

## Tutorial class 2. Set theory & Countability. (corresponding to Lec 3 + 4)

### 1. Set property

During lectures we've proved the following propositions.

**Prop 1**: If  $A \subseteq B$  and  $C \subseteq D$ . then  $A \cap C \subseteq B \cap D$

**Prop 2**:  $(A \cup B) \setminus C = (A \setminus C) \cup (B \setminus C)$ .

Now we give the proof of some others.

**Prop 3**: If  $S_\alpha \subseteq T_\alpha$  for every  $\alpha$ . then

$$\bigcup_{\alpha} S_\alpha \subseteq \bigcup_{\alpha} T_\alpha \quad \text{and} \quad \bigcap_{\alpha} S_\alpha \subseteq \bigcap_{\alpha} T_\alpha$$

Proof:

For the 1st part. consider  $x \in \bigcup_{\alpha} S_\alpha$ , then here  $\exists$  some  $\alpha$  such that  $x \in S_\alpha$ . Since we have  $S_\alpha \subseteq T_\alpha$  for all indice.

we know  $x \in S_\alpha \subseteq T_\alpha \Rightarrow x \in T_\alpha$ . Thus according to the definition.  $x \in \bigcup_{\alpha} T_\alpha$ . which implies  $\bigcup_{\alpha} S_\alpha \subseteq \bigcup_{\alpha} T_\alpha$ .

Similarly. let  $x \in \bigcap_{\alpha} S_\alpha$  (suppose  $\bigcap_{\alpha} S_\alpha$  is nonempty : otherwise the result comes from the trivial fact that  $\emptyset \subseteq \bigcap_{\alpha} T_\alpha$ ). Due to the definition. we know  $x \in S_\alpha$  for all  $\alpha$ ; thus  $x \in T_\alpha$  for all  $\alpha$  (since  $S_\alpha \subseteq T_\alpha$ ), which implies that  $x \in \bigcap_{\alpha} T_\alpha$  and moreover.

$$\bigcap_{\alpha} S_\alpha \subseteq \bigcap_{\alpha} T_\alpha$$

□

**Prop 4.** For  $S \subseteq \mathbb{R}$ ,  $\mathbb{R} \setminus (\mathbb{R} \setminus S) = S$ . If  $A \subseteq B$ . then for every set  $C$ .  $C \setminus B \subseteq C \setminus A$ .

Proof:

Consider the 1st part. We first prove the " $\subseteq$ " relation. For  $b \in R \setminus (R \setminus S)$ , we know  $s \in R$  and  $s \notin R \setminus S$ , which further implies  $s \in S$ . Thus we've reached that  $R \setminus (R \setminus S) \subseteq S$ . For the " $\supseteq$ " relation, suppose  $s \in S$ , then we must have  $s \notin R \setminus S$ ; on the other hand, since  $s \in S \subseteq R$ , according to the definition, we have  $s \in R \setminus (R \setminus S)$ , which ensures  $R \setminus (R \setminus S) \supseteq S$ . Collapse all results above, we attain  $R \setminus (R \setminus S) = S$ .

For the second part, for any  $x \in G \setminus B$ , we have  $x \in C$  and  $x \notin B$ , then actually,  $x \in C$  and  $x \notin A$  (otherwise,  $x \in A \subset B \Rightarrow x \in B$ , a contradiction). Which implies  $x \in G \setminus A$  and  $G \setminus B \subseteq C \setminus A$ , as promised.  $\square$

Prop 5. (Distributive law)

$$A \cap (\bigcup_{\alpha} S_{\alpha}) = \bigcup_{\alpha} (A \cap S_{\alpha}), \quad A \cup (\bigcap_{\alpha} S_{\alpha}) = \bigcap_{\alpha} (A \cup S_{\alpha})$$

Proof: (For your convenience, I write blue for 1st part and red for 2nd part).

Let me show  $A \cap (\bigcup_{\alpha} S_{\alpha}) \subseteq \bigcup_{\alpha} (A \cap S_{\alpha})$  first. Suppose  $x \in V$ , then  $x \in A$  and  $x \in \bigcup_{\alpha} S_{\alpha} \Rightarrow x \in A$  and  $x \in S_{\alpha}$  for some  $\alpha$   
 $\Rightarrow x \in (A \cap S_{\alpha})$  for some  $\alpha$   
 $\Rightarrow x \in \bigcup_{\alpha} (A \cap S_{\alpha})$   
 $\Rightarrow V \subseteq W$ .

For the converse direction, suppose  $x \in \bigcup_{\alpha} (A \cap S_{\alpha})$ , then we

have  $x \in A \cap S_\alpha$  for some  $\alpha$

$$\Rightarrow x \in A \text{ and } x \in S_\alpha \text{ for some } \alpha$$

$$\Rightarrow x \in A \text{ and } x \in \bigcup S_\alpha$$

$$\Rightarrow x \in A \cap (\bigcup S_\alpha)$$

$$\Rightarrow V \supseteq W$$

In conclusion, we have  $V = W \Leftrightarrow A \cap (\bigcup S_\alpha) = \bigcup (A \cap S_\alpha)$ .

Now we comes to the 2nd law. As usual. Let me show

$A \vee (\bigcap S_\alpha) \subseteq \bigcap (A \vee S_\alpha)$ . Suppose  $x \in V$ , then we have

$$\Downarrow \qquad \Downarrow$$

$x \in A$  or  $x \in \bigcap S_\alpha$ . If  $x \in A$ , then  $x \in A \vee S_\alpha$  for all  $\alpha$ .

which promises  $x \in \bigcap (A \vee S_\alpha) \Rightarrow V \subseteq W$ ; if  $x \in \bigcap S_\alpha$ , then  $x \in S_\alpha$  for all  $\alpha \Rightarrow x \in A \vee S_\alpha$  for all  $\alpha$

$$\Rightarrow x \in \bigcap (A \vee S_\alpha) \Rightarrow V \subseteq W.$$

Consider the converse direction. Suppose  $x \in W$ , then

$x \in A \vee S_\alpha$  for all  $\alpha$ . If  $x \in A$ ,  $x \in A \vee (\bigcap S_\alpha) \Rightarrow V \supseteq W$ ;

If  $x \notin A$ , we must have  $x \in S_\alpha$  for all  $\alpha \Rightarrow x \in \bigcap S_\alpha$

$\Rightarrow x \in A \vee (\bigcap S_\alpha) \Rightarrow V \supseteq W$ .

In a word.  $V = W \Leftrightarrow A \vee (\bigcap S_\alpha) = \bigcap (A \vee S_\alpha)$ . □

**Prop 6:** (de Morgan's Law)

$$A \setminus (\bigcup S_\alpha) = \bigcap (A \setminus S_\alpha), \quad A \setminus (\bigcap S_\alpha) = \bigcup (A \setminus S_\alpha)$$

Proof:

Suppose  $x \in A \setminus (\bigcup_{\alpha} S_{\alpha})$ , then we have  $x \in A$  and  $x \notin \bigcup_{\alpha} S_{\alpha}$

$\Rightarrow x \in A$  and  $x \notin S_{\alpha}$  for all  $\alpha$

$\Rightarrow x \in (A \setminus S_{\alpha})$  for all  $\alpha$

$\Rightarrow x \in \bigcap_{\alpha} (A \setminus S_{\alpha}) \Rightarrow A \setminus (\bigcup_{\alpha} S_{\alpha}) \subseteq \bigcap_{\alpha} (A \setminus S_{\alpha})$

On the other hand, suppose  $x \in \bigcap_{\alpha} (A \setminus S_{\alpha})$ , then

$x \in A \setminus S_{\alpha}$  for all  $\alpha \Rightarrow x \in A$  and  $x \notin S_{\alpha}$  for all  $\alpha$

$\Rightarrow x \in A$  and  $x \notin \bigcup_{\alpha} S_{\alpha}$

$\Rightarrow x \in A \setminus (\bigcup_{\alpha} S_{\alpha}) \Rightarrow A \setminus (\bigcup_{\alpha} S_{\alpha}) \supseteq \bigcap_{\alpha} (A \setminus S_{\alpha})$

In conclusion, we've reached that  $A \setminus (\bigcup_{\alpha} S_{\alpha}) = \bigcap_{\alpha} (A \setminus S_{\alpha})$ .

Exercise 1: Try to prove the other law. i.e.

$$A \setminus (\bigcap_{\alpha} S_{\alpha}) = \bigcup_{\alpha} (A \setminus S_{\alpha}).$$

## II. Functions.

We've already introduced the definition of a function.  
Now list some fundamental types of functions.

- ① Identity function.
- ② Composite function.
- ③ Restriction function.
- ④ Surjective function.
- ⑤ Injective function.
- ⑥ Inverse function.
- ⑦ Bijective function.

Now we prove a theorem left in our lecture.

- Thm 1**
- ①  $f: A \rightarrow B$  is bijective iff.  $\exists g: B \rightarrow A$  such that  $g \circ f = I_A$  and  $f \circ g = I_B$ .
  - ②  $f: A \rightarrow B$  and  $g: B \rightarrow C$  are two bijections. then  $g \circ f: A \rightarrow C$  is a bijection.
  - ③ Let  $A, B \subseteq \mathbb{R}$  and  $f: A \rightarrow B$  be a function. If  $y \in B$ , the horizontal line  $y=b$  intersects the graph of  $f$  exactly once. then  $f$  is a bijection.

Proof:

- ①  $\Rightarrow$ : Since  $f$  is bijective. for each  $b \in B$ . there

$\exists! a \in A$  such that  $f(a) = b$ . Hence, we can define a function  $f^{-1}: B \rightarrow A$  by

$$f^{-1}: b \mapsto a \quad (f(a) = b)$$

which is well-defined. Now we check that.

1) for  $\forall a \in A$ .  $f^{-1}(f(a)) = f^{-1}(b) = a$   
 $\Rightarrow f^{-1} \circ f = \text{id}_A$

2) for  $\forall b \in B$ .  $f(f^{-1}(b)) = f(a) = b$ .  
 $\Rightarrow f \circ f^{-1} = \text{id}_B$

So we've already constructed the  $g = f^{-1}$  which we need.

$\Leftarrow$ : we check that:

1)  $f$  is injective.

For  $a_1, a_2 \in A$ . if  $f(a_1) = f(a_2)$ , then apply  $g$  at RHS of this Eq. then:

$$a_1 = g(f(a_1)) = g(f(a_2)) = a_2$$

which confirms that  $f$  is injective.

2)  $f$  is surjective.

For each  $b \in B$ . let  $a = g(b) \in A$ . Then we have

$$f(a) = f(g(b)) = b$$

which implies that  $f$  is surjective.

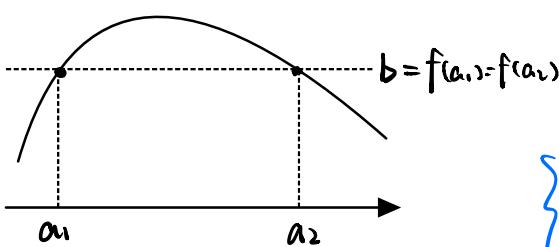
Hence,  $f$  is bijective.

② Exercise 2. Try to check it!

③ As usual, we check that:

i)  $f$  is injective:

Suppose that  $f(a_1) = f(a_2)$  for  $a_1, a_2 \in A$ . Can  $a_1 \neq a_2$  happen? If so, the piece of graph will look like:



Intuition: intersect 2 times?

Formally speaking:

since  $a_1 \neq a_2$  and  $f(a_1) = f(a_2) \stackrel{?}{=} b$ .

$\{(a_1, b), (a_2, b)\} \in G(f)$

(Recall  $G(f) := \{(x, y) : y = f(x)\}$ )

$\{(a_1, b), (a_2, b)\} \cap Y_b \stackrel{?}{=} \{(x, y) : y = b\}$

↳ horizontal line with height b

$\Rightarrow (a_1, b), (a_2, b) \in Y_b \cap f(f)$

but they  
are different!  $\Rightarrow$  contradiction.

$\Rightarrow f$  is surjective.

It's also obvious if you imagine the graph. Formally speaking, for  $b \in B$ ,  $Y_b \cap G(f) \neq \emptyset$ . Select a point  $(x, b) \in Y_b \cap G(f)$ , then we have  $f(x) = b$ ; in other words,  $\exists x \in A$  s.t.  $f(x) = b$ , which confirms that  $f$  is surjective.

Collapse 1). 2). we reach the proof.  $\square$

### III. Equivalent relation.

Def: A relation on a set  $S$  is a subset of  $S \times S$ .

A relation  $R$  on a set  $S$  is an equivalent relation iff

①  $\forall x \in S, (x, x) \in R$  (reflexive)

②  $(x, y) \in R \Rightarrow (y, x) \in R$  (symmetry)

③  $(x, y) \in R, (y, z) \in R \Rightarrow (x, z) \in R$  (transitive)

notation: write  $x \sim y$  when  $(x, y) \in R$ .

$[x] \triangleq \{y \in S : x \sim y\}$   
equivalence class containing (represented by)  $x$ .

Some propositions:

①  $S$  can be decomposed with respect to equivalence, say.

$$S = \bigcup_{x \in S} [x]$$

② The representative element is unambiguous, say

$$x \sim y \Rightarrow [x] = \{z : z \sim x\} = \{z : z \sim y\} = [y]$$

② The . say

$$x \not\sim y \Rightarrow [x] \cap [y] = \emptyset$$

(otherwise,  $\exists z \in [x] \cap [y] \Rightarrow z \sim x, z \sim y \Rightarrow x \sim y$ ).

#### IV. Cardinality.

Based on previous discussions, we know the bijective function can produce an equivalent relation. say :

Sets  $S_1 \sim S_2$  Iff  $\exists$  a bijection  $f: S_1 \rightarrow S_2$ .

Then, we say  $S_1$  and  $S_2$  have the same cardinality if  $(S_1 \sim S_2)$ . Call the number of elements in  $S$  the (cardinal number) of  $S$ , denoted by Card  $S$  or  $|S|$ .

**Remark:**  $|S|$  may be infinite, but still can be compared!

Exercise 3: Does  $\mathbb{Z} \sim \text{Even } \mathbb{Z}$ ?  
all even integers.

Example: We can define a equivalence relation on  $\mathbb{R}$  by setting

$$x \sim y \text{ iff } x-y \in \mathbb{Z}$$

Check that

① Reflexivity

$$\text{For } \forall x \in \mathbb{R}, x-x=0 \in \mathbb{Z} \Leftrightarrow x \sim x$$

② Symmetry

If  $x \sim y$ , we have  $x-y \in \mathbb{Z}$ .

$$\Rightarrow y-x = -z \in \mathbb{Z} \Leftrightarrow y \sim x$$

③ Transitivity

Suppose  $x \sim y, y \sim z$ . then  $\begin{cases} x-y = a \in \mathbb{Z} \\ y-z = b \in \mathbb{Z} \end{cases}$

$$\Rightarrow x-z = (x-y) + (y-z) = a+b \in \mathbb{Z} \Leftrightarrow x \sim z$$

Rank: If we let  $R/\mathbb{Z} \triangleq \{[x] : x \in R\}$  (the collection of all equivalence classes), then  $R/\mathbb{Z} \sim [0, 1)$ !

### Ch3. Countability

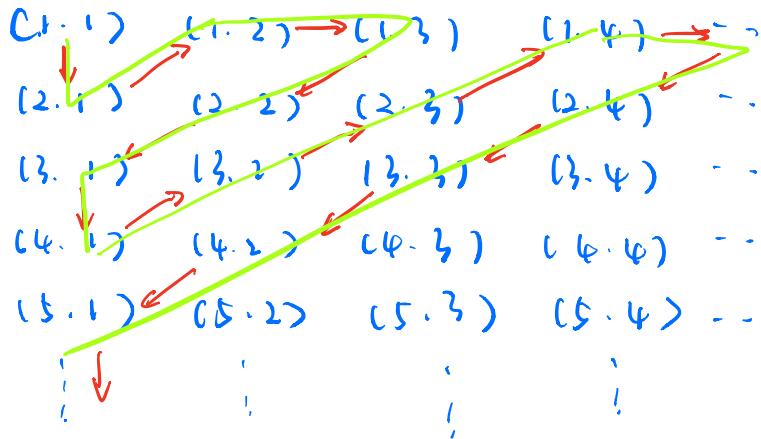
Def: ① A set  $S$  is called countably infinite iff  $S \sim \mathbb{N}$ . (i.e.  $\text{Card } S = \text{Card } \mathbb{N} = \aleph_0$ )

② A set  $S$  is countable iff  $S$  is finite or countably infinite.

**Theorem 2**  $S$  is countable  $\Leftrightarrow S$  can be listed as  
 $S = \{s_1, s_2, \dots\} (s_i \neq s_j \text{ if } i \neq j)$

**Theorem 3** (Bijection theorem) Let  $g: S \rightarrow T$  be a bijection  
 $S$  is countable  $\Leftrightarrow T$  is countable

**Example 1:**  $\mathbb{N} \times \mathbb{N} = \{(m, n) : m, n \in \mathbb{N}\}$  is countably infinite: **diagonal scheme!**



**Example 2:** open interval (e.g.  $(0, 1)$ ) is uncountable.

(idea: construct an element which does not belong to  $(0, 1)$ )

**Exercise 3:** Try to construct a bijection between  $(0, 1)$  and  $[0, 1]$

## V. Countable subset theorem

**Theorem 4:** Let  $A \subseteq B$ . if  $B$  is countable . then  $A$  is countable. (its contrapositive statement is also important)

## VI. Countable union theorem.

**Theorem 5:**  $\bigcup_{n \in \mathbb{N}} A_n$  is countable if  $A_n$  is countable for  $n \geq 1$   
nothing but the diagonal scheme!

## VII. Product theorem.

**Theorem 6:** Any finite product of countable sets is still countable.  
also use the diagonal scheme.

Exercise 4: What happened to the infinite case?

Can you give an uncountable example?

## VIII. Injection theorem.

**Theorem 7** Let  $f: A \rightarrow B$  be injective , if  $B$  is countable.

then  $A$  is countable.

Consequence: (0.1) uncountable  $\Rightarrow A$  uncountable.

Rmk: if  $f: A \rightarrow B$  is injective. we can say  $A$  is embedded in  $B$  by  $f$ .

## IX. Surjection theorem.

**Theorem 8:** Let  $g: A \rightarrow B$  be surjective. If  $A$  is countable. then  $B$  is countable.

Just write  $B = \bigcup_{x \in A} \{f(x)\}$ .

Example:  $\mathbb{Q}$  is countable.

2 ways to show that.

1)  $\mathbb{Q}$  can be seen as a subset of  $\mathbb{Z} \times \mathbb{Z}$

$f: q = \frac{m}{n} \rightarrow (m, n)$  injection

$\Rightarrow \mathbb{Q} = \bigcup_{n \in \mathbb{Z}} S_n = \bigcup_{n \in \mathbb{Z}} \left\{ \frac{n}{m} : m \in \mathbb{Z}, m \neq 0 \right\}$

Obviously  $S_n$  is countable.

**Exercise 5:** Show that (the set of all finite subsets of  $\mathbb{N}$ )  
countable.

Example:  $D = \{x \in \mathbb{R} : x^7 + x^6 + x^5 + \dots + x + 1 \in \mathbb{Q}\}$   
Countable or not?

Seems tricky. but just let:

$$D_r \triangleq \{x \in \mathbb{R} : x^7 + x^6 + \dots + x + 1 = r\} \quad (r \in \mathbb{Q})$$

Then  $D_r$  contains at most 7 elements.

$\Rightarrow D = \bigcup_{r \in \mathbb{Q}} D_r$  is countable (a <sup>countable</sup> union of countable sets is still countable).

Exercise 6. Show that.

$$F \triangleq \{a \in \mathbb{R} : x^5 + ax^3 + 1 = 0 \text{ has a rational root}\}.$$

Is countable or not.