#### Datomic for the 96 Percent

# Cognitect



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# challenges with SQL

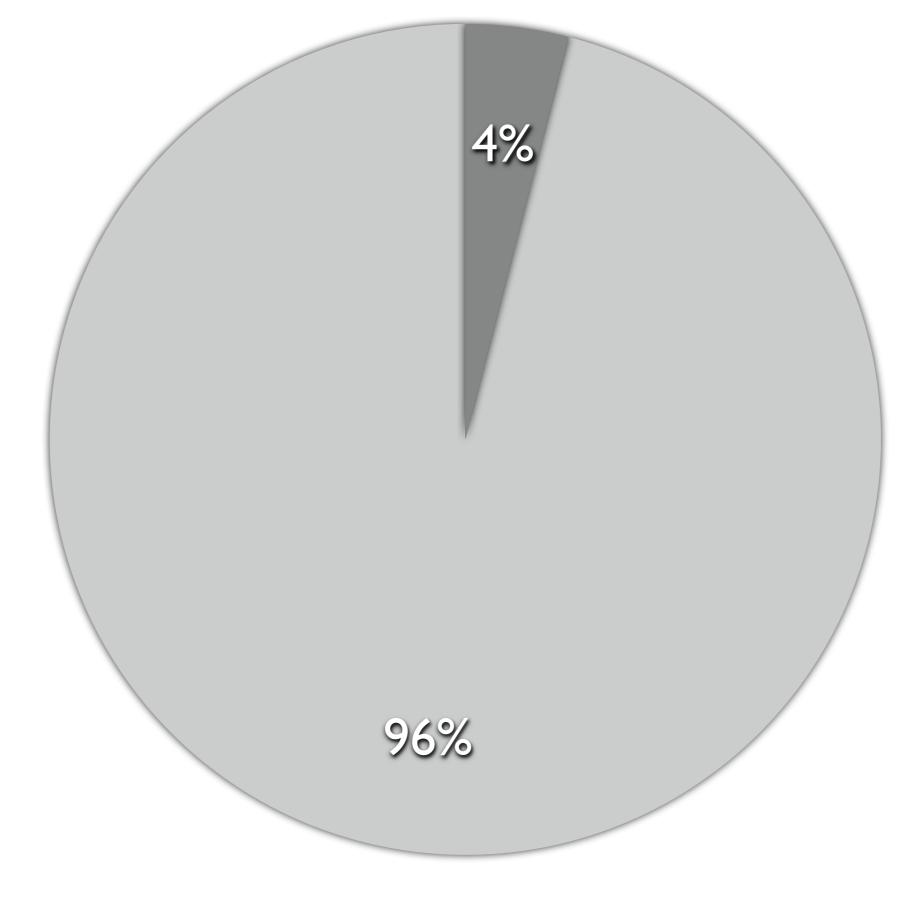
write volumes

read volumes

deployment rigidity

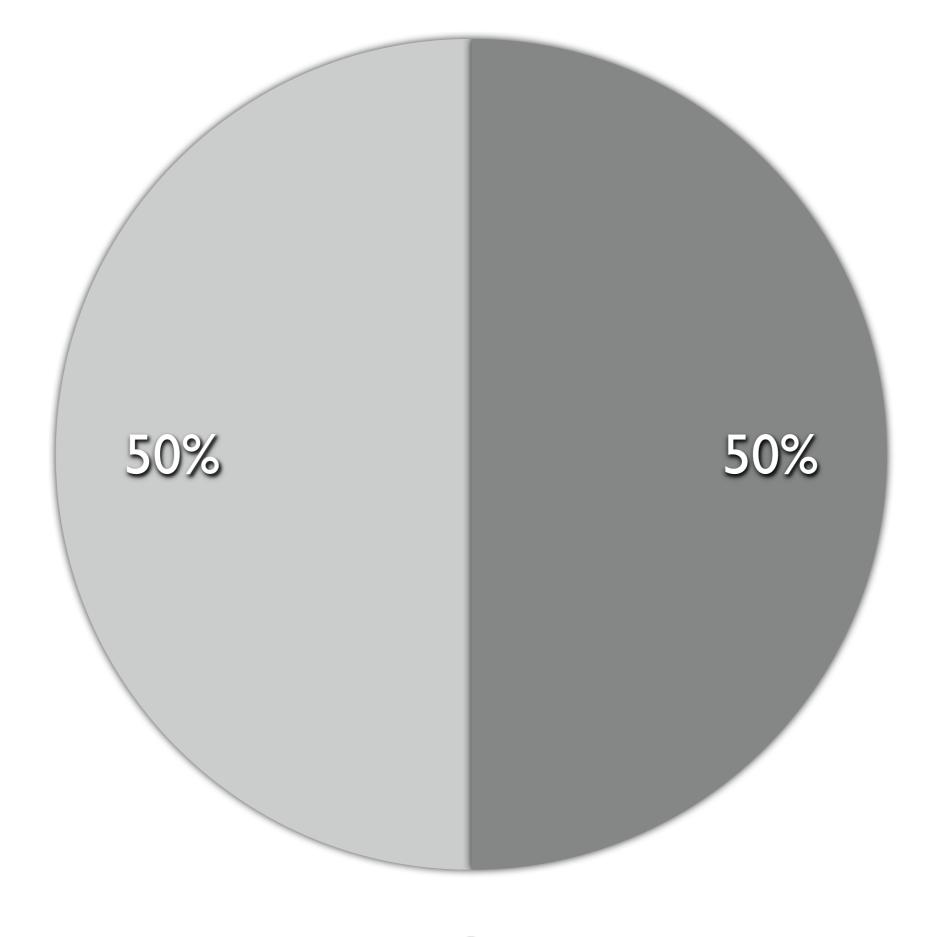
model rigidity

update-in-place



Web Scale

Not So Much



Accurate

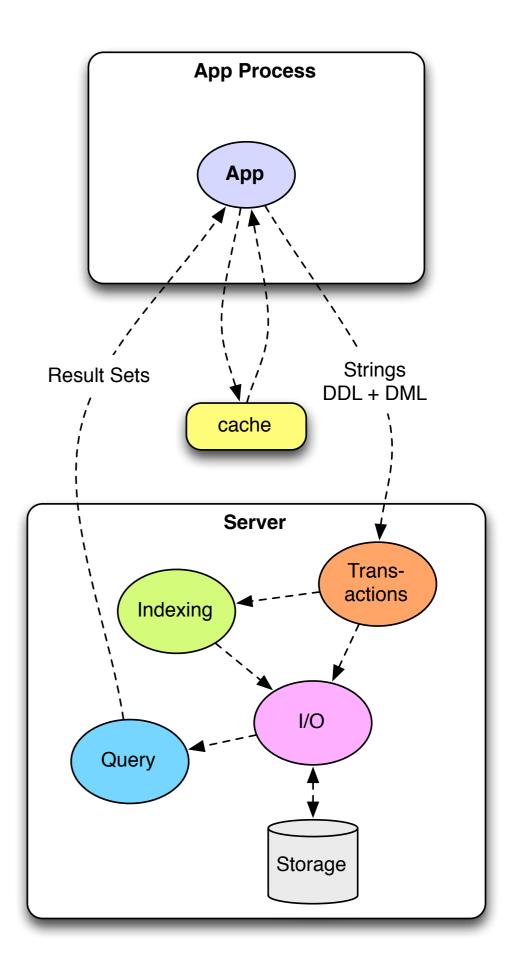
Entirely Ficititious

# eventual consistency



http://en.wikipedia.org/wiki/Everclear\_(alcohol)

# rigid deployment



# rigid models

"People can belong to multiple clubs"

join table person table

club table

id key in person table
person key in join table
club key in join table
id key in club table

#### the laws

memory is expensive

storage is expensive

machines are precious

resources are dedicated

# update in place = transience

characteristic	transient structure
sharing	difficult
distribution	difficult
concurrent access	difficult
access pattern	eager
caching	difficult
examples	Java and .NET collections relational databases NoSQL databases



# answering the challenge

functional

radical deployment, e.g. local logic query, laziness

isolated writes, serialized ACID

time-aware

elastic read scaling

flexible, universal attribute schema

programmable

# transience vs. persistence

characteristic	transient	persistent
sharing	difficult	trivial
distribution	difficult	easy
concurrent access	difficult	trivial
access pattern	eager	eager or lazy
caching	difficult	easy
Java, .NET collections examples relational databases  NoSQL databases		Clojure, F# collections Datomic database

```
Connection conn =
connect("datomic:ddb://us-east-1/mb/mbrainz");
Database db = conn.db();
Set results = q(..., db);
Set crossDbResults = q(..., db1, db2);
Entity e = db.entity(42);
```

```
pluggable storage
protocol
Connection conn =
connect("datomic:ddb://us-east-1/mb/mbrainz");
Database db = conn.db();
Set results = q(..., db);
Set crossDbResults = q(..., db1, db2);
Entity e = db.entity(42);
```

```
Connection conn =
connect("datomic:ddb://us-east-1/mb/mbrainz");
Database db = conn.db(); database is a lazily
                                  realized value, available
                                     to all peers equally
Set results = q(..., db);
Set crossDbResults = q(..., db1, db2);
Entity e = db.entity(42);
```

```
Connection conn =
connect("datomic:ddb://us-east-1/mb/mbrainz");
Database db = conn.db();
                                  query databases,
Set results = q(..., db)
                                   not connections
Set crossDbResults = q(..., db1, db2);
Entity e = db.entity(42);
```

```
Connection conn =
connect("datomic:ddb://us-east-1/mb/mbrainz");
Database db = conn.db();
Set results = q(..., db);
Set crossDbResults = q(..., db1, db2);
Entity e = db.entity(42);
                            join across databases,
                        systems, in-memory collections
```

```
Connection conn =
connect("datomic:ddb://us-east-1/mb/mbrainz");
Database db = conn.db();
Set results = q(..., db);
Set crossDbResults = q(..., db1, db2);
Entity e = db.entity(42);
                            lazy, associative
                             navigable value
```

```
List newData = ...;
Future<Map> f = conn.transactAsync(list);
dbBefore = conn.db.asOf(time);
possibleFuture = db.with(...);
allTime = db.history();
BlockingQueue<Map> queue = conn.txReportQueue();
Log log = conn.log();
Iterable<Map> it = log.txRange(startOfMonth, null);
```

```
List newData = ...;
Future<Map> f = conn.transactAsync(list);
dbBefore = conn.db.asOf(time);
                                         information in
possibleFuture = db.with(...);
                                      generic data structures
allTime = db.history();
BlockingQueue<Map> queue = conn.txReportQueue();
Log log = conn.log();
Iterable<Map> it = log.txRange(startOfMonth, null);
```

```
contains old db, new db, change
List newData 7 ....
Future<Map> f = conn.transactAsync(list);
dbBefore = conn.db.asOf(time);
possibleFuture = db.with(...);
allTime = db.history();
BlockingQueue<Map> queue = conn.txReportQueue();
Log log = conn.log();
Iterable<Map> it = log.txRange(startOfMonth, null);
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List newData = ...;
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BlockingQueue<Map> queue = conn.txReportQueue();
Log log = conn.log();
Iterable<Map> it = log.txRange(startOfMonth, null);
```

```
List newData = ...;
Future<Map> f = conn.transactAsync(list);
dbBefore = conn.db.asOf(time);
possibleFuture = db.with(...);

one possible future
allTime = db.history();
BlockingQueue<Map> queue = conn.txReportQueue();
Log log = conn.log();
Iterable<Map> it = log.txRange(startOfMonth, null);
```

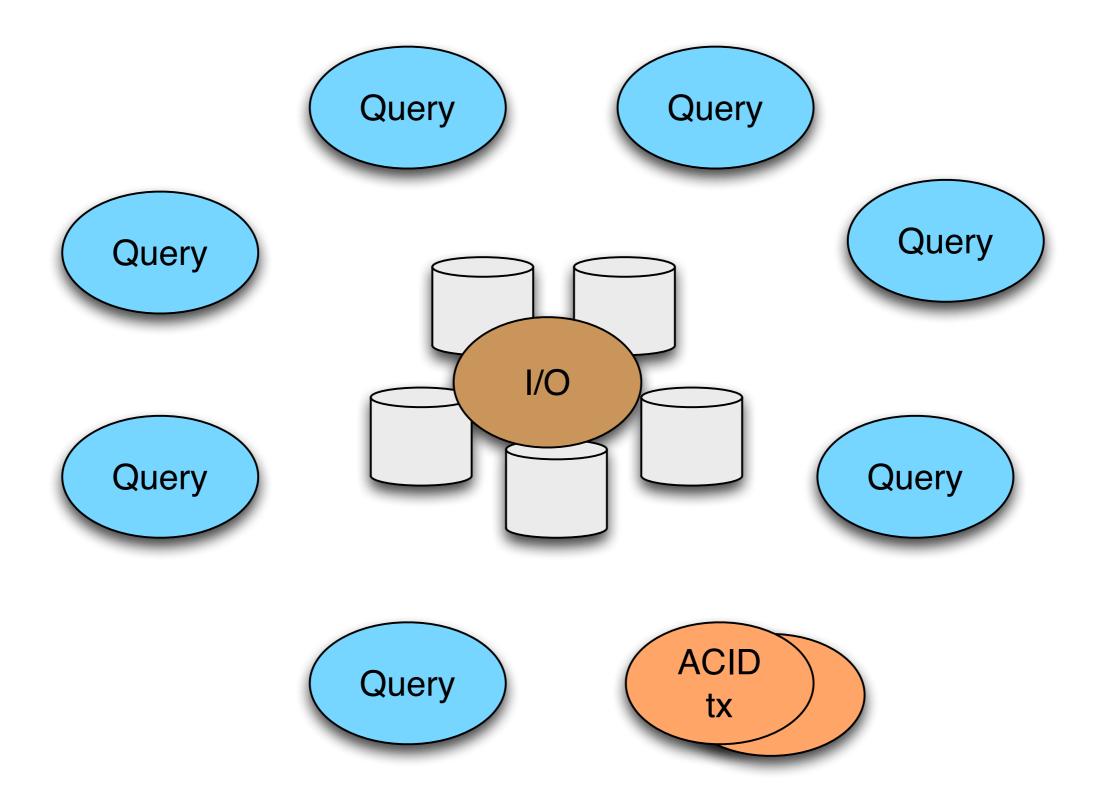
```
List newData = ...;
Future<Map> f = conn.transactAsync(list);
dbBefore = conn.db.asOf(time);
possibleFuture = db.with(...);
allTime = db.history();

all history, overlapped
BlockingQueue<Map> queue = conn.txReportQueue();
Log log = conn.log();
Iterable<Map> it = log.txRange(startOfMonth, null);
```

```
List newData = ...;
Future<Map> f = conn.transactAsync(list);
dbBefore = conn.db.asOf(time);
possibleFuture = db.with(...);
                                 monitor all change
                                   from any peer
allTime = db.history();
BlockingQueue<Map> queue = conn.txReportQueue();
Log log = conn.log();
Iterable<Map> it = log.txRange(startOfMonth, null);
```

```
List newData = ...;
Future<Map> f = conn.transactAsync(list);
dbBefore = conn.db.asOf(time);
possibleFuture = db.with(...);
allTime = db.history();
BlockingQueue<Map> queue = conn.txReportQueue();
Log log = conn.log();
Iterable<Map> it  log.txRange(startOfMonth, null);
                      review any
                       time range
```

# elastic query scaling



# query

# why datalog?

Equivalent to Relational Model + Recursion

Better fit than Prolog for query

No clause order dependency

Guaranteed termination

Pattern-matching style easy to learn

# Example Database

entity	attribute	value
42	:email	jdoe@example.com
43	:email	jane@example.com
42	:orders	107
42	:orders	141

Constrains the results returned, binds variables

[?customer :email ?email]

Constrains the results returned, binds variables

Constrains the results returned, binds variables

```
constant

[?customer :email ?email]
```

Constrains the results returned, binds variables

```
variable variable

| (?customer :email ?email)
```

entity	attribute	value
42	:email	jdoe@example.com
43	:email	jane@example.com
42	:orders	107
42	:orders	141

[?customer :email ?email]

# Constants Anywhere

"Find a particular customer's email"

```
[42 :email ?email]
```

entity	attribute	value
42	:email	jdoe@example.com
43	:email	jane@example.com
42	:orders	107
42	:orders	141

[42 :email ?email]

### Variables Anywhere

"What attributes does customer 42 have?

[42 ?attribute]

entity	attribute	value
42	:email	jdoe@example.com
43	:email	jane@example.com
42	:orders	107
42	:orders	141

#### [42 ?attribute]

### Variables Anywhere

"What attributes and values does customer 42 have?

[42 ?attribute ?value]

entity	attribute	value
42	:email	jdoe@example.com
43	:email	jane@example.com
42	:orders	107
42	:orders	141

#### [42 ?attribute ?value]

#### Where Clause

```
data pattern

[:find ?customer :email]]
```

#### Find Clause

```
variable to return

[:find ?customer
:where [?customer :email]]
```

### Implicit Join

"Find all the customers who have placed orders."

#### **API**

#### q

### Query

### Input(s)

#### In Clause

Names inputs so you can refer to them elsewhere in the query

:in \$database ?email

### Parameterized Query

```
q([:find ?customer
    :in $database ?email
    :where [$database ?customer :email ?email]],
    db,
    "jdoe@example.com");
```

### First Input

```
q([:find ?customer
    :in $database ?email
    :where [$database ?customer :email ?email]],
    db,
    "jdoe@example.com");
```

### Second Input

```
q([:find ?customer
    :in $database ?email
    :where [$database ?customer :email ?email]],
    db,
    "jdoe@example.com");
```

#### Verbose?

```
q([:find ?customer
    :in $database ?email
    :where [$database ?customer :email ?email]],
    db,
    "jdoe@example.com");
```

#### Shortest Name Possible

```
q([:find ?customer
    :in $ ?email
    :where [$ ?customer :email ?email]],
    db,
    "jdoe@example.com");
```

### Elide \$ in Where

```
q([:find ?customer
    :in $ ?email
    :where [ ?customer :email ?email]],
    db,
    "jdoe@example.com");
    no need to
        specify $
```

#### Predicates

Functional constraints that can appear in a :where clause

### Adding a Predicate

"Find the expensive items"

#### **Functions**

Take bound variables as inputs and bind variables with output

```
[(shipping ?zip ?weight) ?cost]
```

# Function Args

```
[(shipping ?zip ?weight) ?cost]

bound inputs

[(shipping ?zip ?weight) ?cost]
```

#### Function Returns

```
[(shipping ?zip ?weight) ?cost]

bind return
values
```

#### **BYO Functions**

Functions can be plain JVM code.

```
public class Shipping {
  public static BigDecimal
  estimate(String zip1, int pounds);
}
```

#### entities

 maplike, point-in-time view of datoms sharing a common e

```
{:db/id 42
:likes "pizza"
:firstName "John"
:lastName "Doe"}
entity

datoms

[42 :likes "pizza"]
       [42 :firstName "John"]
       [42 :lastName "Doe"]
```

#### entities

transformation is purely mechanical

```
{:db/id 42

:likes "pizza"

:firstName "John"

:lastName "Doe"}
```

```
[42 :likes "pizza"]
[42 :firstName "John"]
[42 :lastName "Doe"]
```

### one database, many indexes

structure	attribute
row	EAVT
column	AEVT
document	EAVT, partitions, components
graph	VAET

### transactions

# ids and partitions

# built-in partitions

partition	usage
:db.part/db	schema entities
:db.part/tx	transaction entities
:db.part/user	user entities

### create your own partitions

group related entities in a partition

coarser granularity than e.g. tables

partition is a hint to indexing

group these things together

can help locality

does not affect semantics

#### creating partitions

```
[{:db.install/_partition :db.part/db,
    :db/id #db/id[:db.part/db],
    :db/ident :inventory}
{:db.install/_partition :db.part/db,
    :db/id #db/id[:db.part/db],
    :db/ident :customers}]
```

# uniqueness

# uniqueness

requirement	model with	value types
db-relative opaque id	entity id	opaque (long)
external id	:db.unique/identity attribute	string, uuid, uri
global opaque id	:db.unique/identity squuid	uuid
programmatic name	:db/ident	keyword

## squuids

semi-sequential UUIDs

do not fragment indexes

```
public class Peer;
    public static UUID squuid();
    public static long squuidTimeMillis(UUID squuid);
    // other methods elided for brevity
}
```

## transaction functions

#### add and retract

```
[[:db/add john :likes pizza]
[:db/retract john :likes iceCream]]
```

#### what about update?

```
[[:db/add john :likes pizza]
[:db/retract john :likes iceCream]
[:db/add john :balance 110?]]
```

#### atomic increment

```
[[:db/add john :likes pizza]
[:db/retract john :likes iceCream]
[:inc john :account 10]]
```

#### transaction fns

subset of data fns

run inside transactions

have access to in-tx value of database

as first argument

#### tx function expansion

```
[[:db/add john :likes pizza]
[:db/retract john :likes iceCream]
[:inc john :balance 10]]

[[:db/add john :likes pizza]
[:db/retract john :likes iceCream]
[:db/add john :balance 110]]
```

# lookup the function

```
[:inc john :balance 10]

value

inc = db.entity("inc").get("db/fn");
```

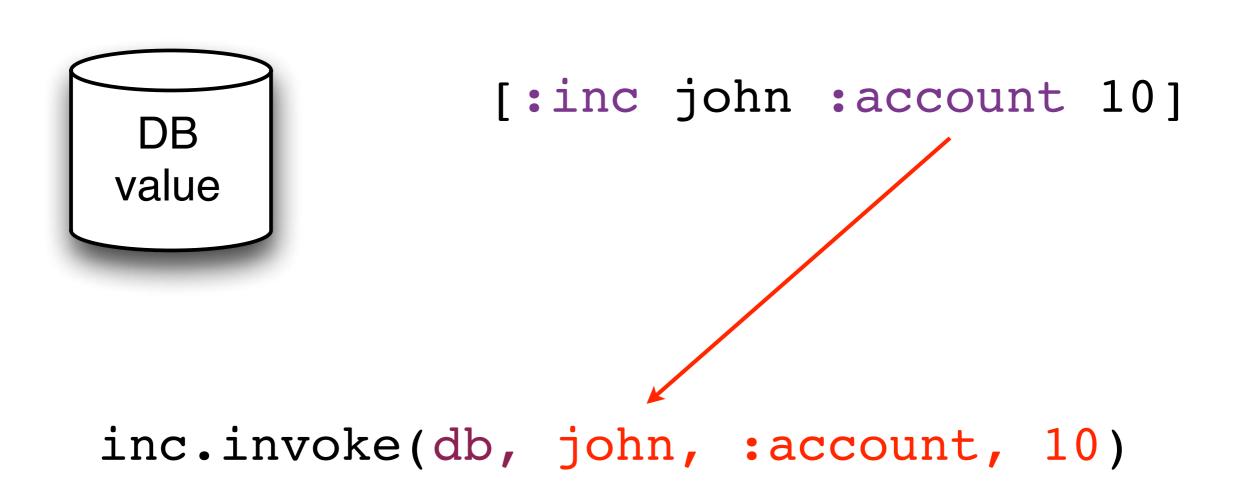
### pass in current db

```
[:inc john :account 10]

value

inc.invoke(db, ...);
```

### pass in args



#### data out

```
DB value [:inc john :account 10]
```

#### inc

```
public static Object inc(Object db, Object e, Object amount)
{
    // lookup entity
    // calculate new balance
    // create assertion
    // return list containing assertion
}
```

#### inc

```
public static Object inc(Object db, Object e, Object a, Object amount) {
    Entity ent = ((Database)db).entity(e);
    Long balance = (Long) ent.get(a) + (Long) amount;
    List updated = list("db/add", e, a, balance);
    return list(updated);
}
```

# modeling

#### modeling rigiditiy

"People can belong to multiple clubs"

join table
person table
club table

id key in person table
person key in join table
club key in join table
id key in club table

#### universal relation

"People can belong to multiple clubs"

[?person :club ?club]

#### stories

attribute	type	cardinality
story/title	string	
story/url	string	
story/slug	string	
news/comments	ref	many

### schema is plain old data

attribute	type	card
story/title	string	I
story/url	string	I
story/slug	string	I
news/comments	ref	many

```
{:db/id #db/id[:db.part/db]
  :db/ident :story/url
  :db/valueType :db.type/string
  :db/cardinality :db.cardinality/one
  :db.install/_attribute :db.part/db}
```

#### users

attribute	type	cardinality
user/firstName	string	
user/lastName	string	
user/email*	string	
user/upVotes	ref	many

\*unique

#### cardinality many

```
[:db/add 42 :upvotes 11]
[:db/add 42 :upvotes 12]
```

#### entities

```
john = db.entity(42);
john.get("user/upVotes").size();
```

#### comments

attribute	type	cardinality
comment/body	string	
comment/author	ref	
news/comments	ref	many

### types do not dictate attrs

attribute	type	cardinality
story/title	string	
story/url	string	I
story/slug	string	I
news/comments	ref	many

attribute	type	cardinality
comment/body	string	l
comment/author	ref	I
news/comments	ref	many

#### relation direction

```
// get child comments
comment.get("news/comments");
```

#### reversing direction

```
;; base case
[(story-comment ?story ?comment)
  [?story :story/title]
  [?story :new/comments ?comment]]

it is a story comment if...
```

```
;; base case
[(story-comment ?story ?comment)
    [?story :story/title]
    [?story :new/comments ?comment]]
    ... there is a story ...
```

```
;; base case
[(story-comment ?story ?comment)
  [?story :story/title]
  [?story :new/comments ?comment]]
  ... with a comment
```

```
;; recursion
[(story-comment ?story ?comment)
  [?parent :news/comments ?comment)
  (story-comment ?story ?parent)]
```

or, it is a story comment if...

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```
;; recursion
[(story-comment ?story ?comment)
  [?parent :news/comments ?comment]
  (story-comment ?story ?parent)]

... it has a parent comment ...
```

# documents

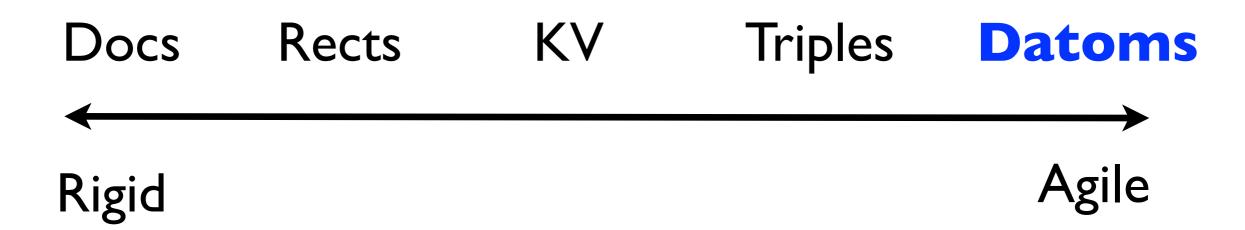
# activity "document"

```
// get upvotes
john.get("user/upVotes");
// get title of an upvoted story
anUpvote.get("story/title");
// get John's comments
john.get("comment/_author");
```

## profile "document"

```
// get facts about John
john.get("user/email");
john.get("user/firstName");
```

# agility



# leverage

KV Docs Rects Triples Datoms
Low
High

# complexities mitigated

Lost Data Managing Time Eventual Consistency

Log Analysis Test Setup DAOs

Defensive Copying DTOs

**ORM** 

Join Tables Objects

Inheritance Relationship Direction Strings

Structural Rigidity Logic Duplication String Injection

Model Caching Imperative Code Data Duplication

Read Transactions Isolation Levels
App Caching

Denormalization

## programmability

Make a column name variable?

Make a table name variable?

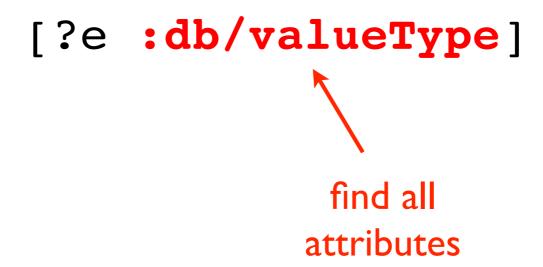
Treat metadata as first-class data?

### first-class attributes

[?person ?attr ?value]

attribute slot
isn't special

### schema made of ordinary data



### user stories





http://thinkrelevance.com/blog/2013/06/12/kurt-zimmer-of-room-key-podcast-episode-033 https://github.com/candera/strangeloop-2013-datomic/blob/master/slides.org

"We use Datomic as an event-source data store and it works wonderfully!"



"Because of the elasticity of Datomic, we were able to reduce our hosting fees by a factor of 10 when we moved off of [a popular NoSQL]."

<redacted>

### resources

#### **Datomic**

http://www.datomic.com/http://blog.datomic.com/2013/06/using-datomic-from-groovy-part-1.htmlhttp://blog.datomic.com/2013/05/a-whirlwind-tour-of-datomic-query\_16.htmlhttps://github.com/datomic/day-of-datomic

#### **Stuart Halloway**

https://github.com/stuarthalloway/presentations/wiki. http://www.linkedin.com/pub/stu-halloway/0/110/543/ https://twitter.com/stuarthalloway mailto:stu@cognitect.com