assembly extern System.Runtime
{
...
}
.assembly extern System.Console
{
...
}
.assembly ConsoleApp
{
...
}
.module ConsoleApp.dll
... types/code goes here ...
```

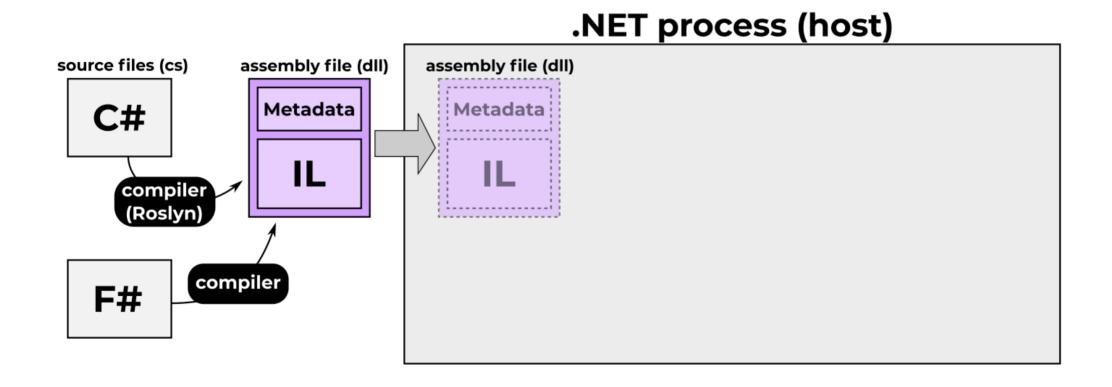
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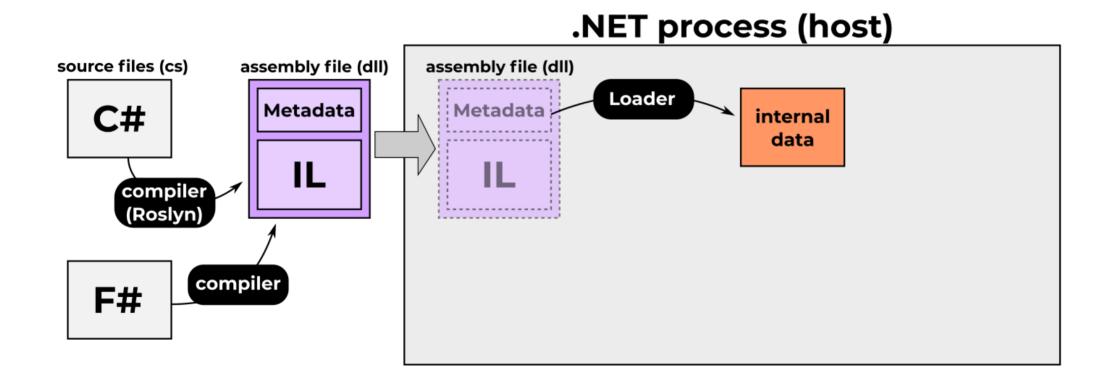
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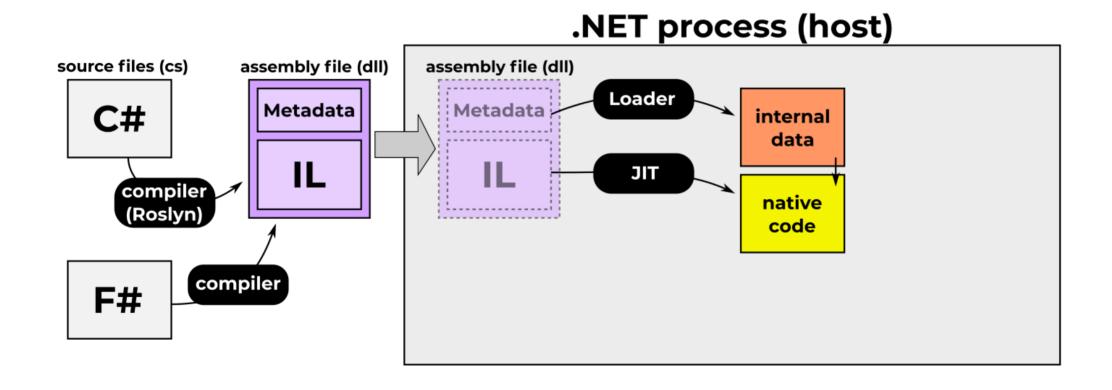
Sidenote: How to: Build a multifile assembly - you can build "multifile assembly" with the command line compiler (csc.exe) and Assembly Linker (al.exe), but for .NET Framework only.

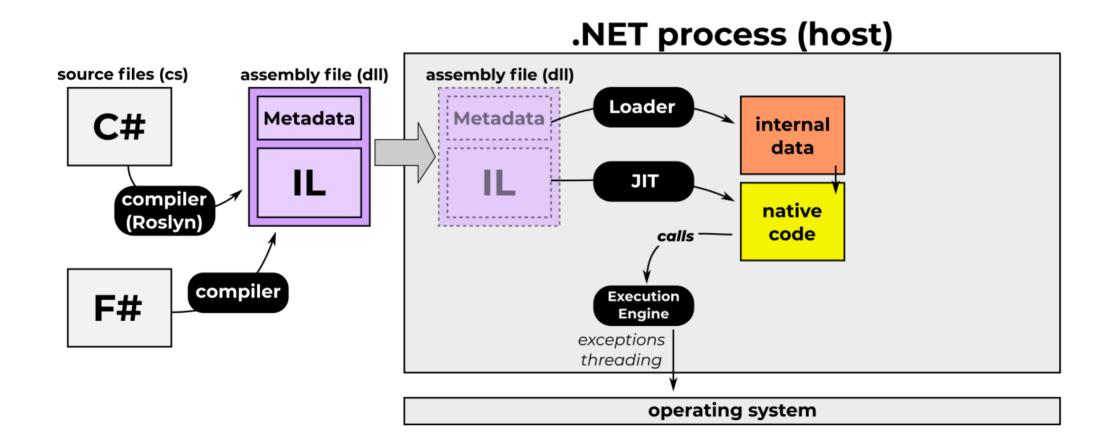


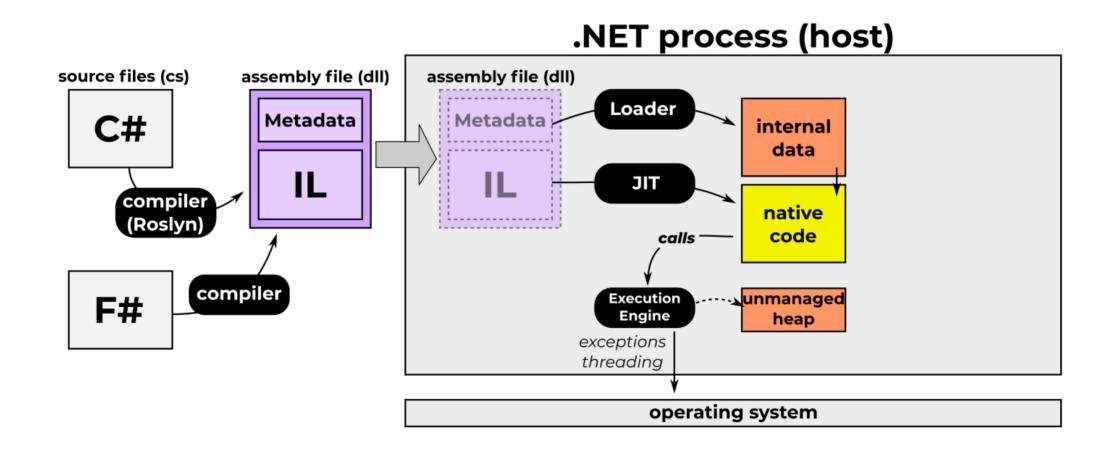


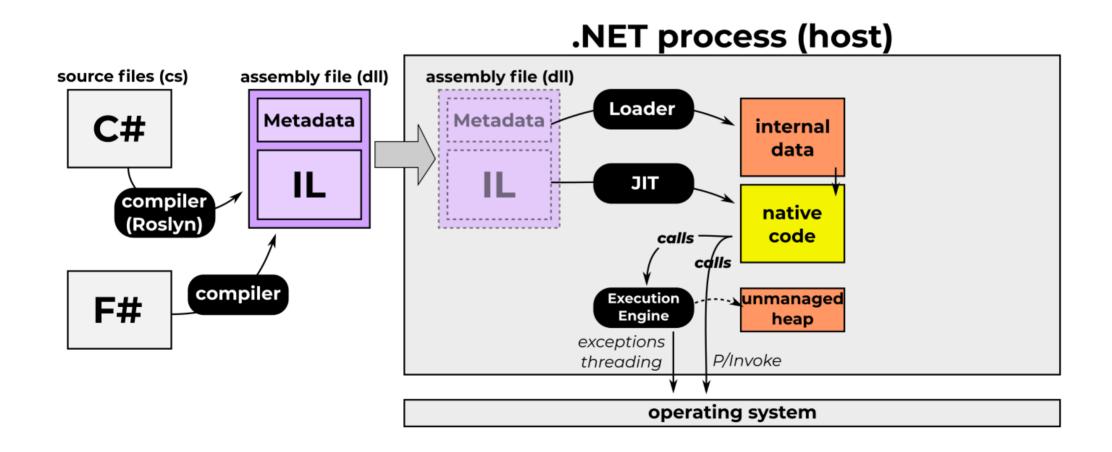


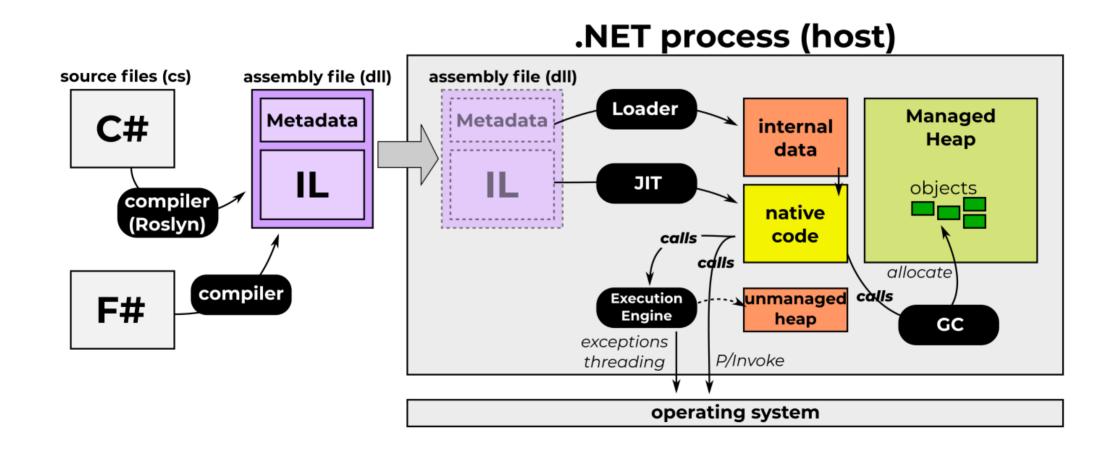


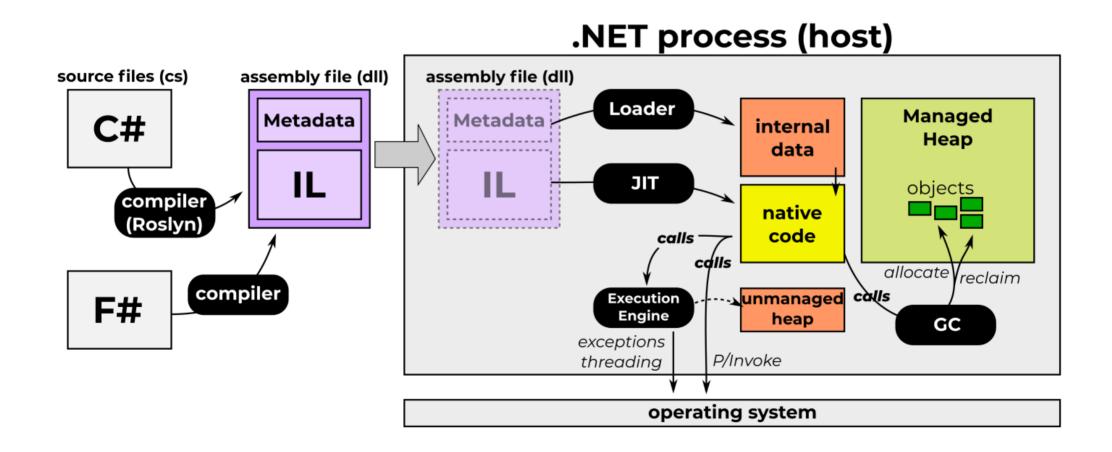


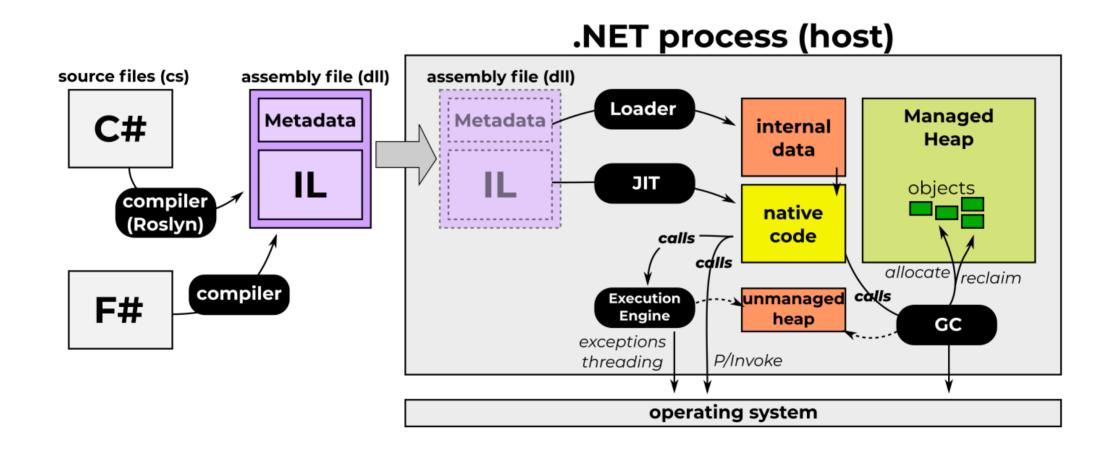


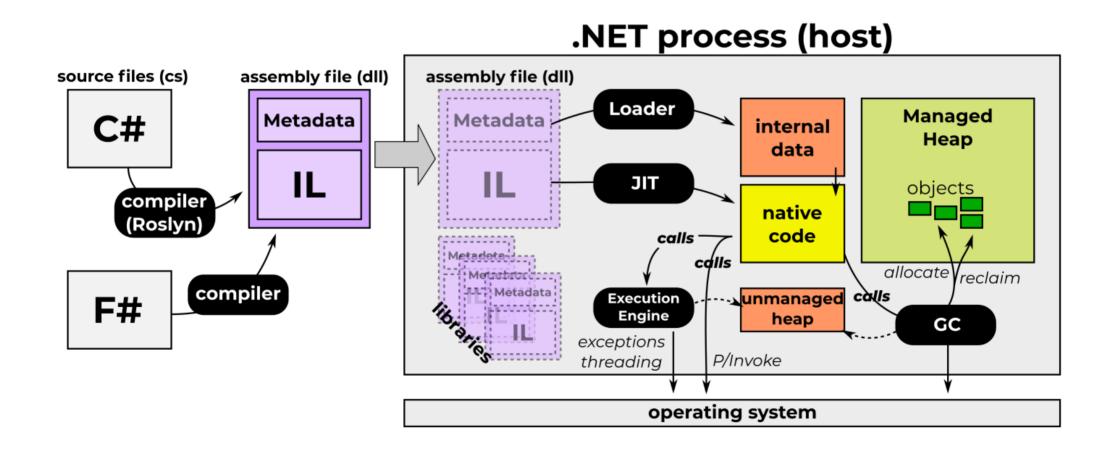


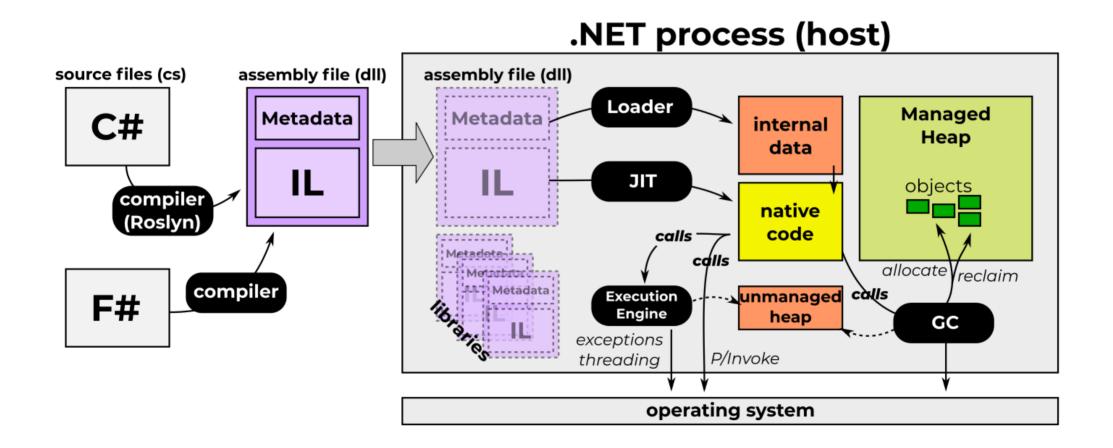






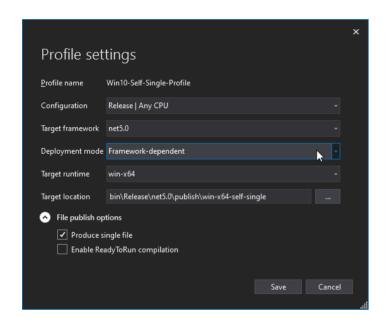




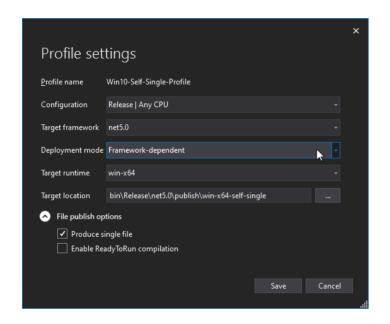


**Notice**: every running .NET app loads **its own, separate .NET runtime** (and... every time it runs).

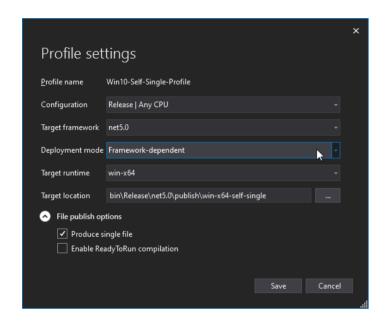
Sidenote: We will return to the topic of unloadable assemblies and AssemblyLoadContext	



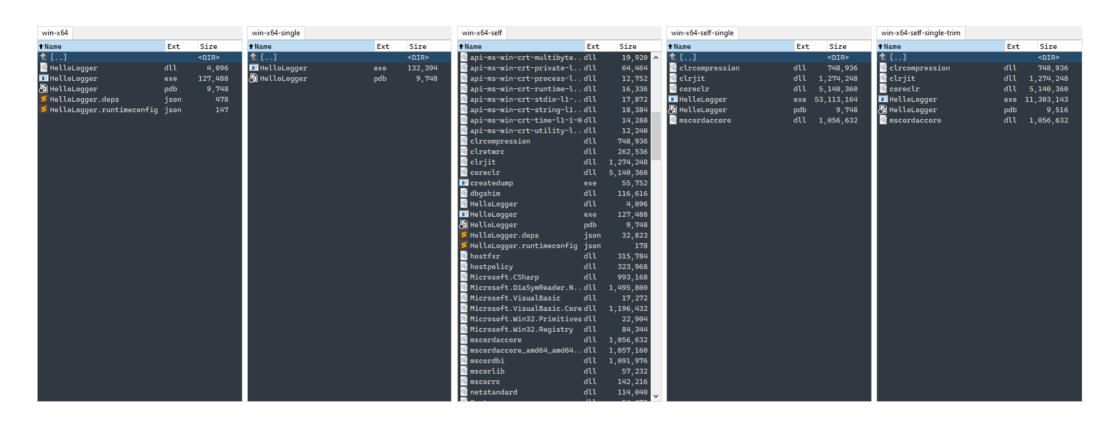
- **Target runtime** Windows 32/64bit, Linux 64bit/ARM, MacOS 64bit, ...
- (optional) **Single file** mode all/most dependencies (libraries, runtime) are packed into single executable file



- Deployment mode:
  - Framework-dependent it will require .NET framework/runtime installed on the target machine. Contains:
    - in normal mode:
      - assembly in PE/COFF format the same dll for all OS
      - host file exe, ELF, Mach-O...
    - in Single file mode:
      - host file with all assemblies inside exe, ELF, Mach-O...



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    - in Single file mode:
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  - Self-contained shipping all runtime and dependencies.
     Contains:
    - in normal mode:
      - a LOT of files all framework libraries, runtime files, assemblies
    - in single file mode:
      - host file with all assemblies/runtime inside exe, ELF, Mach-O...
      - (optionally) some dependencies
- Additionally, in Single file mode for self-contained mode we can enable (experimental) **Trimming**



## Framework-depentent, no single file mode



- HelloLogger.dll assembly containing our app (IL, metadata)
- HelloLogger.exe win-x64 host
- HelloLogger.pdb symbol files
- **HelloLogger.runtimeconfig.json** as it is framework-dependent, it defines the required runtime (and its options):

```
{
"runtimeOptions": {
    "framework": {
        "name": "Microsoft.NETCore.App",
        "version": "5.0.0"
    }
}
```

It triggers looking for shared runtime version at:

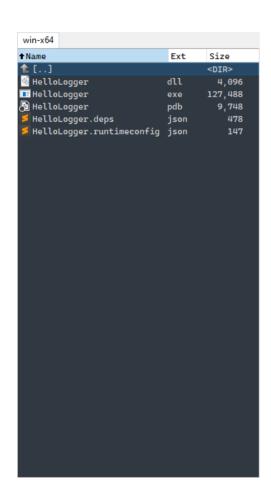
- c:\Program Files\dotnet\shared\Microsoft.NETCore.App\(Windows)
- o /home/{user}/share/dotnet/shared/ (Linux)
- /usr/local/share/dotnet/shared/ (macOS)
- HelloLogger.deps dependency manifest

#### **Assembly anatomy**

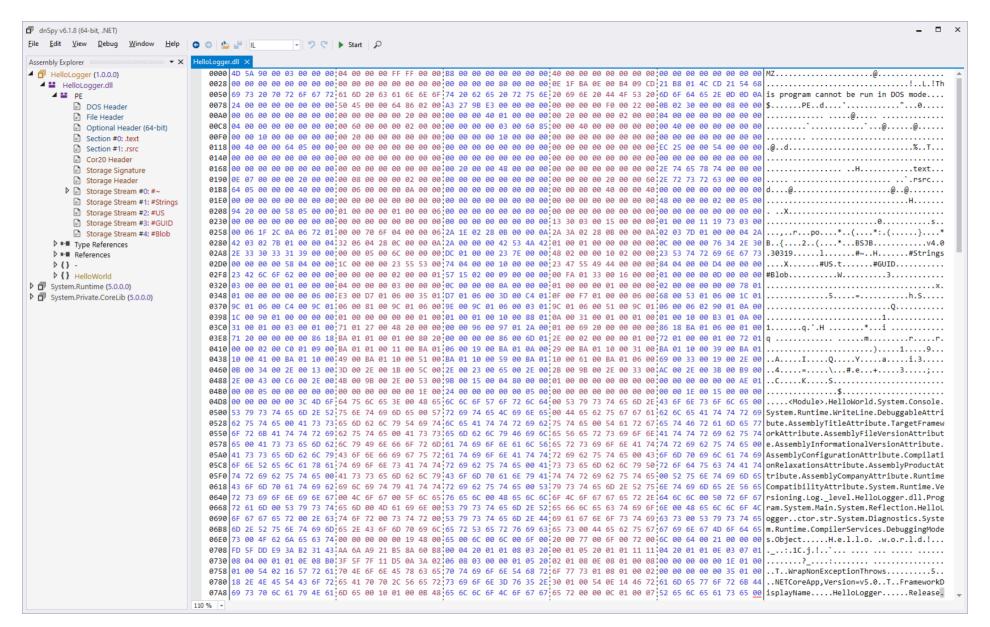
Imagine we compile such trivial C# program into HelloLogger.dll:

```
using System;
public class Program {
    public static void Main() {
        var logger = new Logger(3);
        var level = 44;
        logger.Log(level, "Hello world!");
    }
}
public class Logger
{
    private int _level;
    public Logger(int level) => _level = level;
    public void Log(int level, string str)
    {
        if (level >= _level) Console.WriteLine(str);
    }
}
```

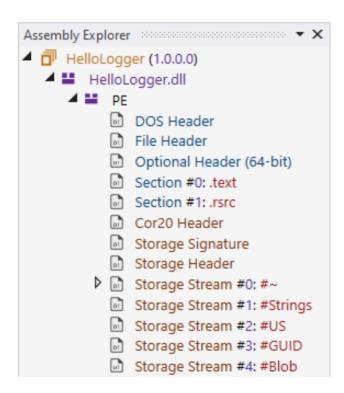
#### **Assembly anatomy**



So, we are interested in analyzing **HelloLogger.dll** file - assembly containing our application (types metadata, all members/methods code in CIL)



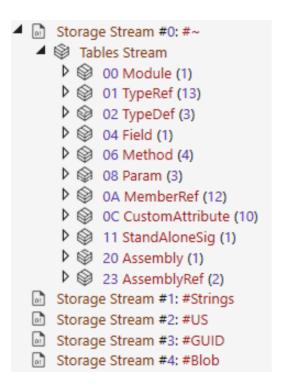
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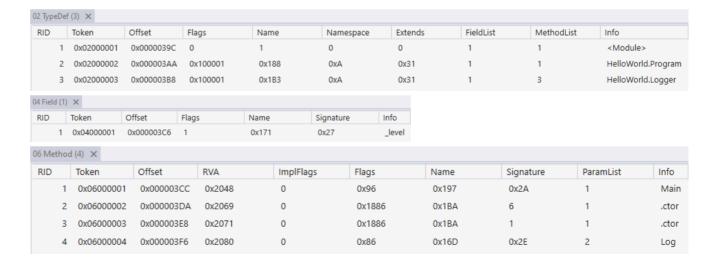
#### Inside HelloLogger.dll:

- standard Portable Executable (PE) & Common Object File Format (COFF) headers:
  - DOS Header
  - File Header
  - Optional Header
- text section "Executable code (free format)" as PE documentation says - reused to contain all the CIL code and metadata!
  - Cor20 header basic information about .NET assembly (expected runtime version, entry point)
  - storage streams
- .rsrc section unmanaged resources

#### **Assembly anatomy - storage streams**



- #~ all metadata about types:
  - Module module defintion
  - TypeRef referenced types (used by dynamic resolve)
  - TypeDef defined types
  - Field fields from all types
  - Method methods froms all types
  - o ...



#### **Assembly anatomy - storage streams**

```
▲ Storage Stream #0: #~

▲ S Tables Stream

     01 TypeRef (13)
    D 🔘 04 Field (1)
    D  OC CustomAttribute (10)
   D 3 AssemblyRef (2)
   Storage Stream #1: #Strings
   Storage Stream #2: #US
   Storage Stream #3: #GUID
 Storage Stream #4: #Blob
```

- #Strings required built-in strings (like type/method names)
- **#US** user-defined strings (yes, **"Hello world!"**) (...we will return to that next weeks)
- #GUID defined Guid values
- #Blob all other inlined data (arrays etc.)

(...we will return to that next weeks)

#### **Assembly anatomy - storage streams**

So, when runtime wants to execute Main method - it finds it, typically by Token, and looks for its RVA (relative virtual address - address relative to the base address of image loaded into memory):



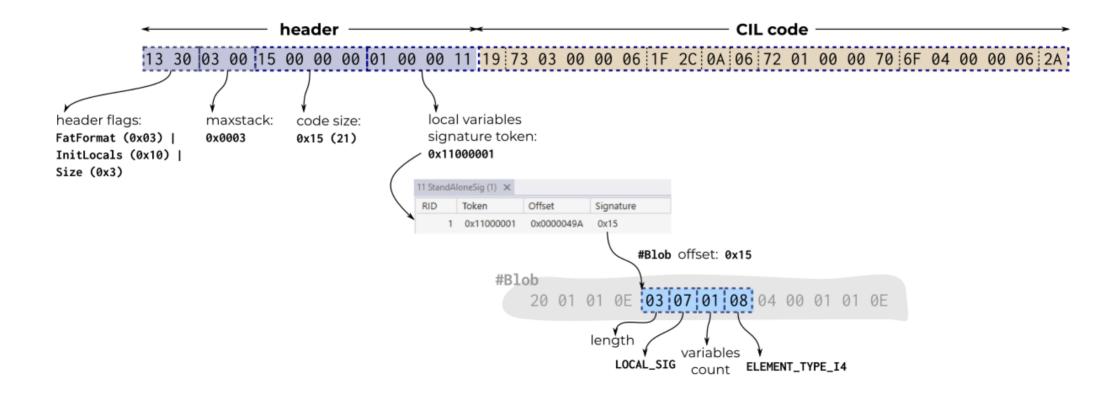
And then it know where the binary CIL code is (33 bytes!):



And as we see, its Signature & ParamList (arguments and return type) may be decoded.

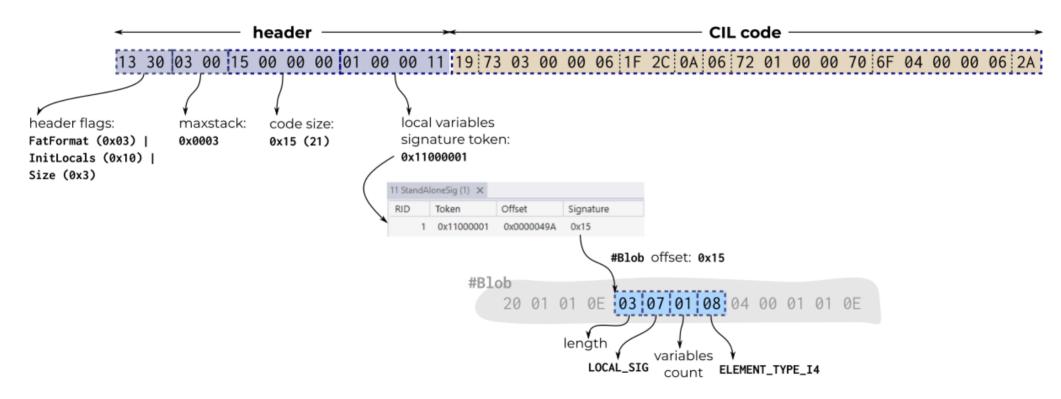
#### **Assembly anatomy - method storage**

So, in the end Main method is represented by the following structure:

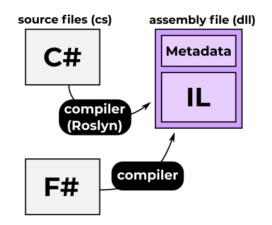


#### **Assembly anatomy - method storage**

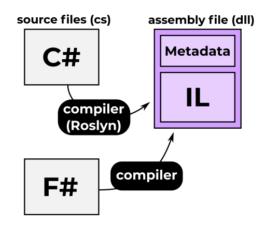
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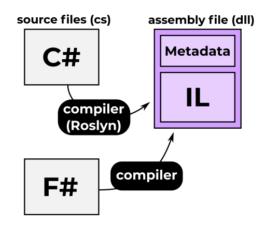
Refer to ECMA-335 22. section for more details.



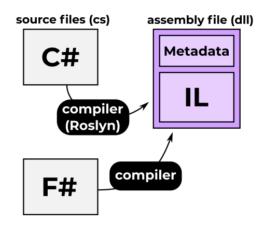
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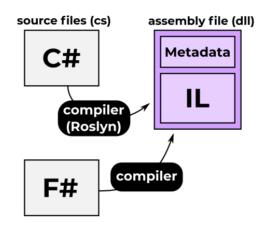
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  - in assembly it is stored in binary format (as we saw)



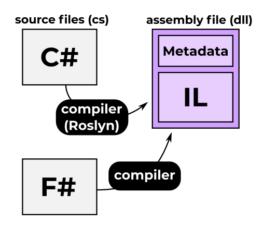
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  - it IS just some abstract stack used by this abstract machine
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Summary: CIL describes in abstract way what your application is doing.

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Summary: CIL describes **in abstract way** what your application is doing. Then .NET runtime uses this information to execute it - typically by Just-in-time compilation.

Let's see how sample C# method is compiled into CIL:

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public class C {
    public int DoCalc(int x, int y)
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        if (x > y)
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```
.class public auto ansi beforefieldinit C
    extends [System.Private.CoreLib]System.Object
    .method public hidebysig
        instance int32 DoCalc (
            int32 x,
            int32 v
        ) cil managed
         ldarg.1
         ldarg.2
         ble.s IL 0006
         ldarg.1
         ret
IL_0006: ldarg.2
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As we see it is **object-based** (classes, methods) and uses built-in types (CTS).

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ECMA-335 3.38:

#### "ldarg.<length> - load argument onto the stack

Format	Assembly Format	Description
FE 09 <unsigned int16=""></unsigned>	ldarg num	Load argument numbered <i>num</i> onto stack.
02	ldarg.0	Load argument 0 onto stack
03	ldarg.1	Load argument 1 onto stack

**Stack Transition: ...** → **...**, value"

#### Notes:

- **Idarg.1** load argument 1 (x) onto the evaluation stack
- **DoCalc** is an instance method (no static) argument 0 is the class **C** instance reference

Evaluation stack transition:  $\emptyset \rightarrow x$ 

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03	ldarg.1	Load argument 1 onto stack
04	ldarg.2	Load argument 2 onto stack

**Stack Transition: ...** → **...**, value"

#### Notes:

• **1darg.2** - load argument 2 (y) onto the evaluation stack

Evaluation stack transition:  $x \rightarrow x$ , y

```
public class C {
    public int DoCalc(int x, int y)
    {
       if (x > y) return x; else return x + y;
    }
}
```

ECMA-335 3.10:

#### "ble.<length> - branch on less than or equal to

Format	Assembly Format	Description
3E <int32></int32>	ble target	branch to target if less than or equal to
31 < <b>int8</b> >	ble.s target	branch to target if less than or equal to, short form

**Stack Transition: ...**, value1, value2 → ... "

#### Notes:

- takes out two topmost values from the evaluation stack (value1, value2)
- transfers control to target (instruction labelled as IL\_0006) if value1 ≤ value2
- target is represented as a signed offset (.s is short form for smaller offsets)

Evaluation stack transition:  $x, y \rightarrow \emptyset$ 

```
public class C {
    public int DoCalc(int x, int y)
    {
        if (x > y) return x; else return x + y;
     }
}
```

ECMA-335 3.38:

#### "ldarg.<length> - load argument onto the stack

Format	Assembly Format	Description
FE 09 <unsigned int16=""></unsigned>	ldarg num	Load argument numbered <i>num</i> onto stack.
02	ldarg.0	Load argument 0 onto stack
03	ldarg.1	Load argument 1 onto stack

**Stack Transition: ...** → **...**, value"

#### Notes:

- we are in case x > y
- **ldarg.1** load argument 1 (x) onto the evaluation stack

Evaluation stack transition:  $\emptyset \rightarrow \mathbf{x}$ 

```
public class C {
    public int DoCalc(int x, int y)
    {
       if (x > y) return x; else return x + y;
    }
}
```

#### ECMA-335 3.57:

#### "ret - return from method

Format	Assembly Format	Description
2A	Ret	Return from method, possibly returning a value

# Stack Transition: value (callee) → ..., value (caller)"

#### Notes:

- takes out the last value (x) from evaluation stack of the calee
- returns from the current method
- pushes the value on the evaluation stack of the caller

Evaluation stack transition: x (callee)  $\rightarrow ..., x$  (caller)

```
public class C {
    public int DoCalc(int x, int y)
    {
       if (x > y) return x; else return x + y;
    }
}
```

ECMA-335 3.38:

#### "ldarg.<length> - load argument onto the stack

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FE 09 <unsigned int16=""></unsigned>	ldarg num	Load argument numbered <i>num</i> onto stack.
02	ldarg.0	Load argument 0 onto stack
03	ldarg.1	Load argument 1 onto stack

**Stack Transition: ...** → **...**, value"

#### Notes:

- we are in case  $x \le y$
- **ldarg.1** load argument 1 (x) onto the evaluation stack
- we've updated offset for ble.s

Evaluation stack transition:  $\emptyset \rightarrow x$ 

```
public class C {
    public int DoCalc(int x, int y)
    {
       if (x > y) return x; else return x + y;
    }
}
```

ECMA-335 3.38:

#### "ldarg.<length> - load argument onto the stack

Format	Assembly Format	Description
		-
FE 09 <unsigned< td=""><td>ldarg num</td><td>Load argument numbered <i>num</i> onto stack.</td></unsigned<>	ldarg num	Load argument numbered <i>num</i> onto stack.
int16>		
	11	T - 1
02	ldarg.0	Load argument 0 onto stack
0.2	111	T - 1 1 t 1
03	ldarg.1	Load argument 1 onto stack
04	ldarg.2	Load argument 2 onto stack

**Stack Transition: ...** → **...**, value"

#### Notes:

• **1darg.2** - load argument 2 (y) onto the evaluation stack

Evaluation stack transition:  $x \rightarrow x$ , y

```
public class C {
    public int DoCalc(int x, int y)
    {
       if (x > y) return x; else return x + y;
    }
}
```

#### ECMA-335 3.1:

#### "add - add numeric values

Format	Assembly Format	Description
58	add	Add two values, returning a new value

**Stack Transition:** ..., value1, value2 → ..., result"

#### Notes:

- takes out two topmost values from the evaluation stack (value1, value2)
- pushes the result of adding them on the evaluation stack

Evaluation stack transition:  $x, y \rightarrow x+y$ 

```
public class C {
    public int DoCalc(int x, int y)
    {
        if (x > y) return x; else return x + y;
    }
}
```

```
.method public hidebysig
instance int32 DoCalc(int32 x, int32 y) cil managed
       ldarg.1
                  /* 03 */
       ldarg.2 /* 04 */
        ble.s IL 0006 /* 31 03 */
       ldarg.1 /* 03 */
                   /* 2A */
        ret
IL_0006: ldarg.1 /* 03 */
                 /* 04 */
       ldarg.2
       add
                   /* 58 */
                    /* 2A */
        ret
```

#### ECMA-335 3.57:

#### "ret - return from method

Format	Assembly Format	Description
2A	Ret	Return from method, possibly returning a value

# Stack Transition: value (callee) → ..., value (caller)"

#### Notes:

- takes out the last value (x+y) from evaluation stack of the calee
- returns from the current method
- pushes the value on the evaluation stack of the caller

Evaluation stack transition: x+y (callee)  $\rightarrow ..., x+y$  (caller)

```
public class C {
    public int DoCalc(int x, int y)
    {
       if (x > y) return x; else return x + y;
    }
}
```

Resulting CIL bytecode:

```
03 04 31 03 03 2A 03 04 58 2A
```

As we see:

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- interpret CIL bytecode as-is implement abstract evaluation stack and execute CIL instruction one after another
  - slow...
  - ... but has some befits (we will return to that)

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Some .NET runtime implementation could:

- interpret CIL bytecode as-is implement abstract evaluation stack and execute CIL instruction one after another
  - o slow...
  - ... but has some befits (we will return to that)
- compile CIL bytecode into native code here we go Just-in-time compiler!
  - JIT happens at the target machine/runtime it will produce x86/x64/ARM code for Linux/Windows/...

```
public class C {
    public int DoCalc(int x, int y)
    {
        if (x > y)
            return x;
        else
            return x + y;
     }
}
```

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public class C {
    public int DoCalc(int x, int y)
    {
        if (x > y)
            return x;
        else
            return x + y;
     }
}
```

```
C.DoCalc(Int32, Int32)
          cmp edx, r8d
          jle short L0008
          mov eax, edx
          ret
L0008: lea eax, [rdx+r8]
          ret
```

So-called *calling convention* on Windows x64:

- for the first fours arguments uses rcx, rdx, r8d and r9d CPU registers
  - (for the rest it uses *thread stack*)
- for the function result uses **rax** CPU register

#### So:

- rcx (hidden) class C instance reference
- rdx x argument
- r8d y argument
- rax for the result

```
public class C {
    public int DoCalc(int x, int y)
    {
        if (x > y)
            return x;
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            return x + y;
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ret

L0008: lea eax, [rdx+r8]

ret
```

- with the help of cmp instruction, compare edx (x argument) with r8d (y argument)
- if  $x \le y$  then jump to L0008 instruction
- **note:** no evaluation stack at all arguments have been translated into CPU registers

```
public class C {
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.method public hidebysig
instance int32 DoCalc(int32 x, int32 y) cil managed
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       ble.s IL_0006
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IL_0006: ldarg.1
       ldarg.2
       add
       ret
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C.DoCalc(Int32, Int32)
cmp edx, r8d
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ret

L0008: lea eax, [rdx+r8]
ret
```

- in case of x > y simply set the result (eax, shortcut of rax) to the edx (which is x) value
- and return

```
public class C {
    public int DoCalc(int x, int y)
    {
        if (x > y)
            return x;
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            return x + y;
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        ret
L0008: lea eax, [rdx+r8]
        ret
```

- in case of x ≤ y use a trick with lea instruction to calculate the sum of arguments (rdx and r8)...
- ...store the result in the eax register...
- and return

### JIT compilation

```
public class C {
    public int DoCalc(int x, int y)
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            return x;
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            return x + y;
     }
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          jle short L0008
          mov eax, edx
          ret
L0008: lea eax, [rdx+r8]
          ret
```

#### Summary:

- no evaluation stack! It's gone:
  - only CPU registers usage
  - thread stack will be used for local variables in more complex methods (but it is **not the same**, remember!)
- it is a pure native code, nothing special about CLR:
  - any programming language, **including C++**, could produce the same result
  - no magic calls into CLR (to check/calculate/validate/...)
  - therefore: this code runs at **native speed** 😂
- so far, we've used optimized Release mode

### **CIL to native code (Debug)**

```
public class C {
    public int DoCalc(int x, int y)
    {
        if (x > y)
            return x;
        else
            return x + y;
     }
}
```

Debug mode is much more verbose:

- much less efficient generated code
- typically stack frame is created to store all local variables
- lifetime of local objects is much longer (we will return to that...)
- contains additional calls into CLR (fe., call 0x00007ffc78028010 here is a "jit helper" CORINFO\_HELP\_DBG\_IS\_JUST\_MY\_CODE which checks if this is "JustMyCode" and needs to be stepped through)

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```
C.DoCalc(Int32, Int32)
    push rbp
    sub rsp. 0x30
    lea rbp, [rsp+0x30]
    xor eax, eax
    mov [rbp-4], eax
    mov [rbp-8], eax
    mov [rbp+0x10], rcx
    mov [rbp+0x18], edx
    mov [rbp+0x20], r8d
    mov rax, 0x7ffc21a0c320
    cmp dword ptr [rax], 0
    ie short L0031
    call 0x00007ffc78028010
L0031: nop
    mov eax, [rbp+0x18]
    cmp eax, [rbp+0x20]
    setg al
    movzx eax, al
    mov [rbp-4], eax
    cmp dword ptr [rbp-4], 0
    ie short L0050
    mov eax, [rbp+0x18]
    mov [rbp-8], eax
    nop
    jmp short L005c
L0050: mov eax, [rbp+0x18]
    add eax. [rbp+0x20]
```

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  - Release mode highly optimized, but it **takes time** to generate
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And the story continues. Non-trivial methods also run on "native speed", like the following  $\pi$  calculation with the help of the Gregory-Leibniz series:

```
// π/4 = 1 - 1/3 + 1/5 - 1/7 + 1/9 - ...
public static double GregoryLeibnizPI(int n)
{
    double sum = 0;
    for (int i = 0; i < n; i++)
    {
        double temp = 4.0 / (1 + 2 * i);
        sum += i % 2 == 0 ? temp : -temp;
    }
    return sum;
}</pre>
```

Again, there are no special calls into CLR etc.

```
GregoryLeibnizGetPI(Int32)
    L0000: vzeroupper
    L0003: vxorps xmm0, xmm0, xmm0
    L0007: xor eax, eax
    L0009: test ecx, ecx
    L000b: ile short L0052
    L000d: lea edx, [rax*2+1]
   L0014: vxorps xmm1, xmm1, xmm1
   L0018: vcvtsi2sd xmm1, xmm1, edx
   L001c: vmovsd xmm2, [C.GregoryLeibnizPI(Int32)]
   L0024: vdivsd xmm1, xmm2, xmm1
   L0028: mov edx, eax
   L002a: shr edx, 0x1f
   L002d: add edx, eax
   L002f: and edx. 0xfffffffe
   L0032: mov r8d, eax
    L0035: sub r8d, edx
    L0038: ie short L0048
    L003a: vmovsd xmm2, [C.GregoryLeibnizPI(Int32)]
    L0042: vxorps xmm1, xmm1, xmm2
    L0046: jmp short L0048
    L0048: vaddsd xmm0, xmm0, xmm1
    L004c: inc eax
    L004e: cmp eax, ecx
    L0050: il short L000d
    L0052: ret
```

An important technique is **inlining**, which is a JIT power to *inline* a method body into another, instead of calling it:

```
public class C {
   public static void M() {
      var rand = new Random();
      N(rand.Next());
   }
   public static void N(int x) {
      Console.WriteLine(x);
   }
}
```

```
.method public hidebysig s
static void M () cil managed
{
   newobj instance void System.Random::.ctor()
   callvirt instance int32 System.Random::Next()
   call void C::N(int32)
   ret
}
```

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   ret
}
```

So, the generated JIT code calls **Console.WriteLine** from the **M** method itself:

```
C.M()
...
mov ecx, esi
call System.Random..ctor(Int32)
mov ecx, esi
call System.Random.InternalSample()
mov ecx, eax
call System.Console.WriteLine(Int32)
...
```

And this happens for *simple* methods and some implementation detail-driven decisions.

Unless, we disable inlining explicitly by MethodImpl attribute (which may be useful mostly in doing some fine-grained benchmarks, as we will see in the course):

```
public class C {
    public static void M() {
        var rand = new Random();
        N(rand.Next());
    }
    [MethodImpl(MethodImplOptions.NoInlining)]
    public static void N(int x) {
        Console.WriteLine(x);
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    public static void N(int x) {
        Console.WriteLine(x);
    }
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.method public hidebysig s
static void M () cil managed
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   callvirt instance int32 System.Random::Next()
   call void C::N(int32)
   ret
}
```

So, the generated JIT code calls **N** from the **M** method, as expected:

```
C.M()
    ...
    mov ecx, esi
    call System.Random..ctor(Int32)
    mov ecx, esi
    call System.Random.InternalSample()
    mov ecx, eax
    call C.N()
    ...
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- summarizing (pseudo-definition):

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  - from 7.9.6.6 Constructors we can read that when an object is created, "space for the new value is allocated in managed memory"
- summarizing (pseudo-definition):
  - .NET uses managed memory (aka Managed Heap) where managed objects are allocated (and optionally, automatically released) by garbage collection

- where's the GC so far? We haven't seen anything related to memory management, yet.
- in ECMA-335 the topic of memory management is little scattered and vague:
  - "managed data", "memory store" and even "garbage collection" mentioned here and there
  - there is no single paragraph/definition that summarizes it all
- mostly:
  - "Managed data data that is allocated and released automatically by the CLI, through a process called garbage collection."
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  - there are no enforced garbage collection implementation details in ECMA
  - only some restrictions to allow, e.g. compacting/moving GC etc.

- putting aside a little vague definitions, in the end we have two CIL instructions for **allocating** new objects:
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- those two are the only ones about "memory management":
  - there is no **delobj** or **delarr** to reclaim memory
  - memory reclamation is automatic, after it is discovered that object is no longer reachable (how it is being discovered will be covered in next modules)

```
using System;
public class Program {
    public static void Main() {
        var logger = new Logger(3);
        var level = 44;
        logger.Log(level, "Hello world!");
    }
}
public class Logger
{
    private int _level;
    public Logger(int level) => _level = level;
    public void Log(int level, string str)
    {
        if (level >= _level) Console.WriteLine(str);
    }
}
```

Main was represented in CIL:

```
19 73 03 00 00 06 1F 2C 0A 06 72 01 00 00 70 6F 04 00 00 06 2A
```

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19 73 03 00 00 06 1F 2C 0A 06 72 01 00 00 70 6F 04 00 00 06 2A
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Which is in textual form:

```
.method public hidebysig static
void Main () cil managed
  .maxstack 3
 .entrypoint
  .locals init ([0] int32 level)
 ldc.i4.3
                         /* 19
 newobj instance void HelloWorld.Logger::.ctor(...
                         /* 7303000006 */
 ldc.i4.s 44
                         /* 1F2C
 stloc.0
                         /* 0A
 ldloc.0
                         /* 06
 ldstr "Hello world!" /* 7201000070 */
 callvirt instance void HelloWorld.Logger::Log(...
                         /* 6F04000006 */
                         /* 2A
 ret
                                      */
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 ldstr
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                          /* 6F04000006 */
 ret
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                                        */
```

**Sidenote:** We see some redundancy here produced by C# compiler:

- stloc.0 pop value from evaluation stack to local variable 0
- ldloc.0 load local variable 0 onto the evaluation stack

Which is no-op and it doesn't make sense 🗑

Sidenote (cont.): Mentally, we can get rid of those two instructions!

```
.method public hidebysig static
void Main () cil managed
{
    .maxstack 3
    .entrypoint
    .locals init ([0] int32 level)

    ldc.i4.3
    newobj instance void HelloWorld.Logger::.ctor(int32)
    ldc.i4.s 44
    stloc.0
    ldloc.0
    ldstr "Hello world!"
    callvirt instance void HelloWorld.Logger::Log(int32, string)
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}
```

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## CIL, memory management and JIT

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```
Program.Main()
    sub rsp, 0x28
    mov rcx, 0x7ffc2265d050
    call 0x00007ffc77ee9b70
    mov dword ptr [rax+8], 3
    mov r8, 0x257c846d0c0
    mov r8, [r8]
    mov rcx, rax
    mov edx, 0x2c
    add rsp, 0x28
    jmp Logger.Log(Int32, System.String)
```

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- again, we don't see anything about evaluation stack only CPU registers usage
- new expression from C# has been translated to newobj and constructor call
- newobj is JITted to a call into CLR allocator
  - now our code is indeed "managed" it uses CLR :)
- constructor call has been inlined
- there is nothing about garbage collection
  - runtime will discover that Logger is not needed we'll cover that!

**Sidenote**: there is also **localloc** CIL instruction, exposed to C# as **stackalloc** expression, which allocates bytes from so-called **local dynamic memory pool**. We'll cover that much later as more advanced use case.

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  - write low-level, *magically* fast code it is not x64 assembly, JIT will do the work (mostly) (as we will see in the demo)

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  - ∘ for... fun 🗑

#### .NET runtimes

- .NFT Framework
- .NET Core 1.0-3.1/.NET 5-6
- Mono (runs Blazor/Xamarin/Unity)
- RIP .NET Compact Framework & .<u>NET Micro Framework</u>
- Rotor/SSCLI Shared Source Common Language Infrastructure
- <u>DotNetAnywhere</u> by Chris Bacon (initial prototype of Blazor)
- DotGNU & Portable .NET (...)
- CoreRT

# .NET runtime & C++

C++ piece of the .NET runtime is not-so-obvious:

- .NET Core more and more parts are being ported to C# (better maintainability, GC pauses/less managed-unmanaged transitions)
  - port JIT and GC to C# issue
  - stdelemref and ldelemaref JIT helpers to C# (dotnet/runtime#32722)
  - moving portions of the unbox helper to C# (dotnet/runtime#32353)
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  - o ...
- CoreRT... many parts rewritten to C# (like type system)
  - still, crucial parts written in C++ GC, JIT, ...

#### **Materials**

- Taking a look at the ECMA-335 Standard for .NET
- Deep-dive into .NET Core primitives: deps.json, runtimeconfig.json, and dll's
- <u>Tiered Compilation</u>
- "Expert .NET 2.0 IL Assembler" by Serge Lidin
- "Compiling for the .Net Common Language Runtime (CLR)" by John Gough

#### **DEMO**