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EXAM 1 STARTS HERE

Q1 The statement is true since the two numbers can always be equal

Q2 "Is child of" does not satisfy reflexivity since a person cannot be their own child. It is also not symmetric or transitive so it is not an equivalence relation

Q3 The proper subset is not a partial order since no set can be a partial order of itself

$$Q4 ((p \rightarrow q) \wedge p) \rightarrow q$$

$$= ((\neg p \vee q) \wedge p) \rightarrow q = (\neg p \wedge p) \vee (q \wedge p) \rightarrow q$$

$$= \neg(q \wedge p) \vee q = \neg q \vee \neg p \vee q = \text{True}$$

Q5 $f(x) = \lceil x \rceil$ is surjective but not injective because there exists at least one $x \in \mathbb{R}$ such that $\lceil x \rceil = z \quad \forall z \in \mathbb{Z}$. Not injective since multiple reals can map to the same integer.

Q6 $\forall s \exists c H(s, c)$ means for all students there exists a computer such that the student s owns the computer c

Q7 The sets $\{b\}$ and \emptyset are lower bounds for X since it is less than or equal to every element in X according to the partial order.

Q8 The empty set is a member of $\{\emptyset\}$.
The empty set is a subset of all sets.
1 and 2 both true.

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- Q9 Assume the opposite of what he are trying to prove:
"There exist some integers m and n such that mn is even but m is odd and n is odd" is the negation
Since $\neg \forall$ is \exists , or gets negated to and, and use the opposite of the predicates.
- Q10 $f(x) = x - 42$ is bijective since there is only output per input and covers the entire codomain
- Q11 $(1,4)$ is not part of the cover relation since 1 is not directly covered by 4
- Q12 Anti-symmetric - two sets cannot be proper subsets of each other
Reflexive - no set is a proper subset of itself
Transitive - If $A \subset B$, $B \subset C$ then $A \subset C$
Asymmetric - $A \subset B$ means $B \subset A$ is impossible
- Q13 $\{a\}$ is not in the power set: $\{\{a\}\}$ is
- Q14 \mathbb{N} is the set of integers ≤ 2 since k starts at 1 and even
 A_k includes integers less than or equal to $2k$
- Q15 This is modus Tollens since Tim having the password implies he can log in. Since he cannot log in, he has no password.
- Q16 False since there is no single positive integer such that it is equal to all positive integers

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Q17 \mathbb{N}_0 and \mathbb{Z} have the same cardinality, so there's a bijective function between them. Real numbers between any interval is an uncountable infinity

Q18 Negate both the subject and predicate, so it becomes "There exist some integers m and n such that mn is even but m is odd and n is odd."

Q19 False. Counterexample: $\lfloor 1.4 + 0.4 \rfloor \neq \lfloor 1.4 \rfloor + \lfloor 0.4 \rfloor$
 $\lfloor 1.8 \rfloor \neq 1 + 0$

Q20 Negate the hypothesis: the equation has an integer solution means there exist integers m_0, n_0 where $42m_0 + 70n_0 = 1000$

Q21 The equivalence $A \leftrightarrow B$ means A is true iff B is true and B is true iff A is true. Symbolically it means $A \rightarrow B$ and $B \rightarrow A$, so A and B must have the same value for $A \leftrightarrow B$ to hold.

Q22 $\neg(\neg(A \leftrightarrow B)) \vee C$ is equivalent to $(\neg A \leftrightarrow B) \rightarrow C$ because $p \rightarrow q = \neg p \vee q$ and $\neg(A \leftrightarrow B) = (\neg A \leftrightarrow B)$

Q23 Numbers in the equivalence class $[-1]$ divide by 3 to have a remainder of -1 . $31 \div 3$ leaves a remainder of 2 so it is not in the equivalence class

Q24 $\neg(A \rightarrow B) = \neg(\neg A \vee B) = A \wedge \neg B$

Q25 $\neg \exists x \forall y (P(x, y) \rightarrow Q(x, y)) = \forall x \exists y \neg (P(x, y) \rightarrow Q(x, y))$
 $= \forall x \exists y (P(x, y) \wedge \neg Q(x, y))$

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Q26 The natural order \leq on the set of integers is a total order since it is comparable on the entire set

Q27 There exists an open but not closed knight tour on a 3×4 board

| | | | |
|----|----|---|----|
| 8 | 11 | 6 | 3 |
| 1 | 4 | 9 | 12 |
| 10 | 7 | 2 | 5 |

EXAM 2 STARTS HERE

Q1 Taking the logarithm of both functions:

$$\ln (\ln n)^n = n \ln (\ln n)$$

$$\ln (n^{\ln n}) = (\ln n)(\ln n) = (\ln n)^2$$

$$\lim_{n \rightarrow \infty} \frac{n \ln (\ln n)}{(\ln n)^2} < \lim_{n \rightarrow \infty} \frac{n}{\ln n} \quad \text{since } \ln (\ln n) < \ln n$$

$$\lim_{n \rightarrow \infty} \frac{n}{\ln n} = \lim_{n \rightarrow \infty} \frac{1}{\frac{1}{n}} = \infty \quad \text{therefore } (\ln n)^n \notin O(n^{\ln n})$$

Q2 True since $\lim_{n \rightarrow \infty} \frac{11n^2 + 7n \log n}{n^3} = 0$

Q3 True since $(-1)^n + 5$ alternates between 4 and 6 which are the bounds in the definition of Θ , and $n^3 \in \Theta(n^3)$

Q4 True since $\lim_{n \rightarrow \infty} \frac{41n^2 + 13n \log n}{n^2 \log n} = \lim_{n \rightarrow \infty} \frac{41}{\log n} + \frac{13}{n} = 0$

Q5 True since $\lim_{n \rightarrow \infty} \frac{1}{\frac{1}{n}} = \infty$

Q6 True since $\lim_{n \rightarrow \infty} \frac{n \log n}{5n^2 + 4n \log n} = \lim_{n \rightarrow \infty} \frac{n \log n}{5n^2} = \lim_{n \rightarrow \infty} \frac{\log n}{5n} = 0$

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Q7 $g(n)$ is an upper and lower bound for $f(n)$ since both are polynomials of highest degree 2. They are not strict bounds since you can choose constants such that $c|g(n)| \leq f(n)$ and $|f(n)| \leq c|g(n)|$. Since $g(n)$ is both an upper and lower bound, it is also a strict bound.

Q8 int i starts at 1 and doubles until $i \leq 99$.
 i goes in the order of $1 \rightarrow 2 \rightarrow 4 \rightarrow 8 \rightarrow 16 \rightarrow 32 \rightarrow 64$
so the block is executed 7 times

Q9 Θ estimate for $f(n)$ is $\Theta(\log_2 n)$ since the index i doubles with each iteration which is consistent with a base 2 logarithm. The logarithm operates on the input size.

Q10 The outer loop runs 99 times and the inner loop runs $a+1$ times for each iteration of the outer loop.
That means the block is run $\sum_{a=1}^{99} (a+1) = \frac{99 \times 100}{2} = 5049$ times

Q11 $f(n)$ is in $\Theta(n^2)$ since it is the sum of the first n natural numbers which is $\frac{n(n+1)}{2}$

Q12 The number of multiplications required is given by $(1+(k-1)) + (1+(k-2)) + \dots + 0$ for each term of $p(x)$.
this is the sum of the first n natural numbers which is $\frac{n(n+1)}{2}$ which is $\Theta(n^2)$

Q13 In this algorithm, only one multiplication is done per coefficient, so it is $\Theta(n)$

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Q14 Each of the 10 elements in the domain has two options,
so there are $2^{10} = 1024$ functions

Q15 Let $b_0 = 1$, $b_1 = 2$, $b_2 = 3$ and $b_n = b_{n-1} + b_{n-2} + b_{n-3}$ for all integers $n \geq 3$.
We want to prove using strong induction that $b_n \leq 2^n$.

Base cases: $b_0 = 1 \leq 2^0 = 1$ $b_1 = 2 \leq 2^1 = 2$ $b_2 = 3 \leq 2^2 = 4$

Now assume $b_k \leq 2^k$ holds for all k $0 \leq k < n$, $n \geq 3$

$$\begin{aligned} b_n &= b_{n-1} + b_{n-2} + b_{n-3} && \text{by definition} \\ &\leq 2^{n-1} + 2^{n-2} + 2^{n-3} && \text{by induction hypothesis} \\ &= 2^n 2^{-1} + 2^n 2^{-2} + 2^n 2^{-3} \\ &= 2^n (2^{-1} + 2^{-2} + 2^{-3}) \end{aligned}$$

$$\text{Therefore } b_n \leq 2^n (2^{-1} + 2^{-2} + 2^{-3}) < 2^n$$

Q16 We want to prove by induction that the sum of squares of the first n positive integers is given by

$$\sum_{j=1}^n (2j-1)^2 = \frac{1}{3} (4n^3 - n) \text{ for all positive integers } n$$

$$\text{Base case: } \sum_{j=1}^1 (2j-1)^2 = 1^2 = \frac{1}{3} (4 \cdot 1^3 - 1)$$

Now assume $\sum_{j=1}^n (2j-1)^2 = \frac{1}{3} (4n^3 - n)$ holds

$$\text{Then } \sum_{j=1}^{n+1} (2j-1)^2 = \sum_{j=1}^n (2j-1)^2 + (2(n+1)-1)^2 = \sum_{j=1}^n (2j-1)^2 + (2n+1)^2$$

$$= \frac{1}{3} (4n^3 - n) + (2n+1)^2 \quad \text{by the induction hypothesis}$$

$$= \frac{1}{3} (4(n+1)^3 - (n+1))$$

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Q17 Using the pigeonhole principle, we have 240000 pigeonholes and 540000 pigeons, so there will be at least $\lceil \frac{540000}{240000} \rceil = 3$ people in Asyeland who have the same number of hairs on their head.

Q18 The coefficient is given by $\binom{7}{3} = \frac{7!}{3!(7-3)!} = 35$

Q19 The number of bit strings of length 10 that contain exactly one 0 or exactly three 0s is given by

$$\binom{10}{1} + \binom{10}{3} = 10 + 120 = 130 \text{ bit strings}$$

Q20 Binomial theorem states that $(X+Y)^n = \sum_{k=0}^n \binom{n}{k} X^{n-k} Y^k$,
which means $\sum_{k=0}^n \binom{n}{k} 2^{n-k} = (2+1)^n = 3^n$

Q21 For the domain of the set with three elements, each element must be mapped to a distinct element in the codomain, so there are $10 \times 9 \times 8$ options for the three elements in the domain.

Q22 The number of bit strings that start with two ones is 2^6 since 6 digits have two options each.
Bit strings that end with three zeros is 2^5 .
Since now there are 5 free digits.

However, some strings can satisfy both conditions, which is given by 2^3 since 5 digits are predetermined.

Therefore there are $2^6 + 2^5 - 2^3 = 88$ bitstrings that start with two ones or end with three zeros.

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Q23 Using the pigeonhole principle, we want to find n pigeons for 12 pigeonholes such that 5 will be in the same hole.
 $\lceil \frac{n}{12} \rceil = 5$ the smallest value of n is 49, since $\frac{48}{12} = 4$

Q24 To count the number of permutations of ABCDIEF that contain the substring ACD, we can treat the substring ACD as a single letter since it must stay together in that order. There are now three other elements, so the permutations is given by $4! = 24$