#### UC Computer Science and Software Engineering

## COSC362 Data and Network Security Semester Spring, 2021

#### **Lab 5: Number Theory**

Exercises from Lecture 10.

#### **QUESTION 1**

If possible (you need to check!), solve for x using the Chinese Remainder Theorem (CRT):

- (a)  $x \equiv 5 \pmod{7}$  and  $x \equiv 7 \pmod{10}$
- (b)  $x \equiv 3 \pmod{7}$  and  $x \equiv 7 \pmod{14}$
- (c)  $x \equiv 2 \pmod{6}$  and  $x \equiv 3 \pmod{11}$

**Methodology:** First find the GCD of p and q. If the GCD is equal to 1 (i.e. p and q are relatively prime) then a solution must exist; otherwise, there is no solution and CRT cannot be applied. If a solution exists, then use the CRT to find x.

More precisely, when gcd(p, q) = 1, we apply CRT as follows. Let p, q be relatively prime. Let n = pq be the modulus. Given integers  $c_1, c_2$ , there exists a unique integer  $x, 0 \le x < n$ , s.t.:

$$x \equiv c_1 \pmod{p}$$
$$x \equiv c_2 \pmod{q}$$

Therefore, the CRT tells us that  $x \equiv \frac{n}{p}y_1c_1 + \frac{n}{q}y_2c_2 \pmod{n}$  where:

$$y_1 \equiv \left(\frac{n}{p}\right)^{-1} \pmod{p}$$
  
 $y_1 \equiv q^{-1} \pmod{p}$   
 $qy_1 \equiv 1 \pmod{p}$  [another way to write the above line]

and

$$y_2 \equiv \left(\frac{n}{q}\right)^{-1} \pmod{q}$$
  
 $y_2 \equiv p^{-1} \pmod{q}$   
 $py_2 \equiv 1 \pmod{q}$  [another way to write the above line]

We observe that  $y_1$  is the inverse of q modulo p. Similarly,  $y_2$  is the inverse of p modulo q.

See in Lecture 3, the slide entitled "Modular Inverses using the Euclidean Algorithm". There exist 2 integers  $k_1, k_2$  such that:

$$qy_1 + pk_1 = 1$$
$$py_2 + qk_2 = 1$$

Observe that the unknowns are  $y_1, y_2, k_1, k_2$ . Then using back substitution, we find those values.

We replace  $y_1$  and  $y_2$  with the values that we have just found in:

$$x \equiv \frac{n}{p} y_1 c_1 + \frac{n}{q} y_2 c_2 \pmod{n}$$
$$\equiv q y_1 c_1 + p y_2 c_2 \pmod{n}$$

and we find x (do not forget to reduce modulo n!).

**Example:** Solve for x using the Chinese Remainder Theorem (CRT):  $x \equiv 2 \pmod{5}$  and  $x \equiv 3 \pmod{7}$ .

Let p=5 and q=7, and the modulus n=pq=35. p and q are 2 prime numbers, so they are relatively prime, so a solution x must exist such that:

$$x \equiv 2 \pmod{5}$$
$$x \equiv 3 \pmod{7}$$

Here,  $c_1 = 2$  and  $c_2 = 3$ .

We now find  $y_1$  and  $y_2$  such that:

$$qy_1 \equiv 1 \pmod{p}$$
$$7y_1 \equiv 1 \pmod{5}$$

and

$$py_2 \equiv 1 \pmod{q}$$
$$5y_2 \equiv 1 \pmod{7}$$

Let  $k_1, k_2$  be 2 integers such that  $7y_1 + 5k_1 = 1$  and  $5y_2 + 7k_2 = 1$ . We find that  $y_1 = 3$  and  $k_1 = -4$ :

$$7y_{1} + 5k_{1} = 1$$

$$y_{1} = \frac{1 - 5k_{1}}{7}$$

$$= \frac{1 - 5 \times (-4)}{7}$$

$$= \frac{1 + 20}{7}$$

$$= \frac{21}{7}$$

$$= 3$$

and that  $y_2 = 3$  and  $k_2 = -2$ :

$$5y_{2} + 7k_{2} = 1$$

$$y_{2} = \frac{1 - 7k_{2}}{5}$$

$$= \frac{1 - 7 \times (-2)}{5}$$

$$= \frac{1 + 14}{5}$$

$$= \frac{15}{5}$$

$$= 3$$

Now we use the values found for  $y_1$  and  $y_2$  in:

$$x \equiv \frac{n}{p} y_1 c_1 + \frac{n}{q} y_2 c_2 \pmod{n}$$

$$\equiv q y_1 c_1 + p y_2 c_2 \pmod{n}$$

$$\equiv (7 \times 3 \times 2) + (5 \times 3 \times 3) \pmod{35}$$

$$\equiv 42 + 45 \pmod{35}$$

$$\equiv 17 \pmod{35}$$

### **QUESTION 2**

Find  $\phi(20)$ ,  $\phi(21)$ ,  $\phi(22)$ ,  $\phi(23)$ ,  $\phi(24)$ ,  $\phi(25)$ .

From the lecture slides:

- $\phi$  is the Euler function.
- $\phi(p) = p 1$  where p is prime.

- $\phi(pq) = (p-1)(q-1)$  where p, q are distinct primes.
- $n = p_1^{e_1} \cdots p_t^{e_t}$  where  $p_i$  are distinct primes, then:

$$\phi(n) = \prod_{i=1}^{t} p_i^{e_i - 1} (p_i - 1)$$

## **QUESTION 3**

Find the discrete logarithm of the number 3 with regard to base 2 for:

- (a) modulus p = 5
- (b) modulus p = 11
- (c) modulus p = 29

In other words, we need to find the value x such that  $2^x = 3 \mod p$  for the above values of p. To do so, we calculate  $2^1 \mod p$ ,  $2^2 \mod p$ ,  $2^3 \mod p$ ,  $2^4 \mod p$ , etc. until finding x such that  $2^x = 3 \mod p$ . More information can be found on slides 35 and 36 of Lecture 10.

### **QUESTION 4**

Use the Fermat test to check whether the following numbers are prime or not:

- 979
- 983

Run the test **at most** 4 times with base values a equal to 2, 3, 11, 17. In particular, we check whether  $a^{979-1} \mod 979$  is equal to 1 or not, for a = 2, 3, 11, 17. Similarly, we check whether  $a^{983-1} \mod 983$  is equal to 1 or not, for a = 2, 3, 11, 17. Check slides 13 and 14 of Lecture 10.

Note that these base values a=2,3,11,17 are not random, and in practice, fixed bases are usually applied.

**Hint:**  $ab \mod n = (a \mod n)(b \mod n) \mod n$ , and in particular,  $(a^m)^k \mod n = (a^m \mod n)^k \mod n$ .

<sup>&</sup>lt;sup>1</sup>See for instance https://www.khanacademy.org/computing/computer-science/cryptography/modarithmetic/a/fast-modular-exponentiation.

# **QUESTION 5**

We first recall the Miller-Rabin algorithm. Let n and u be odd, and v s.t.  $n-1=2^vu$ :

- (a) Pick a at random s.t. 1 < a < n 1
- (b) Set  $b = a^u \mod n$
- (c) If b=1 then return probable prime
- (d) For j = 0 to v 1:
  - If b = -1 then return probable prime
  - Else set  $b = b^2 \mod n$
- (e) Return composite

Use the Miller-Rabin algorithm for:

(a) n = 17.

We can easily see that n is prime. Let us see what the Miller-Rabin test tells us.

(b) n = 15.

We know that  $15 = 3 \times 5$ , hence n is **not** prime. Let us see what the Miller-Rabin test tells us.