

DESIGN & ANALYSIS OF ALGORITHM

CONTINUOUS ASSESSMENT - III

1. a) $n = 5$

$$\text{faces} = 5$$

$$\begin{aligned}\text{vertices} &= 2n - 5 \\ &= 2(5) - 5 \\ &= \boxed{5 \text{ vertices}}\end{aligned}$$

$$\begin{aligned}\text{edges} &= 3n - 6 \\ &= 3(5) - 6 \\ &= \boxed{9 \text{ edges}}\end{aligned}$$

b) ~~Fortune's~~ ^{Fortune's} algorithm is a sweep line algorithm for generating a Voronoi.

b. Space Complexity by Voronoi:

- * The Number of Voronoi vertices is $O(n)$ (i.e. $2n - 5$)
- * The Number of Voronoi edges is $O(n)$ (i.e. $3n - 6$)
- * Size of Voronoi Diagram is $O(n)$

Time complexity of Voronoi is $\boxed{O(n \log n)}$

* Formatti's algorithm is a sweep line algorithm for generating a Voronoi diagram from a set of points in a plane using

$O(n \log n)$ and space is $O(n)$

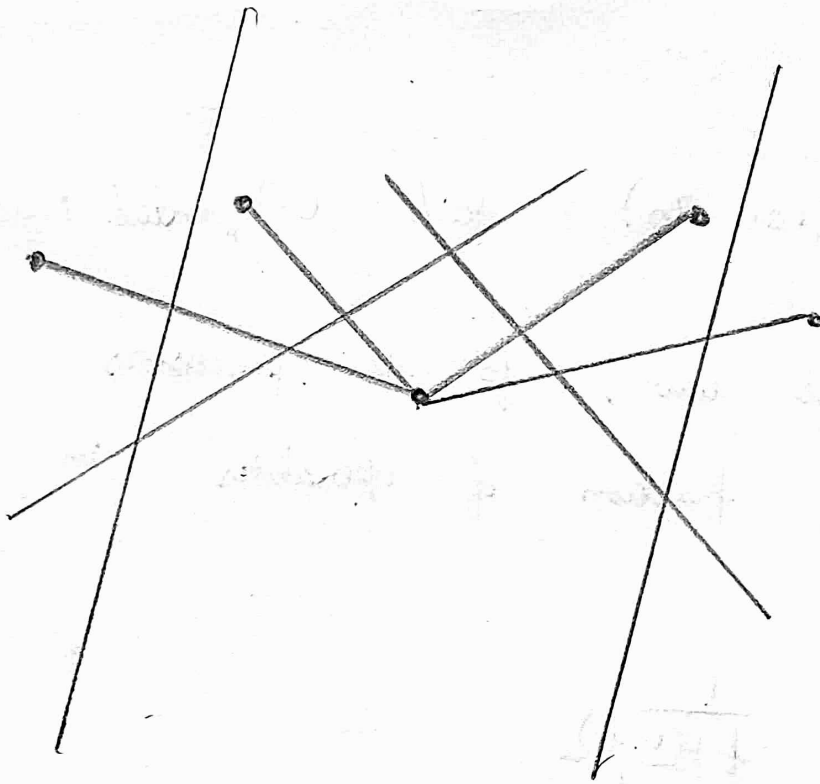
* The number of vertices, edges, faces of the 3d Voronoi - Diagram of n point can be as worst $O(n^2)$

* ~~The Naive~~ In Naive Algorithm point, the complexity is $O(n^2 \log n)$ i.e. half plane of the points $p_j \neq p_i$

* The instrumental time complexity is $O(mn)$

* m is no. of levels and n is no. of sites and space complexity is $O(n)$

1c. for $n = 5$



2. Given $p = 12$

$$f = (100 - 70) = 30\% \quad (\text{Sequential Fraction})$$

By Amdahl's law, for p processors and ~~fractions~~ fraction of operations in sequential computation

$$a) \quad \text{speedup} = \frac{1}{f + \frac{(1-f)}{p}}$$

max speedup when $p \rightarrow \infty$

$$\text{Here max speedup} = \frac{1}{f} = \frac{1}{0.3}$$

$$\boxed{\text{Speedup} = 3.33}$$

b) Speedup for $p = 12$:

$$\Rightarrow \frac{1}{f + \frac{1-f}{p}} = \frac{1}{0.30 + \frac{0.70}{12}}$$

$$= \frac{1}{0.358}$$

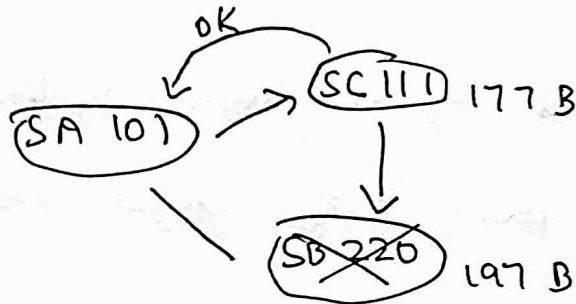
$$\Rightarrow \boxed{2.79}$$

3.

(SA 101) → active

Let us assume SA 101 detects the failure (active)

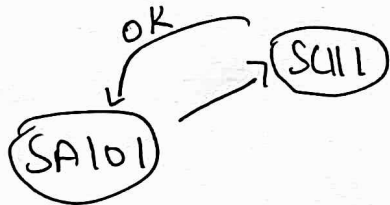
SA 101 sends message SC111 and SB 220



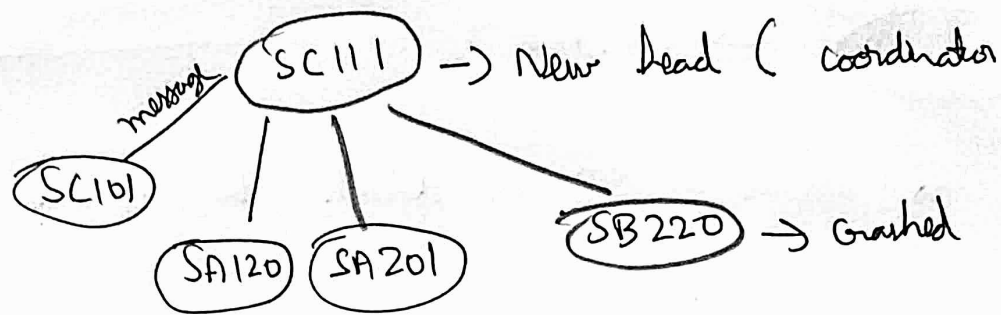
→ SB 220 is crashed

It sends a reply to SC111

→ SC111 sends a response to SA101



→ SC111 becomes the leader and sends a message to all server



a) Worst Case:

* $N-1$ processes all together begin elections and sending messages to process with higher ids

$$\text{message overheads} = O(N^2)$$

* Turn around time is approximately 5 message transmission time

b) Best case:

• $N-2$ coordinator message are sent

• Turn around time is one message transmission

time. Time complexity: $O(n)$

* Bully Algorithm:

When a process finds the coordinator has failed, if it knows its id is the highest, it elects itself as a coordinator, then sends a coordinator message to all processes with lower identifiers than itself.