1. Case Study

1.1. Performance Optimization

We show several examples to demonstrate how EDITLORD improve performance optimization.

Given slow code

```
#include <bits/stdc++.h>
using namespace std;
long long n; long long k;
vector<long long> vec;
long long meme[100007];
long long solve(long long index) {
  if(index==n-1){
    return 0;
  if (meme[index]!=-1) {
    return meme[index];
  long long ret = 10000000000000000;
  long long sol = 100000000000000000000;
  for(long long i = 1 ;i<=k ; i++) {</pre>
    if (index+i<n) {</pre>
      ret = abs (vec[index]-

    vec[index+i]) +

      ⇔ solve(index+i);
    sol = min(sol, ret);
  return meme[index] = sol;
int main(){
  cin>>n;
  cin>>k;
  memset (meme, -1, sizeof (meme));
  long long num;
  for (long long i = 0; i < n; i++) {
    cin>>num;
    vec.push_back(num);
  cout << solve(0) << endl;</pre>
  return 0;
```

```
The given code describe the following
\rightarrow problem: Find the minimum cost to reach
\rightarrow the last index of a sequence of integers,

→ difference between consecutive integers.

The input specification is: The programs take
\rightarrow two integers n and k as input, where n is
   the number of integers in the sequence
   and k is the maximum jump length. Then, n

→ sequence.

The output specification is: Output the
\rightarrow minimum cost to reach the last index from
\hookrightarrow the first index.
Following editing rules should be applied:
switch from multiple nested loops to single

→ linear iterations

switch from recursive function calls to
\rightarrow optimized iterative data handling methods
The corresponding fast code is:
[FAST CODE]
#include<bits/stdc++.h>
using namespace std;
long long int dp[100007],a[100007];
int main(){
 long long int n,k,i,j,x,y,z;
  cin>>n>>k;
 for (i=0; i<n; i++) {
   cin>>a[i];
 dp[0]=0;
  for (i=1; i<n; i++) {</pre>
   x=10000000000007;
   for(j=i-1; j>=i-k&&j>=0; j--) {
      x=min(x,dp[j]+abs(a[i]-a[j]));
   dp[i]=x;
  cout << dp[n-1];
[/FAST CODE]
```

Figure 1: Performance optimization example 1.

Given slow code

```
#include <bits/stdc++.h>
using namespace std;
int n ,k;
vector<int>vi;
long long mem[200005];
bool vis[200005];
long long solve(int indx ) {
  if(indx == n-1)
   return 0;
  if(vis[indx])
   return mem[indx];
  vis[indx] = true;
  mem[indx] = 1e13;
  for(int i = indx+1; i <=min(indx+k)</pre>
  \hookrightarrow , n-1); i++){
    mem[indx] = min(mem[indx],
    \hookrightarrow solve(i) + abs(vi[indx] -
     \hookrightarrow vi[i]));
  return mem[indx];
int main(){
  cin>>n>>k;
  vi.resize(n);
  for(int i = 0; i <n; i++)</pre>
   cin>>vi[i];
  cout << solve (0);
  return 0;
```

EDITLORD output

```
The given code describe the following
\rightarrow problem: Calculate the minimum cost to
    reach the last index in a list of
\hookrightarrow integers, where each step can jump up to
\hookrightarrow k positions.
The input specification is: The programs take
\rightarrow two integers n and k as input, where n is
\,\hookrightarrow\, the number of integers in the list and k
\hookrightarrow is the maximum jump length. Then, n \hookrightarrow integers are read, representing the cost
\hookrightarrow associated with each index.
The output specification is: Output the
\,\,\,\hookrightarrow\,\, minimum cost to reach the last index from
\hookrightarrow the first index.
Following editing rules should be applied:
switch from multiple nested loops for
\hookrightarrow counting elements to direct arithmetic
\hookrightarrow calculation
switch from multiple nested loops to a

→ compact function

switch from cout to printf
switch from cin to scanf
switch from dynamic memory allocation to
\hookrightarrow static memory allocation
The corresponding fast code is:
[FAST CODE]
#include<bits/stdc++.h>
using namespace std;
int n, k;
long long dp[100005];
long long arr[100005];
int main(){
  scanf("%d%d", &n, &k);
  for (int i=1; i<=n; i++)</pre>

    scanf("%lld", &arr[i]);

  for (int i=2;i<=n;i++) {</pre>
    dp[i]=1e18;
    for (int j=max(i-k,1); j<i; j++) {</pre>
       dp[i]=min(dp[i],

    dp[j]+abs(arr[i]-arr[j]));
  printf("%lld",dp[n]);
[/FAST CODE]
```

Figure 2: Performance optimization example 2.

1.2. Decompilation

We show several examples to demonstrate how EDITLORD improves the readability of machine-decompiled code.

Given machine decompiled code

```
void func0(char *param_1,long
\rightarrow param_2)
 char cVar1;
  long lVar2;
 ulong uVar3;
 char *pcVar4;
 int iVar5;
 bool bVar6;
 1Var2 = -1;
 pcVar4 = param_1;
 do {
    if (1Var2 == 0) break;
    1Var2 = 1Var2 + -1;
   cVar1 = *pcVar4;
   pcVar4 = pcVar4 + 1;
  } while (cVar1 != '\0');
  iVar5 = (uint) 1Var2 - 1;
  if (0 < iVar5) {
    uVar3 = 0;
    do {
      cVar1 = param_1[uVar3];
      if ((byte)(cVar1 + 0x9fU) <</pre>
      \rightarrow 0x1a) {
        cVar1 = cVar1 + -0x20;
      else if ((byte)(cVar1 + 0xbfU)
      \hookrightarrow < 0x1a) {
        cVar1 = cVar1 + ' ';
      *(char *)(param_2 + uVar3) =
      bVar6 = uVar3 != ~(uint) lVar2 -

→ 2;

      uVar3 = uVar3 + 1;
    } while (bVar6);
  *(undefined *)(param_2 + iVar5) =

    ○;
  return;
```

```
The given code describe the following
\rightarrow problem: Convert a string from one case
\hookrightarrow (upper or lower) to another, handling

→ special characters appropriately.

The input specification is: The programs take
\hookrightarrow a string as input, which may contain
\,\hookrightarrow\, uppercase and lowercase letters, and a
\rightarrow buffer to store the converted string.
The output specification is: Output the
\rightarrow converted string, where all letters are

→ converted to lowercase, while preserving

Following editing rules should be applied:
switch from arbitrary pointer checks to

→ clearer termination conditions

switch from complex pointer arithmetic to

→ clear variable assignments

switch from indirect variable assignments to
\hookrightarrow direct return statements
switch from complex multi-step declarations
\hookrightarrow to clear single declarations
switch from complex logic to simple loops
switch from complex complex if-else structure
switch from implicit return values to

→ explicit return values

switch from cryptic variable names to

→ descriptive variable names

The corresponding source code is:
[ORIGINAL SOURCE CODE]
static void
func0 (char *str, char *buf)
  int i, len = strlen (str);
  for (i = 0; i < len; i++) {</pre>
   if (str[i] >= 'A' && str[i] <= 'Z')</pre>
     buf[i] = str[i] + ' ';
    else if (str[i] >= 'a' && str[i] <= 'z')
     buf[i] = str[i] - ' ';
    else
      buf[i] = str[i];
 buf[len] = ' \setminus 0';
[/ORIGINAL SOURCE CODE]
```

Figure 3: Decomilation example 1.

Given machine decompiled code

```
undefined8 func0(char *param_1,char
\rightarrow *param_2)
  char cVar1;
 int iVar2;
 uint uVar3;
 char *pcVar4;
  char *pcVar5;
 long lVar6;
  size_t __n;
 bool bVar7;
 1Var6 = -1;
  pcVar4 = param_2;
  do {
    if (lVar6 == 0) break;
    1Var6 = 1Var6 + -1;
    cVar1 = *pcVar4;
   pcVar4 = pcVar4 + 1;
  } while (cVar1 != '\0');
  uVar3 = (uint) lVar6;
  iVar2 = uVar3 - 1;
  pcVar4 = (char

    *)malloc((long)(int)(uVar3 +
      iVar2));
  if (0 < iVar2) {
    lVar6 = (long)iVar2;
      _n = 0;
    do {
      strncpy(pcVar4,param_2 +
      \rightarrow __n,lVar6 - __n);
      strncpy(pcVar4 + (1Var6 -
      \rightarrow __n),param_2,__n);
      pcVar4[lVar6] = '\0';
      pcVar5 =

    strstr(param_1,pcVar4);

      if (pcVar5 != (char *) 0x0) {
        free (pcVar4);
        return 1;
      bVar7 = \underline{\quad } n != uVar3 - 2;
       \underline{n} = \underline{n} + 1;
    } while (bVar7);
  free (pcVar4);
  return 0;
```

```
The given code describe the following
→ problem: Determine if a string is a
→ rotation of another string by checking if
\rightarrow any substring of the second string is a
\rightarrow substring of the first string.
The input specification is: The programs take
\hookrightarrow two strings as input: the first string is
\hookrightarrow the original string, and the second \hookrightarrow string is the potential rotation.
The output specification is: Return a boolean
\hookrightarrow value indicating whether the second
\rightarrow string is a rotation of the first string.
Following editing rules should be applied:
switch from ambiguous function signatures to

→ clear function signatures

switch from complex pointer arithmetic to

→ clear variable assignments

switch from indirect variable assignments to

→ direct return statements

switch from complex logic to simple loops
switch from complex loops to clear
\hookrightarrow single-responsibility sections
switch from complex nested conditionals to

→ clearer sequential logic

switch from complex variable operations to
\hookrightarrow direct assignments
switch from redundant checks to

→ straightforward boolean comparisons

The corresponding source code is:
[ORIGINAL SOURCE CODE]
static int
func0(const char *s1, const char *s2)
  char *s;
  size_t len, i;
  len = strlen(s2);
  s = malloc(len * 2 + 1);
  if (s == NULL) return (0);
  for (i = 0; i < len; i++) {</pre>
    strncpy(s, s2 + i, len - i);
    strncpy(s + len - i, s2, i);
    s[len] = '\0';
    if (strstr(s1, s) != NULL) {
      free(s);
      return (1);
    }
  }
  free(s);
  return (0);
[/ORIGINAL SOURCE CODE]
```

Figure 4: Decomilation example 2.

1.3. Security Hardening

We show several examples to demonstrate how EDITLORD hardens the vulnerable code.

Given vulnerable code

```
The program defines a function that allocates
\hookrightarrow initializes its contents by writing 'A'
\hookrightarrow characters at both the start and end of
\hookrightarrow the buffer. The goal is to ensure that
\rightarrow the buffer is filled correctly without
\rightarrow exceeding its allocated size, thereby
\rightarrow preventing potential buffer overflows.
The input specification is: The function

→ takes a single input parameter, buf_size,

   which specifies the size of the buffer to
\rightarrow be allocated. This size determines how
\rightarrow much memory will be reserved for the
\hookrightarrow buffer.
The output specification is: The function
\rightarrow returns a pointer to the allocated and
\,\,\hookrightarrow\,\, initialized buffer. If the allocation is
\hookrightarrow successful, the buffer will contain 'A'
   characters at the specified positions;
   otherwise, it may return NULL if memory
\hookrightarrow allocation fails.
Following editing rules should be applied:
switch from no checks on function return
\rightarrow values to check function return value
The corresponding secure code is:
[SECURE CODE]
char* allocate_buffer_and_write(size_t
→ buf size) {
    char* buf = (char*)malloc(buf_size);
    if (!buf) {
        return 0;
    memset(buf, 'A', 10);
    memset(buf + buf_size - 10, 'A', 10);
    return buf;
[/SECURE CODE]
```

Figure 5: Secure hardening example 1.

```
#include <stdio.h>
#include <stdlib.h>
#include <sqlite3.h>
#include <string.h>
char *get_user(sqlite3 *db, const
\hookrightarrow char *username, const char
  *password) {
    sqlite3_stmt *stmt;
    char query[256];
    char *userInfo = NULL;
    snprintf(query, sizeof(query),
    → "SELECT userInfo FROM users

→ WHERE username = '%s' AND

→ password = '%s'", username,
    → password);
    if (sqlite3_prepare_v2(db, query,
    \hookrightarrow -1, &stmt, NULL) !=

    SQLITE_OK) {

        fprintf(stderr, "Failed to
         → prepare statement: %s\n",

    sqlite3_errmsg(db));

        return NULL;
    if (sqlite3_step(stmt) ==

    SQLITE_ROW) {

        const unsigned char *result
         \hookrightarrow =

    sqlite3_column_text(stmt,

         \rightarrow 0);
        userInfo = strdup((const char

    *)result);
    sqlite3_finalize(stmt);
    return userInfo;
```

```
The program aims to retrieve user information
\hookrightarrow from a SQLite database based on a
\rightarrow provided username and password. It
\rightarrow connects to the database, executes a
\rightarrow query to find the user, and returns the
\rightarrow user's information if found. The program
appropriately to prevent SQL injection
→ vulnerabilities.
The input specification is: The program takes
\hookrightarrow three command-line arguments: the
\hookrightarrow database filename, a username, and a

→ password. These inputs are used to query

→ the database for matching user

\hookrightarrow credentials.
The output specification is: The program
\hookrightarrow outputs the user's information if a \hookrightarrow matching user is found in the database,
\rightarrow or an error message if no user matches
\rightarrow the provided credentials or if an error
\rightarrow occurs during database operations.
Following editing rules should be applied:
switch from direct SQL string interpolation
\hookrightarrow to use of parameterized logic
The corresponding secure code is:
[SECURE CODE]
char *get_user(sqlite3 *db, const char
\rightarrow *username, const char *password) {
 sqlite3_stmt *stmt;
 char *userInfo = NULL;
  const char *query = "SELECT userInfo FROM
  → users WHERE username = ? AND password =
  if (sqlite3_prepare_v2(db, query, -1,
  fprintf(stderr, "Failed to prepare

    statement: %s\n",

    sqlite3_errmsg(db));
   return NULL;
  sqlite3_bind_text(stmt, 1, username, -1,

    SQLITE_STATIC);

 sqlite3_bind_text(stmt, 2, password, -1,

    SQLITE_STATIC);
  if (sqlite3_step(stmt) == SQLITE_ROW) {
    const unsigned char *result =

    sqlite3_column_text(stmt, 0);

   userInfo = strdup((const char *)result);
  sqlite3_finalize(stmt);
 return userInfo;
[/SECURE CODE]
```

Figure 6: Secure hardening example 2.

```
#include <string.h>

void trim_spaces(char* str) {
   int len = strlen(str);
   while (str[len - 1] == ' ') {
      str[len - 1] = '\0';
      len--;
   }
}
```

```
The program defines a function that trims
\rightarrow trailing spaces from a given string. It
\rightarrow modifies the string in place, ensuring
\hookrightarrow that any whitespace characters at the end
\,\hookrightarrow\, are removed, leaving the rest of the
\hookrightarrow string intact.
The input specification is: The function
\hookrightarrow takes a single string input, which is a
\hookrightarrow character array that may contain trailing
\rightarrow spaces that need to be removed.
The output specification is: The function
\rightarrow does not return any value; instead, it
\rightarrow modifies the input string by removing any

    → trailing spaces, effectively shortening

   the string if necessary.
Following editing rules should be applied:
switch from unvalidated input handling to
\rightarrow check for NULL pointers before access
switch from unvalidated input handling to
\rightarrow check for buffer overflows on memory
\hookrightarrow accesses
The corresponding secure code is:
[SECURE CODE]
void trim_spaces(char* str) {
    int len = strlen(str);
    while (len > 0 && str[len - 1] == ' ') {
         str[len - 1] = ' \ 0';
         len--;
[/SECURE CODE]
```

Figure 7: Secure hardening example 3.

```
The program aims to read a line of input from
\rightarrow the standard input (stdin) into a
\rightarrow dynamically allocated buffer, ensuring
\rightarrow that the input does not exceed the
\rightarrow specified buffer size. It is designed to
\rightarrow prevent buffer overflow vulnerabilities
   by limiting the number of characters read
\rightarrow based on the buffer's capacity.
The input specification is: The function
\rightarrow takes an integer input representing the
\rightarrow maximum size of the buffer, which
\rightarrow dictates how many characters can be
\hookrightarrow safely read from stdin.
The output specification is: The function
\,\hookrightarrow\, containing the user input, ensuring that
   the input is properly null-terminated and

→ does not exceed the allocated buffer

   size.
Following editing rules should be applied:
switch from unvalidated input handling to
\rightarrow check for buffer overflows on memory
\hookrightarrow accesses
switch from no error handling to check for
\hookrightarrow NULL pointers after memory allocation
The corresponding secure code is:
[SECURE CODE]
char* read_stdin_to_buffer(int buf_size) {
    char* buf = (char*) malloc (buf_size);
    fgets(buf, buf_size, stdin);
    return buf;
[/SECURE CODE]
```

Figure 8: Secure hardening example 4.

```
#include <stdlib.h>
int get_price_at_index(int* prices,
    int size, char* index_str) {
    int index = atoi(index_str);
    return prices[index];
}
```

```
The program defines a function that retrieves
\rightarrow the price of an item at a specified index
\rightarrow from an array of prices. The index is
→ provided as a string, which is converted
\rightarrow price at the corresponding index if it's
   valid, or an error if the index is out of
\hookrightarrow bounds.
The input specification is: The function
\rightarrow takes three inputs: an array of integers
\rightarrow representing prices, the size of the

→ array, and a string representing the

\rightarrow index of the price to retrieve.
The output specification is: The function
\,\hookrightarrow\, returns the price at the specified index
\rightarrow if it is valid; otherwise, it returns an
\hookrightarrow error message indicating that the index \hookrightarrow is out of bounds.
Following editing rules should be applied:
switch from direct parsing to implement
\rightarrow comprehensive validation checks for
The corresponding secure code is:
[SECURE CODE]
int get_price_at_index(int* prices, int size,
⇔ char* index_str) {
    int index = atoi(index_str);
    if (index < 0 || index >= size) {
        return -1; // Error: index out of
        → bounds
    return prices[index];
[/SECURE CODE]
```

Figure 9: Secure hardening example 5.