

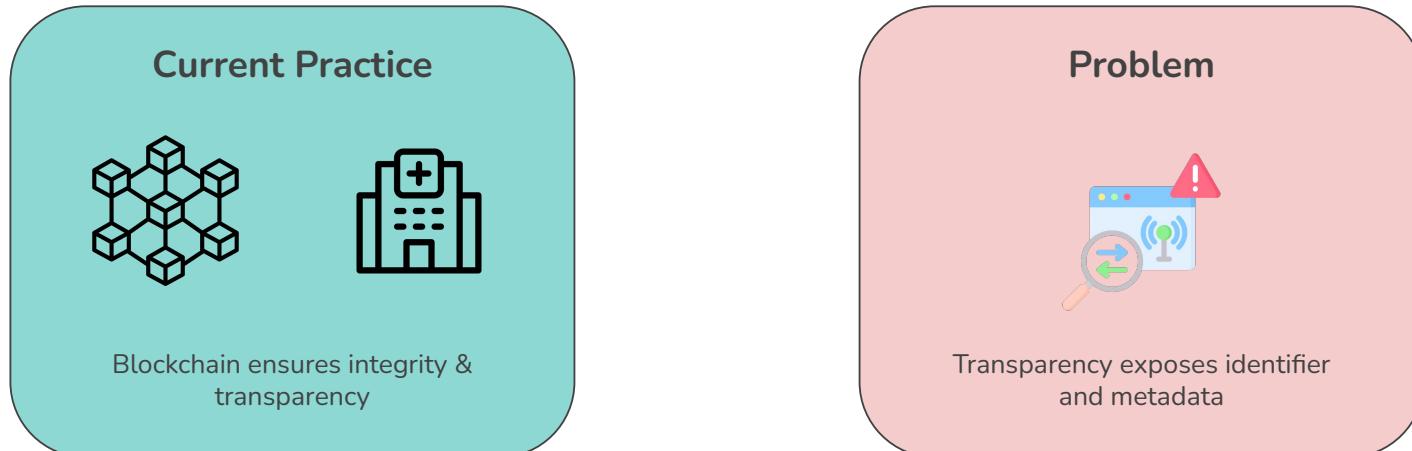
# **Identity-Private Healthcare Data Sharing on Blockchain via Zero-Knowledge Proofs and Account Abstraction**

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# Why Identity Privacy Matters?



Auditability

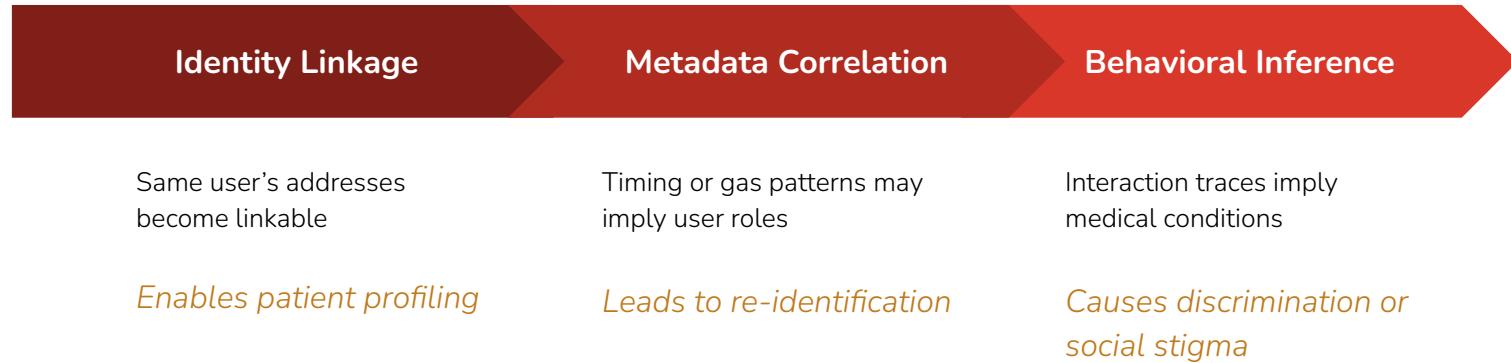


Identity Privacy

We aim to achieve identity privacy without sacrificing auditability



# What Happens When Identities Are Exposed?



These issues undermine trust in blockchain-based healthcare systems



# Research Question

How can blockchain-based healthcare data sharing preserve user identity privacy without compromising transparency and auditability?

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- Blockchain transparency
- Public auditability



- Identity privacy



# Building Blocks for Identity Privacy

## Account Abstraction (AA)

Smart Contract Accounts  
(SCAs) Replace EOAs

Enables programmable identity logic

## Zero-Knowledge Proofs

Prove correctness without revealing secrets

Verify users privately

## Commitments

Commit now, prove later  
Link users and proofs privately

AA defines who can act; ZKPs prove validity; Commitments protect linkage – Together, they enable unlinkable yet verifiable identities



# Privacy Gap in Existing Blockchain-based Healthcare Solutions

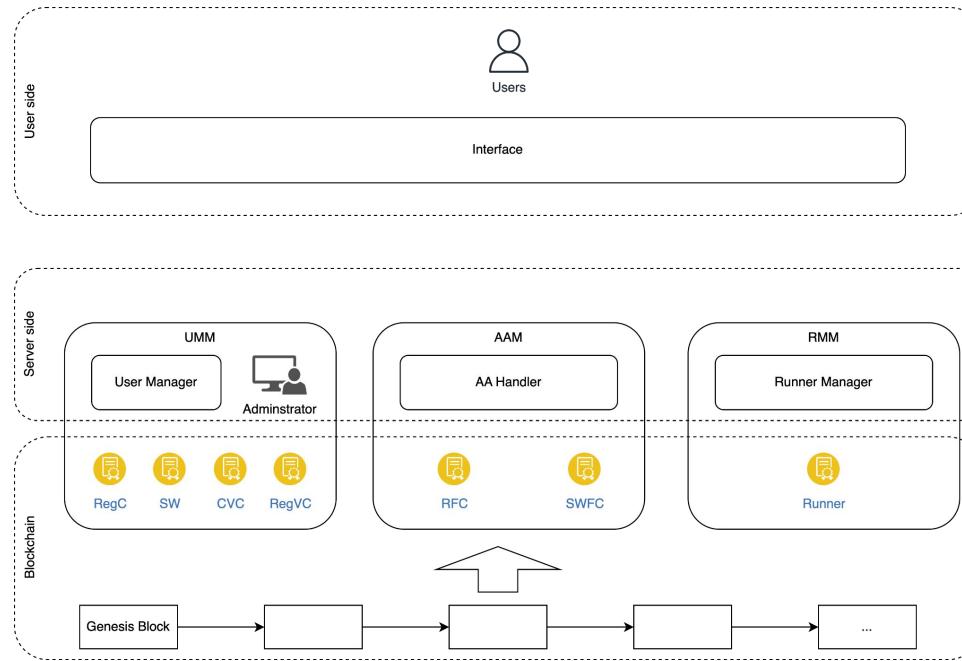
Most blockchain-based healthcare systems secure data content and enforce auditability - but identity privacy and behavioral unlinkability remain unsolved.

	Content Privacy	Identity Privacy	Auditability	Behavioral Unlinkability
CP-ABE/PRE*	✓	✗	✓	✗
ZK-based approaches	✓	✗	✓	✗
EOA-based Systems	✗	✗	✓	✗
Proposed Approach	✓	✓	✓	✓

\*CP-ABE: Ciphertext-Policy Attribute-Based Encryption; PRE: Proxy Re-Encryption

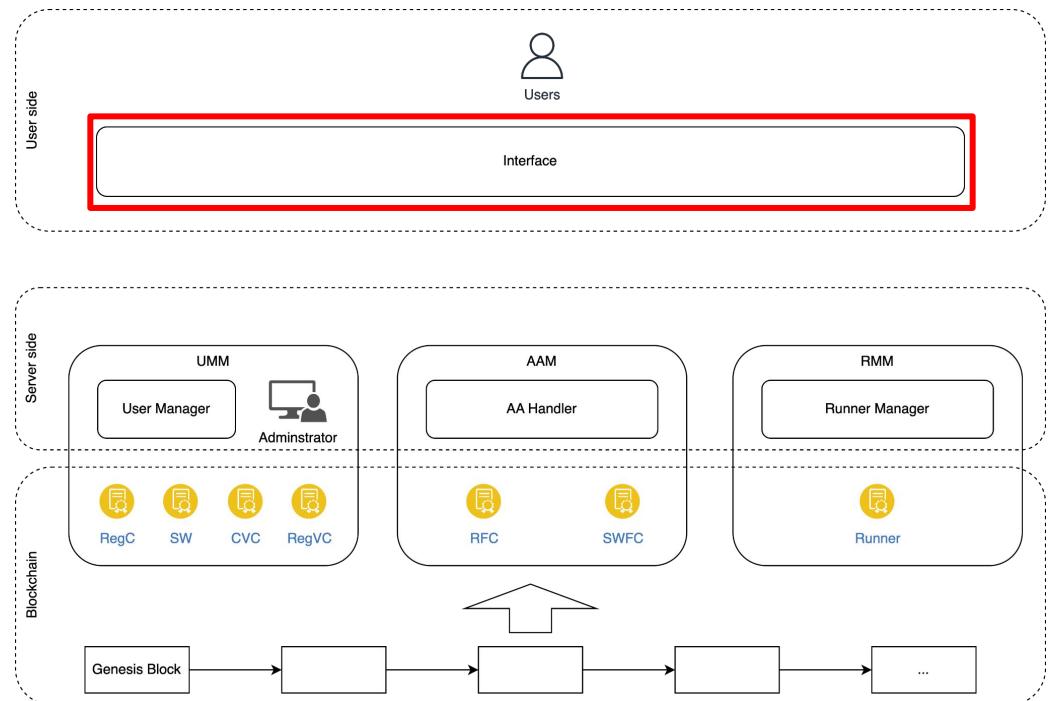
# Proposed Approach Overview

Our approach integrates AA, ZKP and Pedersen Commitments to achieve identity privacy in healthcare data sharing.



# Interface

- **Function:** Handles user input, commitment generation, and ZK proof generation
- **Role:** a bridge between user and blockchain
- **Interaction:** Sends commitments and ZK proofs to User Management Module; submits UserOp to Account Abstraction Module



# User Management Module (UMM)

## Function:

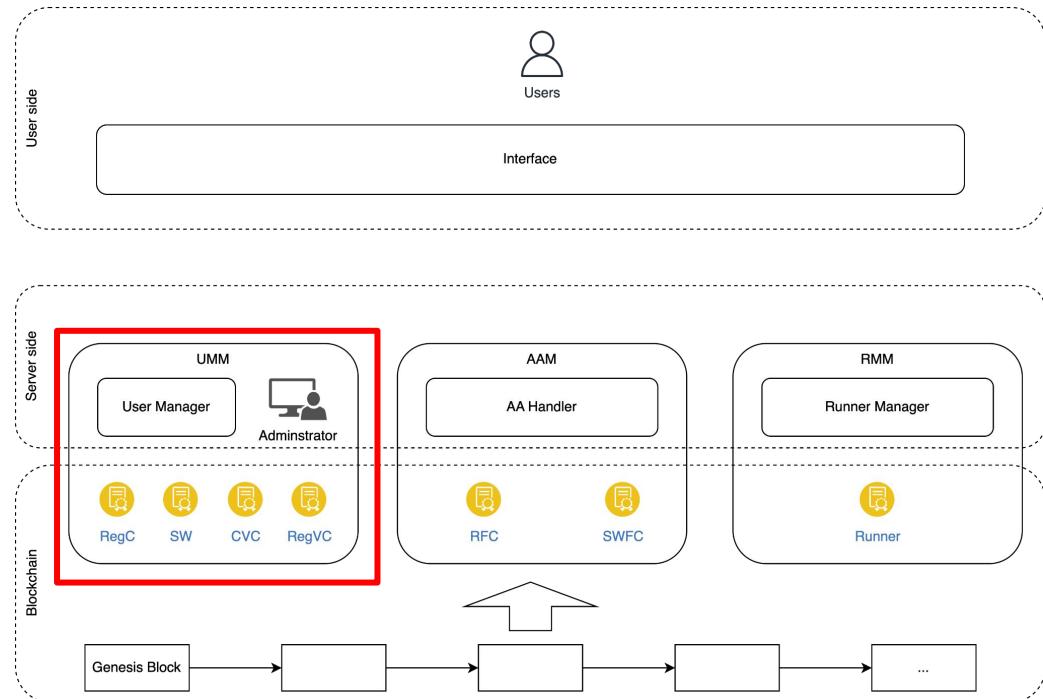
- Manages user registration and proof verification on-chain
- Stores registration commitments and verifies ZK proofs

## Role:

- Serves as the trust anchor in the system - records valid users without exposing identities

## Interaction:

- Receives commitments and ZK proofs from the Interface
- Prepares verified user references to AAM and RMM for execution



# Account Abstraction Module (AAM)

## Function:

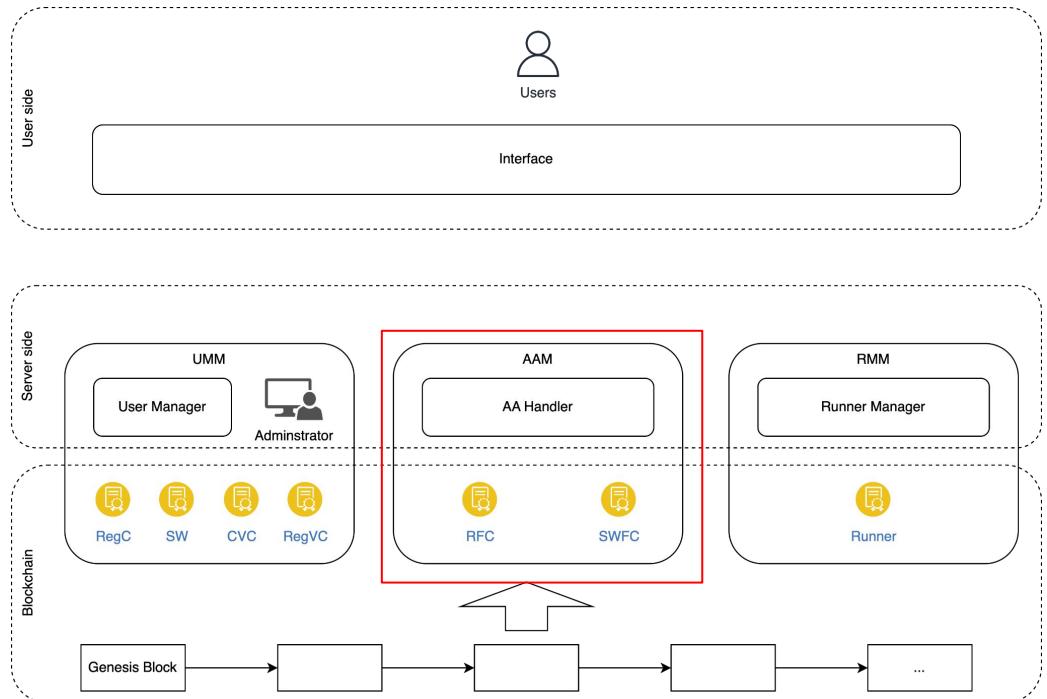
- Create and initializes SCAs and Runner
- Handles UserOp within the AA framework

## Role:

- Acts as the execution layer of the system
- Provides unlinkable identities by decoupling user logic from EOAs

## Interaction:

- Deploys and initializes user SCAs based on commitments
- Receives verified user references from the UMM
- Interacts with the RMM to execute UserOps



# Runner Management Module (RMM)

## Function:

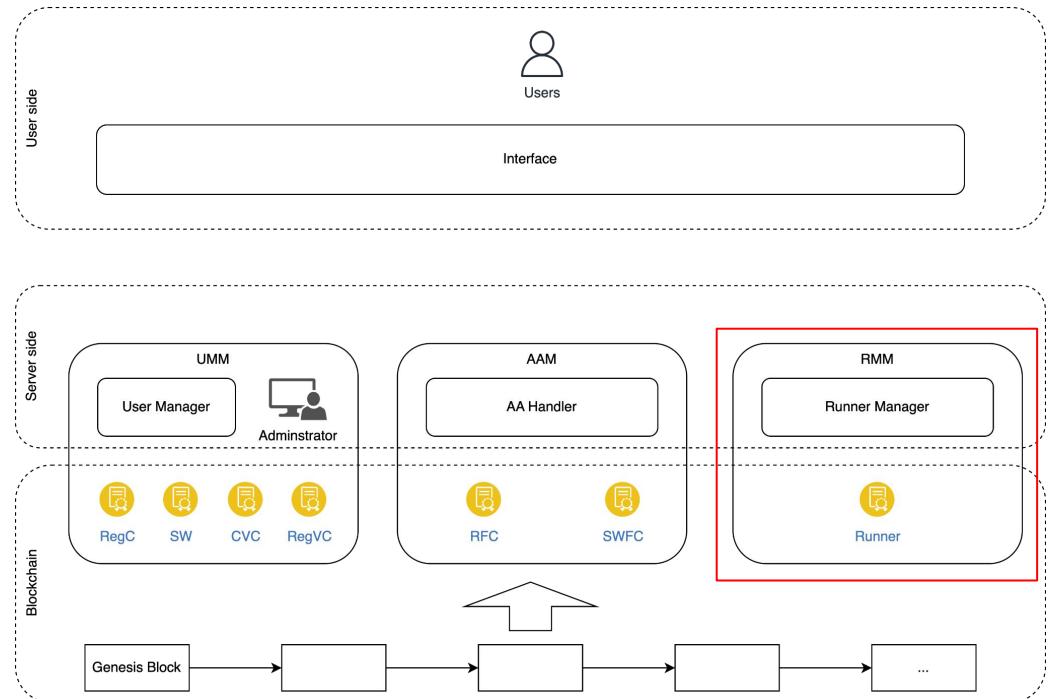
- Manage execution of UserOps via the Runner Contract

## Role:

- Decouples user actions from observable on-chain identities

## Interaction:

- Receives verified user operations from the AAM
- Executes these operations on-chain





# Registration, Verification and Execution

## Registration

User commits  $c_u = \text{Com}(s, r)$  and  $c_r = \text{Com}(\text{Addr}_{SCA}, r)$ ;

Only  $\text{Addr}_{SCA}$  appears on-chain

## Verification

User proves  $c_r$  exists via ZK prof  
(Merkle membership)

No identifier revealed

## Execution

Runner execute User Operation  
(UserOp) via AA framework.

User SCA acts unlinkably

Throughout the process, identity is verified but never revealed

# Registration – Commit Without Revealing Identity

## 1. Commit off-chain

User generates  $c_u = \text{Com}(s, r)$  and create a SCA with  $c_u$

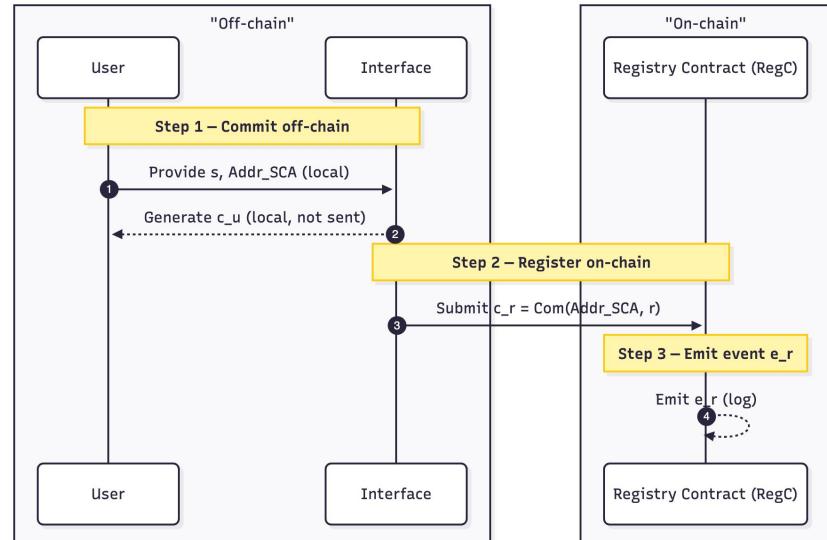
## 2. Register on-chain

User sends  $c_r = \text{Com}(\text{Addr}_{SCA}, r)$  to Registry Contract

## 3. Emit event

Registry confirms  $c_r$  and emits  $e_r$

**Only  $\text{Addr}_{SCA}$  appears on chain**



# Verification – Proof of Membership without Revealing Identity

## 1. Proof Generation

User generates zero-knowledge proof  $p_{mm}$  locally using Merkle path.  $(s, r)$  remain private off-chain.

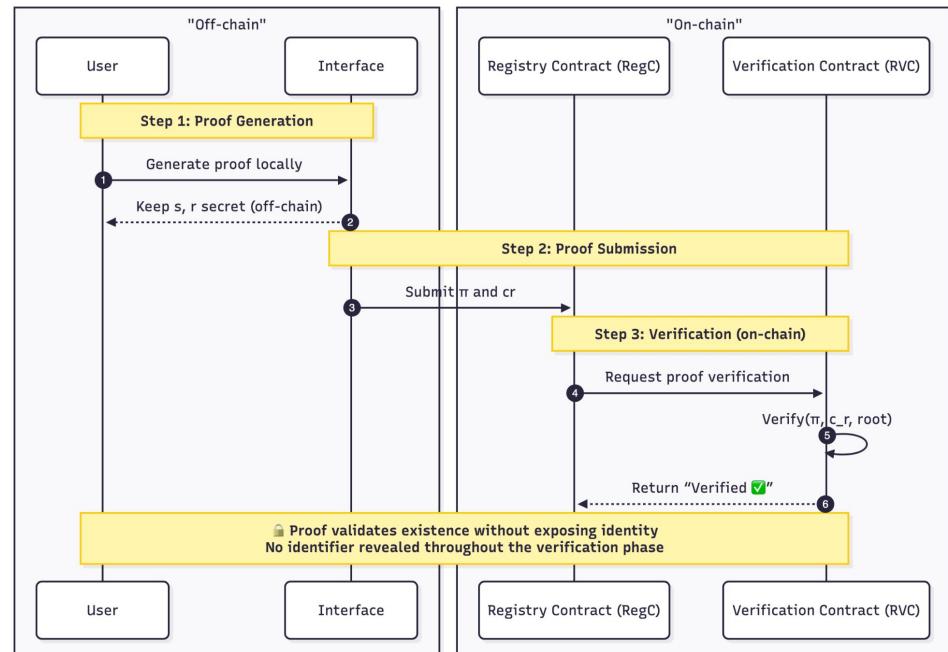
## 2. Proof Submission

User submits  $p_{mm}$  and  $c_r$  to Registry Contract

## 3. Verification

RVC verifies proof correctness and confirms membership

**No identifier or secret is ever revealed.**



# Execution – Unlinkable Operation within AA Framework

## 1. UserOp Preparation

User constructs UserOp and attaches ZK proof  $p_{mm}$ . All secrets  $(s, r)$  stay off-chain.

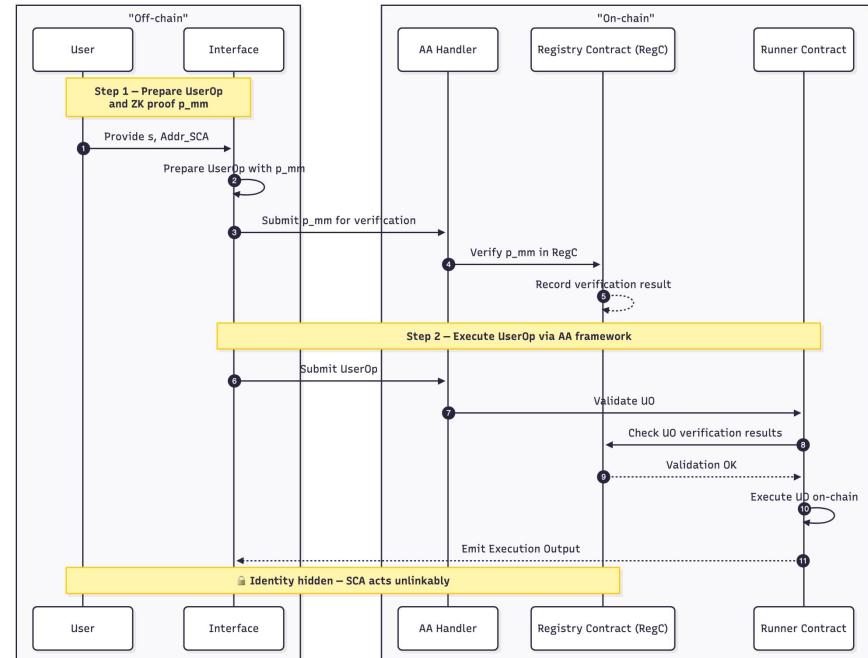
## 2. Verification & Validation

$p_{mm}$  is checked by UMM (RegC) through AA handler. Runner queries verification state.

## 3. Execution

Runner executes UserOp on-chain on behalf of user's SCA. Identity remains unlinkable throughout.

**Identity unobservable**





# Evaluation Scenarios

## Experiment Setup

- All experiments were conducted on the Polygon Amoy Testnet
- Smart contracts include Registry (RegC), Runner (RC), and Verifier (VC, exported),
- Implemented using Solidity 0.8.27, Hardhat 2.26.3, Circom 2.0, and snarkjs 0.7.5.

## Evaluation Scenarios

1. ERC-4337 vs. Proposed Approach  
Attacker knows full (user, SCA) mapping.
2. Account Shuffling vs. Proposed Approach  
Attacker knows one SCA per user and infers others via clustering.

# Attacker Model

We adopt the taxonomy of Wagner & Eckhoff\* and assume a **global, passive, external attacker** with bounded computation.

## Global

- Full visibility of ledger (transactions, state, events etc.)

## Passive

- Only observes; no injection / modification / blocking

## External

- Not registered; holds no secrets; cannot produce ZK witnesses

## Static

- Use fixed inference heuristics (non-adaptive)

## Prior knowledge

### ERC-4337 setting

- Attacker knows  $\langle \text{user}, \text{SCA} \rangle$  mapping
- Knows total number of users

### Account shuffling setting

- Attacker knows one SCA per user & total user count
- Must infer extra addresses via clustering

\*Wagner, Isabel, and David Eckhoff. "Technical privacy metrics: a systematic survey." *ACM Computing Surveys (Csur)* 51.3 (2018): 1-38.



# Evaluation Metrics

01

Entropy

- Measures identity uncertainty – How unpredictable user-address mapping is for the attacker
- **Higher = Better privacy**

02

Anonymity Set Size

- Measures how many users are indistinguishable from each other
- **Higher = Better privacy**

03

Adjusted Rand Index (ARI)

- Measures clustering accuracy between predicted and true user identities
- **Lower = Better privacy**

04

Normalized Mutual Information (NMI)

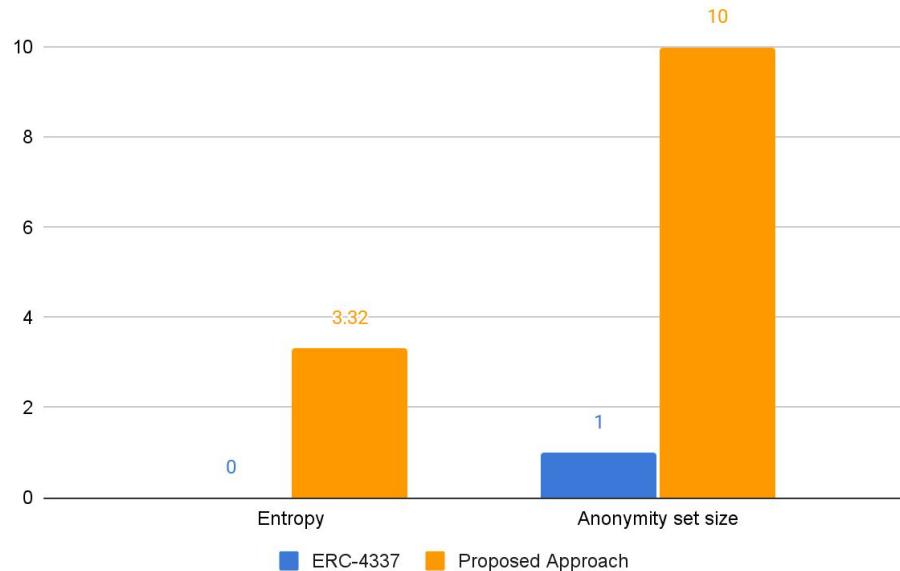
- Measures information overlap between attacker's inference and ground truth
- **Lower = Better privacy**



# Identity Privacy Improvements over ERC-4337

- Higher Entropy: Attacker prediction becomes almost indistinguishable from random guessing
- Larger Anonymity Set: Each user is hiding among more indistinguishable users, making re-identification harder

**Metrics collectively show a clear privacy advantage over ERC-4337**

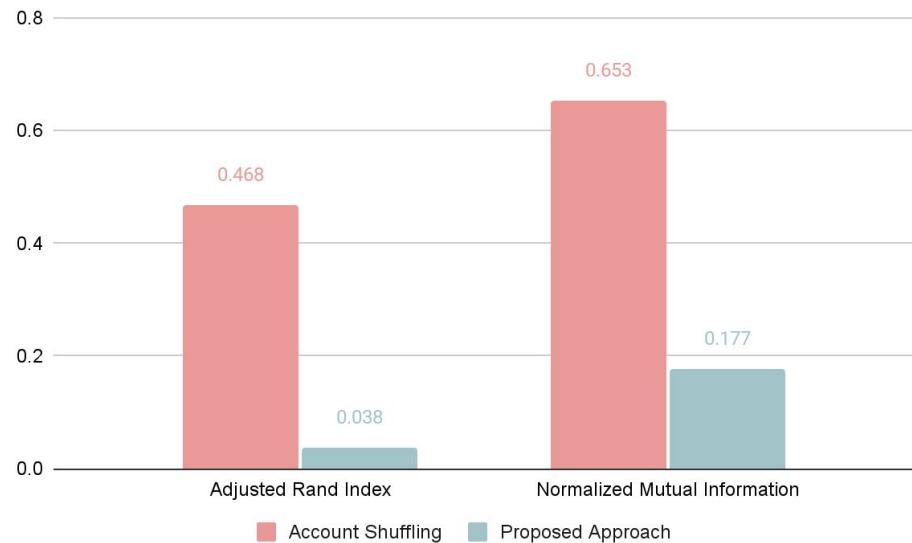




# Identity Privacy Improvement over Account Shuffling

- Attacker inference: Clustering via privileged operation heuristic
- Lower ARI and NMI: attacker cannot reliably cluster user addresses, making identity inference ineffective

**Clustering attacks fail, proving unlinkability beyond shuffling-based defenses.**



\*Privileged operation heuristic: If an account performs a privileged operation, on a particular smart contract, it is assumed to be controlled by the owner of that contract



# Discussion

Achievements	Limitations	Future Works
<ul style="list-style-type: none"><li>• Identity Unlinkability</li><li>• Verifiable Execution</li><li>• Quantitative Privacy Gains</li></ul>	<ul style="list-style-type: none"><li>• Trusted setup (Groth 16)</li><li>• No dynamic credential management</li><li>• Limited evaluation (gas/scalability).</li><li>• Simplified attacker model.</li></ul>	<ul style="list-style-type: none"><li>• Trying transparent proof (e.g., Plonk, Halo2).</li><li>• Add credential revocation/delegation</li><li>• Evaluate scalability and gas costs</li><li>• Map to GDPR/HIPAA compliance</li></ul>



# Conclusion: Toward Identity-Private Blockchain Healthcare

1. We built a privacy-preserving healthcare data sharing system that protects *user identities* on blockchain.
2. Our approach integrates Account Abstraction, Zero-Knowledge Proofs, and Pedersen Commitments into a deployable framework.
3. It achieves unlinkability and verifiability.
4. Experiments show higher entropy and lower inference accuracy, confirming strong resistance to identity inference attacks.

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*Transparency and auditability can coexist with privacy*



# Q&A

Thank you for your attention

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