

# Natural Language Processing

Lecture 5: Introduction to Syntax and  
Formal Languages.

10/1/2021

COMS W4705  
Yassine Benajiba

# Sentences:

## the good, the bad, and the ugly

- Some good sentences:
  - *the boy likes a girl*
  - *the small girl likes a big girl*
  - *a very small nice boy sees a very nice boy*
- Some bad sentences:
  - *the boy the girl likes*
  - *small boy likes nice girl*
- Ugly word salad: *very like nice the girl boy*

# Syntax as an Interface

- Syntax can be seen as the interface between morphology (structure of words) and semantics.
- Why treat syntax separately from semantics?
- Can judge if a sentence is grammatical or not, even if it doesn't make sense semantically.

*Colorless green ideas sleep furiously.*

*\*Sleep ideas furiously colorless green.*

# Key Concepts of Syntax

- Constituency and Recursion.
- Dependency.
- Grammatical Relations.
- Subcategorization.
- Long-distance dependencies.

# Constituents

- A constituent is a group of words that behave as a single unit (within a hierarchical structure).
- Noun-Phrase examples:
  - [they], [the woman], [three parties from Brooklyn],  
[a high-class spot such as Mindy's], [the horse raced past the barn]
  - Noun phrases can appear before verbs (among other things) and they must be complete:
    - \**from arrive...*
    - \**the is ....*
    - \**spot sat....*

# Constituency Tests

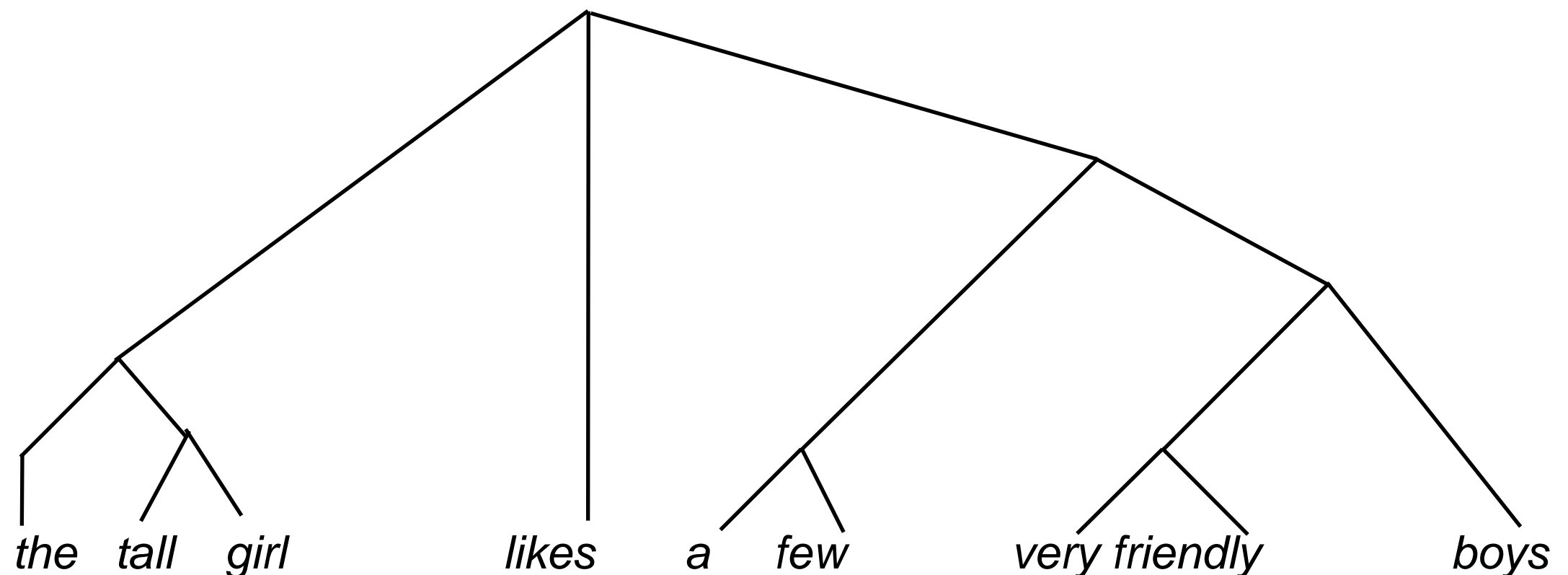
- ***On September seventeenth, I'd like to fly to New York.***
- ***I'd like to fly to New York on September seventeenth.***
- ***I'd like to fly on September seventeenth to New York.***
- ***\*On I'd like to fly to New York September seventeenth.***
- ***\*On September I'd like to fly seventeenth to New York.***

# More Constituency Tests

- There is a great number of constituency tests. They typically involve moving constituents around or replacing them.
- Topicalization:
  - *I won't eat **that pizza**      **That pizza**, I won't eat    \***pizza** I won't eat that*
- Pro-form Substitution:
  - *I don't know **the man who sent flowers**.      I don't know **him**.  
\*I don't know **him** flowers.*
- Wh-question test.
  - ***Where** would you like to fly on September seventeenth?*
  - ***When** would you like to fly to New York?*

# Sentence Structure as Trees

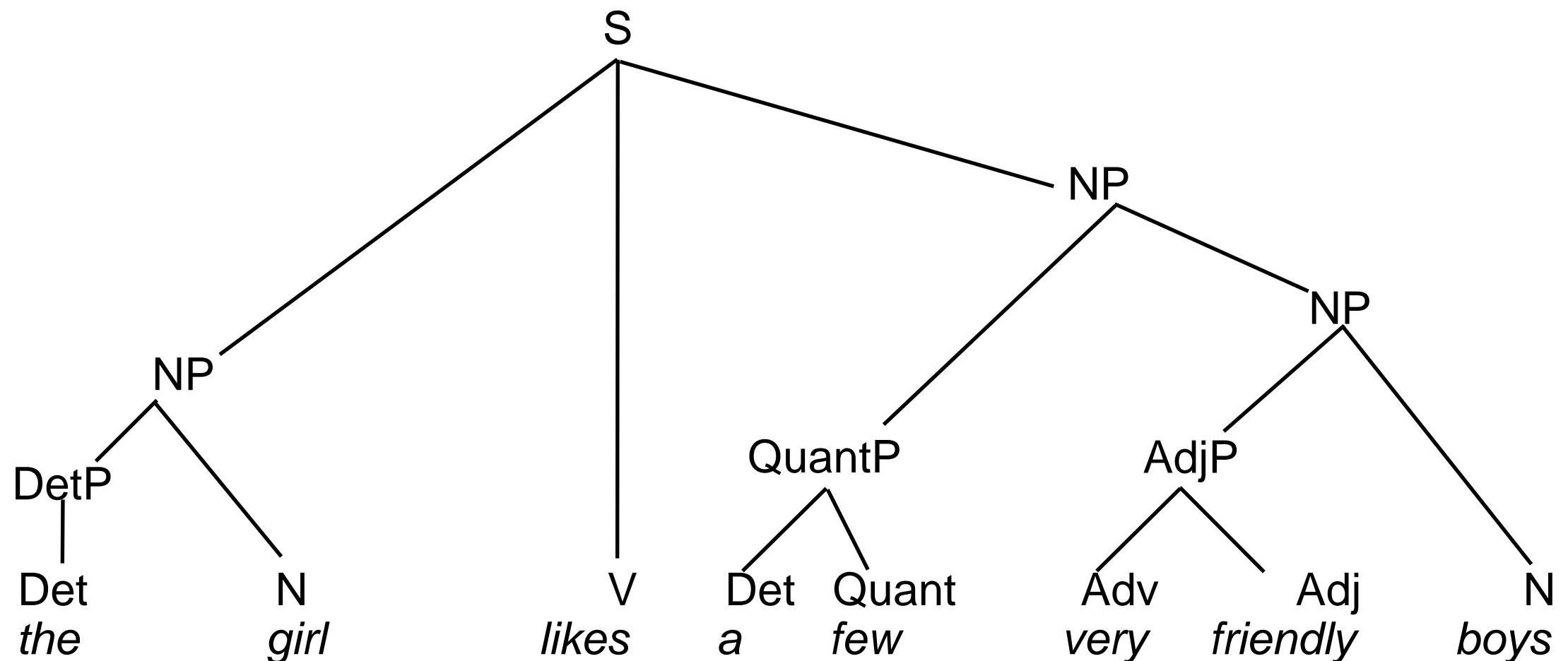
- [the tall girl likes a few very friendly boys]
- [[the tall girl] likes [a few very friendly boys]]
- [[the] tall girl] likes [[a few] [very friendly] boys]]





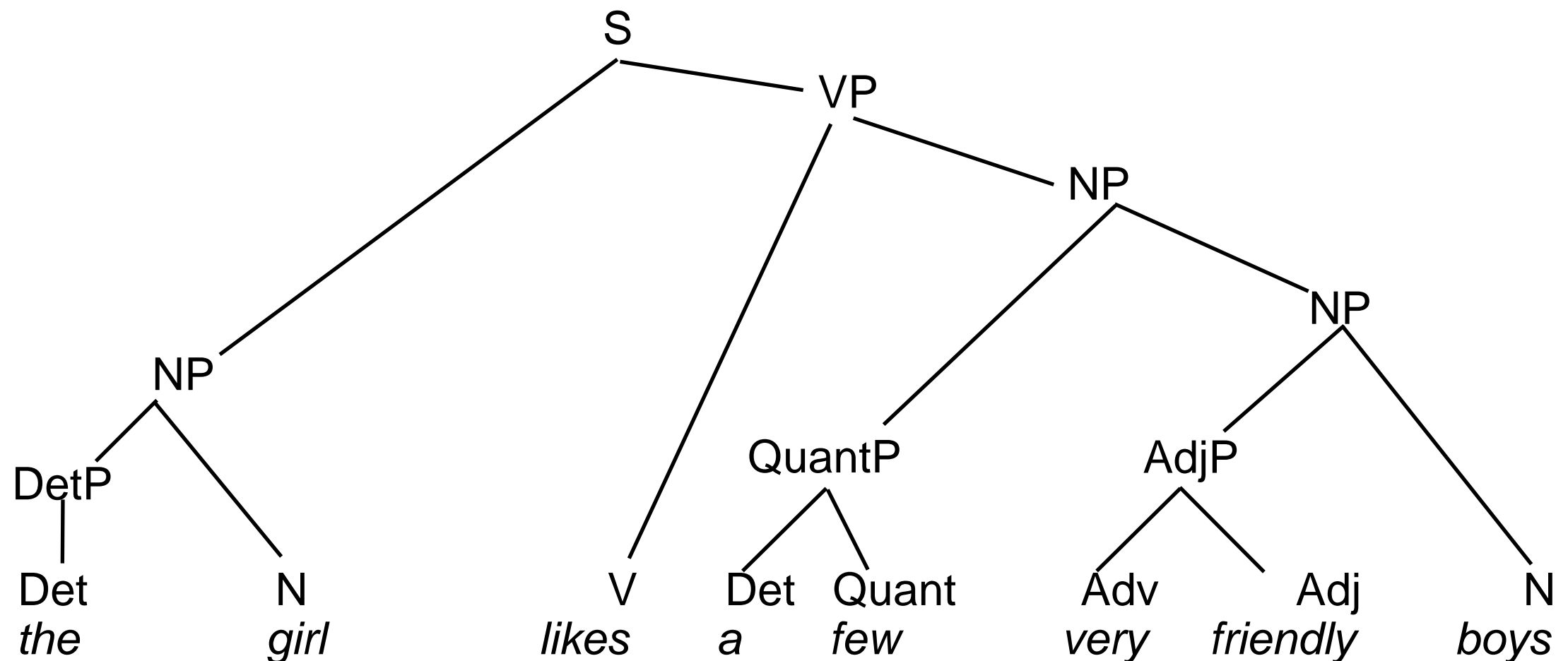
# Constituent Labels

- Choose constituents so each one has one non-bracketed word: the **head**.
- Category of Constituent: XP, where X is the part-of-speech of the head  
NP, VP, AdjP, AdvP, DetP



# Constituent Labels

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# Review: Constituency

*The students easily completed the difficult NLP homework.*

Which constituents can you identify? What tests could you use?

# Recursion in Language

- One of the most important attributes of Natural Languages is that they are **recursive**.
  - *He made pie  
[with apples [from the orchard [near the farm [in ...]]]]*
  - *[The mouse [the cat [the dog chased]] ate] died.*
- There are infinitely many sentences in a language, but in predictable structures.
- How do we model the set of sentences in a language and their structure?

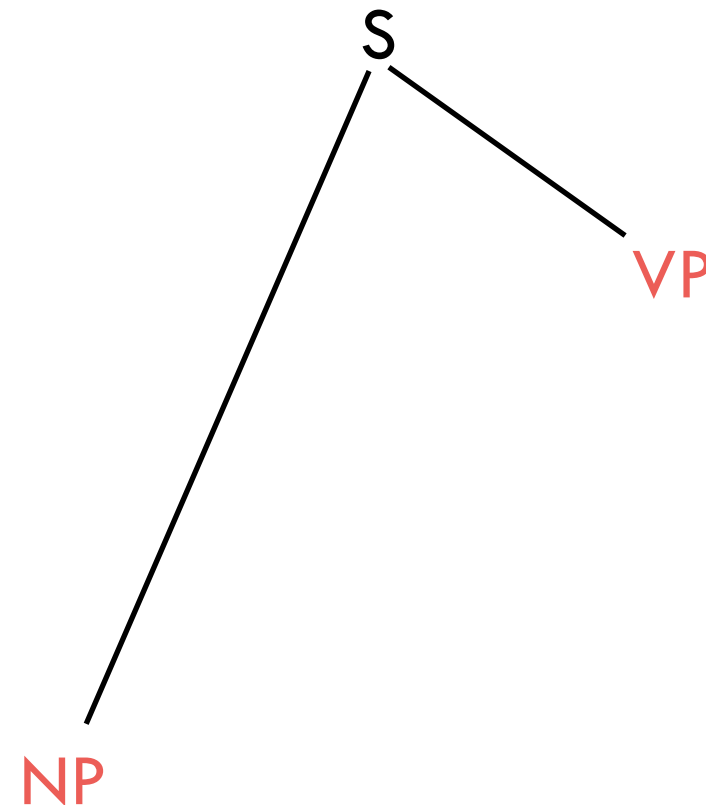
# Context Free Grammars (CFG)

$S \rightarrow NP VP$	$V \rightarrow \text{ saw }$
$VP \rightarrow V NP$	$P \rightarrow \text{ with }$
$VP \rightarrow VP PP$	$D \rightarrow \text{ the }$
$PP \rightarrow P NP$	$N \rightarrow \text{ cat }$
$NP \rightarrow D N$	$N \rightarrow \text{ tail }$
$NP \rightarrow NP PP$	$N \rightarrow \text{ student }$

S

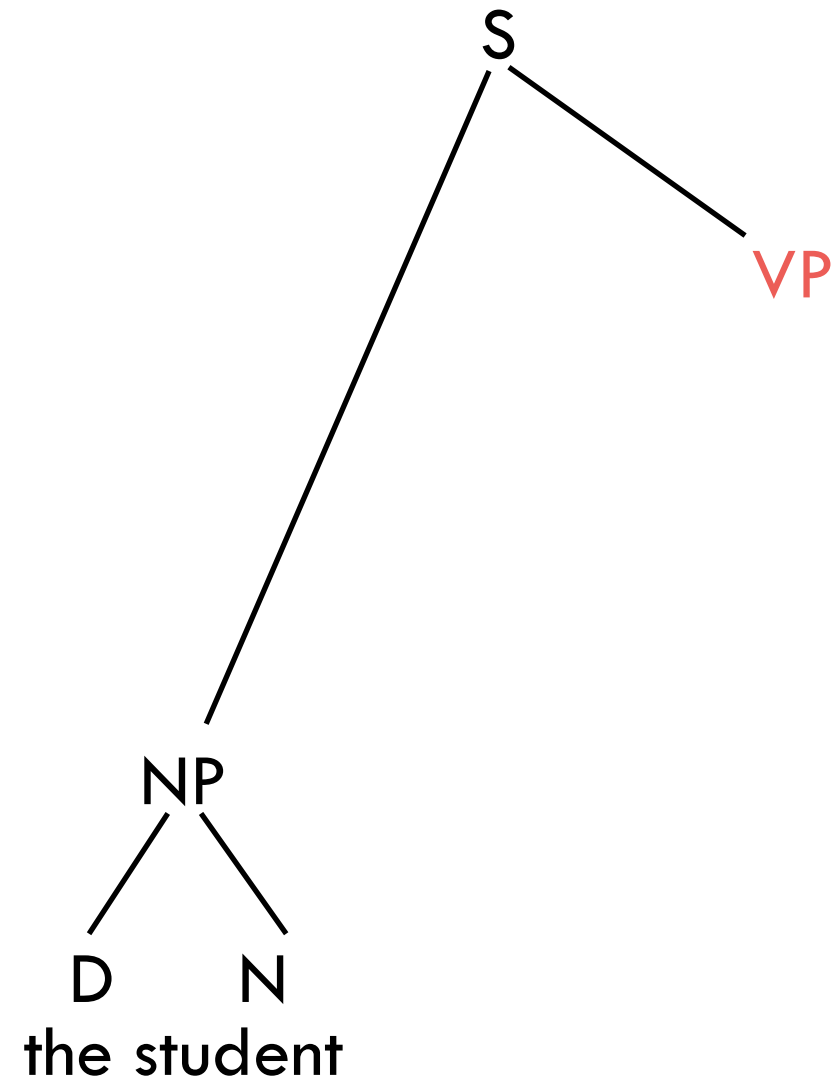
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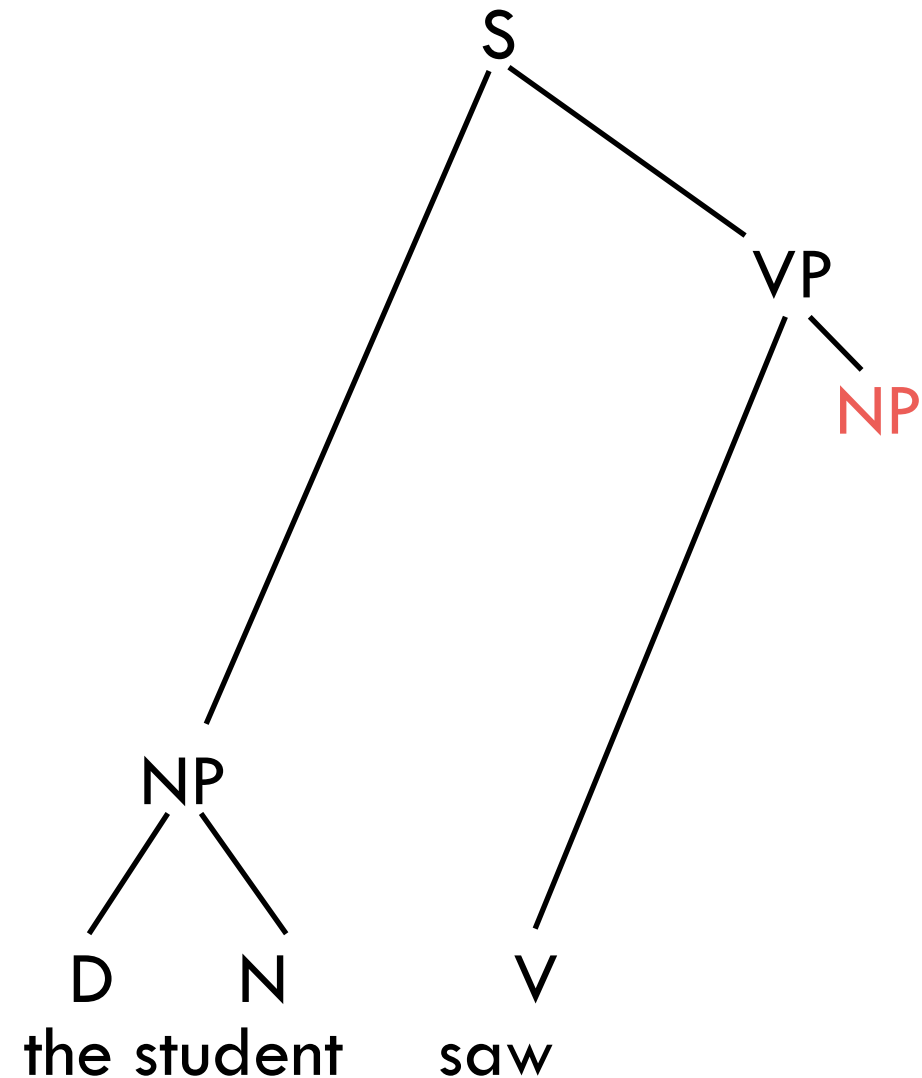
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VP $\rightarrow$ V NP	P $\rightarrow$ with
VP $\rightarrow$ VP PP	D $\rightarrow$ the
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NP $\rightarrow$ NP PP	N $\rightarrow$ student



# Context Free Grammars (CFG)

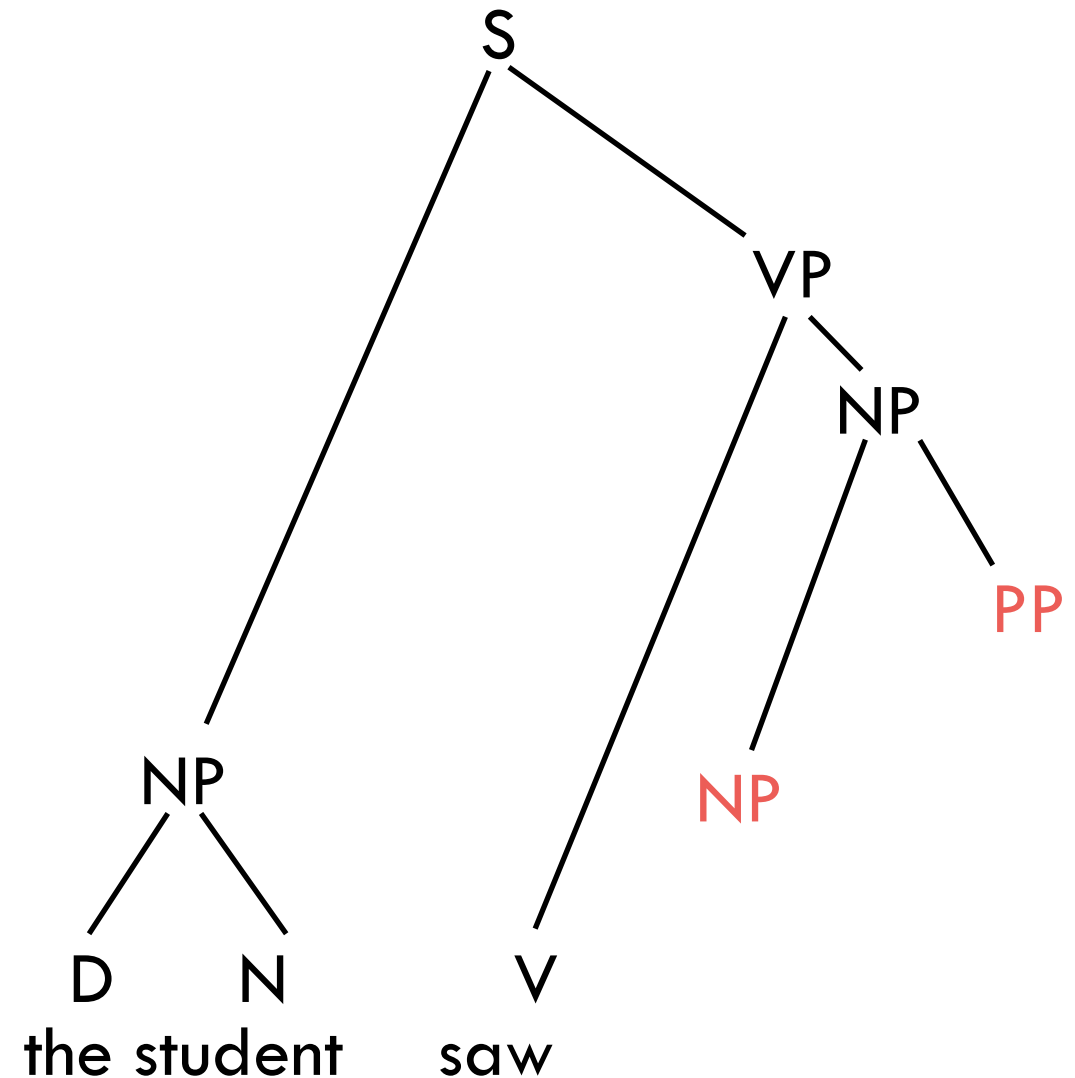
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VP	→	V	NP		P	→	with
VP	→	VP	PP		D	→	the
PP	→	P	NP		N	→	cat
NP	→	D	N		N	→	tail
NP	→	NP	PP		N	→	student





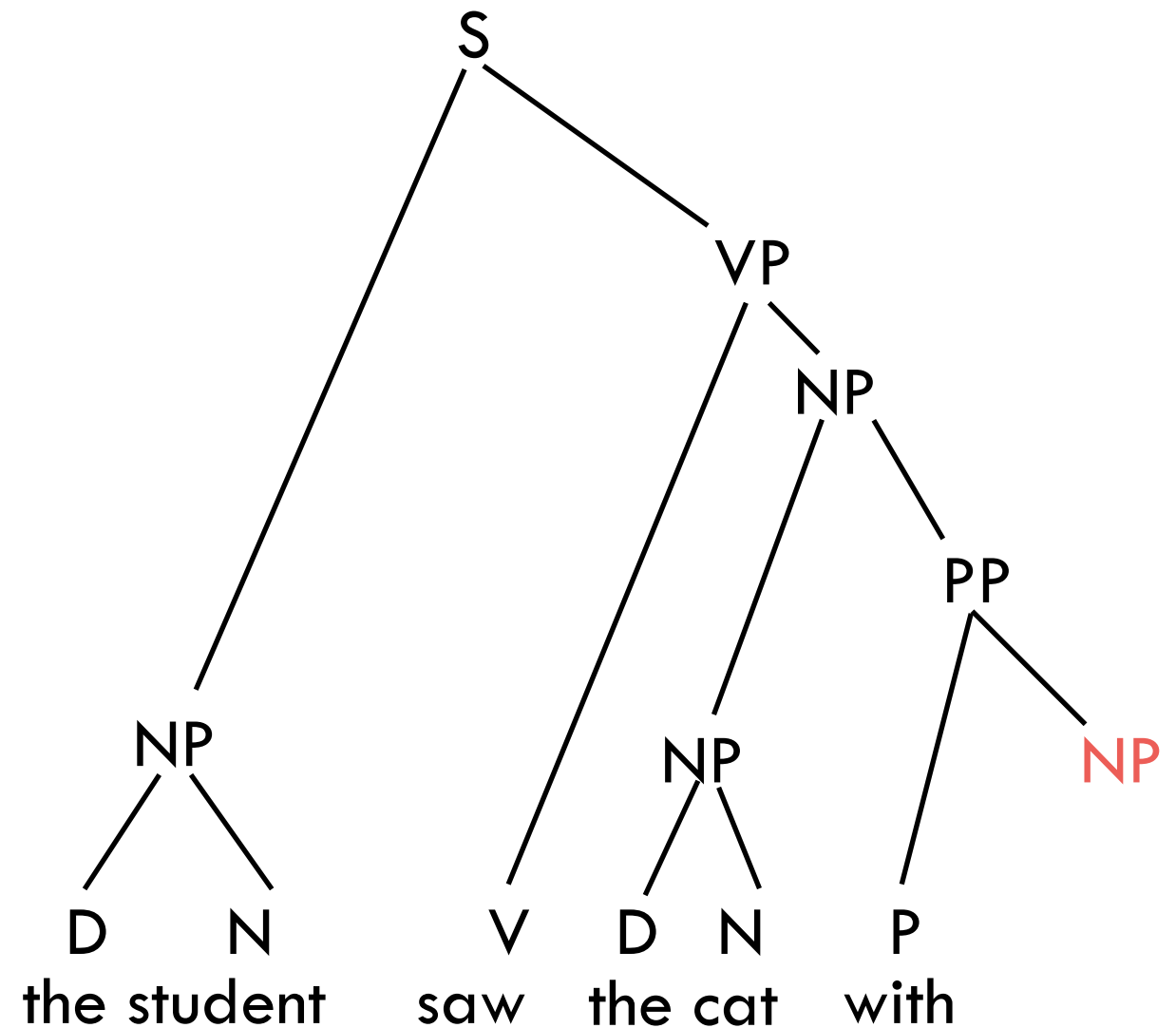
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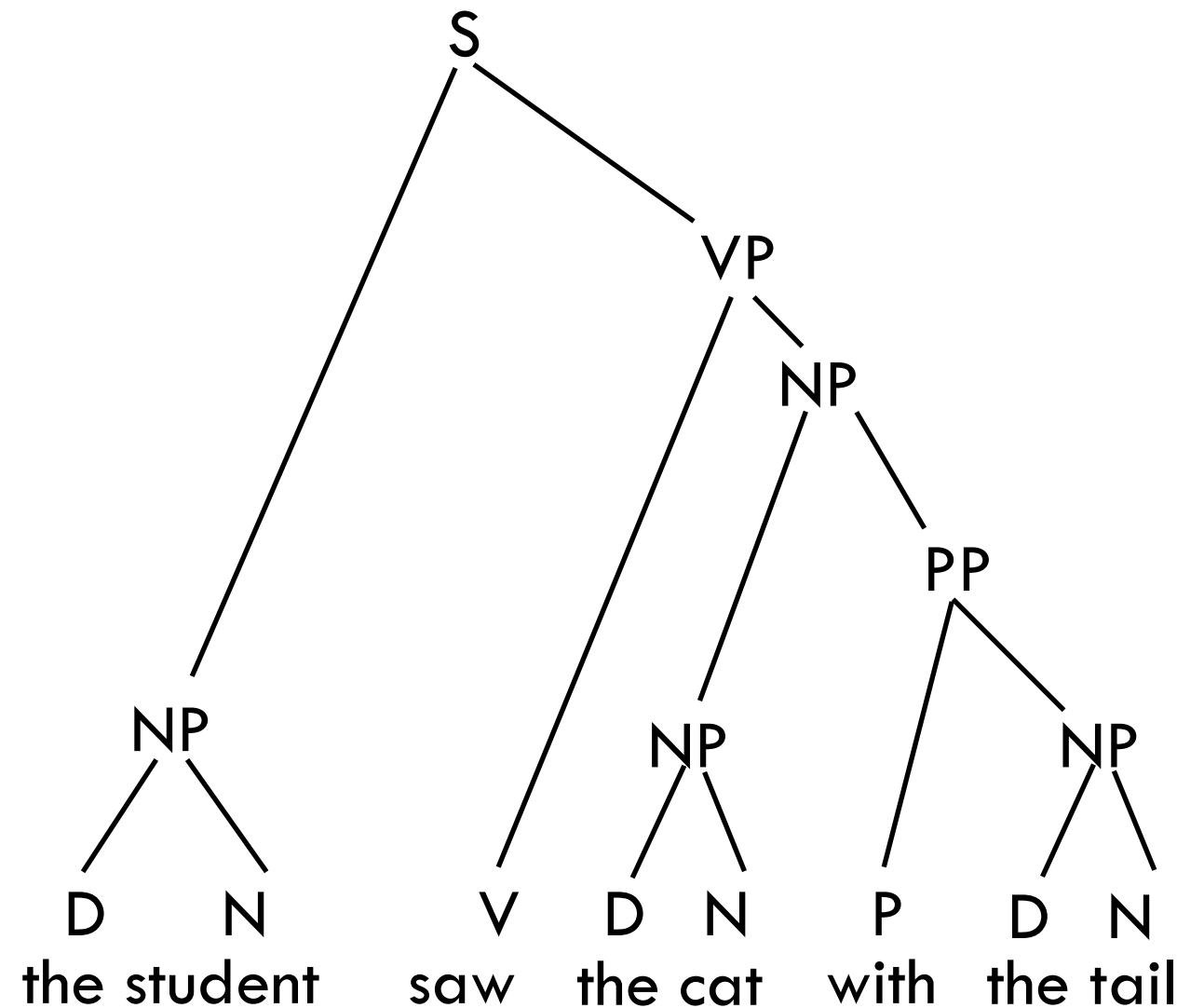
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# Context Free Grammars

- A context free grammar is defined by:
  - Set of **terminal symbols**  $\Sigma$ .
  - Set of **non-terminal symbols**  $N$ .
  - A **start symbol**  $S \in N$ .
  - Set  $R$  of **productions** of the form  $A \rightarrow \beta$ ,  
where  $A \in N$  and  $\beta \in (\Sigma \cup N)^*$ , i.e.  $\beta$  is a string of  
terminals and non-terminals.

# Language of a CFG

- Given a CFG  $G=(N, \Sigma, R, S)$ :
  - Given a string  $\alpha A \gamma$ , where  $A \in N$ , we can derive  $\alpha \beta \gamma$  if there is a production  $A \rightarrow \beta \in R$ .
  - $\alpha \Rightarrow \beta$  means that  $G$  can derive  $\beta$  from  $\alpha$  in a single step.
  - $\alpha \Rightarrow^* \beta$  means that  $G$  can derive  $\beta$  from  $\alpha$  in a finite number of steps.
  - The **language of  $G$**  is defined as the set of all terminal strings that can be derived from the start symbol.

$$L(G) = \{\beta \in \Sigma^*, \text{ s.t. } S \Rightarrow^* \beta\}$$

# Derivations and Derived Strings

- CFG is a string rewriting formalism, so the **derived objects** are strings.
- A derivation is a sequence of rewriting steps.
- CFGs are **context free**: applicability of a rule depends only on the nonterminal symbol, not on its context.
  - Therefore, the order in which multiple non-terminals in a partially derived string are replaced does not matter.  
We can represent identical derivations in a **derivation tree**.
- The derivation tree implies a parse tree.

# Recursion in CFGs

$S \rightarrow NP VP$	$V \rightarrow \text{ saw }$
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**Parse Tree:**

NP

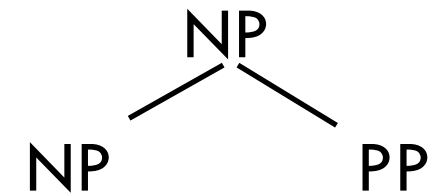
**Derived String:**

NP

# Recursion in CFGs

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Parse Tree:



Derived String:

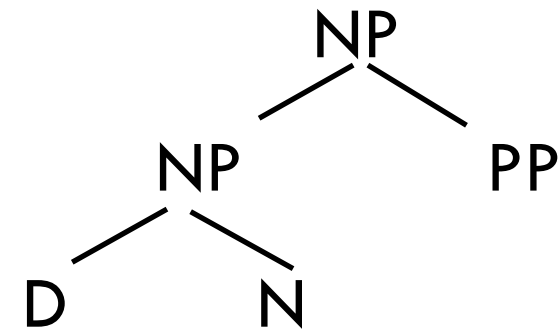
NP PP



# Recursion in CFGs

$S \rightarrow NP VP$	$V \rightarrow \text{say}$
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Parse Tree:



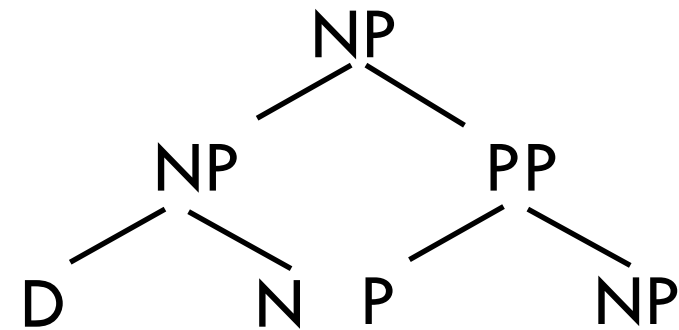
Derived String:

the student PP

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Parse Tree:



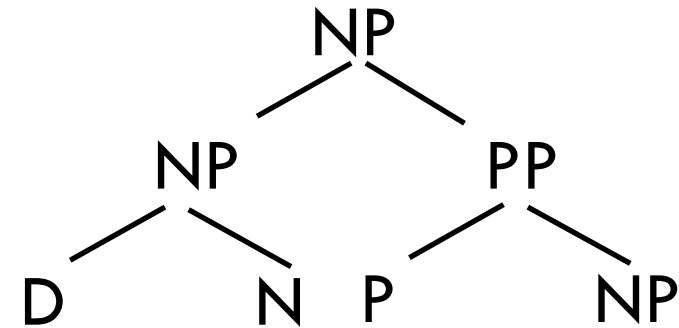
Derived String:

the student P NP

# Recursion in CFGs

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Parse Tree:



Derived String:

the student with NP

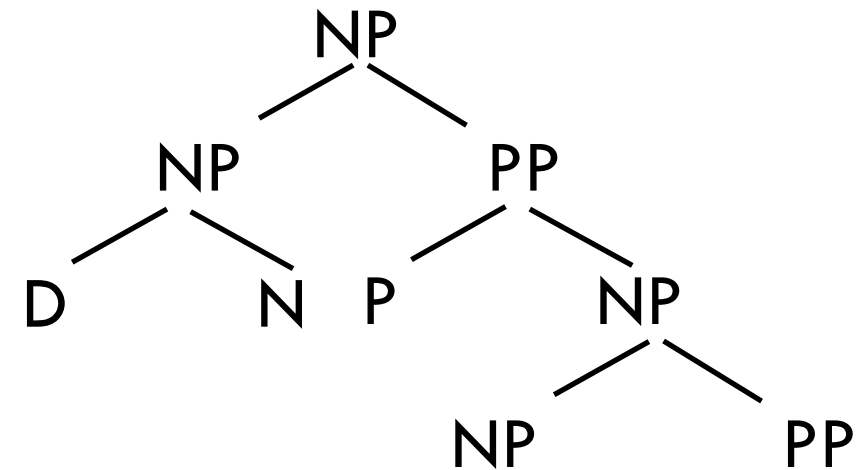
# Recursion in CFGs

S	→	NP VP	V	→	saw
VP	→	V NP	P	→	with
VP	→	VP PP	D	→	the
PP	→	P NP	N	→	cat
NP	→	D N	N	→	tail
NP	→	NP PP	N	→	student

## Derived String:

the student with NP PP

## Parse Tree:



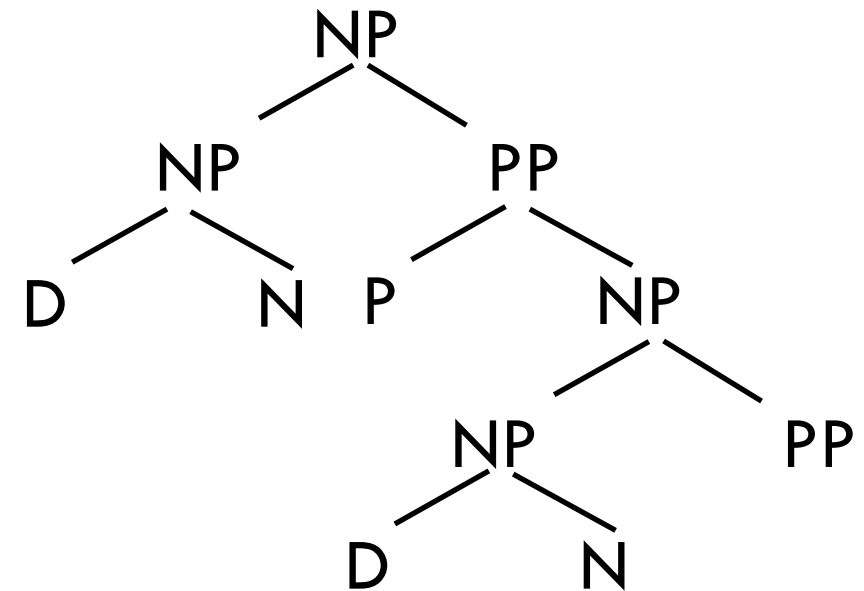
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**Derived String:**

the student with the cat PP

**Parse Tree:**



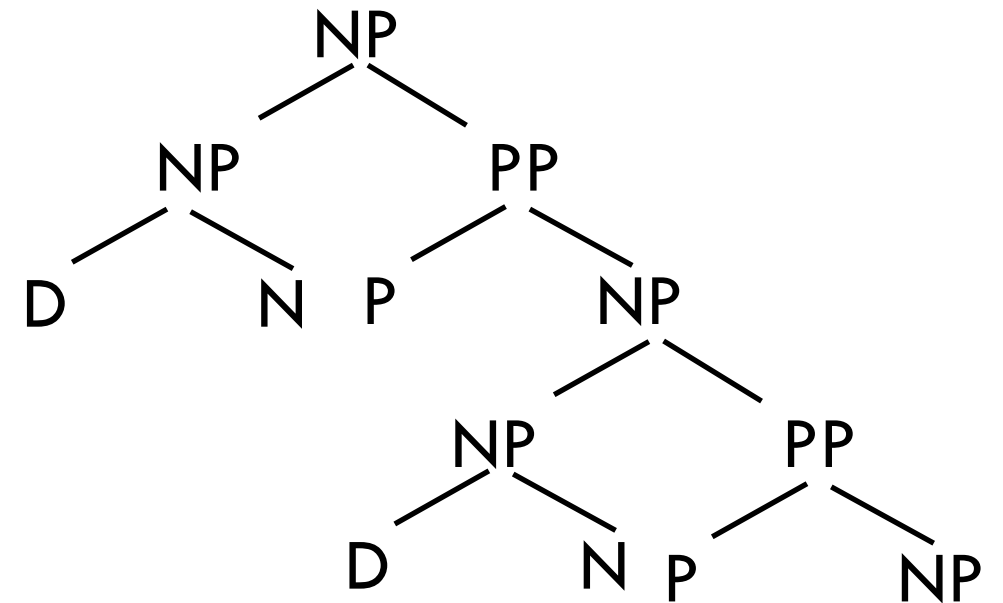
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**Derived String:**

the student with the cat with NP

**Parse Tree:**



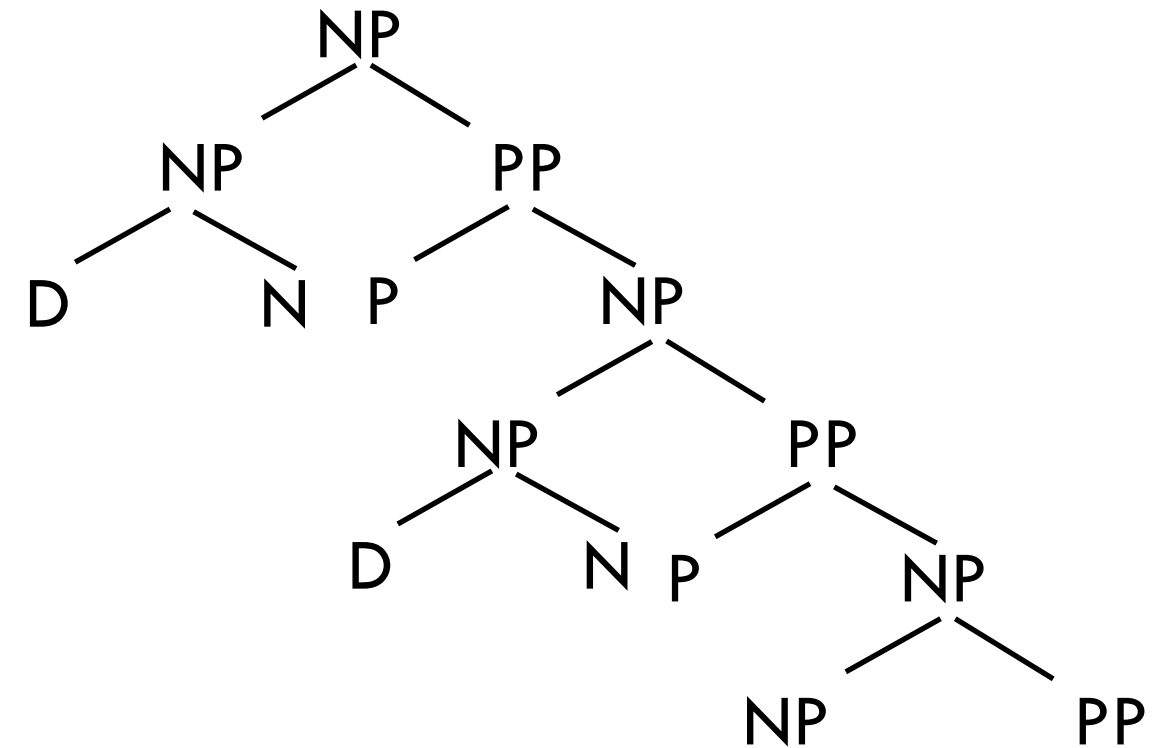
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PP	→	P NP	N	→	cat
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NP	→	NP PP	N	→	student

## Derived String:

the student with the cat with NP PP

## Parse Tree:



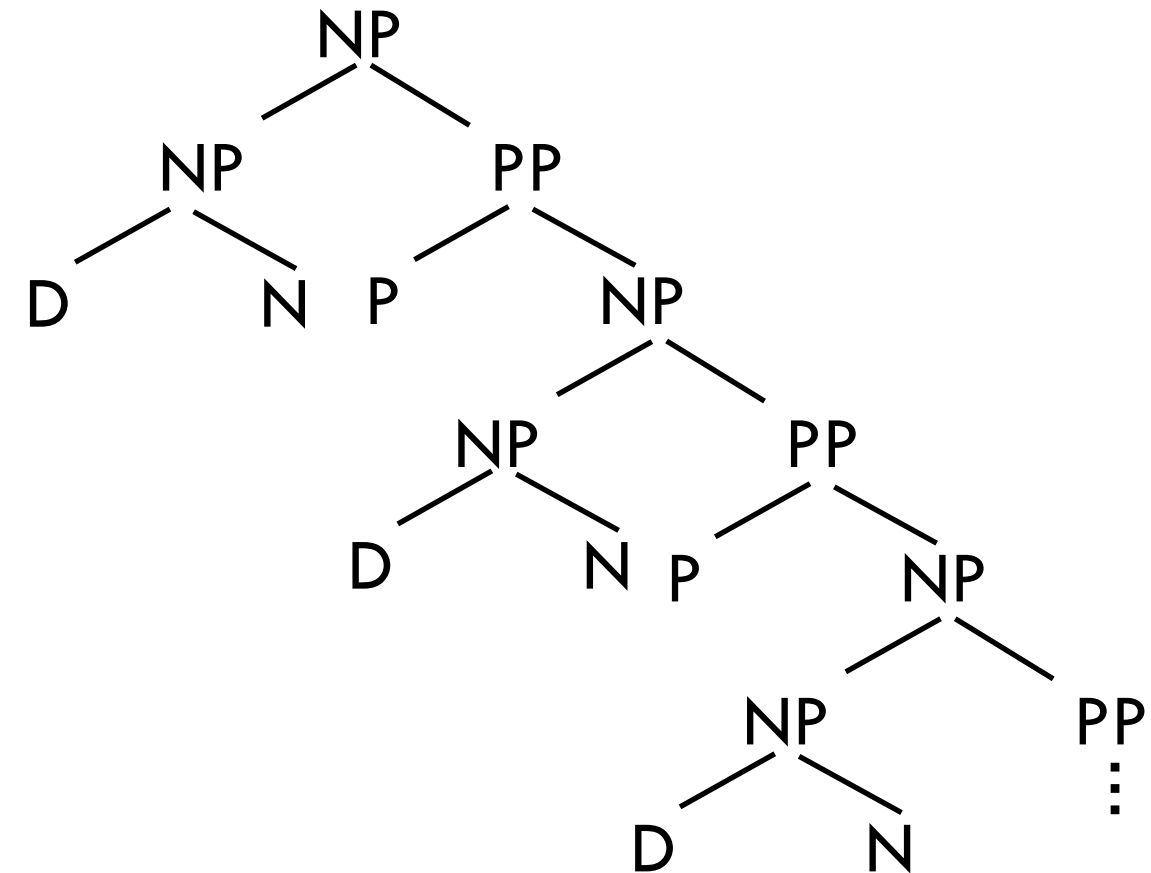
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**Derived String:**

the student with the cat with the tail PP

**Parse Tree:**





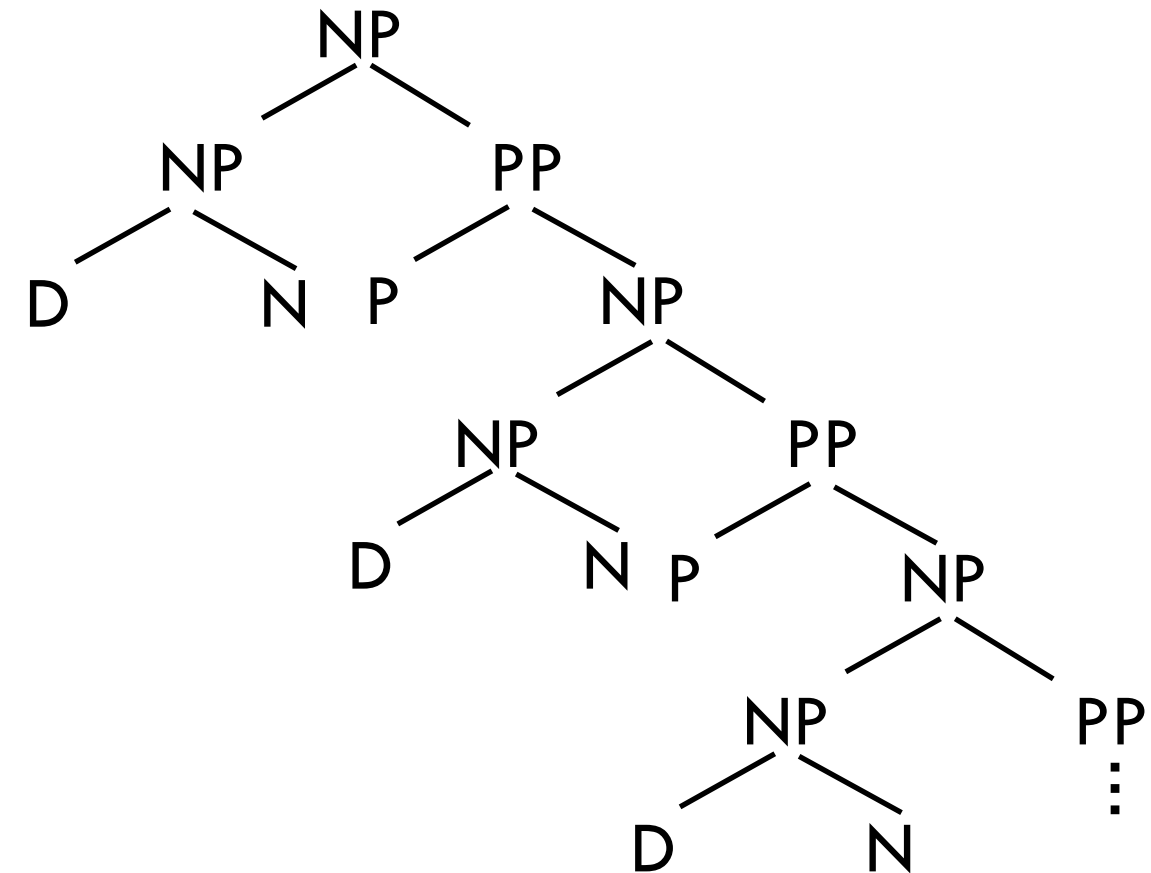
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**Derived String:**

the student with the cat with the tail PP

**Parse Tree:**



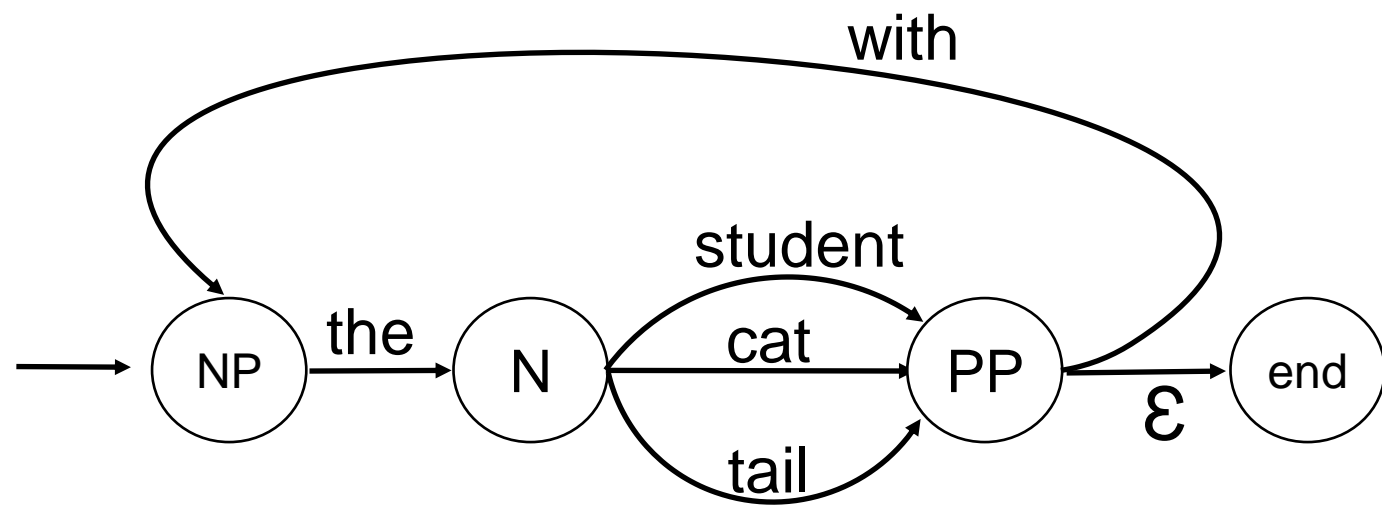
# Regular Grammars

- A regular grammar is defined by:
  - Set of **terminal symbols**  $\Sigma$ .
  - Set of **non-terminal symbols**  $N$ .
  - A **start symbol**  $S \in N$ .
  - Set  $R$  of **productions** of the form  $A \rightarrow aB$ , or  $A \rightarrow a$  where  $A, B \in N$  and  $a \in \Sigma$ .

# Finite State Automata

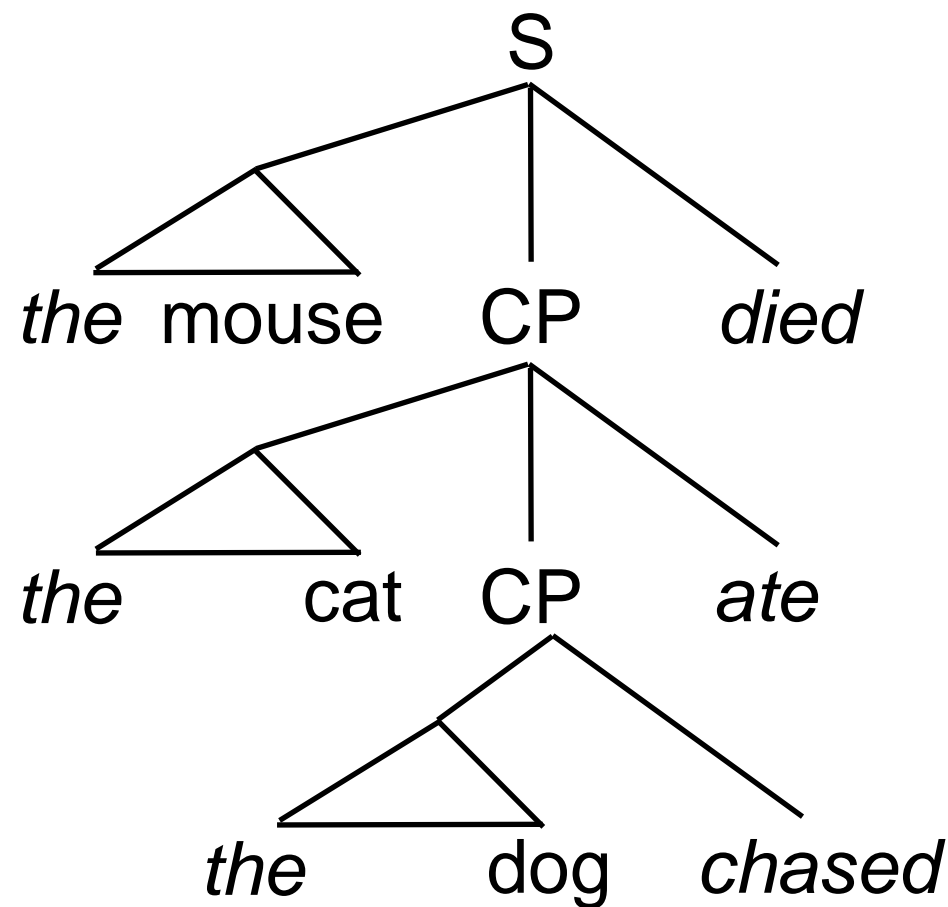
- Regular grammars can be implemented as finite state automata.

NP  $\rightarrow$  the N  
N  $\rightarrow$  student PP  
N  $\rightarrow$  cat PP  
N  $\rightarrow$  tail PP  
PP  $\rightarrow$  with NP  
PP  $\rightarrow \epsilon$



- The set of all regular languages is strictly smaller than the set of context-free languages.

# Center Embeddings

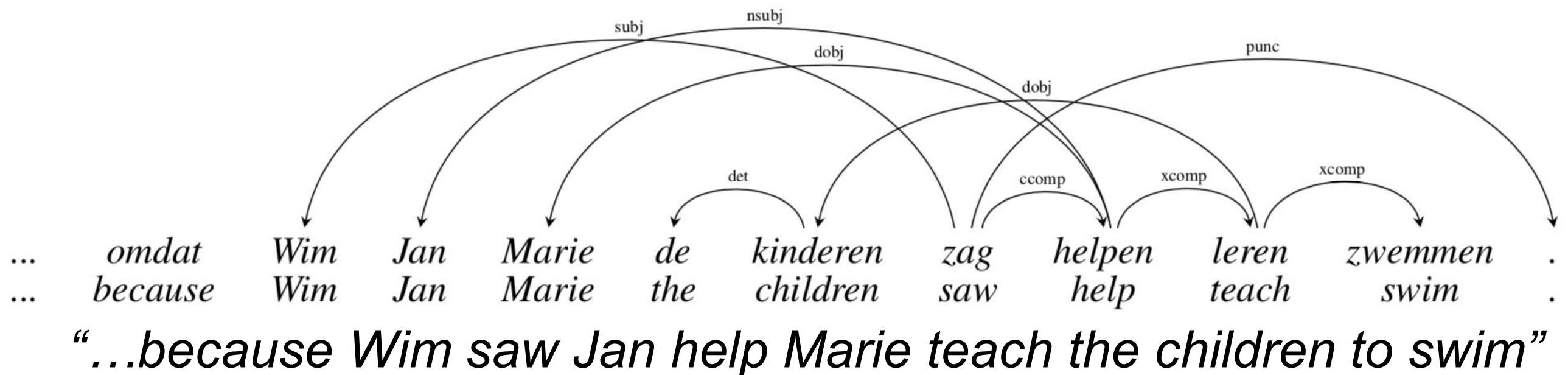


- Problem: Regular grammars cannot capture long-distance dependencies.
- This example follows the pattern  $\mathbf{a^n b^n}$ .  
Can show that is language is not regular  
(using the “pumping lemma”).

Linguistically, this is not a perfect analysis.

# Is Natural Language Context Free?

- Probably not. An example from Dutch:

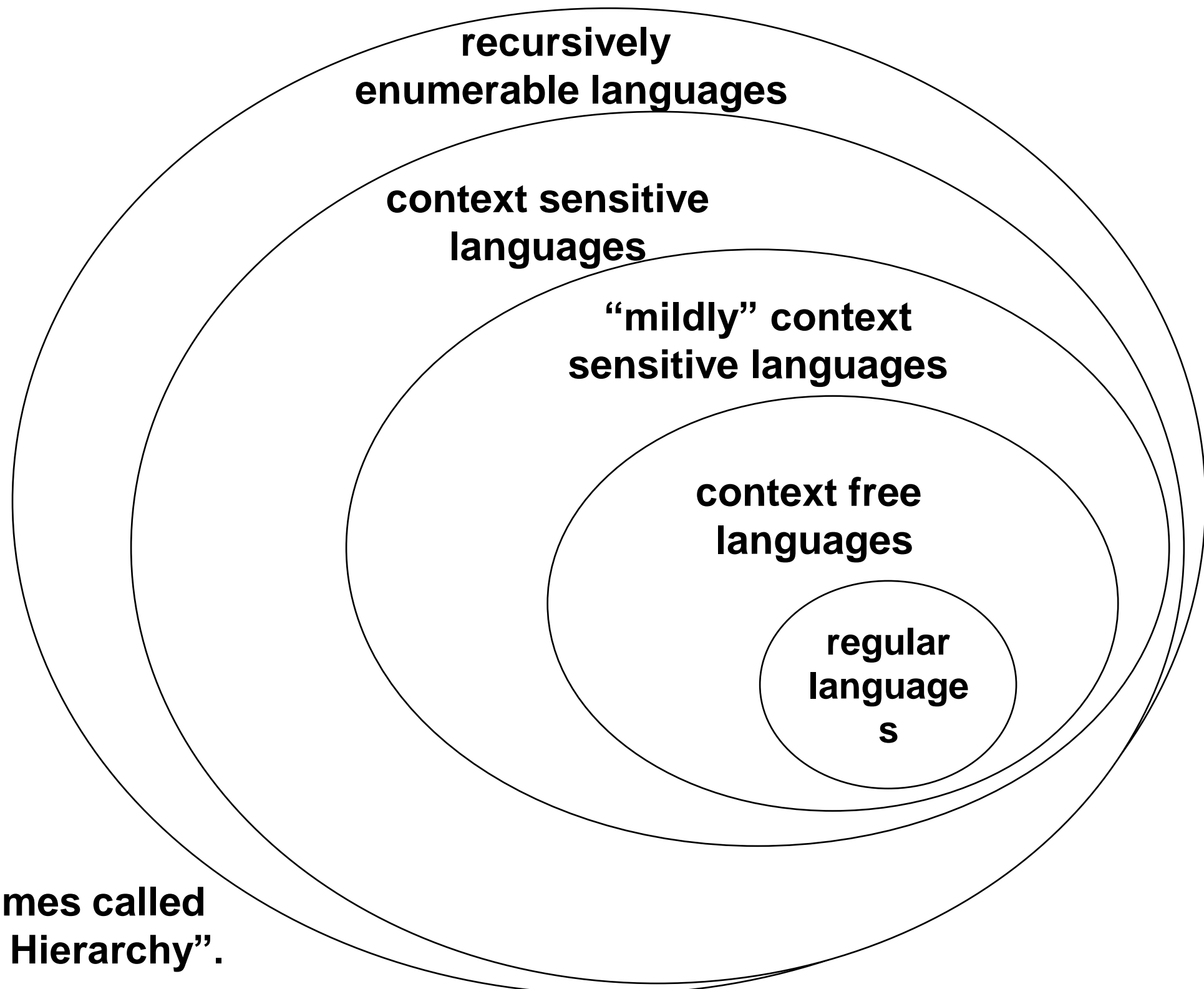


- Context Free Grammars cannot describe crossing dependencies. For example, it can be shown that

$$a^n b^m c^n d^m$$

is not a context free language.

# Complexity Classes



This is sometimes called the “Chomsky Hierarchy”.

# Formal Grammar and Parsing

- Formal Grammars are used in linguistics, NLP, programming languages.
- We want to build a compact model that describes a complete language.
- Need efficient algorithms to determine if a sentence is in the language or not (**recognition problem**).
- We also want to recover the structure imposed by the grammar (**parsing problem**).

# Acknowledgments

- Some slides by Kathy McKeown and Owen Rambow.