

### Example 9

Insert 20, 15, 5, 9, 16, 8, 13, 4 into an initially empty hash table of 4 entries. Use the linear hashing scheme and hash functions  $h_i(K) = K \bmod 2^i$  for  $i \geq 2$ .

0	
1	
2	
3	

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■ Initially, *splitindex* = 0.

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- Initially, *splitindex* = 0.
- $h_i(20) = 0 \geq \textit{splitindex}$ . Insert 20 to  $h_i(20) = 0$ .

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- $h_i(20) = 0 \geq \textit{splitindex}$ . Insert 20 to  $h_i(20) = 0$ .
- Insert 15 and 5 to  $h_i(15) = 3$  and  $h_i(5) = 1$ , respectively.

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- Entry 0 is split, and 20 moves to  $h_{i+1}(20) = 4$ .  $splitindex = 1$ .

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- Entry 0 is split, and 20 moves to  $h_{i+1}(20) = 4$ .  $splitindex = 1$ .
- $h_i(16) = 0 < splitindex$ . Insert 16 to  $h_{i+1}(16) = 0$ .



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- Entry 0 is split, and 20 moves to  $h_{i+1}(20) = 4$ .  $splitindex = 1$ .
- $h_i(16) = 0 < splitindex$ . Insert 16 to  $h_{i+1}(16) = 0$ .
- $h_i(8) = 0 < splitindex$ . Insert 8 to  $h_{i+1}(8) = 0$ . The second collision and split.

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- $h_i(8) = 0 < splitindex$ . Insert 8 to  $h_{i+1}(8) = 0$ . The second collision and split.
- Entry 1 is split, and 5 moves to  $h_{i+1}(5) = 5$ .  $splitindex = 2$ .

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- $h_i(8) = 0 < splitindex$ . Insert 8 to  $h_{i+1}(8) = 0$ . The second collision and split.
- Entry 1 is split, and 5 moves to  $h_{i+1}(5) = 5$ .  $splitindex = 2$ .
- $h_i(13) = 1 < splitindex$ . Insert 13 to  $h_{i+1}(13) = 5$ . The third collision and split.

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1	9
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6	

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- Insert 15 and 5 to  $h_i(15) = 3$  and  $h_i(5) = 1$ , respectively.
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- Entry 1 is split, and 5 moves to  $h_{i+1}(5) = 5$ .  $splitindex = 2$ .
- $h_i(13) = 1 < splitindex$ . Insert 13 to  $h_{i+1}(13) = 5$ . The third collision and split.
- Entry 2 is split, but it is empty.  $splitindex = 3$ .

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- $h_i(13) = 1 < splitindex$ . Insert 13 to  $h_{i+1}(13) = 5$ . The third collision and split.
- Entry 2 is split, but it is empty.  $splitindex = 3$ .
- $h_i(4) = 0 < splitindex$ . Insert 4 to  $h_{i+1}(4) = 4$ . The fourth collision and split.

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- $h_i(13) = 1 < splitindex$ . Insert 13 to  $h_{i+1}(13) = 5$ . The third collision and split.
- Entry 2 is split, but it is empty.  $splitindex = 3$ .
- $h_i(4) = 0 < splitindex$ . Insert 4 to  $h_{i+1}(4) = 4$ . The fourth collision and split.
- Entry 3 is split, and 15 moves to  $h_{i+1}(15) = 7$ .  $splitindex = 0$ .

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- Entry 2 is split, but it is empty.  $splitindex = 3$ .
- $h_i(4) = 0 < splitindex$ . Insert 4 to  $h_{i+1}(4) = 4$ . The fourth collision and split.
- Entry 3 is split, and 15 moves to  $h_{i+1}(15) = 7$ .  $splitindex = 0$ .
- Ready for the next round.