CMSC 245 Wrap-up

This class is about understanding how programs work

To do this, we're going to have to learn how a computer works

Learned a ton in the class

Regexp Lexical vs. Dynamic Scoping Closures Parsing Objects Racket Heaps **Functions** Stacks Method dispatch Classes Garbage collection Calling conventions Assembly

To apologize for making you write so much I wrote 732 lines of C++ yesterday

- Today we're going to design an interpreter
- Our source language will be a subset of Scheme
 - Numbers, variables, if, lambdas, let, begin, set!
- We'll write our own lexer, grammar, and parser
 - Starting from what you already wrote in labs
- Our interpreter will use data structures from the course
- And will include garbage collection under the hood

Raw Text

Lexer

Regex

Parser

CFG

AST

C++ (Sub)classes

Interpreter

Methods on AST

Raw Text

Lexer

scanner. 1 ~20 lines of code

Parser

parser.cc —150 lines of code

AST

interpreter.h —220 lines of code

Interpreter

interpreter.cc — 160 lines of code

Lesson

Sometimes the most best way to do something is to find someone who's already done it for you...

Symbol Table

Garbage Collector

HAMT

Boehm GC

HAMT

Hash Array-Mapped Trie

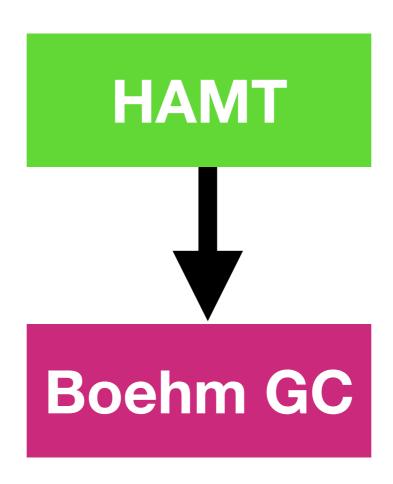
Think of this is as a hash table that is "quick" to copy

Boehm GC

High-performance GC for C

We'll use this to make it so our interpreter is automatically garbage collected

We'll have our hash table use the GC under the hood



So when we put things into HAMT, they are automatically GC'd

The grammar...

(In EBNF, allows E+)

```
START -> E $
E -> number
E -> identifier
E -> ( OP E+ )
E -> ( begin E+ )
E -> ( lambda (ID+) E )
E -> ( set! x E )
E -> ( E+ )
OP --> +|-|*|=
```

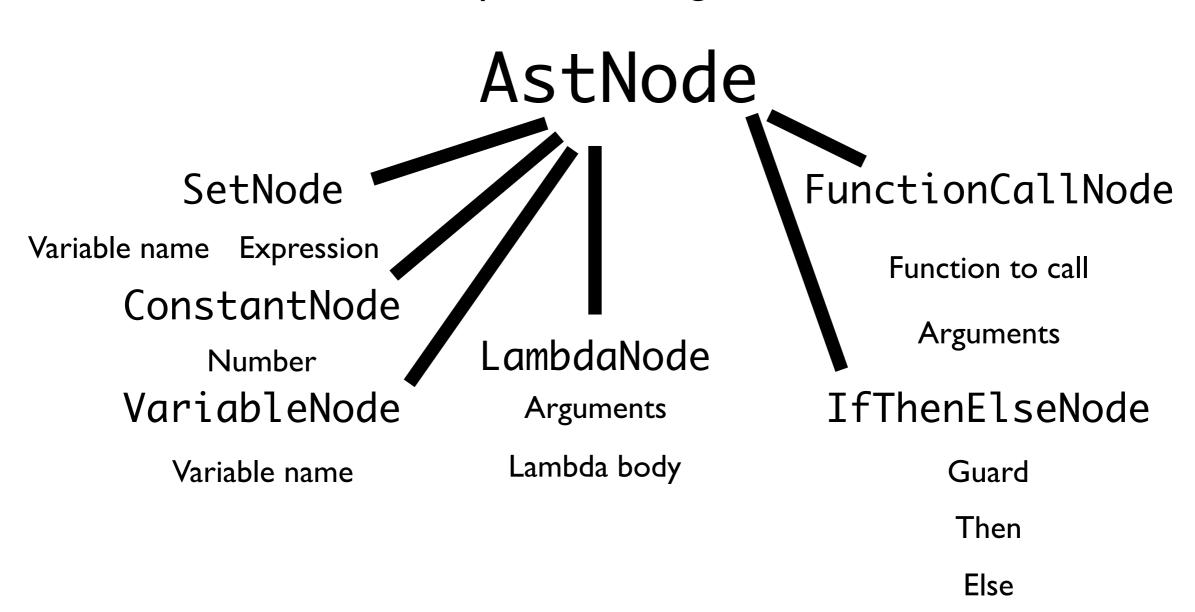
The grammar...

(In EBNF, allows E+)

We'll dig into this in a few mins

AST

Idea: Represent using subclasses



The lexer...

```
{ continue; }
[ \t]
[\n]
              { tokenCount++; return NEWLINE; }
              { continue; }
            { tokenCount++; return LPAREN; }
")"
              { tokenCount++; return RPAREN; }
"+"
              { tokenCount++; return PLUS; }
'' _ ''
              { tokenCount++; return MINUS; }
"*"
              { tokenCount++; return TIMES; }
"lambda"
              { tokenCount++; return LAMBDA; }
"let"
              { tokenCount++; return LET; }
"<E0F>"
              { tokenCount++; return END_OF_INPUT; }
-?{digit}+ { tokenCount++; return INT; }
{identifier} { tokenCount++; return IDENTIFIER; }
              { scannerError(); continue; }
```

The Parser

I started from code we gave you in Lab 5...

But I cheated because it's not LL(I)

See parser.cc

(5 minute tour)

The Symbol Table

Is a dictionary that takes strings to addresses in the heap

Means most things are stored on heap

Necessitates GC (we'll discuss next)

The Symbol Table

Wrapper for strings Representation of pointers

typedef hamt<HashedString, Address> environment;

HAMT is a dictionary

Two methods:

- Get: Takes a dictionary and key, gives us address
 - Which we then look up in heap
- Insert: Takes a dictionary, key, and value
 - Returns a **new** dictionary

The Heap

Stores two possible things:

- Plain old numbers
- Closures
- You could add other things (strings, etc..)

To find X, we look up address in symbol table, then use that address to look up through the heap

Wrapper around std::string

```
Wrapper around std::string
typedef hamt<HashedString, Address>
                 environment;
     Closures
          struct Closure {
            AstNode *function;
            environment *environment;
  Values };
typedef variant<int, Closure> value;
            Variant is new in C++17
Container that allows me to store anything from any set of types
 get<int>(x) // gets the integer value assuming
              // x is an integer
```

```
Wrapper around std::string
typedef hamt<HashedString, Address>
               environment;
    Closures
         struct Closure {
           AstNode *function;
           environment *environment;
  Values };
typedef variant<int, Closure> value;
    Heap
 hamt<Address, value> *heap
         = new hamt<Address, value>();
```

```
Address *putValueInHeap(value v) {
  heapSize++;
  Address* addr =
    new ((Address*)GC_MALLOC(sizeof(Address)))
          Address({heapSize});
  value * val =
      new ((value*)GC_MALLOC(sizeof(value))) value(v));
  heap =
    const_cast<hamt<Address, value> *>
                    (heap->insert(addr,val));
  return addr;
value getValueFromHeap(Address a) {
  return *heap->get(&a);
```

```
Address *putValueInHeap(value v) {
  heapSize++;
  Address* addr =
    new ((Address*)GC_MALLOC(sizeof(Address)))
          Address({heapSize});
  value * val =
      new ((value*)GC_MALLOC(sizeof(value))) value(v));
  heap =
    const_cast<hamt<Address, value> *>
                    (heap->insert(addr,val));
  return addr;
                             Tracks the object with GC
value getValueFromHeap(Address a) {
  return *heap->get(&a);
```

Every AstNode implementation has a method

execute : symbol table -> value

There is a "top level" symbol table where global variables go

(top of interpreter.cc)

```
int main() {
  while (true) {
    cout << "> ";
                                  REPL
    AstNode *AST = parseE();;
    executeToplevelAst(AST);
};
  void executeToplevelAst(AstNode *node) {
    value result = node->execute(globalEnvironment);
    if (holds_alternative<int>(result)) {
      cout << get<int>(result) << endl;</pre>
    } else {
      get<Closure>(result).function->render();
      cout << endl;</pre>
```