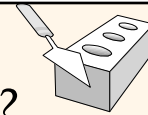


# *The Relational Model*

## Chapter 3



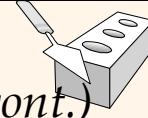
## *Why Study the Relational Model?*

- ❖ Most widely used model.
  - Vendors: IBM, Microsoft, Oracle, etc.
- ❖ “Legacy systems” in older models
  - e.g., IBM’s IMS (Information Management System) – hierarchical model
- ❖ Recent competitor: object-oriented model
  - ObjectStore, Versant, Ontos
  - A synthesis emerging: *object-relational model*
    - Oracle, IBM DB2, MS SQL Server

[https://en.wikipedia.org/wiki/Object\\_database](https://en.wikipedia.org/wiki/Object_database)

<https://en.wikipedia.org/wiki/NoSQL>

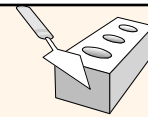
## Why Study the Relational Model? (cont.)



### ❖ Relational model features

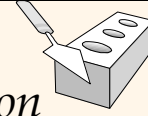
- Very simple and elegant data representation
- Even novice users can understand the contents of a database
- Supports a popular high level query language – SQL
- Complex queries can be easily expressed

## Relational Database: Definitions



- ❖ *Relational database*: a set of *relations* (tables)
- ❖ *Relation*: made up of 2 parts:
  - *Schema* : specifies name of relation, plus name and type of each column.
    - e.g., Students (*sid*: string, *name*: string, *login*: string, *age*: integer, *gpa*: real).
  - *Instance* : a *table*, with rows and columns.  
#Rows = *cardinality*, #fields = *degree*.
- ❖ Can think of a relation as a *set* of rows or *tuples* (i.e., all rows are distinct).

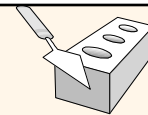
## Example Instance of Students Relation



sid	name	login	age	gpa
53666	Jones	jones@cs	18	3.4
53688	Smith	smith@ee	18	3.2
53650	Smith	smith@math	19	3.8

- ❖ Cardinality = 3, degree = 5, **all rows distinct**
- ❖ Do all columns in a relation instance have to be distinct?

## Relational Query Languages

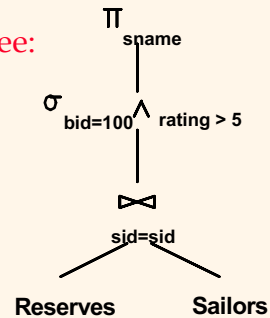


- ❖ A major strength of the relational model: supports simple, powerful *querying* of data.
- ❖ Queries can be written intuitively, and the DBMS is responsible for efficient evaluation.
  - The key: precise semantics for relational queries.
  - Allows the optimizer to extensively re-order operations, and still ensure that the answer does not change (Chapter 12).

## Overview of Query Evaluation

```
SELECT S.sname
FROM Reserves R, Sailors S
WHERE R.sid=S.sid AND
      R.bid=100 AND S.rating>5
```

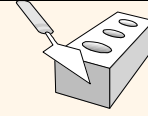
RA Tree:



## The SQL Query Language

- ❖ Developed by IBM (for the pioneering system - System R) in the 1970s
- ❖ Need for a standard since it is used by many vendors
- ❖ Standards:
  - SQL-86
  - SQL-89 (minor revision)
  - SQL-92 (major revision)
  - SQL-1999 (major extensions)
  - SQL-2003 (minor revision)
  - SQL-2008 (minor revision)
  - SQL-2011 (minor revision)
  - SQL-2016 (latest release)

## Creating Relations in SQL



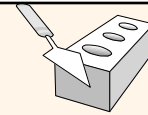
- ❖ Creates the Students relation. Observe that the type (**domain**) of each field is specified, and enforced by the DBMS whenever tuples are added or modified.

```
CREATE TABLE Students
(sid: CHAR(20),
 name: CHAR(20),
 login: CHAR(10),
 age: INTEGER,
 gpa: REAL)
```

- ❖ As another example, the Enrolled table holds information about courses that students take.

```
CREATE TABLE Enrolled
(sid: CHAR(20),
 cid: CHAR(20),
 grade: CHAR(2))
```

## Adding and Deleting Tuples



- ❖ Can insert a single tuple using:

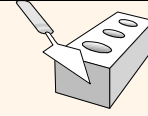
```
INSERT INTO Students (sid, name, login, age, gpa)
VALUES (53688, 'Smith', 'smith@cs', 18, 3.2)
```

- ❖ Can delete all tuples satisfying some condition (e.g., name = Smith):

```
DELETE
FROM Students S
WHERE S.name = 'Smith'
```

➡ *Powerful variants of these commands are available; more later!*

## Update Tuples

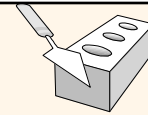


- ❖ Modify the column values in an existing row using the UPDATE command

```
UPDATE Students S
SET S.age = S.age + 1, S.gpa = S.gpa - 1
WHERE S.sid = 53688
```

```
UPDATE Students S
SET S.gpa = S.gpa - 0.1
WHERE S.gpa >= 3.6
```

## The SQL Query Language



- ❖ To find all 18 years old students, we can write:

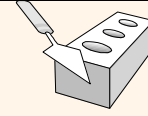
```
SELECT *
FROM Students S
WHERE S.age=18
```

sid	name	login	age	gpa
53666	Jones	jones@cs	18	3.4
53688	Smith	smith@ee	18	3.2

- To find just names and logins, replace the first line:

```
SELECT S.name, S.login
```

## Querying Multiple Relations



- ❖ What does the following query compute?

```
SELECT S.name, E.cid  
FROM Students S, Enrolled E  
WHERE S.sid=E.sid AND E.grade="A"
```

Given the following instances of Enrolled and Students:

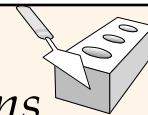
sid	name	login	age	gpa
53666	Jones	jones@cs	18	3.4
53688	Smith	smith@eecs	18	3.2
53650	Smith	smith@math	19	3.8

sid	cid	grade
53831	Carnatic101	C
53831	Reggae203	B
53650	Topology112	A
53666	History105	B

we get:

S.name	E.cid
Smith	Topology112

## Destroying and Altering Relations



```
DROP TABLE Students
```

- ❖ Destroys the relation Students. The **schema information** and the **tuples** are deleted.

```
ALTER TABLE Students  
ADD COLUMN firstYear: integer
```

- ❖ The schema of Students is altered by adding a new field; every tuple in the current instance is extended with a **null** value in the new field.