

CPSC 5031 :

Data Structures & Algorithms

Lecture 7: Heap, Heapsort, Priority Queue
(Levitin, Chapter 6.4)

Overview

- What is a heap? Semiheap?
- How does heapsort work?
- Time efficiency of heapsort
- Priority queues

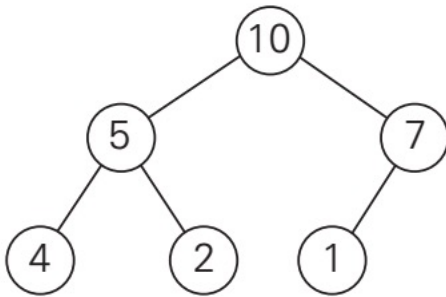
Heaps (max heap)

A **heap** is a binary tree with properties:

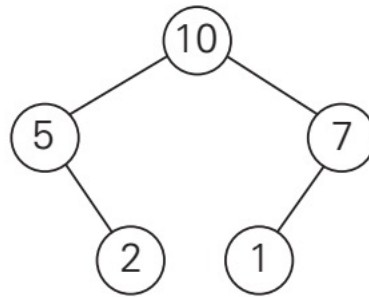
1. It is **complete**
 - Each level of tree completely filled
 - Except possibly bottom level (nodes in left most positions)
2. It satisfies ***heap-order property***
 - Data in each node \geq data in children

Heaps

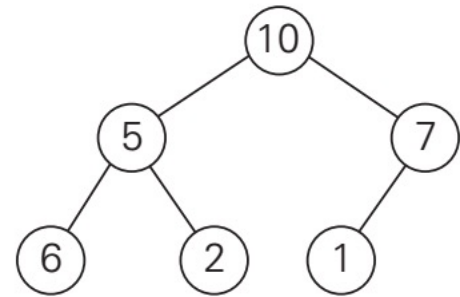
- Which of the following are heaps?



A



B



C

Heaps

- Maxheap?– by default
- Minheap?

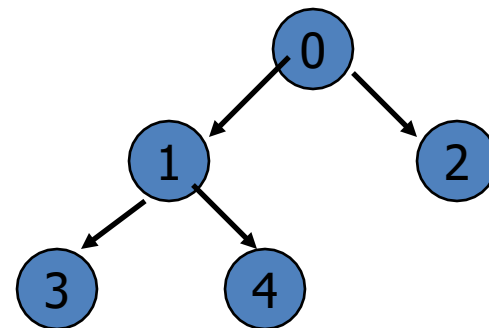
Implementing a Heap

- What data structure is good for its implementation?

Implementing a Heap

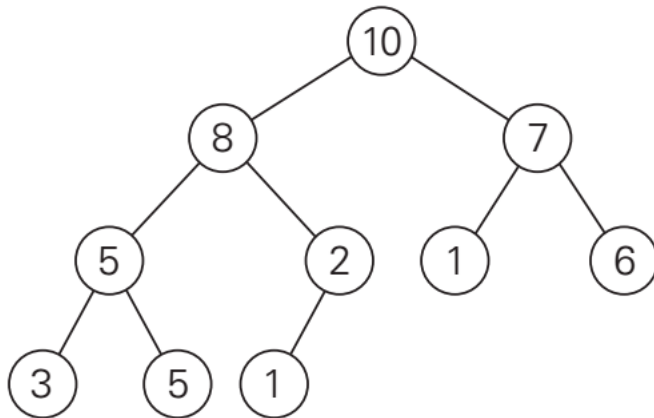
- Use an array or vector, why?
- Number the nodes from top to bottom
 - Number nodes on each row from left to right
- Store data in i^{th} node in i^{th} location of array (vector)

A heap is an array
pretending to be a tree,
NOT the other way
around!



Implementing a Heap

- Note the placement of the nodes in the array



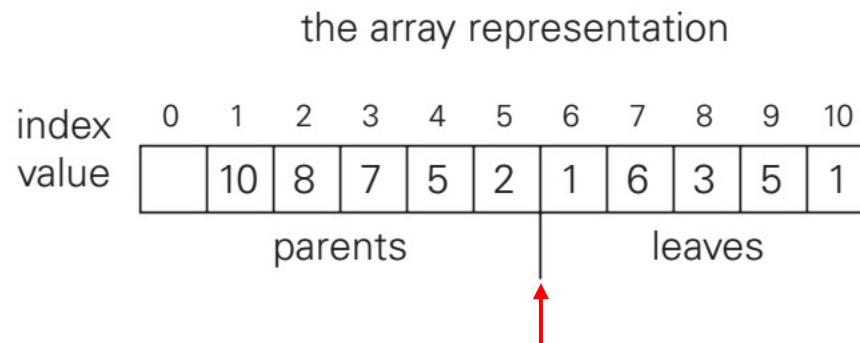
the array representation

index	0	1	2	3	4	5	6	7	8	9	10
value		10	8	7	5	2	1	6	3	5	1

parents | leaves

Implementing a Heap

- Using an array `myArray[]`
 - Children of i^{th} node are at `myArray[2*i+1]` and `myArray[2*i+2]`
 - Parent of the i^{th} node is at `myArray[(i-1)/2]`



For size n , where is this line drawn?

Compare to BST?

- Isn't this the same thing as an array BST?
- NO!
 - Look at the ordering of the array below
 - Heaps are good when you need to know the max/first element, *but nothing past that*
 - Hence very good for prioritization (as opposed to sorting)

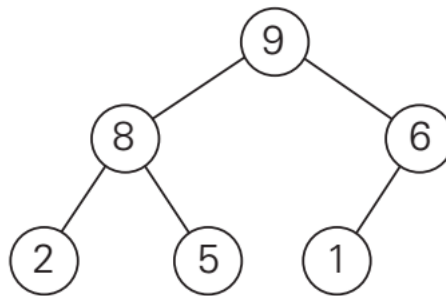
the array representation

	0	1	2	3	4	5	6	7	8	9	10
index											
value		10	8	7	5	2	1	6	3	5	1
		parents						leaves			

Basic Heap Operations

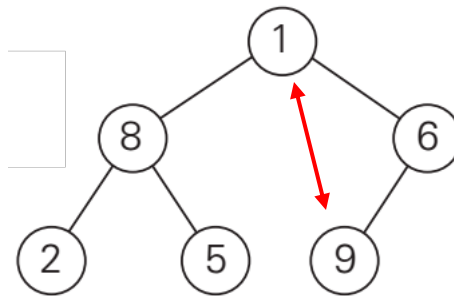
- Constructor
 - Set `mySize` to 0, allocate array
- Empty
 - Check value of `mySize`
- Retrieve the max item
 - Return root of the binary tree, `myArray[0]`
- How about delete max item?
 - Think about it? 🤔

Delete Max Item?



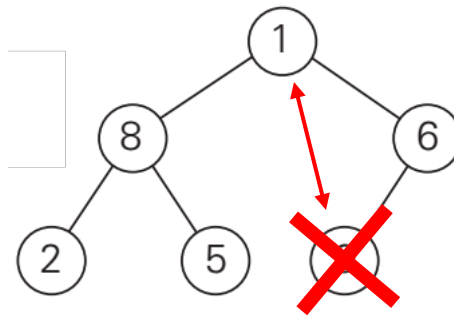
Basic Heap Operations

- Delete max item
 - Max item is the root, replace with last node in tree



Basic Heap Operations

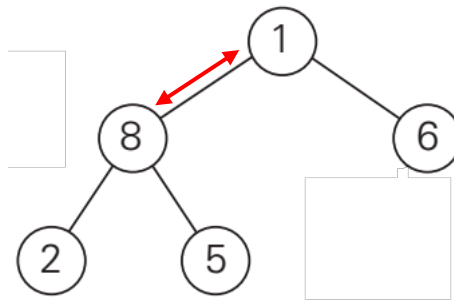
- Delete max item
 - Max item is the root, replace with last node in tree



- Take the last element

Basic Heap Operations

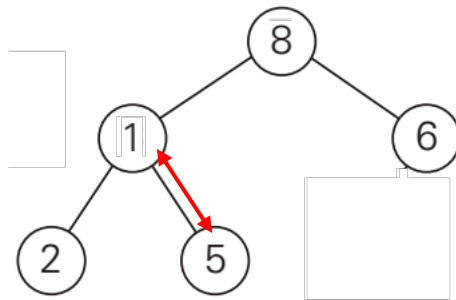
- Delete max item
 - Max item is the root, replace with last node in tree



- Take the last element
- Then interchange root with larger of two children
- Continue this with the resulting sub-tree(s) →
percolate down

Basic Heap Operations

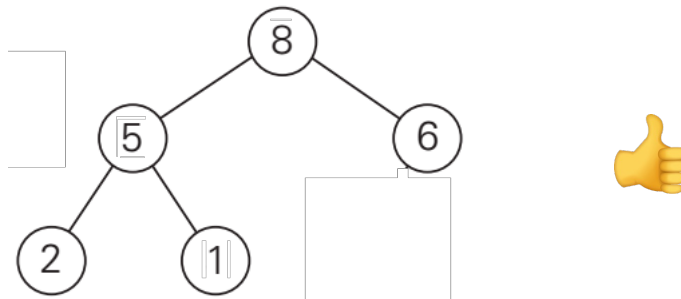
- Delete max item
 - Max item is the root, replace with last node in tree



- Take the last element
- Then interchange root with larger of two children
- Continue this with the resulting sub-tree(s) →
percolate down

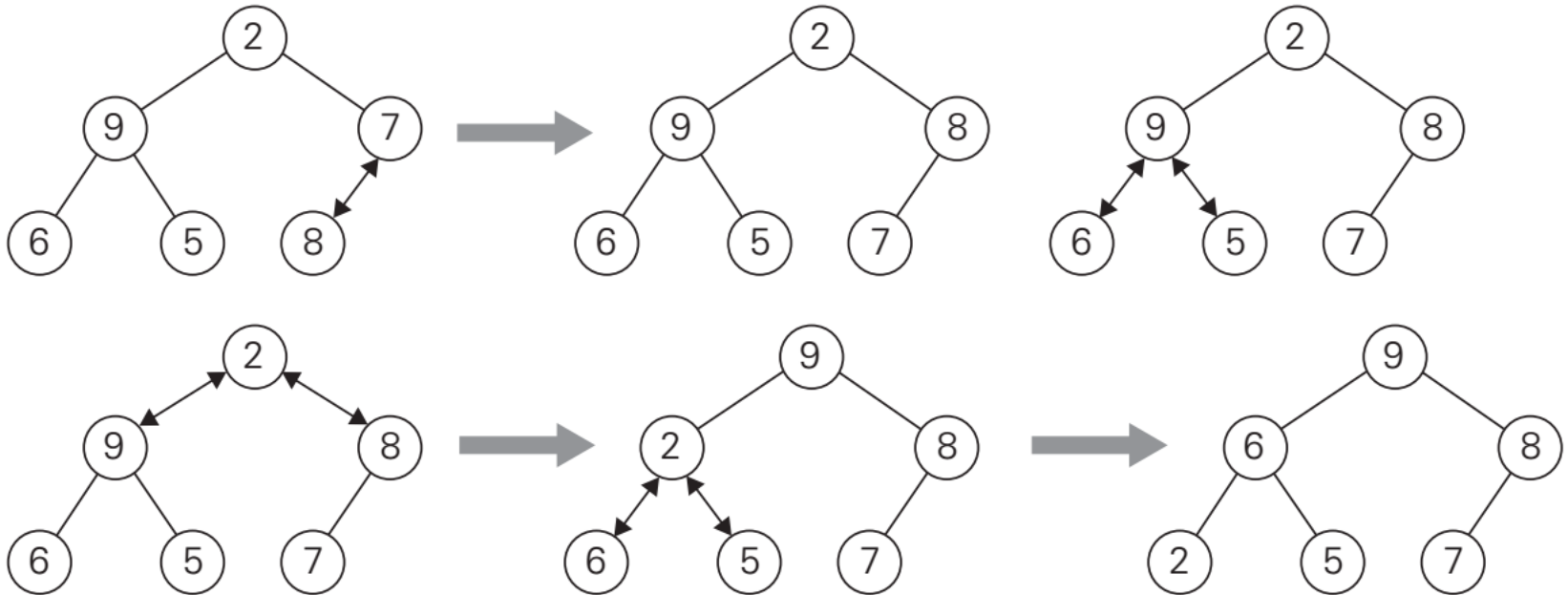
Basic Heap Operations

- Delete max item
 - Max item is the root, replace with last node in tree



- Take the last element
- Then interchange root with larger of two children
- Continue this with the resulting sub-tree(s) →
percolate down

Percolate Down Algorithm



Percolate Down Algorithm

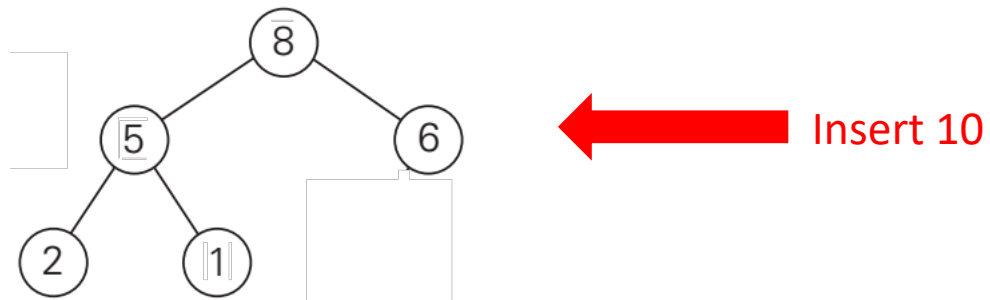
```
//Constructs a heap from elements of a given array
// by the bottom-up algorithm
//Input: An array  $H[1..n]$  of orderable items
//Output: A heap  $H[1..n]$ 
for  $i \leftarrow \lfloor n/2 \rfloor$  downto 1 do
     $k \leftarrow i$ ;  $v \leftarrow H[k]$ 
     $heap \leftarrow \text{false}$ 
    while not  $heap$  and  $2 * k \leq n$  do
         $j \leftarrow 2 * k$ 
        if  $j < n$  //there are two children
            if  $H[j] < H[j + 1]$   $j \leftarrow j + 1$ 
        if  $v \geq H[j]$ 
             $heap \leftarrow \text{true}$ 
        else  $H[k] \leftarrow H[j]$ ;  $k \leftarrow j$ 
     $H[k] \leftarrow v$ 
```

Insert an item into the heap

- Anybody have an idea?

Basic Heap Operations

- Insert an item
 - Amounts to a percolate up routine
 - Place new item at end of array



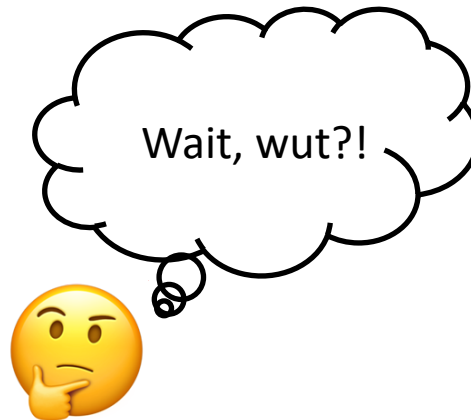
- Interchange with parent so long as it is greater than its parent

Percolate Up Algorithm

- Why percolate up?
 - When to terminate the up process?
- `void Heap::percolate_up()`
- `void Heap::insert(const T& item)`

Heapsort

- Two-stage algorithm.
- Given a list of integers...
 - “Heapify” the list
 - Then delete them all.



Heapsort - Heapify

Stage 1 (heap construction)

2 9 **7** 6 5 8

2 **9** 8 6 5 7

2 9 8 6 5 7

9 **2** 8 6 5 7

9 6 8 2 5 7

(Hint: it helps to draw the line showing where leaves are...)

Heapsort – “Delete” all nodes

Stage 2 (maximum deletions)

9 6 8 2 5 7

7 6 8 2 5 | **9**

8 6 7 2 5

5 6 7 2 | **8**

7 6 5 2

2 6 5 | **7**

6 2 5

5 2 | **6**

5 2

2 | **5**

2

(Remember: deleting just moves the root to the very end!)

Summary of HeapSort

- Given an array...
- Stage 1: Heapify this array into a heap (percolate up)
- Stage 2: Exchange the root node with the last element and shrink the list by pruning the last element (percolate down). Repeat on $|n-1|$ array.

Heapsort complexity

- $T(n) = O(n \log n)$
 - Stage 1: $O(n)$
 - Stage 2: $C(n) = 2n - 2 \log n$, so $O(n \log n)$
- $O(1)$ additional space requirement!

Priority Queue

- A collection of data elements...
 - Items stored in order by priority
 - Higher priority items removed ahead of lower
 - Do we care about sort order other than priority?

- Operations

- Constructor
- Insert
- Find, remove smallest/largest (priority) element
- Replace
- Change priority
- Delete an item
- Join two priority queues into a larger one

“Multi-set” vs.
strictly ordered array



Priority Queue

- Implementation possibilities
 - As a list (array, vector, linked list)
 - $T(n)$ for search, removeMax, insert operations?
 - As an ordered list
 - $T(n)$ for search, removeMax, insert operations?
 - Best is to use a heap
 - $T(n)$ for basic operations?
 - Why?

Should we use a PQ?

Given a list of m words, I'd like to query for the n th most frequent word(s). (e.g. Given 1000 words, tell me the 5th most frequent word(s), or the 10th most frequent word(s) etc). Would like to query multiple times for different n 's.

An approach...

1. Count word freq
2. Insert into a max-heap
3. Delete $n-1$ times
4. Take the top!

...but...

Priority Queues are heaps
...and PQ delete is destructive!

In C#

```
var freq = rawWords.Split(' ')
    .GroupBy(w=>w)
    .ToDictionary(w=>w.Key, w=>w.Count())
    .ToLookup(kvp=>kvp.Value, kvp=>kvp.Key)
    .OrderByDescending(kvp=>kvp.Key)
    .ToList();
```

```
var fifth = freq[4].Value;
```

...terrible perf, but short code...? 🤔