CPU card parsing and operation instructions



Category Column: NFC technology

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Part 1 CPU Basics

1. Why use CPU cards?

IC cards can be divided into contact IC cards, contactless IC cards and composite cards in terms of interface mode. In terms of device technology, it can be divided into unencrypted memory cards, encrypted memory cards and CPU cards. The non-encrypted card has no security, and the data in the card can be arbitrarily rewritten, and the encrypted memory card adds a logical encryption circuit on the basis of the ordinary memory card to become an encrypted memory card. Because the logic encryption memory card uses password control logic to control the access and rewrite of EEPROM, it needs to verify the password before it can be written, so it is safe for the chip itself, but it is not safe in the application. It has the following unsafe factors:

- 1. The password is transmitted in plaintext on the line and is easy to be intercepted;
- 2. For system vendors, passwords and encryption algorithms are transparent.
- 3. The logical encryption card cannot authenticate whether the application is legitimate. For example, if someone fakes an ATM and you have no way of knowing its legitimacy, when you insert a credit card and enter your PIN, the password of the credit card is intercepted. Another example is INTENET online shopping, if the card is encrypted logically, the shopper is also unable to determine the legitimacy of the online store.

It is precisely because of the insecurity factors in the use of logical encryption cards that the development of CPU cards has been promoted. The CPU card can achieve the legitimacy authentication of the person, the card, and the system.

2. Three types of authentication for CPU cards

CPU cards have three authentication methods:

Cardholder legitimacy authentication - PIN verification

Card legitimacy certification - internal certification

System legitimacy certification - external certification

Cardholder Legitimacy Certification:

The process of verification by entering a personal order by the cardholder.

System legitimacy certification (external certification) process:

system card,

Send a random number X

Encrypt random numbers [with specified algorithm and key].

[Use the specified algorithm and key] to decrypt Y and get the result Z Compare X, Z, if the same means that the system is legitimate; Card legitimacy verification (internal authentication) process: System card

Send a random number X

Encrypt random numbers with the specified algorithm and key].

[Use the specified algorithm and key] to decrypt Y and get the result Z

Compare X, Z, if the same means that the card is legitimate;

In the above authentication process, the key does not appear in plaintext on the line, it is encrypted by random numbers every time it is sent, and because there is a random number to participate, it is ensured that the content of each transmission is different. If intercepted, it doesn't make any sense. This is not only a password-to-password authentication, but also a method authentication method, just like the cipher telegram used in the early days in the army, the sender encrypts the message into ciphertext according to a certain method and sends it, and then the receiver receives it and decrypts the ciphertext according to a certain method.

Through this authentication method, there is no attack point on the line, and the card can also verify the legitimacy of the application.

However, because the key and algorithm used by the system side for authentication are in the application, it still cannot remove the aggressiveness of the system vendor.

Here, we introduce the concept of SAM card.

SAM card is a CPU card with special performance, which is used to store keys and encryption algorithms, and can complete mutual authentication, password verification, encryption and decryption operations in transactions, and is generally used as an identity mark.

Thanks to the advent of SAM cards, we have a more complete system solution.

When issuing a card, we store the master key into the SAM card, and then the master key in the SAM card encrypts the characteristic byte of the user card (such as the application serial number) to generate a sub-key, and injects the sub-key into the user card. Due to the uniqueness of the application serial number, the sub-key within each user card is different

Once the key is injected into the card, it does not appear outside the card. When in use, the master key of the SAM card generates a sub-key and stores it in the RAM area to encrypt and decrypt data.

The above certification process takes the following form:

System legitimacy certification (external certification) process:

SAM Card System card Send a random number X

The SAM card generates a sub-key to encrypt the random number

Decrypt Y to get result Z

Compare X, Z, if the same means that the system is legitimate;

Card legitimacy verification (internal authentication) process:

SAM Card System card

Send a random number X

Encrypt random numbers with the specified algorithm and key].

SAM card decrypts Y, and results Z

Compare X, Z, if the same means that the card is legitimate;

In this way, the key in the application is transferred to the SAM card, and the authentication becomes the card-card authentication, and the system vendor no longer has responsibility.

3. Line protection

When the card transmits data to the outside world, if it is transmitted in clear text, the data can be easily loaded and analyzed. At the same time, it is also possible to tamper with the transmitted data, and to solve this problem, the CPU card provides a line protection function.

There are two types of line protection, one is to encrypt the transmitted data in DES and transmit it in ciphertext form to prevent interception and analysis. The second is to attach a MAC address (security message authentication code) to the transmitted data, which the receiver will first verify after receiving it, and receive it only after the verification is correct, so as to ensure the authenticity and integrity of the data.

Fourth, the hardware structure diagram

EEPROM is used to store user data; The ROM is used to store the COS operating system, and the RAM area is used to store the intermediate variables of the COS runtime. COS (chip operation system) is a chip operating system, similar to DOS and WINDWOS, and a CPU card without COS is as unusable as a PC without DOS and WINDOWS. COS is solidified into the ROM area by the chip supplier when the chip leaves the factory, and this process is called masking.

COS is the core part of the CPU card, and it implements the security of the CPU card together with the hardware. Part 1 CPU Basic Knowledge

1. Why use CPU card

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Cardholder legitimacy authentication

The process of verification by entering a personal order by the cardholder.

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System card,

send random number X

[with the specified algorithm and key] to encrypt

the random number

[with the specified algorithm and key] to decrypt Y, and obtain the result Z

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Card legitimacy authentication (internal authentication) process:

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In the above authentication process, the key does not appear in plaintext on the line, it is encrypted by random numbers every time it is sent, and because there is a random number to participate, it is ensured that the content of each transmission is different. If intercepted, it doesn't make any sense. This is not only a password-to-password authentication, but also a method authentication method, just like the cipher telegram used in the early days in the army, the sender encrypts the message into ciphertext according to a certain method and sends it, and then the receiver receives it and decrypts the ciphertext according to a certain method.

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The above authentication process becomes the following form:

System legitimacy authentication (external authentication) process:

SAM card

system card sends random number X

SAM card generates a sub-key to encrypt random number

Decrypt Y, and get the result Z

Compare X, Z, if the same means that the system is legitimate;

Card legitimacy authentication (internal authentication) process:

SAM card system

card

sends random number X

with the specified algorithm, key] to encrypt the random number

SAM card decrypts Y, and the result Z

compares X, Z, if it is the same, it means that the card is legitimate;

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COS is the core part of the CPU card, and it implements the security of the CPU card together with the hardware.

Part 2 Introduction to

SmartCOS SmartCOS is a chip operating system independently developed by Minghua Company, which passed the certification of the People's Bank of China in June 1999. COS is mainly divided into four parts:

1. SMARTCOS FILE SYSTEM

CPU CARD IS MANAGED IN THE FORM OF FILESSMARTCOS SUPPORTS THE FOLLOWING FILE SYSTEMS.

1. Files can be divided into MF files, DF files, and EF files

MF: The master file is the root of the entire file system, which is unique and equivalent to the root directory;

DF: A dedicated file, equivalent to a subdirectory, which can be used to store all the files of an application. A DF can be an application, or multiple DFs can be used for the same application.

EF: A basic file that stores a variety of application data and management information.

2. EF is divided into two types from the storage content:

Security basic file: used to store the key, each directory can only create a security basic file, the key file can not be selected by file selection, the key content can not be read, but it can be used and modified when the conditions are met.

Working basic file: used to store the actual data of the application, the number and size are limited by space. Reads and writes when conditions are met.

3. Basic file structure The structure of

basic files can be divided into the following four types:

Binary files

Data is read and written in bytes, and the length of each read and write cannot exceed 110 bytes; Can be used to store out-of-order data.

Linear fixed-length record file:

Each record is of fixed length, and the record can be accessed by record number, and the record range does not exceed 254;

The length of each record does not exceed 110 bytes, and the key file is a linear fixed-length record file, and the length of each record is fixed at 25 bytes. It can be used to store regular, fixed-length data.

Linear Variable Length Record Files:

The length of each record can vary, but the maximum length cannot exceed 110

bytes, which can be accessed by the record number.

Cyclic fixed-length record file structure:

It is equivalent to a ring record queue, which is stored according to the first-in, first-out principle, the latest record number is 1, the last written record number is 2, and so on, the earliest record is automatically overwritten after the record is full.

4. File structure diagram

EF and DF can be established under MF;

DF can no longer be created under DF, only EF can be created;

KEY file: used to control the creation, reading and writing of files under MF

Data file

Key file:

Used to control the creation and access of files under DF

Data files (such as wallet files, etc.

).

5. Calculation of file space The header file length of

MF is 10 bytes + the length of the file name (5-16 bytes)

The header file length of DF is 10 bytes + the length of the file name The space occupied by

the EF file:

Define the space of the record and loop record file =File header space (10 bytes) + number of records * record length The space of

the variable-length record structure file = the file header space (10 bytes) + the space requested when creating The space occupied by the

key file = the file header space (10 bytes) + the number of keys \ast 25 bytes The space of

the wallet file = the file header (10 bytes) + the file body (17 bytes

The space of the passbook file = file header (10 bytes) + file body (20 bytes)

The file system is established, so how to ensure the security of the file, the following describes the security system.

2. SMARTCOS security system

1. State machine is the security state:

refers to a security level that the card is currently in, with (0---F) 16 security states. Automatically set to 0 after reset, and the security status of the current application will be automatically cleared to 0 after being successfully selected or reset. A change in security status must be achieved through key authentication.

**Only PIN reconciliation and external authentication in the current directory can change the security status.

2. Security attributes, i.e., access

permissions, are specified when creating files.

It is a compartmental concept, for example, if the read permission of a file is XY, it means that the state machine (security state) M of the current application must meet the requirements of X= \(\lambda M = \lambda Y \). If the read permission is set to 2F, the current state machine (security state) reaches 2 or higher, and the file can be read.

3. The relationship between the key and the security state

Each key defines the subsequent state when it is established, that is, the security state that can be achieved after passing the key authentication.

4

. The relationship between security status and access rights

(security status)

The use of permissions ------ keys----- the subsequent status ----- file read and write limits

5. The security status is only valid in the current directory, once you select another application, the state machine will automatically jump to the minimum permission 0. Security is independent between directories.

4. Reset Response

Symbol Byte Content Content Explanation

TS 3B Forward Convention

TO 6C TB1 and TC1 exist, and the historical characters are 12

TB1 00 No additional programming voltage is required

TC1 02 requires 2 additional protection times

T1-TC XX Historical Characters

Specific Meanings of SMARTCOS Historical Characters:

Symbol Byte Content Content Interpretation

T1 XX The version number of SMARTCOS

T2 XX Card Status Byte

T3 86 Minghua IC Card Manufacturing Organization Identification Number

T4 38

T5-TC XX Card Unique Serial Number

Card Status Byte is described as follows:

B7 B6 B5 B4 B3 B2 B1 B0 Status

011 XXX XXX 001 XXX XXX XXX The card has been initialized and successful, the card is not initialized, the card initialization process is locked,

00000XX the card is not personalized

0 0 1 0 X X X X The card personalization is not complete

0 1 1 0 X X X X The card is personalization successful

0 0 0 1 X X X X The card was not personalized, and the card was locked

0 1 1 1 1 X X X X The card was personalized, and the card was locked

There are concepts about initialization and personalization in the card state byte, and we look at these two concepts from the life cycle of the card, which generally includes the following parts:

chip chip manufacturer

Mask COS Chip Manufacturer

Encapsulated into a card Card Producer

Card Initialization (COS Enabled) Card Manufacturer Transfer Password Protection

Card Personalization Card

Use of Cards Users of Cards

Recycling of

Cards Card Issuers

The

initialization process is to activate COS, define the COS version, and then use COS after this process.

The personalization process refers to the creation of files, the writing of user data, etc.

The card manufacturer is protected by a transmission password when it is initialized and delivered to the card issuer.

The IC card must support either T=0 or T=1 protocols, but not both, while the terminal must support both T=0 and bothT=1 protocol.

T=0 communication protocol is an asynchronous half-duplex character transmission protocol;

T=1 communication protocol is an asynchronous half-duplex block transmission protocol;

In the ISO7816-3 standard, these two protocols are specified;

The protocol used by the IC card is specified in TD1, and if TD1 is not in the reset response message, it means that the protocol with T=0 is used for communication. After the reset answer, the IC card and the terminal communicate with the protocol specified by the IC card.

5. Instruction Analysis

Here, we take an e-wallet application as an example to explain the instructions of SmartCOS.

1. File structure:

The following file structure is defined in the PBoC specification:

1) Basic data file of public application of electronic passbook ED/e-wallet EP application

File structure:

File Identifier (SFI) '21' (decimal)

File Type Transparent

File Size 30

File Access Control Read = Free Overwrite = Security Information

Required Bytes Data Element Length

1-8 Issuer ID 8

9 App Type ID 1

10 App Version 1

11-20 App Serial Number 10

21-24 App Enabled Date 4

25-28 App Expiration Date 4

29-30 Issuer Custom FCI Data 2

2) Cardholder basic data file

for e-passbook ED/e-wallet EP application

File Identifier (SFI) '22' (decimal)

File type Transparent

File size 39

File access control Read = free to rewrite = security information

required Byte Data Element Length

1 Card Type Identifier 1

2 Bank Employee Identifier 1

3-22 Cardholder Name 20

23-38 Cardholder ID Number 16

39 Cardholder ID Type 1

3) Electronic Passbook ED Transaction Details File

Document Identifier (SFI) '24' (Decimal)

File Type Loop

File Access Control Read = PIN Protection

Overwrite = Not Allowed

Record Size

Bytes Element Length

1-2 ED or EP Online or Offline Transaction Serial Number

2 3-5 Overdraft Limit 3

6-9 Transaction Amount 4

10 Transaction Type Identifier 1

11-16 Terminal Number 6

17-20 Transaction Date (Terminal) 4

21-23 Transaction Time (Terminal) 3

- 2. The security design is as follows:
- 1) External authentication can be carried out after checking the password;
- 2) The external authentication key No. 01 is used to control the storage of the e-wallet;

- 3) The No. 02 external authentication key is used to control the modification of the basic file and the modification of the key;
- 4) You can make purchases after checking the password.

The following keys are installed in the KEY file:

KEY Type Identification KID Subsequent status of usage rights The function description of the KEY

OB 01 OF 1 Personal PIN for PIN verification

08 01 11 2 External authentication key, used for e-wallet storage

08 02 1F F External authentication key, used for modifying basic files and keys

01 01 22 No Loop Key, used to generate Loop MAC

00 02 01 No e-wallet consumption key, used to generate wallet consumption MAC

02 01 03 No TAC key, used to generate Cache, Consumption, Withdrawal, and TAC

for modifying overdraft limits 05 01 33 No application maintenance key, used to generate MAC for application lock, application unlock, card lock and read, update binary, and record commands

3. Instruction sequence:

1) The card issuance process, creating a file on the card and installing the key

[Create MF] APDU command: 80 (CLA) e0 (INS) 00 (P1) 00 (P2) 18 (Lc) ff ff (8 byte transfer code) ff (Establish security attributes of files under MF) 01 (short file identifier for directory files) 31 50 41 59 2e 53 59 53 2e 44 44 46 30 31 (name of the file created).

[Create DF] APDU commands: 80 (CLA), E0 (INS), 01 (P1), 00 (P2) 0D (LC message length), 2F 01 (file identifier), ff (create file permissions), 00 (COS reserved), A0 00 00 00 03 86 98 07 01 (ADF name).

Create File under DF APDU command: 80 (CLA) E0 (INS) 02 (P1) 00 (P2) 07 (File Information Length) 6F 02 (Key File File Identifier) 05 (File Type) FF (Add Permission for New Key) 00 (COS Retention) 09 (Number of Records) 19 (Record Length):

Install PIN [Write Key] APDU command: 80 (CLA), E8 (INS) 00 (P1) 00 (P2) 0A (key information length), 01 (key identifier), 01 (key version number), 00 (algorithm identification), 0B (key type), 0F (usage permission), 01 (subsequent status), 2F (modify permission), 33 (error counter), 12 34 (personal password).

Install the external authentication key (DEAK) [Write Key] APDU command: 80 (CLA), E8 (INS) 00 (P1) 00 (P2) 18 (key information length) 01 (key identifier) 01 (key version number) 00 (algorithm identifier) 08 (key type) 11 (usage rights) 02 (subsequent status) FF (modify permissions) 33 (error counter) XX XX XX (Key Content)

Install External Authentication Key (DEAK) [Write Key] APDU Command: 80 (CLA), E8 (INS) 00 (P1) 00 (P2) 18 (Key Information Length) 02 (Key Identifier) 01 (Key Version Number) 00 (Algorithm Identification) 08 (Key Type) 1F (Usage Permission) 0F (Subsequent Status) FF (Modify Permission) 33 (Error counter) XX XX (Key Contents

). Install the electronicsThe consumption key DPK of the wallet EP. [Write Key] APDU commands: 80 (CLA), E8 (INS) 00 (P1) 00 (P2) 18 (key information length) 02 (key identifier) 01 (key version number) 00 (algorithm identifier) 00 (key type) 01 (usage rights) 00 (subsequent status) FF (modify permissions) 00 (error counter) XX XX (KEY CONTENT).

DTK [Write Key] APDU command: 80 (CLA), E8 (INS) 00 (P1) 00 (P2) 18 (key information length) 01 (key identifier) 01 (key version number) 00 (algorithm identifier) 07 (key type) 0F (usage permission) 00 (subsequent status) FF (modify permission) 00 (error counter) XX XX XX XX (Key Content)

Installing the Application Maintenance Key DAMK[Write Key] APDU command: 80 (CLA) E8 (INS) 00 (P1) 00 (P2) 18 (Key Information Length) 01 (Key Identifier) 01 (Key Version Number) 00 (Algorithm Identification) 05 (Key Type) 0F (Permissions) 00 (Subsequent Status) FF (Modify Permissions) 00 (Error Counter) XX XX XX (Key Contents).

[Create File] APDU command: 80 (CLA) E0 (INS) 02 (P1) 00 (P2) 07 (LC message length) 00 15 (file identifier) 00 (binary file type) 0F (read permission) FF (update permission) 00 1e (file length).

Create File APDU Command: 80 (CLA) E0 (INS) 02 (P1) 00 (P2) 07 (File Info Length) 00 16 (File Identifier) 00 (Binary File Type) 0F (Read Permission) FF (Update Permission) 00 27 (File Length).

Create File APDU command: 80 (CLA) E0 (INS) 02 (P1) 00 (P2) 07 (File Information Length) 00 18 (File Identifier) 03 (Circular Record File Type) 1F (Read Permission) 10 (Update Permission) 0a 17 (File Length).

Write a public application basic data file [Update Binary] APDU command: 00 (CLA), D6 (INS), 95 (P1), 00 (P2), 1E (message length), A0 00 00 00 00, 03 00 00 01 (issuer ID), 03 (application type ID), 01 (application version), 00 00 19 98 08, 15 00 00 00 01 (application serial number), 20 00 10 01 (application activation date), 20 02 12 31 (application expiration date) 55 66 (issuer-defined FCI data).

Create Wallet EP File [Create File] APDU Command: 80 (CLA) E0 (INS) 02 (P1) 00 (P2) 07 (Lc) 00 01 (File Identifier) 06 (File Type) 00 (Permission 1) 00 (Permission 2) 00 (LEN1)

00 (LEN2)

[Create End] APDU command: 80 (CLA) E0 (INS) 01 (P1) 01 (P2) 02 (File Identifier Length) 2F 01 (File Identifier).

MF [Create End] APDU Command: 80 (CLA) E0 (INS) 00 (P1) 01 (P2) 02 (File Identifier Length) 3F 00 (File Identifier)

2) Transaction Process

Here, let's take the captive process and consumption process as examples:

A. Trapping -- Deposit the amount into the

card Before depositing, you must proofread the personal population order and the external certification No. 01;

Select Application[Select File]APDU Command: 00 (CLA) A4 (INS) 00 (P1) 00 (P2) 02 (Length) 01 (File Identifier)

Verify PIN [Verify] APDU Command: 00 (CLA) 20 (INS) 00 (P1) 00(P2) 02(Length) 12 34(PIN)

External Authentication: [

Get Challenge] APDU Command: 00 (CLA) 84 (INS) 00 (P1) 00 (P2) 08

3DES encryption of random numbers with No. 01 external authentication key;

[External Authentication] APDU command: 00(CLA) 82(INS) 00(P1) 01

(P2) 08 (length) XX XX XX (encrypted random number)

Trapping:[

Initializing Trap [Initalize For Load] APDU Command: 80(CLA) 50(INS) 00(P1) 02(P2) 0B(Length) 01 (Key Identifier) 00 00 10 00 (Transaction Amount) 00 00 00 00 00 01

(Terminal

Number)

If the initialization of the trap is successful, the content of the reply data field is as follows:

Description Length (bytes)

Old balance of e-passbook or e-wallet 4

E-passbook or e-wallet online transaction serial number 2

Key

Algorithm identification 1

Pseudo-random number ICC 4

MAC1 4

3DES encryption is generated

by using the corresponding content of the circle (4-byte pseudo-random number ICC + 2-byte electronic passbook or e-wallet online transaction serial number +80 00).process key;

Trap [Credit For Load] APDU command: 80 (CLA) 52 (INS) 00 (P1) 00 (P2) 0B (LC)

YY YY MM DD (transaction date) HH MM SS (Trading time) XX XX XX XX (MAC2).

Calculation of MAC2:

Initial value: 00 00 00 00 00 00 00 00 00

Key: The process key generated above Data

generated for MAC2: 00 00 10 00 (4-byte transaction amount) 02 (transaction type ID) 00 00 00 01 (6-byte terminal number) YY YY MM DD (Transaction Date) HH MM SS (Trading Hours)

If the deposit transaction is successful, the online transaction serial number of the e-wallet file is added to 1, the transaction amount is added to the balance of the e-wallet, and a record is added to the transaction details file.

B. Consumption process:

Verify PIN before consumption:

Verify PIN [Verify] APDU command: 00 (CLA) 20 (INS) 00 (P1) 00 (P2) 02 (length) 12 34 (PIN).

[Initialize For Purchase] APDU command: 80(CLA) 50(INS) 01(P1) 02(P2) 0B (length) 01 (consumption key identifier) 00 00 00 01 (consumption amount) 00 00 00 00 (terminal number)

If the initialization consumption is successful, the data field content is answered:

Description Length (bytes)

E-passbook or e-wallet old balance 4

E-passbook or e-wallet online transaction serial number 2

Overdraft limit 3

Key version number 1 Algorithm identifier 1

Pseudo-random number ICC 4

Consumption[Debit For Purchase]APDU Command: 80 (CLA) 54 (INS) 01 (P1) 00 (P2) 0F (LC)

YY YY YY YY (Terminal Transaction Serial Number) YY YY MM DD (Terminal Transaction Date) HH MM SS (Trading Hours) XX XX (MAC1)

Use the consumption/withdrawal key pair (4-byte pseudo-random number ICC + 2-byte electronic passbook or e-wallet online transaction serial number + the last 2 bytes of the terminal transaction serial number) to encrypt the process key;

Calculation of MAC1: Initial

value: 00 00 00 00 00 00 00 00 00 (8 bytes)

Key: Process key Data

to generate MAC code: 4 bytes of transaction amount + 1 byte of transaction type ID + 6 bytes of terminal number + 4 bytes of terminal transaction date +3 bytes terminal transaction time

If the consumption transaction is successful, the online transaction serial number of the e-wallet file is added by 1, the e-wallet subtracts the transaction amount, and a record is added to the transaction detail file.

6. Analysis of other instructions The

following instructions have many questions in the process of use, which are explained as follows: 1.

The ciphertext installation process of the key:

In version 3.2, the ciphertext installation of the key will be different, and its process is as follows:

1) Create a key file:

When creating a key file, the relevant file information is as follows:

LC File information

07 File identifier (2 bytes) File type (1 byte) Permission 1 (1 byte) Permission 2 (1 byte) Len1 (1 byte) Len2 (1 byte).

Permission 1 indicates the permission to add a new key, and permission 2 indicates how to install the key. Permission 2 is set to 80h, which means that the installation is in ciphertext.

- 2) Use the application master key of the previous layer (SmartCOS 3.2 stipulates that the external authentication key with key ID 01 is the application master key) to encrypt the key information (the key information is: plaintext key information length + key information + 80 00 00 00 00 00 00, and the encryption method of the key information is standard Triple DES or Single DES, if the key is 16 bytes, use Triple DES encryption, if it is 8 bytes, it is encrypted with Single DES.
- 3) Generate MAC code, the initial value is: 4 bytes of random number + 00 00 00, the data generated MAC code is: 5 command headers + encrypted key information.
- 4) The ciphertext installation of other keys is encrypted with the application master key.
- 5) When installing the master key of the application under MF in ciphertext, use the transmission key of the card to install it.
- 2. Use security packets to write binaries
- 1) When creating binaries, the fourth digit of the file type is 1, which means that security packets are used.
- 2) LC is the number of bytes written + 4
- 3) Calculation of security packets (MAC): The

initial value is: 4 bytes of random number + '00 00 00 00'

The key is: the code value of the application maintenance key

CLA 04

INS D6

P1 XX P2 XX

Lc Write Data Length + 4

DATA Write Data + 4 Bytes MAC 8

and PIN The

unblock PIN unlock command packet is encoded as follows:

Code value CLA 84

INS 24

P1 00

P2 01 - Unlock PIN

Lc 00

DATA encrypts the PIN data element + Packet Authentication Code (MAC) data element, using the PIN unlock key.

Wherein:

Encrypted personal password data element:

PIN unlock key pair PIN plaintext length + PIN + 80 00 (make up for the multiple of 8) to encrypt the

generation of MAC:

Initial value: 4 bytes random number + 00 00

Password: PIN unlock key Data

generated MAC code: CLA+INS+P1+P2+LC+encrypted PIN (8 bytes)

6. Use

of CPUDEMO

7. Instructions for

using RD reader libraries (1) int cpu_reset(int icdev,unsigned char *data_buffer);

NOTE: CPU Card Reset Operation Parameters: icdev: Device identifier

dat returned by initializationa_buffer: String returned by the reset response

Returns: = 0 Correct < 0 Error (see error code)

(2) int cpu_protocol(int icdev,int len, unsigned char * send_cmd,

unsigned char *receive data)

Note: Send a command string to the reader according to the transmission protocol of the CPU card T=1

Parameters: icdev: device identifier returned by initialization send_cmd: command string for card operations (T=1 format)

len: length of the command string

receive_data: data string returned by the CPU card

Return: = 0 Correct < 0 Error (see Error Code).

(3) int cpu_comres(int icdev,int len, unsigned char *send_cmd, unsigned char *receive_data)

Notes:

1) Whether it is a T=0 or T=1 card, RD readers transmit commands in T=1 format.

The command string send_cmd sent includes:

NAD+PCB+LEN+COMMAND+BCC

For T=0 cards, NAD and PCB can be filled with '00';

The structure of the return value receive_data is the same as that of Send_cmd.

2) cpu_comres is different from cpu_protocol()

a. When the upper-layer library receives the state byte SW1SW=61XX, it Cpu_protocol () function to do further processing, and sends a Get Respond command to the reader, and returns the data and the state byte together; The cpu_comres() function does nothing and returns the state byte SW1SW2 directly.

b. When the upper-layer library receives the state byte SW1SW=6CXX, it Cpu_protocol () function for further processing, assigns XX to le, and then resends the last instruction, and returns the data and the state byte together; The cpu_comres() function does nothing and returns the state byte SW1SW2 directly.

Part 2 Introduction to SmartCOS

SmartCOS is a chip operating system independently developed by Minghua Company, which passed the People's Bank of China certification in June 1999.

COS is divided into four parts:

1. The file system of SMARTCOS

The CPU card is managed in file mode, and SmartCOS supports the following file systems:

1. Files can be divided into MF files, DF files, and EF files

MF: Master file, which is the root of the entire file system and is unique, equivalent to the root directory.

DF: A dedicated file, equivalent to a subdirectory, which can be used to store all the files of an application. A DF can be an application, or multiple DFs can be used for the same application.

EF: A basic file that stores a variety of application data and management information.

2. EF is divided into two types in terms of storage content:

Security basic file: used to store the key, each directory can only create a security basic file, the key file can not be selected by file selection, the key content can not be read, but can be used and modified when the conditions are met.

Working basic file: used to store the actual data of the application, the number and size are limited by space. Reads and writes when conditions are met.

3. Basic file structure

The structure of a basic document can be divided into the following four types:

Binary file

Data is read and written in bytes, and the length of each read and write cannot exceed 110 bytes. Can be used to store out-of-order data.

Linear Fixed-Length Record Files:

Each record is of fixed length, and records can be accessed by record number, and the record range does not exceed 254;

The length of each record does not exceed 110 bytes, and the key file is a linear fixed-length record file, and the length of each record is fixed at 25 bytes. It can be used to store regular, fixed-length data.

Linear Variable Length Record File:

The length of each record can vary, but the maximum length cannot exceed 110

bytes, which can be accessed by the record number.

Structure of the cyclic fixed-length record file:

It is equivalent to a ring record queue, which is stored according to the first-in, first-out principle, the latest record number is 1, the last written record number is 2, and so on, the earliest record is automatically overwritten after the record is full.

4. File structure diagram

EF and DF can be created under MF;

DF can no longer be created under DF, only EF can be created;

KEY file: used to control the creation, reading, and writing of files in MF

Data files

Key file:

It is used to control the creation and access of files under DF

Data files (e.g. wallet files, etc.)

5. Calculation of file space

MF header file length is 10 bytes + file name length (5-16 bytes)

The length of the DF header file is 10 bytes + the length of the file name

EF file footprint:

Define the space of the record and loop record files = file header space (10 bytes) + number of records * record length

The space of the variable-length record structure file = the header space (10 bytes) + the space requested when it is created

Space occupied by the key file = Header space (10 bytes) + Number of keys × 25 bytes

Wallet file space = file header (10 bytes) + file body (17 bytes)

Space of passbook file = File header (10 bytes) + File body (20 bytes)

The file system is established, so how to ensure the security of the file, the following describes the security system.

Second, the security system of SMARTCOS

1. The state machine is the safe state:

Refers to a security level that the card is currently in, with (0---F) 16 security states. Automatically set to 0 after reset, and the security status of the current application will be automatically cleared to 0 after being successfully selected or reset. A change in security status must be achieved through key authentication.

**Only PIN reconciliation and external authentication in the current directory can change the security status.

2. Security attributes are access rights

Access permissions are specified when the file is created.

It is a compartmental concept, for example, if the read permission of a file is XY, it means that the state machine (security state) M of the current application must meet the requirements of X= $\langle M = \langle Y | M = \langle Y | M = \langle M =$

3. The relationship between the key and the security state

Each key defines a subsequent state when it is established, that is, the security state that can be achieved after passing the key authentication, and among various keys, only PIN authentication and external authentication can change the security status of the current directory.

4. The relationship between security status and access rights

(Security Status)

Usage permissions----- key usage ----- subsequent status----- file read and write retrieval limits

5. The security status is only valid in the current directory, once you select another application, the state machine will automatically jump to the minimum permission 0. Security is independent between directories.

Fourth, reset response

The symbol byte content is explained

TS 3B Forward Convention

TO 6C TB1 and TC1 exist, and the historical character is 12

The TB1 00 requires no additional programming voltage

The TC1 02 requires 2 additional protection periods

T1-TC XX historical characters

Specific meanings of SMARTCOS historical characters:

The symbol byte content is explained

The version number of the T1 XX SMARTCOS

T2 XX card status bytes

T3 86 Minghua IC card manufacturing organization identification number

T5-TC XX card unique serial number

The card status bytes are described as follows:

B7 B6 B5 B4 B3 B2 B1 B0 Status

011 XXX XXX 001 XXX XXX The card has been initialized and successfully, the card has not been initialized, and the card initialization process is locked

00000X The card is not personalized

0 0 1 0 X X X X The personalization of the card is not finished

0 1 1 0 X X X X X The card is personalized

0001 X X X X The card personalization was not successful and the card was locked

0 1 1 1 1 X X X X X The card is successfully personalized and the card is locked

0

There are concepts about initialization and personalization in the card state byte, and we look at these two concepts from the life cycle of the card, which generally includes the following parts:

Chips chip producers

Mask COS chip manufacturer

Encapsulated into a card card manufacturer

Initialization of the card (COS enabled) by the card producer

Transmission password protection

Personalization of the card card issuer

Use of the card The consumer of the card

Recycling of cards by card issuers

The initialization process is to activate COS and define the COS version, after which COS can be used.

The personalization process refers to the creation of files, the writing of user data, etc.

The card manufacturer is protected by a transmission password when it is initialized and delivered to the card issuer.

The IC card must support either T=0 or T=1 protocols, but not both, while the terminal must support both T=0 and T=1 protocols.

T=0 communication protocol is an asynchronous half-duplex character transmission protocol;

T=1 communication protocol is an asynchronous half-duplex block transmission protocol;

In the ISO7816-3 standard, these two protocols are specified;

The protocol used by the IC card is specified in TD1, and if TD1 is not in the reset response message, it means that the protocol with T=0 is used for communication. After the reset answer, the IC card and the terminal communicate with the protocol specified by the IC card.

5. Instruction Analysis

Here, we take an e-wallet application as an example to explain the instructions of SmartCOS.

1. File structure:

The following file structure is defined in the pedestrian specification:

1) The basic data file of the public application of the e-passbook ED/e-wallet EP application

File structure:

File Identifier (SFI) '21' (decimal)

File Type: Transparent

File size 30

File Access Control Read = Free Rewrite = Security Information Required

Byte data object length

1-8 Issuer ID 8

9 App Type Identifier 1

10 App Version 1

11-20 Apply serial number 10

21-24 App Enabled Date 4

25-28 Apply effective date 4

29-30 Issuer-defined FCI data 2

2) Cardholder basic data file of e-passbook ED/e-wallet EP application

File Identifier (SFI) '22' (decimal)

File Type: Transparent

File size 39

File Access Control Read = Free Rewrite = Security Information Required

Byte data object length

- 1 Card Type Identification 1
- 2 Employee Identification 1 of the Bank
- 3-22 Cardholder Name 20
- 23-38 Cardholder ID number 16
- 39 Cardholder ID Type 1

3) Electronic passbook ED transaction details file

File Identifier (SFI) '24' (decimal)

File Type: Loop

File Access Control Read = PIN Protection

Rewrite = Not allowed

Record size 23

Byte data object length

- 1-2 ED or EP online or offline transaction serial number 2
- 3-5 Overdraft Limit 3
- 6-9 Transaction amount 4
- 10 Transaction Type Identifier 1
- 11-16 Terminal No. 6
- 17-20 Trading Date (Terminal) 4
- 21-23 Trading hours (terminal) 3
- 2. The safety design is as follows:
- 1) External authentication can be carried out after checking the password;
- 2) The external authentication key No. 01 is used to control the storage of the e-wallet;
- 3) The No. 02 external authentication key is used to control the modification of the basic file and the modification of the key;
- 4) You can make purchases after checking the password.

The following keys are installed in the KEY file:

KEY TYPE IDENTIFIES THE KID STATUS OF THE PERMISSION TO USE THE KEY AND DESCRIBES THE FUNCTION OF THE KEY

OB 01 OF 1 PIN for PIN verification

08 01 11 2 External authentication key, used for wallet storage

08 02 1F F external authentication key, which is used to modify basic files and keys

01 01 22 No lapse key, used to generate a lapse MAC

00 02 01 No wallet consumption key, used to generate wallet consumption MAC

02 01 03 No TAC key, which is used to generate TAC for deposit, consumption, cash withdrawal, and modification of overdraft limit

05 01 33 No Application Maintenance Key, used to generate MAC for App Lock, App Unlock, Card Lock, Card Lock, and Read, Update Binary, and Log commands

3. Instruction sequence:

1) During the card issuance process, create a file and install the key on the card

[Create MF] APDU command: 80 (CLA) e0 (INS) 00 (P1) 00 (P2) 18 (Lc) ff ff (8 byte transfer code) ff (Establish security attributes of files under MF) 01 (short file identifier for directory files) 31 50 41 59 2e 53 59 53 2e 44 44 46 30 31 (name of the file created)

[Create DF] APDU Command: 80 (CLA) E0 (INS) 01 (P1) 00 (P2) 0D (LC Message Length) 2F 01 (File Identifier) FF (Establish File Permissions) 00 (COS Reserved) A0 00 00 00 03 86 98 07 01 (ADF Name)

Create File in DF APDU Command: 80 (CLA) E0(INS)02(P1) 00(P2) 07 (File Information Length) 6F 02 (Key File File Identifier) 05 (File Type) FF (Add Permission for New Key) 00 (COS Reservation) 09 (Number of Records) 19 (Record Length):

Install PIN [Write Key] APDU command: 80 (CLA) E8 (INS) 00 (P1) 00 (P2) 0A (key information length) 01 (key identifier) 01 (key version number) 00 (algorithm identification) 0B (key type) 0F (usage permission) 01 (subsequent status) 2F (modify permission) 33 (error counter) 12 34 (personal password)

Install the external authentication key (DEAK) [Write Key] APDU command: 80 (CLA), E8 (INS) 00 (P1) 00 (P2) 18 (key information length) 01 (key identifier) 01 (key version number) 00 (algorithm identifier) 08 (key type) 11 (usage rights) 02 (subsequent status) FF (modify permissions) 33 (error counter) XX XX XX (Key Contents)

Install the external authentication key (DEAK) [Write Key] APDU command: 80 (CLA), E8 (INS), 00 (P1), 00 (P2), 18 (key information length), 02 (key identifier), 01 (key version number), 00 (algorithm identifier), 08 (key type), 1F (usage permission), 0F (subsequent status), FF (modify permission), 33 (error counter), xx xx xx XX (Key Contents)

. Install the consumption key DPK of the e-wallet EP. [Write Key] APDU Command: 80 (CLA) E8 (INS) 00 (P1) 00 (P2) 18 (Key Information Length) 02 (Key Identifier) 01 (Key Version Number) 00 (Algorithm Identification) 00 (Key Type) 01 (Permission to Use) 00 (Subsequent Status) FF (Modify Permission) 00 (Error Counter) XX XX

Install the Wallet's Trap Key DLK [Write Key] APDU command: 80 (CLA) E8 (INS) 00 (P1) 00 (P2) 18 (Key Information Length) 01 (Key Identifier) 01 (Key Version Number) 00 (Algorithm Identifier) 01 (Key Type) 22 (Usage Permissions) 00 (Subsequent Status) FF (Modify Permissions) 00 (Error Counter) XX XX XX XX (Key Contents)

DTK [Write Key] APDU command: 80 (CLA), E8 (INS) 00 (P1) 00 (P2) 18 (key information length) 01 (key identifier) 01 (key version number) 00 (algorithm identifier) 07 (key type) 0F (usage permission) 00 (subsequent status) FF (modify permission) 00 (error counter) XX XX XX XX XX XX XX XX (key contents)

Install the application maintenance key DAMK[Write Key] APDU command: 80 (CLA), E8 (INS), 00 (P1), 00 (P2), 18 (key information length), 01 (key identifier), 01 (key version number), 00 (algorithm identifier), 05 (key type), 0F (usage permission), 00 (subsequent status), FF (modify permission), 00 (error counter), xx xx xx XX (Key Contents)

Create File APDU Command: 80 (CLA) E0 (INS) 02 (P1) 00 (P2) 07 (Lc message length) 00 15 (File identifier) 00 (binary file type) 0F (read permission) FF (update permission) 00 1e (file length)

Create File [Create File] APDU command: 80 (CLA) E0 (INS) 02 (P1) 00 (P2) 07 (file information length) 00 16 (file identifier) 00 (binary file type) 0F (read permission) FF (update permission) 00 27 (file length)

Create File APDU Command: 80 (CLA) E0 (INS) 02 (P1) 00 (P2) 07 (File Message Length) 00 18 (File Identifier) 03 (Circular Record File Type) 1F (Read Permission) 10 (Update Permission) 0a 17 (File Length)

Write a public application basic data file [Update Binary] APDU command: 00 (CLA), D6 (INS), 95 (P1), 00 (P2), 1E (message length), A0 00 00 00 00, 03 00 00 01 (issuer ID), 03 (application type ID), 01 (application version), 00 00 19 98 08, 15 00 00 00 01 (application serial number), 20 00 10 01 (application activation date), 20 02 12 31 (application expiration date) 55 66 (issuer-defined FCI data)

Create File for EP [Create File] APDU Command: 80 (CLA) E0 (INS) 02 (P1) 00 (P2) 07 (Lc) 00 01 (File Identifier) 06 (File Type) 00 (Permission 1) 00 (Permission 2) 00 (LEN1) 00 (LEN2)

Create End (DF) APDU Command: 80 (CLA) E0 (INS) 01 (P1) 01 (P2) 02 (File Identifier Length) 2F 01 (File Identifier)

MF [Create End] APDU Command: 80 (CLA) E0 (INS) 00 (P1) 01 (P2) 02 (File Identifier Length) 3F 00 (File Identifier)

2) Transaction process

Here, we take the captive process and consumption process as an example:

 $A. \ Circle \ deposit - \ deposit \ the \ amount \ into \ the \ card$

Before captivity, the personal population order and the external certification No. 01 must be proofread;

Select Apply[Select File]APDU Command: 00 (CLA) A4 (INS) 00 (P1) 00 (P2) 02 (Length) 2F 01 (File Identifier)

Verify PIN [Verify] APDU command: 00(CLA) 20(INS) 00(P1) 00(P2) 02(length) 12 34(PIN)

External Certification: [

Take the random number [Get Challenge] APDU command: 00(CLA) 84(INS) 00(P1) 00(P2) 08

3DES encryption of random numbers with No. 01 external authentication key;

External Authentication APDU command: 00(CLA) 82(INS) 00(P1) 01 (P2) 08 (length) XX XX XX (ENCRYPTED RANDOM NUMBER)

Captivity:[

MAC1 4

Initialize Trap [Initalize For Load] APDU command: 80 (CLA) 50 (INS) 00 (P1) 02 (P2) 0B (length) 01 (key identifier) 00 00 10 00 (transaction amount) 00 00 00 00 00 01 (final.) Endpoint number)

If the initialization of the captive is successful, the content of the data field is answered:

Description Length (bytes)
Old balance of e-passbook or e-wallet4
e-Passbook or e-Wallet online transaction serial number 2
Key version number 1
Algorithm ID 1
Pseudo-random number ICC 4

3DES encryption is carried out with the content of the circled response (4-byte pseudo-random number ICC + 2-byte electronic passbook or e-wallet online transaction number +80 00) to generate the process key;

[Credit For Load] APDU Command: 80 (CLA) 52 (INS) 00 (P1) 00 (P2) 0B (LC) YY YY MM DD (Trading Date) HH MM SS (Trading Hours) XX XX (MAC2)

MAC2 CALCULATIONS:

Initial value: 00 00 00 00 00 00 00 00 Key: The process key generated above

Data generated for MAC2: 00 00 10 00 (4-byte transaction amount) 02 (transaction type ID) 00 00 00 00 01 (6-byte terminal number) YY YY MM DD (transaction date) HH

MM SS (transaction time)

If the deposit transaction is successful, the online transaction number of the e-wallet file is added by 1, the transaction amount is added to the balance of the e-wallet, and a record is added to the transaction detail file.

B. Consumption process:

Verify the PIN before making a purchase:

Verify PIN [Verify] APDU command: 00(CLA) 20(INS) 00(P1) 00(P2) 02(length) 12 34(PIN)

Initialize Consumption [Initialize For Purchase] APDU command: 80 (CLA) 50 (INS) 01 (P1) 02 (P2) 0B (length) 01 (consumption key identifier) 00 00 00 01 (consumption amount) 00 00 00 01 (terminal number)

If the initialization consumption is successful, the data field content is answered:

Description Length (bytes)

Old balance of e-passbook or e-wallet4

e-Passbook or e-Wallet online transaction serial number 2

Overdraft Limit 3

Key version number 1

Algorithm ID 1

Pseudo-random number ICC 4

Debit For Purchase APDU Command: 80 (CLA) 54 (INS) 01 (P1) 00 (P2) 0F (LC)

YY YY YY YY (Terminal Transaction Number) YY YY MM DD (Terminal Transaction Date) HH MM SS (Trading Hours) XX XX XX (MAC1)

Use the consumption/withdrawal key pair (4-byte pseudo-random number ICC + 2-byte electronic passbook or e-wallet online transaction serial number + the last 2 bytes of the terminal transaction serial number) to encrypt the process key;

MAC1 Calculation:

Initialization: 00 00 00 00 00 00 00 00 00 (8 bytes)

Key: The process key

The data generated by the MAC code: 4 bytes of transaction amount + 1 byte of transaction type identification + 6 bytes of terminal number + 4 bytes of terminal transaction date + 3 bytes of terminal transaction time

If the transaction is successful, add 1 to the online transaction number of the e-wallet file, minus the transaction amount from the e-wallet, and add a record to the transaction details file.

6. Analysis of other instructions

The following directives have many doubts in the process of use, which are explained as follows:

1. The ciphertext installation process of the key:

In version 3.2, the ciphertext installation of keys will be different, and the process is as follows:

1) Create a key file:

When you create a key file, the following table describes the file information:

LC for file information

07 File Identifier (2 bytes) File Type (1 byte) Permission 1 (1 byte) Permission 2 (1 byte) Len1 (1 byte) Len2 (1 byte)

Permission 1 indicates the permission to add a new key, and permission 2 indicates how to install the key. Permission 2 is set to 80h, which means that the installation is in ciphertext.

- 2) Use the application master key of the previous layer (SmartCOS 3.2 specifies: the external authentication key with the key identifier 01 is the application master key) to encrypt the key information (the key information is: plaintext key information length + key information + 80 00 00 00 00 00 00 00, and the encryption method of the key information is standard Triple DES or Single DES, if the key is 16 bytes, it is encrypted with Triple DES, if it is 8 bytes, It is encrypted with Single DES.
- 3) Generate MAC code, the initial value is: 4 bytes of random number + 00 00 00, the data generated MAC code is: 5 command headers + encrypted key information.
- 4) The ciphertext installation of other keys is encrypted with the application master key.
- 5) When installing the master key of the application under MF in ciphertext, use the transmission key of the card to install it.
- 2. Use security messages to write binary files
- 1) When creating a binary file, the fourth digit of the file type is 1, which indicates that a security packet is used.
- 2) LC is the number of bytes written + 4
- 3) Calculation of Security Messages (MAC):

The initial value is: 4 bytes of random number + '00 00 00 00'

The key is: the application maintenance key

Code value

CLA 04

INS D6

P1 Xx

P2 Xx

LC Write Data Length +4

DATA Write data + 4 bytes MAC

8. Unlock PIN Unblock personal password

The command packet is encoded as follows:

Code value

CLA 84

INS 24

P1 00

P2 01 - Unlock PIN

Lc 0C

DATA encrypts the PIN data element + Packet Authentication Code (MAC) data element, using the PIN unlock key.

Thereinto:

Encrypted PIN Token:

The PIN unlock key is encrypted and generated by PIN plaintext length + PIN + 80 00 (a multiple of 8).

MAC Generation:

Initial value: 4-byte random number + 00 00 00 00

Password: PIN unlock key

MAC address code generated data: CLA+INS+P1+P2+LC+encrypted PIN (8 bytes)

6. Use of CPUDEMO

7. Instructions for using the RD reader library

(1) int cpu_reset(int icdev,unsigned char *data_buffer);

NOTE: CPU card reset operation

Parameters: icdev: The device identifier returned by initialization

data_buffer: The string returned by the reset response

Returns: = 0 correct < 0 error (see error code)

(2) int cpu_protocol(int icdev,int len, unsigned char *send_cmd,

unsigned char *receive_data)

Note: A command string is sent to the reader based on the transmission protocol of the CPU card T=1

Parameters: icdev: The device identifier returned by initialization send_cmd: Command string for card operation (T=1 format)

len: The length of the command string

receive_data: The string of data returned by the CPU card

Returns: = 0 correct < 0 error (see error code)

(3) int cpu_comres(int icdev,int len, unsigned char *send_cmd,

unsigned char *receive_data)

Notes:

1) Regardless of whether it is a T=0 or T=1 card, the RD reader transmits commands in T=1 format.

The send_cmd of the command strings sent includes:

NAD+PCB+LEN+COMMAND+BCC

For T=0 cards, NAD and PCB can be filled with '00';

The return value receive_data structure is the same as Send_cmd.

- 2) cpu_comres is different from cpu_protocol ().
- a. When the upper-layer library receives the status byte SW1SW=61XX, it Cpu_protocol () function for further processing, and sends a Get Respond command to the reader, and returns the data and the status byte together. The cpu_comres() function does nothing and returns the state byte SW1SW2 directly.
- b. When the upper-layer library receives the state byte SW1SW=6CXX, it Cpu_protocol the () function for further processing, assigns XX to le and resends the last instruction, and returns the data and the state byte together; The cpu_comres() function does nothing and returns the state byte SW1SW2 directly. Part 1 CPU Basic Knowledge
- 1. Why use CPU card

IC card can be divided into contact IC card, contactless IC card and composite card from the interface mode. In terms of device technology, it can be divided into unencrypted memory cards, encrypted memory cards and CPU cards. The non-encrypted card has no security, and the data in the card can be arbitrarily rewritten, and the encrypted memory card adds a logical encryption circuit on the basis of the ordinary memory card to become an encrypted memory card. Because the logic encryption memory card

uses password control logic to control the access and rewrite of EEPROM, it needs to verify the password before it can be written, so it is safe for the chip itself, but it is not safe in the application. It has the following insecurity factors:

- 1. The password is transmitted in plaintext on the online road and is easy to be intercepted;
- 2. For system vendors, passwords and encryption algorithms are transparent.
- 3. The logical encryption card cannot authenticate whether the application is legitimate. For example, if someone fakes an ATM and you have no way of knowing its legitimacy, when you insert a credit card and enter your PIN, the password of the credit card is intercepted. Another example is INTENET online shopping, if the card is encrypted logically, the shopper is also unable to determine the legitimacy of the online store.

It is precisely because of the insecurity factors in the use of logical encryption cards that the development of CPU cards has been promoted. The CPU card can achieve the legitimacy authentication of the person, the card, and the system.

2. Three types of authentication of CPU card

CPU card has three authentication methods:

cardholder legitimacy authentication - PIN verification

Card legitimacy authentication - internal authentication

System legitimacy authentication - external authentication

Cardholder legitimacy authentication

The process of verification by entering a personal order by the cardholder.

System legitimacy authentication (external authentication) process:

System card,

send random number X

[with the specified algorithm and key] to encrypt

the random number

[with the specified algorithm and key] to decrypt Y, and obtain the result ${\sf Z}$

Compare X,Z, If it is the same, it means that the system is legitimate;

Card legitimacy authentication (internal authentication) process:

System card

sends random number X

with the specified algorithm and key to encrypt

the random number

[with the specified algorithm and key] to decrypt Y to obtain the result Z

Compare X, Z, if the same means that the card is legitimate;

In the above authentication process, the key does not appear in plaintext on the line, it is encrypted by random numbers every time it is sent, and because there is a random number to participate, it is ensured that the content of each transmission is different. If intercepted, it doesn't make any sense. This is not only a password-to-password authentication, but also a method authentication method, just like the cipher telegram used in the early days in the army, the sender encrypts the message into ciphertext according to a certain method and sends it, and then the receiver receives it and decrypts the ciphertext according to a certain method.

Through this authentication method, there is no attack point on the line, and the card can also verify the legitimacy of the application.

However, because the key and algorithm used by the system side for authentication are in the application, it still cannot remove the aggressiveness of the system vendor.

Here, we introduce the concept of SAM card.

SAM card is a CPU card with special performance, which is used to store keys and encryption algorithms, and can complete mutual authentication, password verification, encryption and decryption operations in transactions, and is generally used as an identity mark.

Thanks to the advent of SAM cards, we have a more complete system solution.

When issuing a card, we store the master key into the SAM card, and then the master key in the SAM card encrypts the characteristic byte of the user card (such as the application serial number) to generate a sub-key, and injects the sub-key into the user card. Due to the uniqueness of the application serial number, the sub-key within each user card is different.

Once the key is injected into the card, it does not appear outside the card. When in use, the master key of the SAM card generates a sub-key and stores it in the RAM area to encrypt and decrypt data.

The above authentication process becomes the following form:

System legitimacy authentication (external authentication) process:

SAM card

system card sends random number X

SAM card generates a sub-key to encrypt random number

Decrypt Y, and get the result Z

Compare X, Z, if the same means that the system is legitimate;

Card legitimacy authentication (internal authentication) process:

SAM card system

card

sends random number X

with the specified algorithm, key] to encrypt

the random number

SAM card decrypts Y, and the result Z

compares X, Z, if it is the same, it means that the card is legitimate;

In this way, the key in the application is transferred to the SAM card, and the authentication becomes the card-card authentication, and the system vendor no longer has responsibility.

3. When the line protection

card transmits data to the outside world, if it is transmitted in plaintext, the data is easy to be loaded and analyzed. At the same time, it is also possible to tamper with the transmitted data, and to solve this problem, the CPU card provides a line protection function.

There are two types of line protection, one is to encrypt the transmitted data in DES and transmit it in ciphertext form to prevent interception and analysis. The second is to attach a MAC address (security message authentication code) to the transmitted data, which the receiver will first verify after receiving it, and receive it only after the verification is correct, so as to ensure the authenticity and integrity of the data.

Fourth, the hardware structure diagram

EEPROM is used to store user data; The ROM is used to store the COS operating system, and the RAM area is used to store the intermediate variables of the COS runtime. COS (chip operation system) is a chip operating system, similar to DOS and WINDWOS, and a CPU card without COS is as unusable as a PC without DOS and WINDOWS. COS is solidified into the ROM area by the chip supplier when the chip leaves the factory, and this process is called masking.

COS is the core part of the CPU card, and it implements the security of the CPU card together with the hardware.

Part 2 Introduction to

SmartCOS SmartCOS is a chip operating system independently developed by Minghua Company, which passed the certification of the People's Bank of China in June 1999. COS is mainly divided into four parts:

1. SMARTCOS FILE SYSTEM

CPU CARD IS MANAGED IN THE FORM OF FILESSMARTCOS SUPPORTS THE FOLLOWING FILE SYSTEMS.

1. Files can be divided into MF files, DF files, and EF files

MF: The master file is the root of the entire file system, which is unique and equivalent to the root directory;

DF: A dedicated file, equivalent to a subdirectory, which can be used to store all the files of an application. A DF can be an application, or multiple DFs can be used for the same application.

EF: A basic file that stores a variety of application data and management information.

2. EF is divided into two types from the storage content:

Security basic file: used to store the key, only one security basic file can be created in each directory, and the key file is not!t can be selected by file selection, and the key content cannot be read, but it can be used and modified when the conditions are met.

Working basic file: used to store the actual data of the application, the number and size are limited by space. Reads and writes when conditions are met.

3. Basic file structure The structure of

basic files can be divided into the following four types:

Binary files

Data is read and written in bytes, and the length of each read and write cannot exceed 110 bytes; Can be used to store out-of-order data.

Linear fixed-length record file:

Each record is of fixed length, and the record can be accessed by record number, and the record range does not exceed 254;

The length of each record does not exceed 110 bytes, and the key file is a linear fixed-length record file, and the length of each record is fixed at 25 bytes. It can be used to store regular, fixed-length data.

Linear Variable Length Record Files:

The length of each record can vary, but the maximum length cannot exceed 110

bytes, which can be accessed by the record number. $% \left(1\right) =\left(1\right) \left(1\right) \left($

Cyclic fixed-length record file structure:

It is equivalent to a ring record queue, which is stored according to the first-in, first-out principle, the latest record number is 1, the last written record number is 2, and so on, the earliest record is automatically overwritten after the record is full.

4. File structure diagram

EF and DF can be established under MF;

DF can no longer be created under DF, only EF can be created;

KEY file: used to control the creation, reading and writing of files under MF

Data file

Key file:

Used to control the creation and access of files under DF

Data files (such as wallet files, etc.

).

5. Calculation of file space The header file length of

MF is 10 bytes + the length of the file name (5-16 bytes)

The header file length of DF is 10 bytes + the length of the file name The space occupied by

the EF file:

Define the space of the record and loop record file =File header space (10 bytes) + number of records * record length The space of

the variable-length record structure file = the file header space (10 bytes) + the space requested when creating The space occupied by the

key file = the file header space (10 bytes) + the number of keys * 25 bytes The space of

the wallet file = the file header (10 bytes) + the file body (17 bytes

The space of the passbook file = file header (10 bytes) + file body (20 bytes)

The file system is established, so how to ensure the security of the file, the following describes the security system.

2. SMARTCOS security system

1. State machine is the security state:

refers to a security level that the card is currently in, with (0---F) 16 security states. Automatically set to 0 after reset, and the security status of the current application will be automatically cleared to 0 after being successfully selected or reset. A change in security status must be achieved through key authentication.

**Only PIN reconciliation and external authentication in the current directory can change the security status.

2. Security attributes, i.e., access

permissions, are specified when creating files.

It is a compartmental concept, for example, if the read permission of a file is XY, it means that the state machine (security state) M of the current application must meet the requirements of X= \(\lambda M = \lambda Y \). If the read permission is set to 2F, the current state machine (security state) reaches 2 or higher, and the file can be read.

3. The relationship between the key and the security state

Each key defines the subsequent state when it is established, that is, the security state that can be achieved after passing the key authentication.

4

. The relationship between security status and access rights

(security status)

The use of permissions ------ keys----- the subsequent status ----- file read and write limits

5. The security status is only valid in the current directory, once you select another application, the state machine will automatically jump to the minimum permission 0. Security is independent between directories.

4. Reset Response

Symbol Byte Content Content Explanation

TS 3B Forward Convention

TO 6C TB1 and TC1 exist, and the historical characters are 12

TB1 00 No additional programming voltage is required

TC1 02 requires 2 additional protection times

T1-TC XX Historical Characters

Specific Meanings of SMARTCOS Historical Characters:

Symbol Byte Content Content Interpretation

T1 XX The version number of SMARTCOS

T2 XX Card Status Byte

T3 86 Minghua IC Card Manufacturing Organization Identification Number

T4 38

T5-TC XX Card Unique Serial Number

Card Status Byte is described as follows:

B7 B6 B5 B4 B3 B2 B1 B0 Status

011 XXX XXX 001 XXX XXX XXX The card has been initialized and successful, the card is not initialized, the card initialization process is locked,

 $0\ 0\ 0\ 0\ 0\ X\ X$ the card is not personalized

0 0 1 0 X X X X The card personalization is not complete

0 1 1 0 X X X X The card is personalization successful

0 0 0 1 X X X X The card was not personalized, and the card was locked

0 1 1 1 1 X X X X The card was personalized, and the card was locked

There are concepts about initialization and personalization in the card state byte, and we look at these two concepts from the life cycle of the card, which generally includes the following parts:

chip chip manufacturer

Mask COS Chip Manufacturer

Encapsulated into a card Card Producer

Card Initialization (COS Enabled) Card Manufacturer Transfer Password Protection Card Personalization Card

Use of Cards Users of Cards

Recycling of

Cards Card Issuers

The

initialization process is to activate COS, define the COS version, and then use COS after this process.

The personalization process refers to the creation of files, the writing of user data, etc.

The card producer delivers the card to the issuer after initializationWhen used, it is protected by a transmission password.

The IC card must support either T=0 or T=1 protocols, but not both, while the terminal must support both T=0 and T=1 protocols.

T=0 communication protocol is an asynchronous half-duplex character transmission protocol;

T=1 communication protocol is an asynchronous half-duplex block transmission protocol;

In the ISO7816-3 standard, these two protocols are specified;

The protocol used by the IC card is specified in TD1, and if TD1 is not in the reset response message, it means that the protocol with T=0 is used for communication. After the reset answer, the IC card and the terminal communicate with the protocol specified by the IC card.

5. Instruction Analysis

Here, we take an e-wallet application as an example to explain the instructions of SmartCOS.

1. File structure:

The following file structure is defined in the PBoC specification:

1) Basic data file of public application of electronic passbook ED/e-wallet EP application

File structure:

File Identifier (SFI) '21' (decimal)

File Type Transparent

File Size 30

File Access Control Read = Free Overwrite = Security Information

Required Bytes Data Element Length

1-8 Issuer ID 8

9 App Type ID 1

10 App Version 1

11-20 App Serial Number 10

21-24 App Enabled Date 4

25-28 App Expiration Date 4

29-30 Issuer Custom FCI Data 2

2) Cardholder basic data file

for e-passbook ED/e-wallet EP application

File Identifier (SFI) '22' (decimal)

File type Transparent

File size 39

File access control Read = free to rewrite = security information

required Byte Data Element Length

1 Card Type Identifier 1

2 Bank Employee Identifier 1

3-22 Cardholder Name 20

23-38 Cardholder ID Number 16

39 Cardholder ID Type 1

3) Electronic Passbook ED Transaction Details File

Document Identifier (SFI) '24' (Decimal)

File Type Loop

File Access Control Read = PIN Protection

Overwrite = Not Allowed

Record Size

Bytes Element Length

- 1-2 ED or EP Online or Offline Transaction Serial Number
- 2 3-5 Overdraft Limit 3
- 6-9 Transaction Amount 4
- 10 Transaction Type Identifier 1
- 11-16 Terminal Number 6
- 17-20 Transaction Date (Terminal) 4
- 21-23 Transaction Time (Terminal) 3
- 2. The security design is as follows:
- 1) External authentication can be carried out after checking the password;
- 2) The external authentication key No. 01 is used to control the storage of the e-wallet;
- 3) The No. 02 external authentication key is used to control the modification of the basic file and the modification of the key;
- 4) You can make purchases after checking the password.

The following keys are installed in the KEY file:

KEY Type Identification KID Subsequent status of usage rights The function description of the KEY

OB 01 OF 1 Personal PIN for PIN verification

08 01 11 2 External authentication key, used for e-wallet storage

08 02 1F F External authentication key, used for modifying basic files and keys

01 01 22 No Loop Key, used to generate Loop MAC

00 02 01 No e-wallet consumption key, used to generate wallet consumption MAC

02 01 03 No TAC key, used to generate Cache, Consumption, Withdrawal, and TAC

for modifying overdraft limits 05 01 33 No application maintenance key, used to generate MAC for application lock, application unlock, card lock and read, update binary, and record commands

- 3. Instruction sequence:
- 1) The card issuance process, creating a file on the card and installing the key

[Create MF] APDU command: 80 (CLA) e0 (INS) 00 (P1) 00 (P2) 18 (Lc) ff ff (8 byte transfer code) ff (Establish security attributes of files under MF) 01 (short file identifier for directory files) 31 50 41 59 2e 53 59 53 2e 44 44 46 30 31 (name of the file created).

[Create DF] APDU commands: 80 (CLA), E0 (INS), 01 (P1), 00 (P2) 0D (LC message length), 2F 01 (file identifier), ff (create file permissions), 00 (COS reserved), A0 00 00 00 03 86 98 07 01 (ADF name).

Create File under DF APDU command: 80 (CLA) E0 (INS) 02 (P1) 00 (P2) 07 (File Information Length) 6F 02 (Key File File Identifier) 05 (File Type) FF (Add Permission for New Key) 00 (COS Retention) 09 (Number of Records) 19 (Record Length):

Install PIN [Write Key] APDU command: 80 (CLA), E8 (INS) 00 (P1) 00 (P2) 0A (key information length), 01 (key identifier), 01 (key version number), 00 (algorithm identification), 0B (key type), 0F (usage permission), 01 (subsequent status), 2F (modify permission), 33 (error counter), 12 34 (personal password).

Install the external authentication key (DEAK) [Write Key] APDU command: 80 (CLA), E8 (INS) 00 (P1) 00 (P2) 18 (key information length) 01 (key identifier) 01 (key version number) 00 (algorithm identifier) 08 (key type) 11 (usage rights) 02 (subsequent status) FF (modify permissions) 33 (error counter) XX XX XX (Key Content)

Install External Authentication Key (DEAK) [Write Key] APDU Command: 80 (CLA), E8 (INS) 00 (P1) 00 (P2) 18 (Key Information Length) 02 (Key Identifier) 01 (Key Version Number) 00 (Algorithm Identification) 08 (Key Type) 1F (Usage Permission) 0F (Subsequent Status) FF (Modify Permission) 33 (Error counter) XX XX XX XX (Key Contents).

. Install the consumption key DPK of the e-wallet EP. [Write Key] APDU commands: 80 (CLA), E8 (INS) 00 (P1) 00 (P2) 18 (key information length) 02 (key identifier) 01 (key version number) 00 (algorithm identifier) 00 (key type) 01 (usage rights) 00 (subsequent status) FF (modify permissions) 00 (error counter) XX XX (KEY CONTENT).

DTK [Write Key] APDU command: 80 (CLA), E8 (INS) 00 (P1) 00 (P2) 18 (key information length) 01 (key identifier) 01 (key version number) 00 (algorithm identifier) 07 (key type) 0F (usage permission) 00 (subsequent status) FF (modify permission) 00 (error counter) XX XX XX XX (Key Content)

Installing the Application Maintenance Key DAMK[Write Key] APDU command: 80 (CLA) E8 (INS) 00 (P1) 00 (P2) 18 (Key Information Length) 01 (Key Identifier) 01 (Key Version Number) 00 (Algorithm Identification) 05 (Key Type) 0F (Permissions) 00 (Subsequent Status) FF (Modify Permissions) 00 (Error Counter) XX XX XX (Key Contents).

[Create File] APDU command: 80 (CLA) E0 (INS) 02 (P1) 00 (P2) 07 (LC message length) 00 15 (file identifier) 00 (binary file type) 0F (read permission) FF (update permission) 00 1e (file length).

Create File APDU Command: 80 (CLA) E0 (INS) 02 (P1) 00 (P2) 07 (File Info Length) 00 16 (File Identifier) 00 (Binary File Type) 0F (Read Permission) FF (Update Permission) 00 27 (File Length).

Create File APDU command: 80 (CLA) E0 (INS) 02 (P1) 00 (P2) 07 (File Information Length) 00 18 (File Identifier) 03 (Circular Record File Type) 1F (Read Permission) 10 (Update Permission) 0a 17 (File Length).

Write a public application basic data file [Update Binary] APDU command: 00 (CLA), D6 (INS), 95 (P1), 00 (P2), 1E (message length), A0 00 00 00 00, 03 00 00 01 (issuer ID), 03 (application type ID), 01 (application version), 00 00 19 98 08, 15 00 00 00 01 (application serial number), 20 00 10 01 (application activation date), 20 02 12 31 (application expiration date) 55 66 (issuer-defined FCI data).

Create Wallet EP File [Create File] APDU Command: 80 (CLA) E0 (INS) 02 (P1) 00 (P2) 07 (Lc) 00 01 (File Identifier) 06 (File Type) 00 (Permission 1) 00 (Permission 2) 00 (LEN1) 00 (LEN2)

[Create End] APDU command: 80 (CLA) E0 (INS) 01 (P1) 01 (P2) 02 (File Identifier Length) 2F 01 (File Identifier).

MF [Create End] APDU Command: 80 (CLA) E0 (INS) 00 (P1) 01 (P2) 02 (File Identifier Length) 3F 00 (File Identifier)

2) Transaction Process

Here, let's take the captive process and consumption process as examples:

A. Trapping—— Deposit the amount into the

card Before depositing, you must proofread the personal population order and the external certification No. 01;

Select Application[Select File]APDU Command: 00 (CLA) A4 (INS) 00 (P1) 00 (P2) 02 (Length) 01 (File Identifier)

Verify PIN [Verify] APDU Command: 00 (CLA) 20 (INS) 00 (P1) 00(P2) 02(Length) 12 34(PIN)

External Authentication: [

Get Challenge] APDU Command: 00 (CLA) 84 (INS) 00 (P1) 00 (P2) 08

3DES encryption of random numbers with No. 01 external authentication key;

[External Authentication] APDU command: 00(CLA) 82(INS) 00(P1) 01

(P2) 08 (length) XX XX XX (encrypted random number)

Trapping:[

Initializing Trap [Initalize For Load] APDU Command: 80(CLA) 50(INS) 00(P1) 02(P2) 0B(Length) 01 (Key Identifier) 00 00 10 00 (Transaction Amount) 00 00 00 00 00 01 (Terminal Number)

If the initialization of the trap is successful, the content of the reply data field is as follows:

Description Length (bytes)

Old balance of e-passbook or e-wallet 4

E-passbook or e-wallet online transaction serial number 2

Key

Algorithm ID 1

Pseudorandom Number ICC 4

MAC1 4

3DES encryption is carried out with the content of the circled response (4-byte pseudo-random number ICC + 2-byte electronic passbook or e-wallet online transaction number +80 00) to generate the process key;

Trap [Credit For Load] APDU command: 80 (CLA) 52 (INS) 00 (P1) 00 (P2) 0B (LC) YY YY MM DD (transaction date) HH MM SS (Trading time) XX XX XX XX (MAC2).

Calculation of MAC2:

Initial value: 00 00 00 00 00 00 00 00 00 Key: The process key generated above Data

generated for MAC2: 00 00 10 00 (4-byte transaction amount) 02 (transaction type ID) 00 00 00 01 (6-byte terminal number) YY YY MM DD (Transaction Date) HH MM SS (Trading Hours)

If the deposit transaction is successful, the online transaction serial number of the e-wallet file is added to 1, the transaction amount is added to the balance of the e-wallet, and a record is added to the transaction details file.

B. Consumption process:

Verify PIN before consumption:

Verify PIN [Verify] APDU command: 00 (CLA) 20 (INS) 00 (P1) 00 (P2) 02 (length) 12 34 (PIN).

[Initialize For Purchase] APDU command: 80(CLA) 50(INS) 01(P1) 02(P2) 0B (length) 01 (consumption key identifier) 00 00 00 01 (consumption amount) 00 00 00 00 01 (terminal number)

If the initialization consumption is successful, the data field content is answered:

Description Length (bytes)

E-passbook or e-wallet old balance 4

E-passbook or e-wallet online transaction serial number 2

Overdraft limit 3

Key version number 1

Algorithm identifier 1

Pseudo-random number ICC 4

Consumption[Debit For Purchase]APDU Command: 80 (CLA) 54 (INS) 01 (P1) 00 (P2) 0F (LC)

YY YY YY YY (Terminal Transaction Serial Number) YY YY MM DD (Terminal Transaction Date) HH MM SS (Trading Hours) XX XX (MAC1)

Use the consumption/withdrawal key pair (4-byte pseudo-random number ICC + 2-byte electronic passbook or e-wallet online transaction serial number + the last 2 bytes of the terminal transaction serial number) to encrypt the process key;

Calculation of MAC1: Initial

value: 00 00 00 00 00 00 00 00 00 (8 bytes)

Key: Process key Data

to generate MAC code: 4 bytes of transaction amount + 1 byte of transaction type ID + 6 bytes of terminal number + 4 bytes of terminal transaction date +3 bytes terminal transaction time

If the consumption transaction is successful, the online transaction serial number of the e-wallet file is added by 1, the e-wallet subtracts the transaction amount, and a record is added to the transaction detail file.

6. Analysis of other instructions The

following instructions have many questions in the process of use, which are explained as follows: 1.

The ciphertext installation process of the key:

In version 3.2, the ciphertext installation of the key will be different, and its process is as follows:

1) Create a key file:

When creating a key file, the relevant file information is as follows:

LC File information

07 File identifier (2 bytes) File type (1 byte) Permission 1 (1 byte) Permission 2 (1 byte) Len1 (1 byte) Len2 (1 byte).

Permission 1 indicates the permission to add a new key, and permission 2 indicates how to install the key. Permission 2 is set to 80h, which means that the installation is in ciphertext.

- 2) Use the application master key of the previous layer (SmartCOS 3.2 stipulates that the external authentication key with key ID 01 is the application master key) to encrypt the key information (the key information is: plaintext key information length + key information + 80 00 00 00 00 00 00 00), and the encryption method of the key information is standard Triple DES or Single DES, if the key is 16 bytes, use Triple DES encryption, if it is 8 bytes, it is encrypted with Single DES.
- 3) Generate MAC code, the initial value is: 4 bytes of random number + 00 00 00, the data generated MAC code is: 5 command headers + encrypted key information.
- 4) The ciphertext installation of other keys is encrypted with the application master key.
- 5) When installing the master key of the application under MF in ciphertext, use the transmission key of the card to install it.
- 2. Use security packets to write binaries
- 1) When creating binaries, the fourth digit of the file type is 1, which means that security packets are used.
- 2) LC is the number of bytes written + 4
- 3) Calculation of security packets (MAC): The

initial value is: 4 bytes of random number + '00 00 00 00'

The key is: the code value of the application maintenance key

CLA 04

INS D6

P1 XX

P2 XX

Lc Write Data Length + 4

DATA Write Data + 4 Bytes MAC 8

and PIN The

unblock PIN unlock command packet is encoded as follows:

Code value CLA 84

INS 24

P1 00

P2 01 - Unlock PIN

Lc 0C

DATA encrypts the PIN data element + Packet Authentication Code (MAC) data element, using the PIN unlock key.

Wherein

Encrypted personal password data element:

PIN unlock key pair PIN plaintext length + PIN + 80 00 (make up for the multiple of 8) to encrypt the

generation of MAC:

Initial value: 4 bytes random number + 00 00

Password: PIN unlock key Data

generated MAC code: CLA+INS+P1+P2+LC+encrypted PIN (8 bytes)

6. Use

of CPUDEMO

7. Instructions for

using RD reader libraries (1) int cpu_reset(int icdev,unsigned char *data_buffer);

NOTE: CPU Card Reset Operation

Parameters: icdev: Device identifier returned by initialization

data_buffer: String returned by reset response

Returned: = 0 Correct < 0 Error (see Error Code).

(2) int cpu_protocol(int icdev, int len, unsigned char *send_cmd,

unsigned char *receive_data)

Note: T = the CPU card1's transport protocol sends a command string to the reader

Parameters: icdev: device identifier returned by initialization send_cmd: command string for card operations (T=1 format)

len: length of command string

receive_data:
Returns: = 0 Correct
< 0 False (see error code)

(3) int cpu_comres(int icdev, int len, unsigned char *send_cmd,

unsigned char *receive_ data)

Precautions:

1) Regardless of whether it is a T=0 or T=1 card, the RD reader transmits commands in T=1 format.

The command string send cmd sent includes:

NAD+PCB+LEN+COMMAND+BCC

For T=0 cards, NAD and PCB can be filled with '00';

The structure of the return value receive_data is the same as that of Send_cmd.

- 2) cpu_comres is different from cpu_protocol()
- a. When the upper-layer library receives the state byte SW1SW=61XX, it Cpu_protocol () function to do further processing, and sends a Get Respond command to the reader, and returns the data and the state byte together; The cpu_comres() function does nothing and returns the state byte SW1SW2 directly.
- b. When the upper-layer library receives the state byte SW1SW=6CXX, it Cpu_protocol () function for further processing, assigns XX to le, and then resends the last instruction, and returns the data and the state byte together; The cpu_comres() function does nothing and returns the state byte SW1SW2 directly.