

Qualifying Exam 2020 Solution

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1a $(d.a.b.0) \setminus \{d, b\} + (b.(*.0 + c.a.0))$

1a - solution 1 $(d.a.b.0) \setminus \{d, b\}$

$$\frac{}{(d.a.b.0) \setminus \{d, b\} + (b.(*.0 + c.a.0))} \xrightarrow{*} (d.a.b.0) \setminus \{d, b\} \text{Choice}_1$$

1a - solution 2 $(b.(*.0 + c.a.0))$

$$\frac{}{(d.a.b.0) \setminus \{d, b\} + (b.(*.0 + c.a.0))} \xrightarrow{*} (b.(*.0 + c.a.0)) \text{Choice}_2$$

1b $(a.b.0 || c.a.0) || c.b.0$

1b - solution 1 $(a.b.0 || c.a.0) || b.0$

$$\frac{\frac{}{c.b.0 \xrightarrow{c} b.0} \text{Prefix}}{(a.b.0 || c.a.0) || c.b.0} \xrightarrow{c} ((a.b.0 || c.a.0) || b.0) \text{Par}_2$$

1b - solution 2 $(b.0 || c.a.0) || c.b.0$

$$\frac{\frac{\frac{}{a.b.0 \xrightarrow{a} b.0} \text{Prefix}}{(a.b.0 || c.a.0) \xrightarrow{a} (b.0 || c.a.0)} \text{Par}_1}{((a.b.0 || c.a.0) || c.b.0) \xrightarrow{a} (b.0 || c.a.0) || c.b.0} \text{Par}_1$$

1b - solution 3 $(a.b.0||a.0)||c.b.0$

$$\frac{\frac{\frac{\text{Prefix}}{c.a.0 \xrightarrow{c} a.0}}{(a.b.0||c.a.0) \xrightarrow{c} (a.b.0||a.0)} \text{Par}_2}{((a.b.0||c.a.0)||c.b.0) \xrightarrow{c} (a.b.0||a.0)||c.b.0} \text{Par}_1$$

1b - solution 4 $(a.b.0||a.0)||b.0$

$$\frac{\frac{\frac{\text{Prefix}}{c.a.0 \xrightarrow{c} a.0}}{(a.b.0||c.a.0) \xrightarrow{c} (a.b.0||a.0)} \text{Par}_2 \quad \frac{\text{Prefix}}{c.b.0 \xrightarrow{c} b.0}}{((a.b.0||c.a.0)||c.b.0) \xrightarrow{*} ((a.b.0||a.0)||b.0)} \text{Par}_3$$

1c $((a.0 + b.0)||c.0)||e.0$

1c - solution 1 $((a.0 + b.0)||c.0)||0$

$$\frac{\frac{\text{Prefix}}{e.0 \xrightarrow{e} 0}}{(((a.0 + b.0)||c.0)||e.0) \xrightarrow{e} (((a.0 + b.0)||c.0)||0)} \text{Par}_2$$

1c - solution 2 $((a.0 + b.0)||0)||e.0$

$$\frac{\frac{\frac{\text{Prefix}}{c.0 \xrightarrow{c} 0}}{((a.0+b.0)||c.0) \xrightarrow{c} ((a.0+b.0)||0)} \text{Par}_2}{(((a.0 + b.0)||c.0)||e.0) \xrightarrow{c} (((a.0 + b.0)||0)||e.0)} \text{Par}_1$$

1c - solution 3 $(a.0||c.0)||e.0$

$$\frac{\frac{\frac{\text{Choice}_1}{(a.0+b.0) \xrightarrow{*} a.0}}{((a.0+b.0)||c.0) \xrightarrow{*} (a.0||c.0)} \text{Par}_1}{(((a.0 + b.0)||c.0)||e.0) \xrightarrow{*} ((a.0||c.0)||e.0)} \text{Par}_1$$

1c - solution 4 $(b.0||c.0)||e.0$

$$\frac{\frac{\frac{}{\text{Choice}_2} \xrightarrow{*} (a.0+b.0)}{(a.0+b.0)||c.0} \xrightarrow{*} (b.0||c.0)}{\text{Par}_1} \xrightarrow{*} ((a.0+b.0)||c.0)||e.0 \xrightarrow{*} ((b.0||c.0)||e.0) \text{Par}_1$$

1d $((b.c.0||a.d.0)||b.0) \setminus \{b, c\}$

1d - solution 1 $((b.c.0||d.0)||b.0) \setminus \{b, c\}$

$$\frac{\frac{\frac{\frac{}{\text{Prefix}} a.d.0 \xrightarrow{a} d.0}{(b.c.0||a.d.0) \xrightarrow{a} (b.c.0||d.0)} \text{Par}_2}{((b.c.0||a.d.0)||b.0) \xrightarrow{a} ((b.c.0||d.0)||b.0)} \text{Par}_1}{((b.c.0||a.d.0)||b.0) \setminus \{b, c\} \xrightarrow{a} ((b.c.0||d.0)||b.0) \setminus \{b, c\}} \text{Restrict}$$

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Proof by induction over Q.

$P = 0$ There are no derivations, this case does not matter.

$P = \alpha.P'$. There is one derivation: P' .

$$\text{sort}(\alpha.P') = (\{\alpha\} - *) \cup \text{sort}(P').$$

We must prove that $\text{sort}(P') \subseteq (\alpha.P')$.

$$\text{sort}(P') \subseteq \{\alpha\} \cup \text{sort}(P').$$

In the case where $\alpha = *$, then the two sets are equal.

Otherwise, this is naturally evident by the laws of sets.

$P = P_1 + P_2$. There are two derivations: P_1 and P_2 .

$$\text{sort}(P_1 + P_2) = \text{sort}(P_1) \cup \text{sort}(P_2).$$

For P_1 :

We must prove that $\text{sort}(P_1) \subseteq \text{sort}(P_1 + P_2)$.

$$\text{sort}(P_1) \subseteq \text{sort}(P_1) \cup \text{sort}(P_2).$$

This is naturally evident by the laws of sets.

The same holds for P_2 .

$P = P_1 || P_2$. There are three derivations: $P'_1 || P_2$, $P_1 || P'_2$, $P'_1 || P'_2$
 The inductive hypothesis is: $sort(P'_1) \subseteq sort(P_1)$ and $sort(P'_2) \subseteq sort(P_2)$
 $sort(P_1 || P_2) = sort(P_1) \cup sort(P_2)$.
 For the first derivation $P'_1 || P_2$:
 We must prove that $sort(P'_1 || P_2) \subseteq sort(P_1 || P_2)$.
 $sort(P'_1) \cup sort(P_2) \subseteq sort(P_1) \cup sort(P_2)$.
 This holds by the laws of sets when you apply the inductive hypothesis.
 The same method applies to the two other derivations.

$P = P' \setminus X$. There is one derivation: $P'' \setminus X$
 The inductive hypothesis is: $sort(P'') \subseteq sort(P')$ and that P'' isn't derived using a label in X .
 $sort(P' \setminus X) = sort(P') - X$
 $sort(P'' \setminus X) = sort(P'') - X$
 We must prove that $sort(P'' \setminus X) \subseteq sort(P' \setminus X)$
 $sort(P'') - X \subseteq sort(P') - X$
 This holds by the inductive hypothesis, and the laws of sets.

□