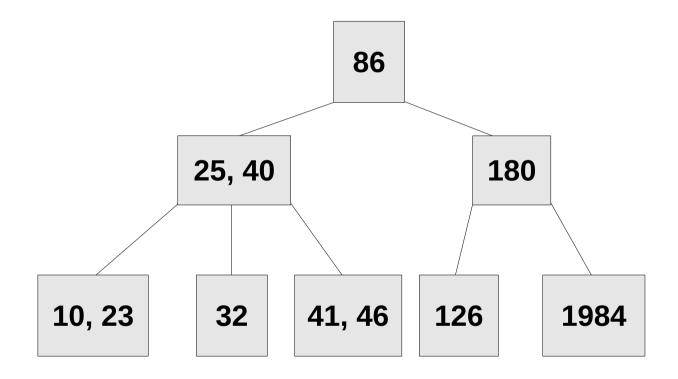
## 2-4 Trees

(Yet another balanced tree structure, but way more fun than AVL trees.)



Slides by **Sean Szumlanski** for **COP 3503**, Computer Science II

(Let's play the game!)

Insert 10 (into an initially empty tree):

(Let's play the game!)

Insert 10 (into an initially empty tree):

**10** 

(Let's play the game!)

Insert 40:

**10** 

(Let's play the game!)

Insert 40:

10, 40

(Let's play the game!)

Insert 180:

10, 40

(Let's play the game!)

Insert 180:

**10**, **40**, **180** 

(Things are getting nice and cozy in there!)

(Let's play the game!)

Insert 25:

10, 40, 180

(Let's play the game!)

Insert 25:

**10**, **25**, **40**, **180** 

Oh no!

The node has too many elements. This is called an overflow.

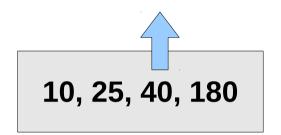
It's time to...

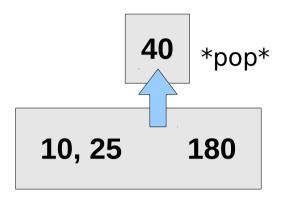
SQUISH UP!

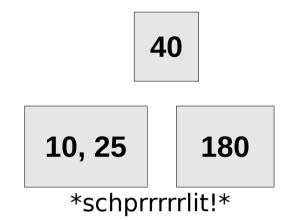
(This is also called "splitting," which is insufferably dull.)

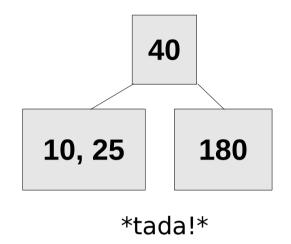
Our contract this semester: Let's agree always to squish up the third element.

That's a bit arbitrary. We could also have chosen the second element. When coding these up, you just have to make a choice and stick with it.

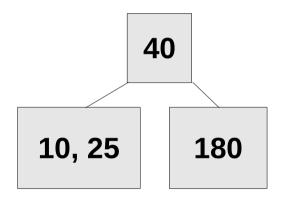




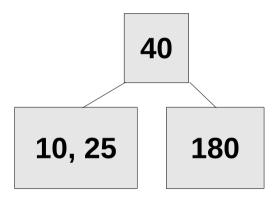




(Let's play the game!)

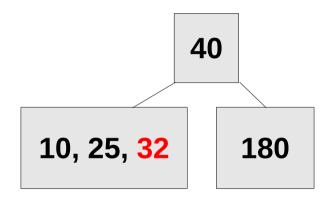


Let's continue inserting...



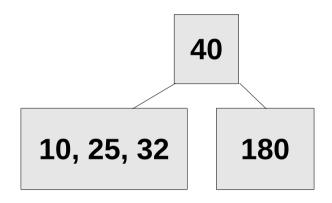
Insert 32

(Let's play the game!)

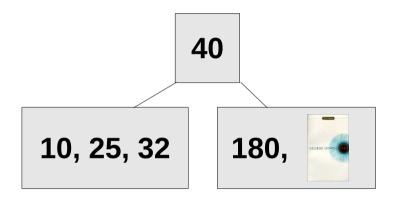


Insert 32

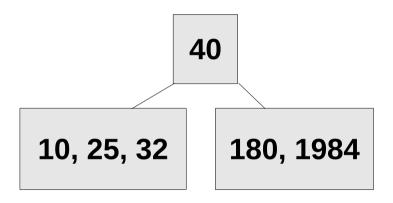
(Always insert at a leaf node.)



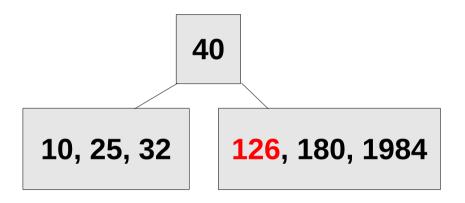
Insert 1984



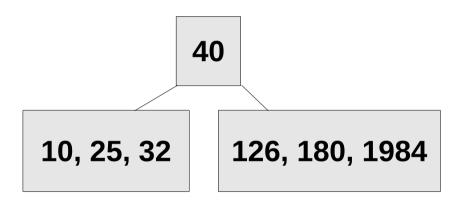
Insert 1984



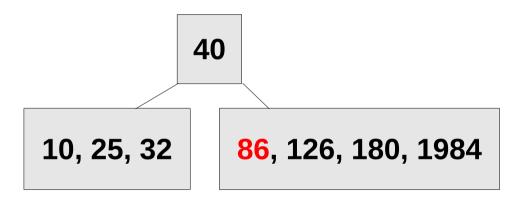
Insert 126



Insert 126

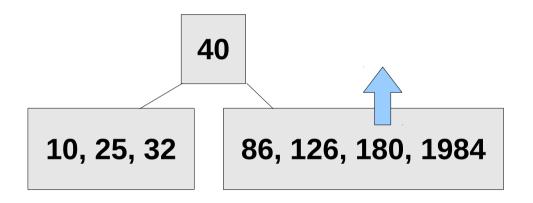


Insert 86



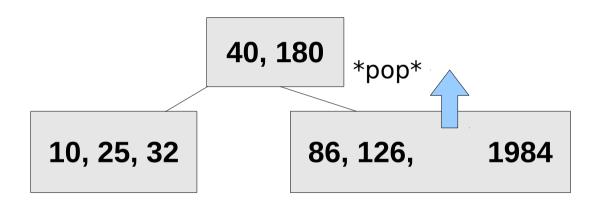
Insert 86

(Let's play the game!)

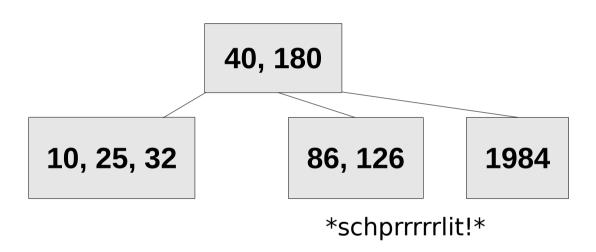


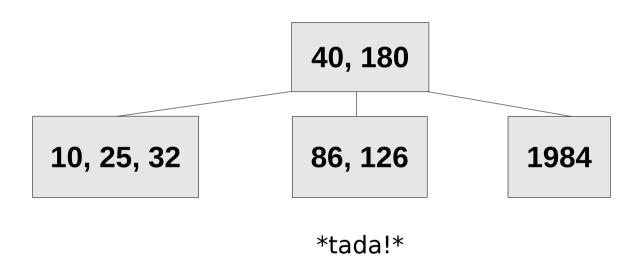
Squish up! (Which element?)

(Let's play the game!)

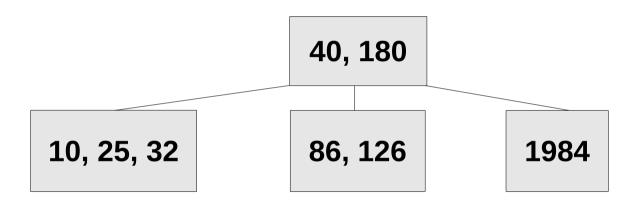


Squish up! (Which element?)

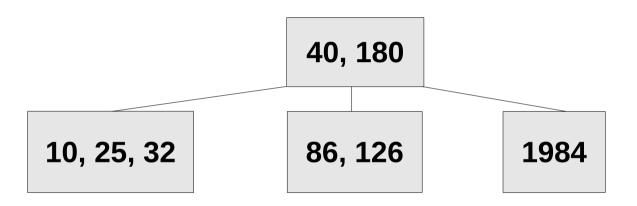




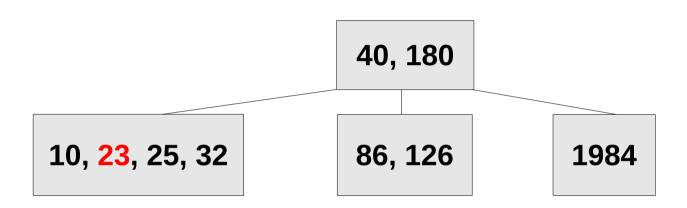
(Let's play the game!)



Let's continue inserting...

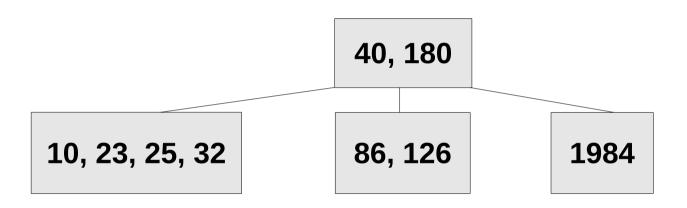


Insert 23



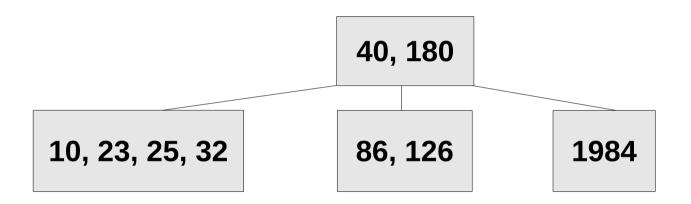
Insert 23

(Let's play the game!)



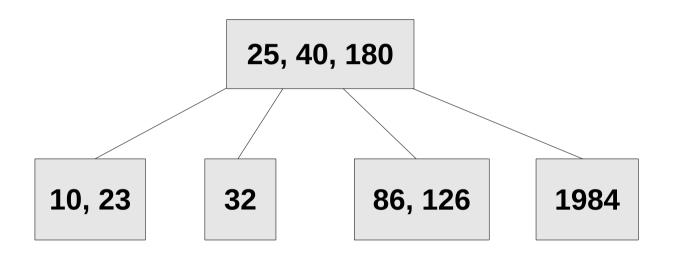
25 gets squished up.

(Let's play the game!)

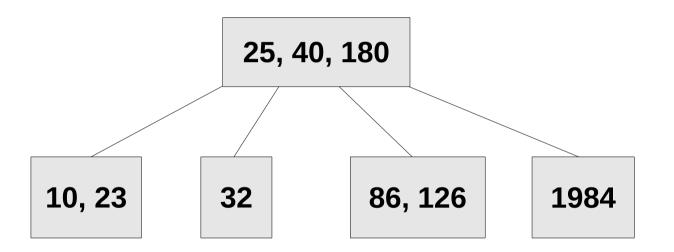


Let's skip the theatrics for the sake of time.

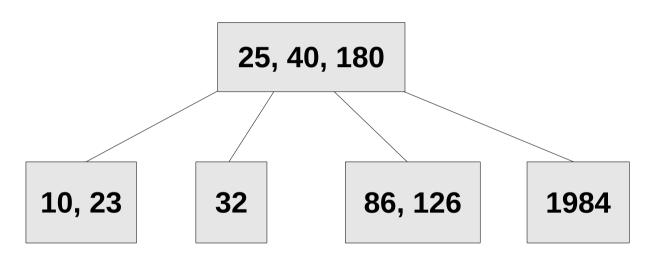
Do you see how the structure will change?



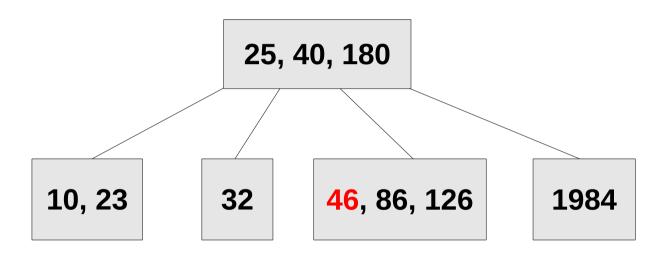
(Let's play the game!)



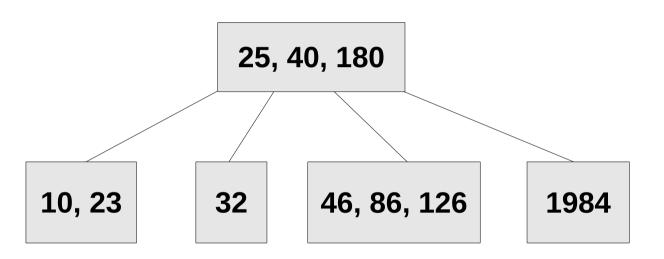
Just two more insertions to make this thing bust...



Insert 46

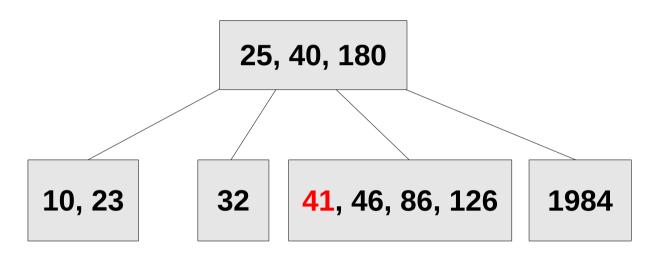


Insert 46



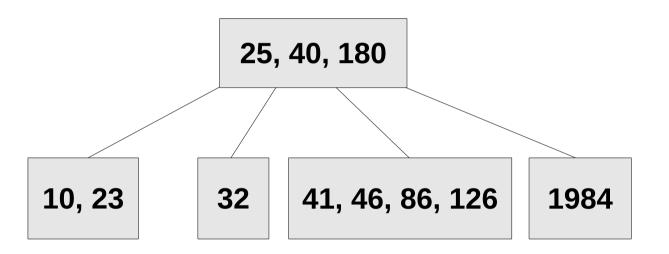
Insert 41

(Let's play the game!)



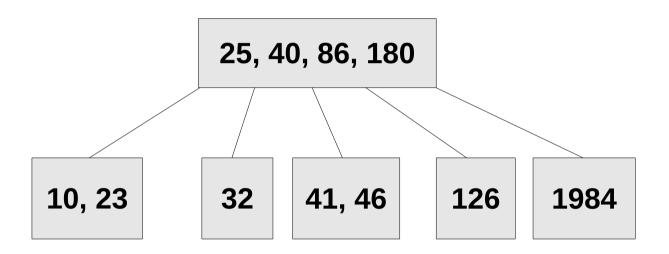
Insert 41

(Let's play the game!)



86 gets squished up

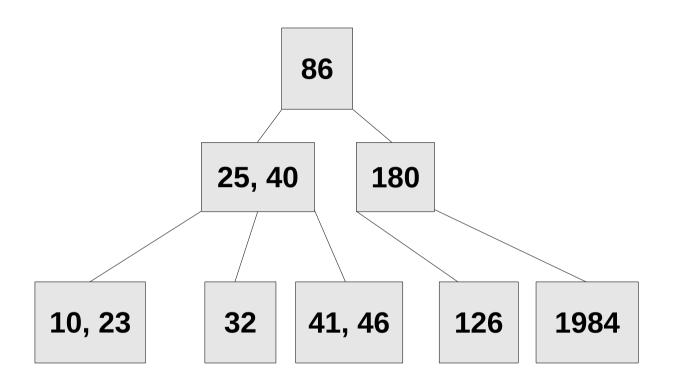
(Let's play the game!)



But now the root node has an overflow!

86 gets squished up again.

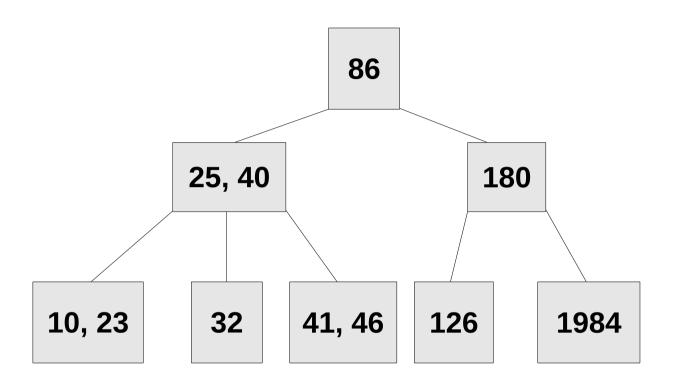
(Let's play the game!)



But now the root node has an overflow!

86 gets squished up again.

(Let's play the game!)

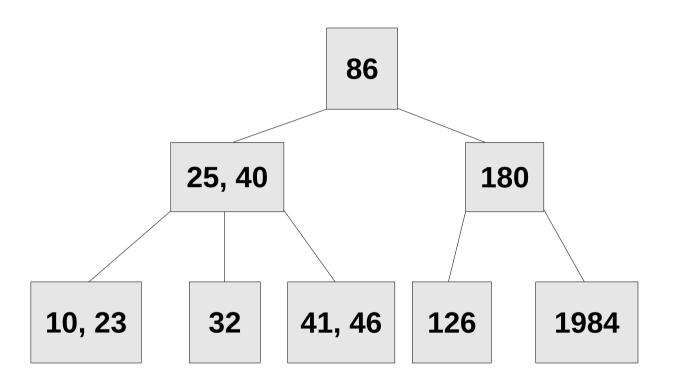


But now the root node has an overflow!

86 gets squished up again.

Let's take a look at deletion now.

(this is going to get crazy)

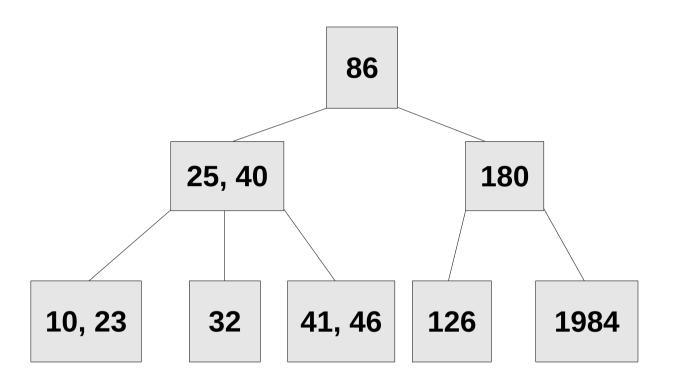


If we delete a non-leaf node, we do it BST style:

replace with largest value in left subtree
--or-replace with smallest value in right subtree

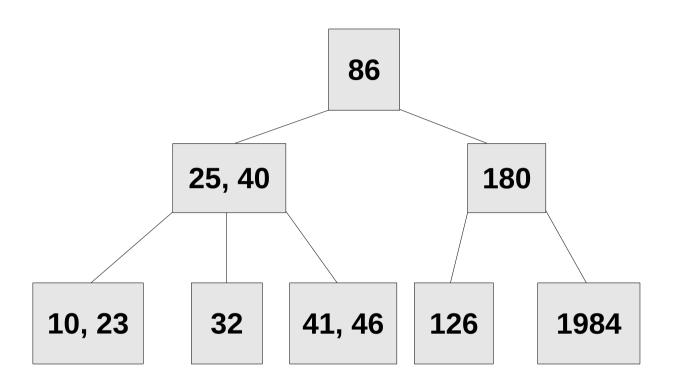
But that means we're always bring up a value from a leaf node!

(this is going to get crazy)

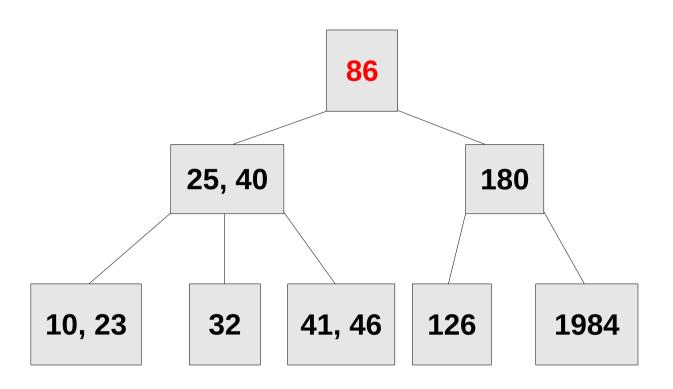


So, leaf node deletions are the only interesting deletions.

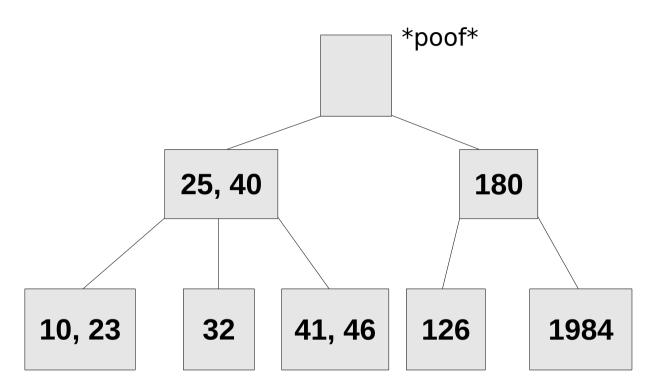
Let's take a look...



Delete 86

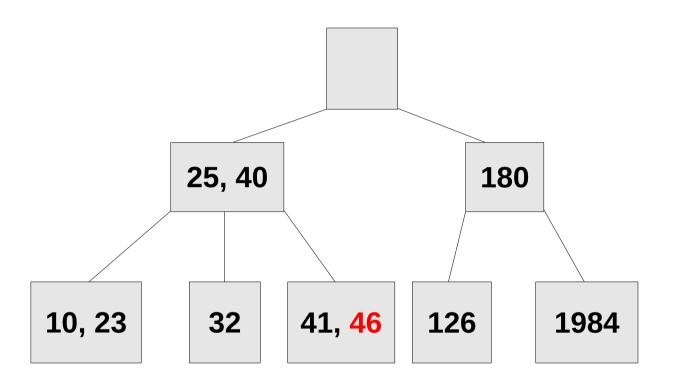


Delete 86



Delete 86
Which element moves up?

(this is going to get crazy)

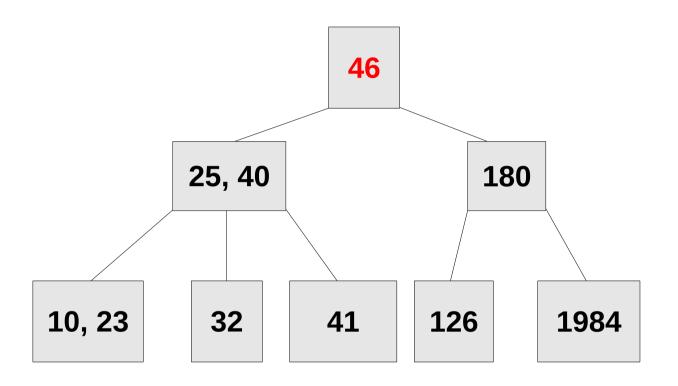


Delete 86
Which element moves up?

If we grab an element from a leaf node with multiple elements, we're okay:

No change in leaf node heights. No children to worry about.

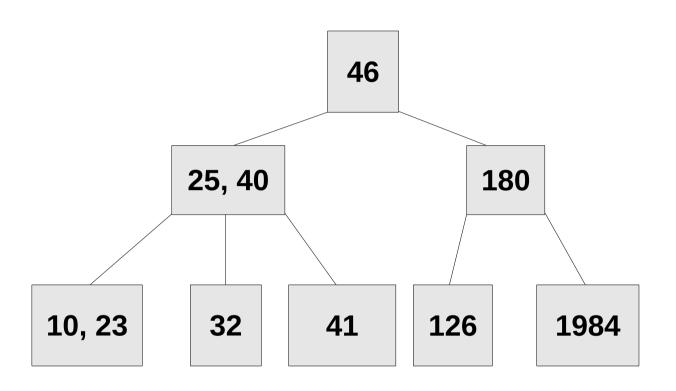
(this is going to get crazy)



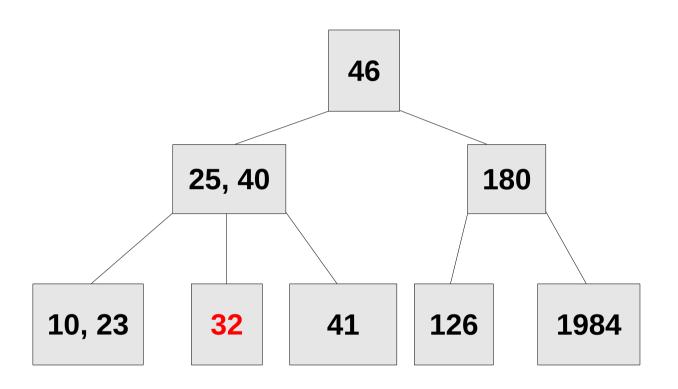
Delete 86
Which element moves up?

If we grab an element from a leaf node with multiple elements, we're okay:

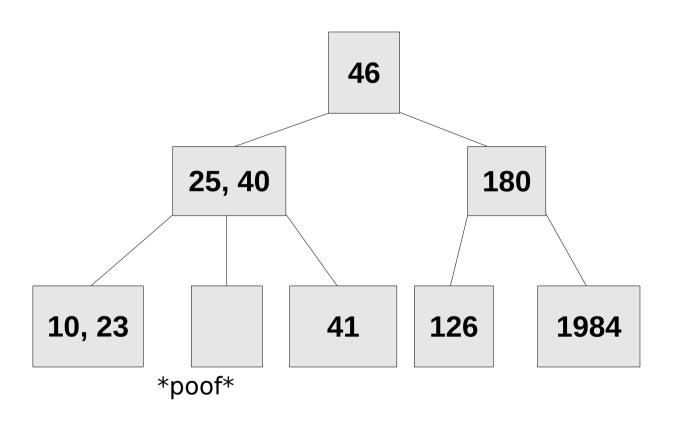
No change in leaf node heights. No children to worry about.



Delete 32

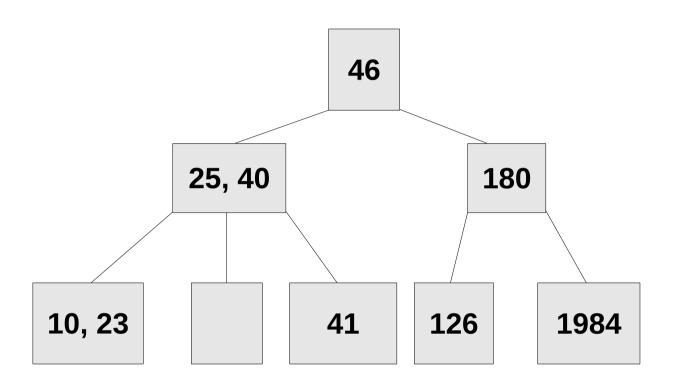


Delete 32



Delete 32

(this is going to get crazy)



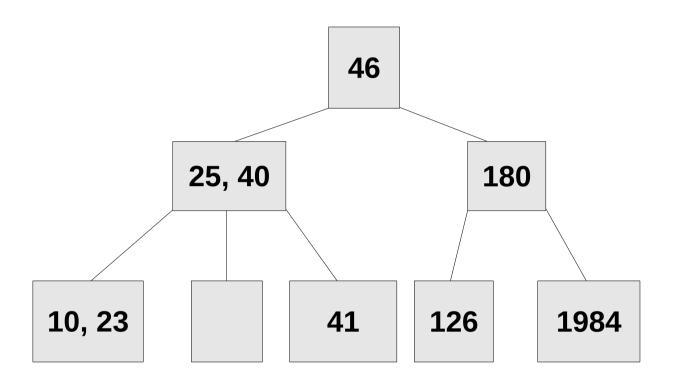
Delete 32

We can't eliminate that node, though! It won't be a valid 2-4 tree if we do.

Instead, we go begging for help from our siblings.

First look left, then look right. Our left sibling has some elements to spare.

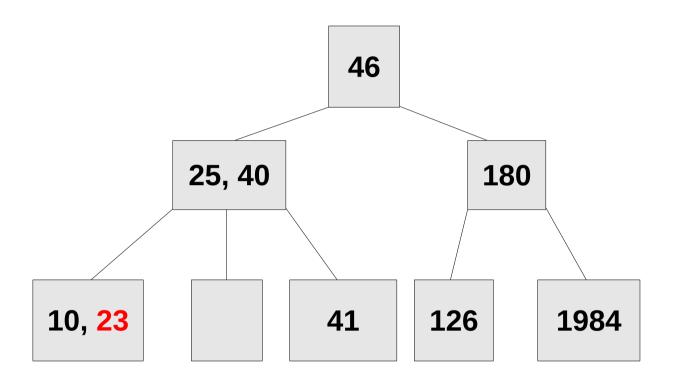
(this is going to get crazy)



Delete 32

We can't just move 23 over. That would violate our ordering property.

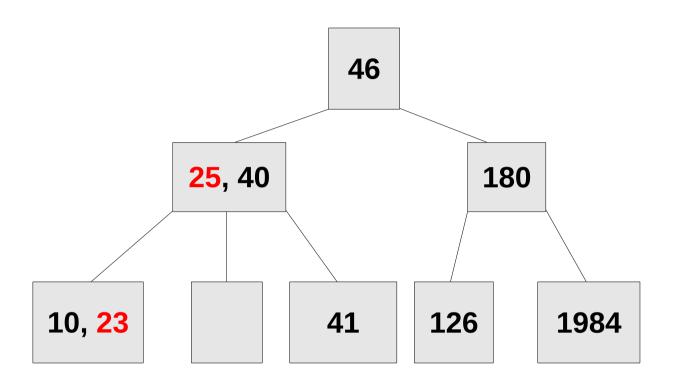
(this is going to get crazy)



Delete 32

We can't just move 23 over. That would violate our ordering property.

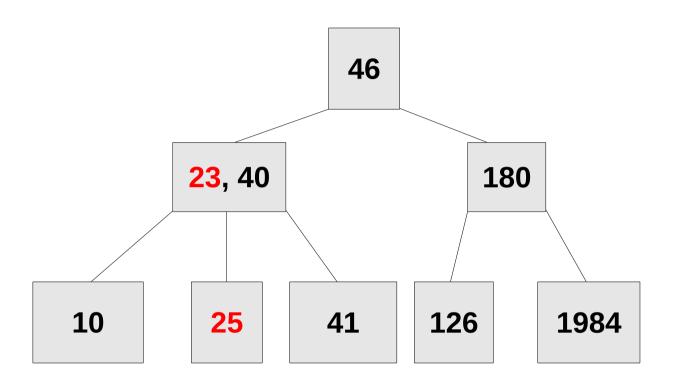
(this is going to get crazy)



Delete 32

We can't just move 23 over. That would violate our ordering property.

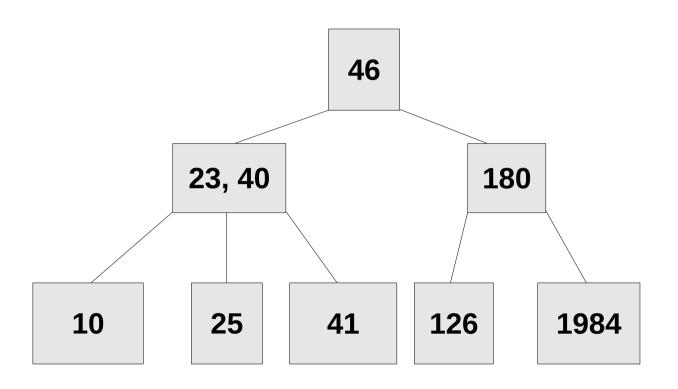
(this is going to get crazy)



Delete 32

We can't just move 23 over. That would violate our ordering property.

(this is going to get crazy)

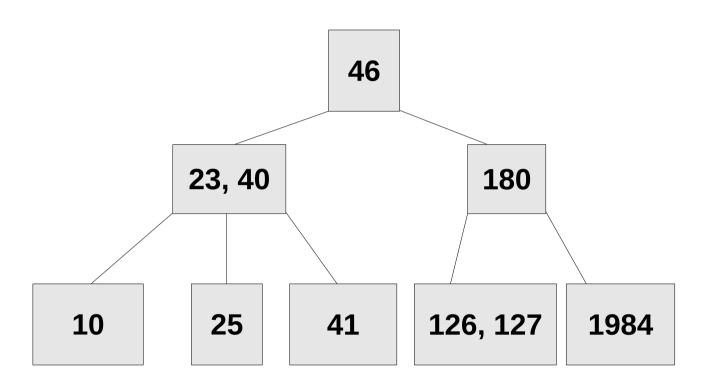


Delete 32

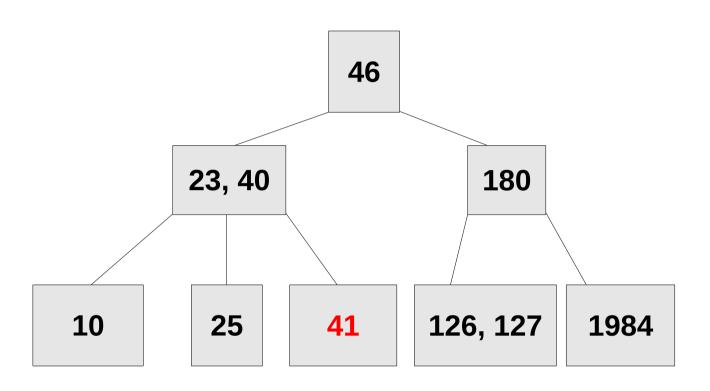
We can't just move 23 over. That would violate our ordering property.

Instead, 23 moves up and knocks its parent down.

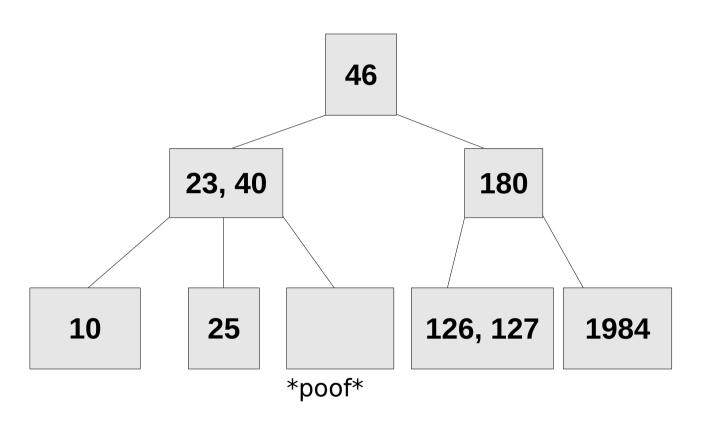
Tada!



Delete 41

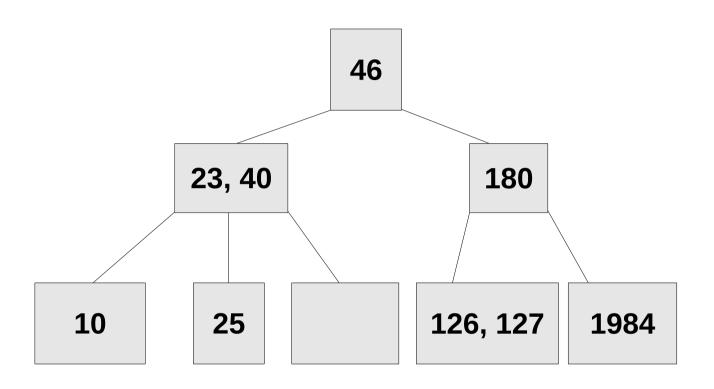


Delete 41



Delete 41

(this is going to get crazy)



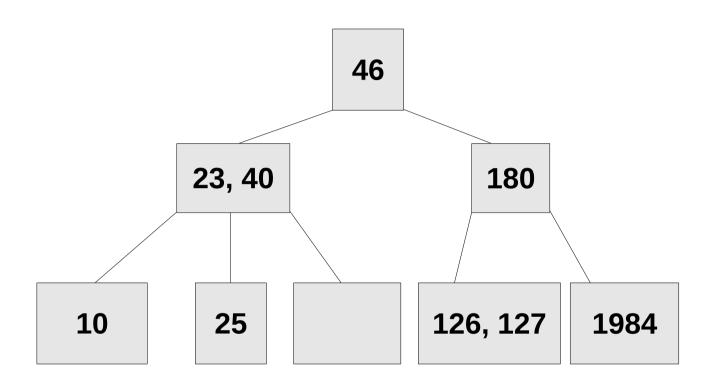
Delete 41

But the empty node has no siblings with extra elements!

Note: (126, 127) is not a sibling of that empty node.

(They have different parents!)

(this is going to get crazy)

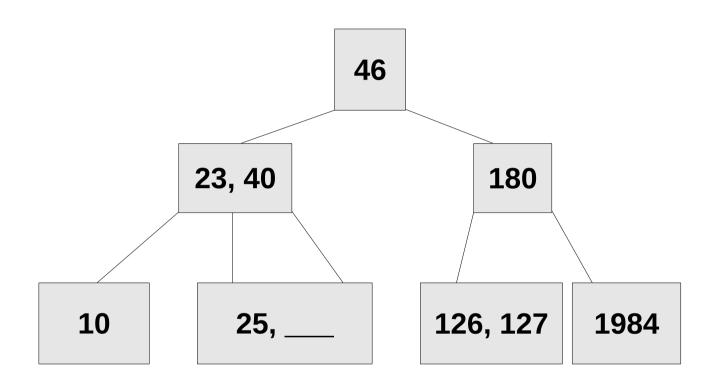


Delete 41

THIS IS A CRISIS. But we have a solution: fuse and drop.

First, fuse the empty node with its left sibling (or right sibling, if it has no left sibling).

(this is going to get crazy)



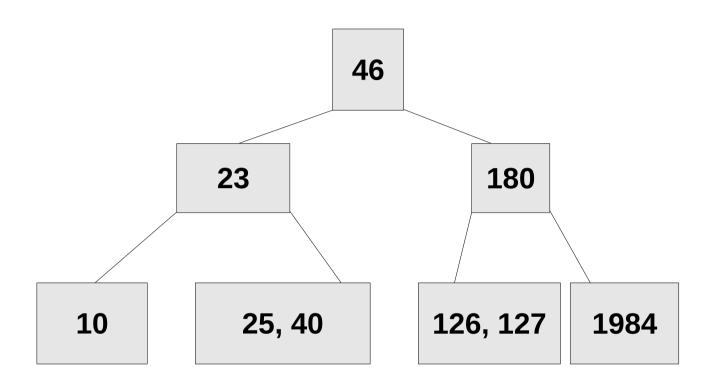
Delete 41

THIS IS A CRISIS. But we have a solution: fuse and drop.

First, fuse the empty node with its left sibling (or right sibling, if it has no left sibling).

Now drop down the parent that was in between the newly fused nodes.

(this is going to get crazy)

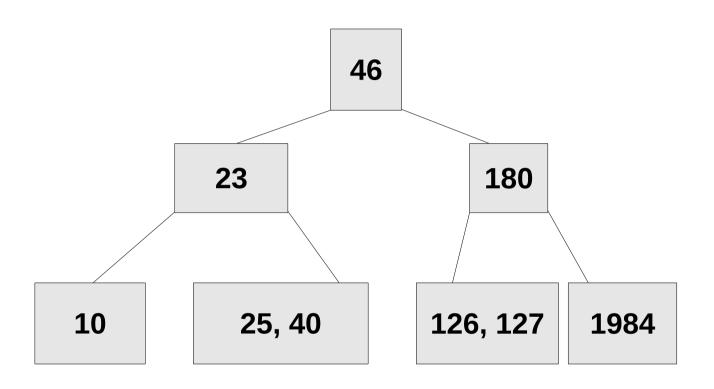


Delete 41

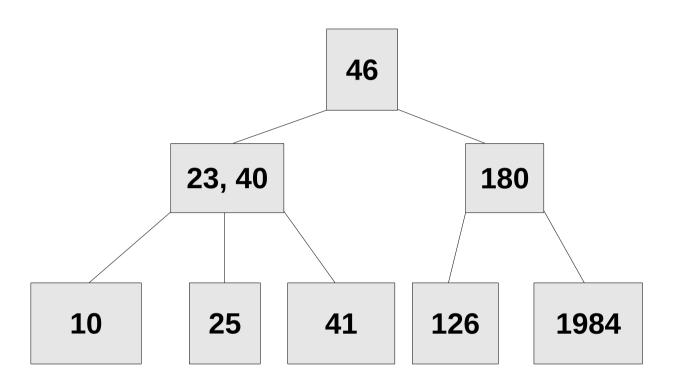
THIS IS A CRISIS. But we have a solution: fuse and drop.

First, fuse the empty node with its left sibling (or right sibling, if it has no left sibling).

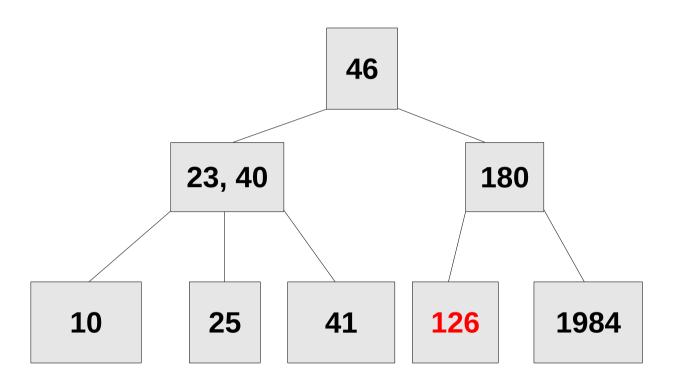
Now drop down the parent that was in between the newly fused nodes.



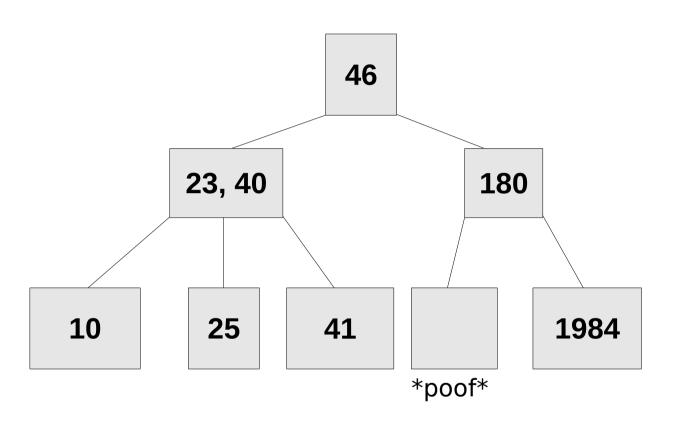
Tada!



Delete 126

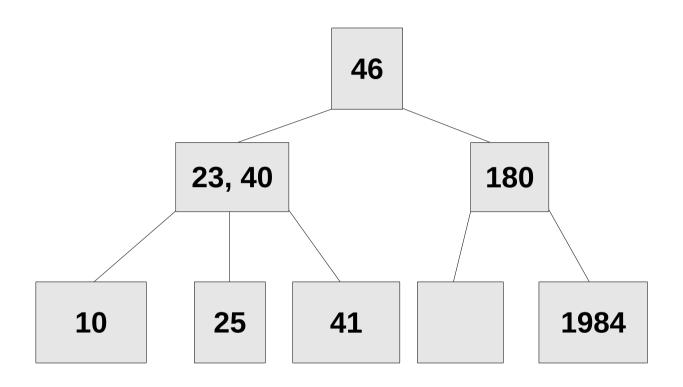


Delete 126



Delete 126

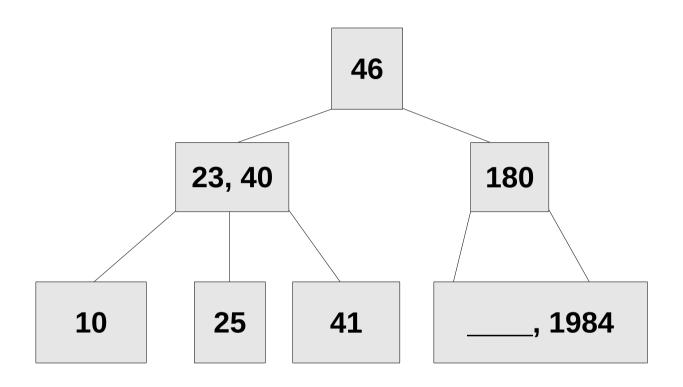
(this is going to get crazy)



Delete 126

Again, the empty node has no siblings with extra elements! So, fuse and drop! First, fuse the empty node with its left sibling (or right sibling, if it has no left sibling).

(this is going to get crazy)



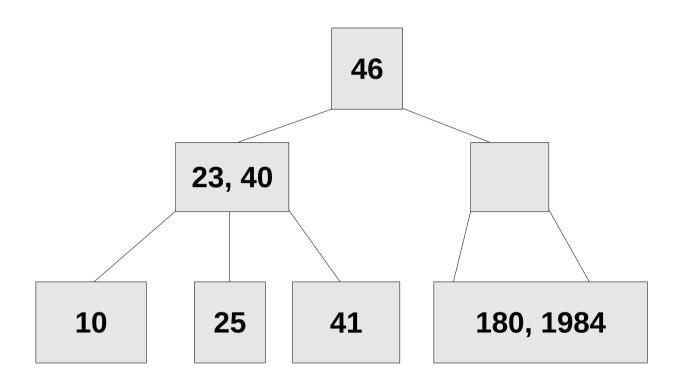
Delete 126

Again, the empty node has no siblings with extra elements! So, fuse and drop!

First, fuse the empty node with its left sibling (or right sibling, if it has no left sibling).

Now drop down the parent that was in between the newly fused nodes.

(this is going to get crazy)



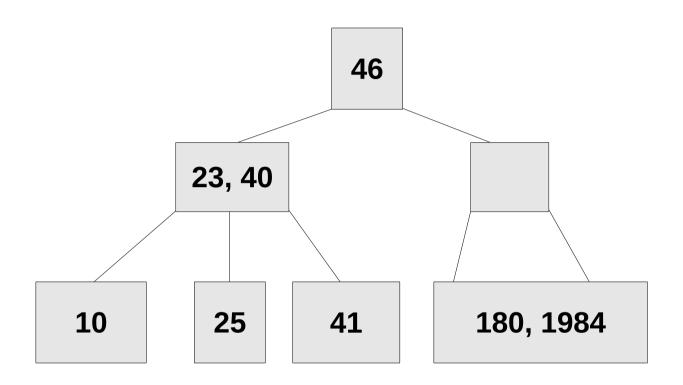
Delete 126

Again, the empty node has no siblings with extra elements! So, fuse and drop!

First, fuse the empty node with its left sibling (or right sibling, if it has no left sibling).

Now drop down the parent that was in between the newly fused nodes.

(this is going to get crazy)

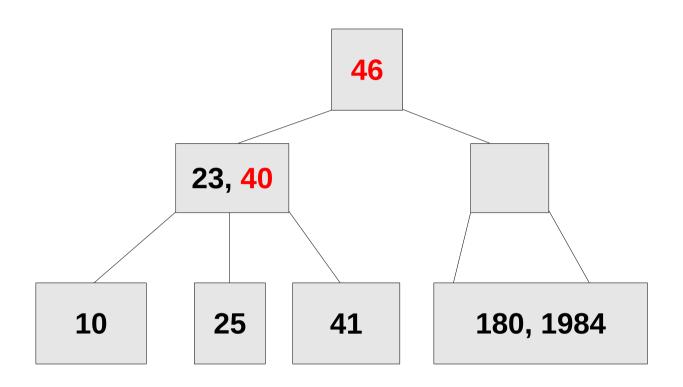


Delete 126

There's now this empty node that we can fix with a transfer.

40 moves up. 46 moves down.

(this is going to get crazy)

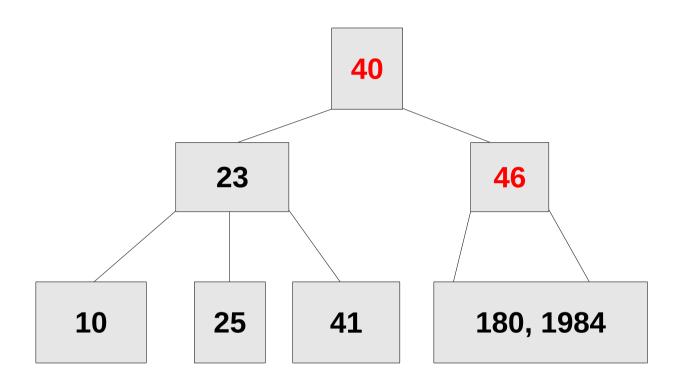


Delete 126

There's now this empty node that we can fix with a transfer.

40 moves up. 46 moves down.

(this is going to get crazy)

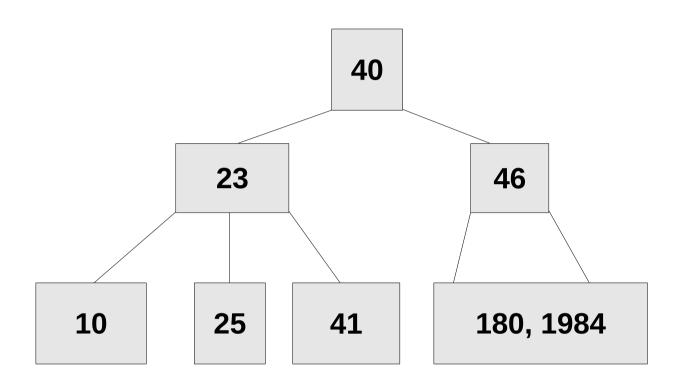


Delete 126

There's now this empty node that we can fix with a transfer.

40 moves up. 46 moves down.

(this is going to get crazy)



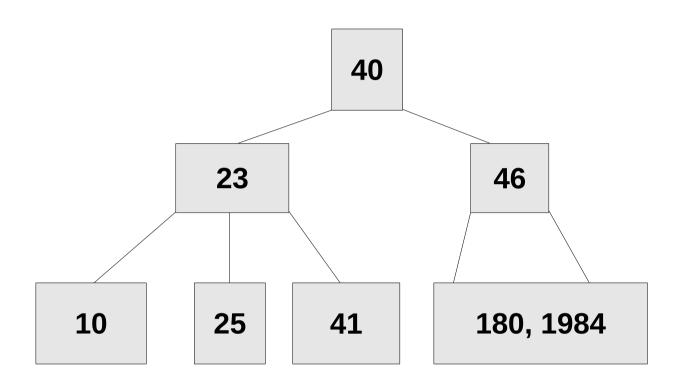
Delete 126

There's now this empty node that we can fix with a transfer.

40 moves up. 46 moves down.

Notice that 41 is now orphaned, but 46 really only has one child.

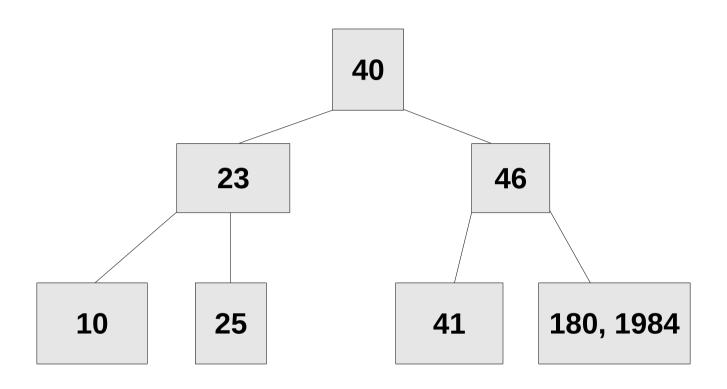
(this is going to get crazy)



Delete 126

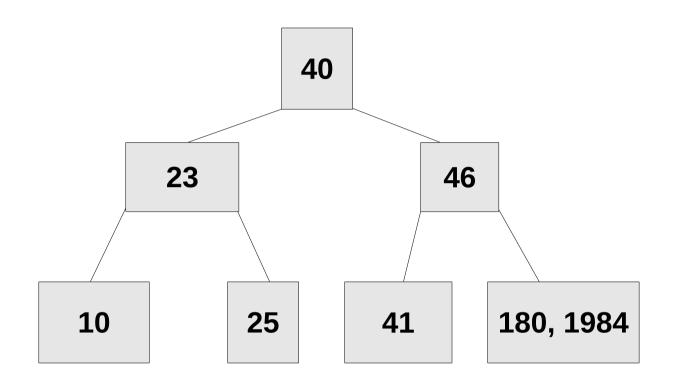
41 moves over to become the left child of 46.

(this is going to get crazy)



Delete 126

41 moves over to become the left child of 46.



Tada!

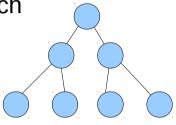
Okay, but what's the runtime for insert, delete, and search?

# 2-4 Tree Height

(which also tells us the runtime...)

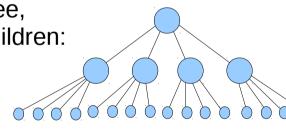
Consider a tree of height *h*.

In the smallest such tree, each node has two children:



$$h = 2$$

# nodes in bottom level: 2h



$$h = 2$$

# nodes in bottom level: 4h

Note: If that binary tree has n nodes, it has a layer of n + 1 null references. So, we have:

$$2^{h+1} \le n+1 \le 4^{h+1}$$

$$(h+1) \le \log_2(n+1) \le 2h+2$$

$$h \le \log(n+1) - 1$$
 and  $h \ge (1/2)\log(n+1) - 1$ 

$$(1/2)\log(n+1) - 1 \le h \le \log(n+1) - 1 => \Theta(\log n)$$

by taking log, of the entire expression

by solving each side separately for h

because math and formal defn. of Big-Theta

