

Assignment 3 – MSCI 240

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1. Introduction: A brief summary of the whole report

This report describes a high-level description and the simplified complexity time analysis for the methods in the LinkedList class.

2. A brief description and overview of the class

Class name: LinkedList

The LinkedList class is a data structure that will manage a linked list of ListNode objects. The class has 1 private field and 1 public field, and different methods that are helpful to manipulate and retrieve information from the linked list. The methods include 4 accessors, 4 mutators, and 1 sorting function.

The bubble sort algorithm was used because is a straightforward and easy to understand algorithm. The implementation is concise and did not require a lot of lines of code. It was also adaptable to sorting linked lists from the array list one.

Class Structure:

1. Constructor:

- **__init__(self):** Creates a new empty linked list and initialized private fields:
 - o self.head (public field) : head of the linked list.
 - o self._length (private field): length of the list.

2. Accessors:

- **def length(self) -> int:** Returns the number of elements stores in the linked list, representing the length of the linked list.
 - o Returns: self._length.
- **def contains(self, value, map_to_key=lambda x: x) -> bool:** This method checks if a specific value or mapped key belongs to the linked list.

This function calls the `find_index_byvalue` method, which returns the index of a value if in the linked list.

- Returns **True** if value is in linked list.
- Returns **False** if value is not in linked list.
- **def find_index_byvalue(self, value, map_to_key=lambda x: x) -> int:** This method iterates through the linked list and returns the index of the value if this or its mapped key is in the linked list.
 - Returns the first index of a provided value or its mapped key in the linked list.
 - Returns -1 if not found.
- **def get(self, k: int) -> object:** This method returns the element stored at an specified index k in the linked list. For this, the method iterates through the linked list by until current reaches the kth node.
 - Raises an **IndexError** if the index is out of range.
 - Returns `current.data` (element at index k)

3. Mutators:

- **def set(self, k: int, value: object):** This method sets the value of the kth node in the linked list to the provided new value.
 - Raises an **IndexError** if the index is out of range.
 - Raises an **IndexError** if k is out of range.
 - Sets the data of the kth node to the value.
- **def append(self, value: object):** Adds a new `ListNode` with the provided (obj) to the end of the linked list, storing the provided value.
 - Creates a new list node: `ListNode(obj)`
 - If the list is empty, it sets the new node as the head of the list.
 - If the list is not empty, it iterates through the list until it reaches the last node and sets the next node to the new node.
 - It increments the length of the list.

- **def insert(self, k: int, value: object):** Inserts a new list node with the provided value:object as the kth node in the linked list, shifting subsequent values rightward.
 - Raises an **IndexError** if the index is out of range.
 - If inserting at index 0, it sets the new node as the head and adjusts the next pointer.
 - If index not at 0, it iterates through the list to the node before the specified index and inserts the new node.
 - Increments the length of the list
- **def remove(self, k: int) -> object:** Removes the kth list node in the linked list and returns its value.
 - Raises an **IndexError** if the index is out of range.
 - If removing at index 0, it sets the head to the next node.
 - If not at index 0, iterates through the list to the node before the specified index and adjusts the next pointer to skip the node being removed.
 - Reduces the length of the list.
 - Returns the removed value

4. Sorting:

- **def sort(self, map_to_key=lambda x: x, descending=False):** Sorts the list in-place using the bubble sort algorithm.
 - Returns sorted list.

3. Class Implementation Description:

a. Class Diagram

| Class | <i>LinkedList</i> | |
|-----------|-------------------|----------------------|
| Fields | Constructor | head |
| | | _length |
| Behaviour | Accessors | length() |
| | | contains() |
| | | find_index_byvalue() |
| | | get() |
| | Mutators | set() |
| | | append() |

| | | |
|--|---------|----------|
| | | insert() |
| | | remove() |
| | Sorting | sort() |

4. Simplified Time Analysis

Constructor:

```
def __init__(self):
    """Create an empty list."""
    self.head = None # empty list with no head
    self._length = 0 # empty list with 0 length
```

- This has 2 statements.
 - 2 assigning variables.
- Total:
 - **$T(n) = 2$**
 - **Big O = $O(1)$**
- The constructor has constant time complexity $O(1)$ since the algorithm is only assigning variables.

Accessors:

- **def length(self) -> int:**

```
def length(self) -> int:
    """Return number of elements stored in the linked list."""
    return self._length
```

- This has only 1 returning statement.
- Total:
 - **$T(n) = 1$**
 - **Big O = $O(1)$**
- The length method has constant time complexity $O(1)$ since the algorithm is only returning the length.

- `def contains(self, value, map_to_key=lambda x: x) -> bool:`

```
def contains(self, value, map_to_key=lambda x : x) -> bool:
    """ contains
        Returns True if value is in the list, False otherwise

        map_to_key: an optional mapping function from the stored
        object type to the key to compare with.
    """
    # Returns True if the value is in the list
    return self.find_index_byvalue(value, map_to_key=map_to_key) >= 0
```

- This has 1 statement:
 - Calls find_index_byvalue method:
 - Big O = $O(n, k) = k * n = O(n^2)$
- Total:
 - $T(n, k) = 1 + (2k+5)n$
 - Big O = $O(n, k) = k * n = O(n^2)$
- The contains method has complexity $O(n^2)$ which is a quadratic time complexity.

- `def find_index_byvalue(self, value, map_to_key=lambda x: x) -> int:`

```
def find_index_byvalue(self, value, map_to_key=lambda x : x) -> int:
    """ find_index_byvalue
        Returns the first index of value, or -1 if not found

        value: the value to find as an object,
        or the key to compare to (if using map_to_key)

        map_to_key: an optional mapping function from the stored
        object type to the key to compare with.
    """
    # for loop to iterate through the linked list
    for i in range(self.length()):
        # checks if the value is in the list
        # compares the value to the key of the current node
        if map_to_key(self.get(i)) == value:
            # returns the index of the value
            return i
    return -1
```

- This has:
 - 1 for loop statement (n):
 - Calls length () method with $T(n) = 1$
 - One if statement that calls the map_to_key method and get accessor:
 - map_to_key: 1 statement
 - get(): $2k + 2$
 - 1 return statement for returning the index
- Total:
 - $T(n, k) = (1+1+2k+2+1)n$
 - $T(n, k) = (2k+5)n$
 - **Big O = $O(n, k) = k*n = O(n^2)$**
- The contains method has complexity **$O(n^2)$** which is a quadratic time complexity given the get method was called inside a for loop.
- **def get(self, k: int) -> object:**

```
def get(self, k : int) -> object:
    """Return element at index k."""
    # check if k is in range
    if not 0 <= k < self._length:
        # raise an error if k is out of range
        raise IndexError("Index k out of range")
    # set current to reference the head of the linked list
    current = self.head
    # iterate through the list until current reaches the kth node
    for i in range (k):
        # check if current is not None and the node is valid at the current position
        if current is not None:
            # set current to reference the next node
            current = current.next
        else:
            # raise an error if current is None
            raise IndexError("Index out of range")
    # returns the data stored in the kth node
    return current.data
```

- This has:

- 1 statement setting current to reference the head
- 1 for loop (k)
 - 1 if statement checking current is not None
 - 1 statement setting current to reference the next node
- 1 statement returning the data stored in kth node
- Total:
 - $T(k) = 1 + 2k + 1$
 - $T(k) = 2k + 2$
 - **Big O = O(k)**
- The get statement has a linear time complexity, where k is a constant.

Mutators:

- **def set(self, k: int, value: object):**

```
def set(self, k, value):
    """Set element at index k"""
    # check if k is in range
    if not 0 <= k < self._length:
        # raise an error if k is out of range
        raise IndexError("Index k out of range")
    # set current to reference the head of the linked list
    current = self.head
    # iterate through the list until current reaches the kth node
    for i in range(k):
        # Update current to point to the next node in the list
        current = current.next
        # check if current index is out of range by checking if it is None
        if current is None:
            raise IndexError("Index out of range")
    # set the data of the kth node to the value
    current.data = value
```

- This has:
 - 1 statement setting current to reference the head
 - 1 for loop (k) to iterate through the list until kth node
 - 1 statement updating current to the next node
 - 1 statement setting the data of the kth node to the value

- Total:
 - $T(k) = 1 + k + 1$
 - $T(k) = k + 2$
 - **Big O = $O(k)$**
- The set statement has a linear time complexity, where k is a constant.
- **def append(self, value: object):**

```
def append(self, obj):
    """Add object to end of the array."""
    # create a new list node with obj
    new_list_node = ListNode(obj)
    if not self.head:
        # if the list is empty, set the new node as the head
        self.head = new_list_node
    else:
        # if the list is not empty, set current to the head
        current = self.head
        while current.next:
            # iterate until the last node is reached
            current = current.next
        # set the next node to reference the new node
        current.next = new_list_node
    # increment the length of the list
    self._length += 1
```

- This has:
 - 1 statement creating a new list node
 - Creates ListNode instance:
 - 2 statements assigning value and reference
 - 1 if statement
 - 1 statement setting the new node as the head
 - 1 else statement
 - 1 statement setting current to the head
 - 1 while statement with $\log_2(n)$
 - 1 statement
 - 1 statement setting the next node to the reference new node
 - 1 statement incrementing length

- length accessor with $T(n) = 1$
- Total:
 - $T(n) = 10 + \log_2(n)$
 - Big O = $O(\log_2(n))$
- The append method has a $O(\log_2(n))$ that represents logarithmic time complexity. This is given for the while loop statement.
- **def insert(self, k: int, value: object):**

```
def insert(self, k, value):
    """Insert value at index k, shifting subsequent values rightward."""
    # NOTE: 0 <= k <= n is okay
    # TODO: stub
    # check if k is in range
    if k < 0 or k > self._length:
        raise IndexError("Index out of range")
    # create a new node
    new_node = ListNode(value)
    # check first index
    if k == 0:
        # sets the new node to reference the head
        new_node.next = self.head
        # adjust the next pointer
        self.head = new_node
    else:
        current = self.head
        for i in range(k - 1):
            # iterate through the list to the node before the kth node
            current = current.next
        # set the next node of the new node to the next node of the current node
        new_node.next = current.next
        # set the next node to the new node
        current.next = new_node
    # increment the length of the list
    self._length += 1
```

- This has:
 - 1 statement creating a new list node
 - Creates ListNode instance:
 - 2 statements assigning value and reference
 - 1 else statement:

- 1 statement setting current to next
- 1 for loop (k-1)
 - 1 statement setting current to next
- 2 statements
 - 1 statement incrementing length
 - 1 statement from length()
- Total:
 - $T(k) = 5 + k + 4$
 - $T(k) = k + 9$
 - **Big O = O(k)**
- The insert statement has a linear time complexity, where k is a constant.

- **def remove(self, k: int) -> object:**

```
def remove(self, k : int) -> object:
    """Remove the value at index k, returning it
    """
    # check if k is in range
    if k < 0 or k >= self._length:
        # raise an error if k is out of range
        raise IndexError("Index out of range")
    if k == 0:
        # update head when removing at index 0
        removed_value = self.head.data
        # sets head to reference the next node
        self.head = self.head.next
    else:
        # set current to the head
        current = self.head
        for i in range(k - 1):
            # iterate through the list to the node before the kth node
            current = current.next
        # set the removed value to the data of the kth node
        removed_value = current.next.data
        current.next = current.next.next
    # reduce the length of the list
    self._length -= 1
    return removed_value
```

- This has:

- 1 else statement:
 - 1 statement setting current to next
 - 1 for loop (k-1)
 - 1 statement setting current to next
 - 2 statements
- 1 statement incrementing length
 - 1 statement from length()
- 1 return statement
- Total:
 - $T(k) = 2 + k + 6$
 - $T(k) = k + 8$
 - **Big O = O(k)**
- The remove statement has a linear time complexity, where k is a constant.

Sorting:

- **def sort(self, map_to_key=lambda x: x, descending=False):**

```
# check if list is length 1 or 0
if self._length <= 1:
    return

    # Iterate through each element in the list
for i in range(self._length):
    current = self.head
    # Iterate through the unsorted part of the list
    for j in range(0, self._length - i - 1):
        # Create key for the current element in the unsorted part
        key1 = map_to_key(self.get(j))
        # Create key for next element in the unsorted part
        key2 = map_to_key(self.get(j+1))
        # Swap if the keys are not in order
        # Compare keys based on the specified order (ascending or descending).
        if ((key1 > key2) if not descending else (key1 < key2)):
            # Swap the elements
            temp = current.data
            current.data = current.next.data
            current.next.data = temp
        current = current.next
```

- This has:
 - 1 for loop (n)
 - 1 statement from length()
 - 1 statement

- 1 for loop (n)
 - 1 statement from length()
 - 1 statement creating key 1
 - 1 statement from map_to_key
 - $4k + 5$ from get()
 - 1 statement creating key 2
 - 1 statement from map_to_key
 - $2k + 2$ from get()
 - 1 if statement
 - 3 statements
- 1 statement assigning current_next to current
- Total:
 - $T(n, k) = 2n + (3+1+4k+5+1+1+4k+5+1+1+1)n$
 - **$T(n, k) = 2n + (19+8k)n$**
 - **Big O = $O(n, k) = O(n^2)$**
- To calculate the time complexity, it will depend on the product of the time complexities of the outer loop, the inner loop, and the key mapping and swap. In the worst-case time complexity of the provided sorting algorithm is **$O(n^2)$**