How Can Organisations Innovate?

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18 June 2018

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What if there are constraints that agents are not aware of?

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Duh!

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- An agent's decision making is constrained by the decisions of other agents (dependency network)

G. J. Baxter et al.

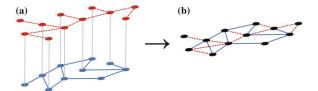


Fig. 2.1 a Two interdependent networks. A vertex in one network has a mutual dependence, represented by grey vertical lines, on zero or one vertex in the other network. b This can be reduced to a multiplex network by merging the mutually dependent vertices, and representing the edges of each network by different kinds or colours of edges

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- Suppose there is a problem in the organisation that can only be solved if *all* agents coordinate.

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- Each agent has a dependency threshold: number of innovating dependency neighbours for it to 'see' the benefit of innovation

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- positive pay-off if it innovates and sufficiently many of its dependency neighbours innovate
- **negative pay-off** if it innovates *but* not enough agents that constrain you decide with innovate
- neutral pay-off if it remains in the default state.

We assume dependency is reflexive; if A influences B then B influences A.

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- Agents have varying 'social awareness': number of 'friends' required to persuade you to adopt new ideas

Agents must be aware and receive a positive pay-off from innovation to change its state from default to innovation.

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 - In random degree distribution networks, same as rewiring edges.

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- 1. Agent randomly samples both thresholds independently from the same normal distribution (vary mean r and fix $\sigma=0.2$) [3]
- 2. Dependency threshold sampled from normal distribution but social threshold = if one friend is innovating agent is aware.

Agents with a *negative* dependency and social thresholds are *innovators* who would innovate no matter what.

social network \otimes social threshold \otimes dependency layer \otimes dependency threshold.

Consider social layer = dependency layer and the thresholds are identical for each node.

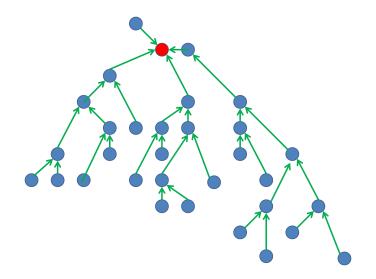
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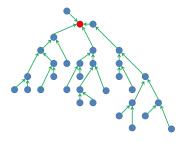
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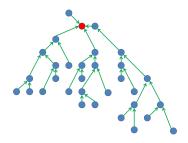
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We can think of this process of synchronously updating the state of nodes as an **infection** from infinitely far away in the network.





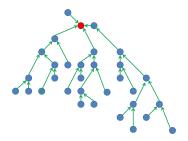
Initially all nodes are in the 0 state. Infection initiated on level n=0 at time t=0.

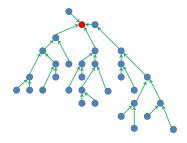


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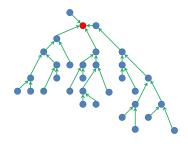
Probability that a node on layer n+1 is infected at time n+1 is dependent on the fraction of nodes infected in layer n:

$$q_{n+1} = g(q_n; r) \tag{1}$$



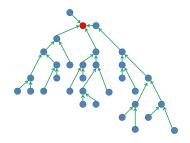


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Final fraction of infected nodes is

$$\rho = \rho(q_{\infty}; r) \tag{2}$$

We look at Erdős-Rényi graphs of varying mean degree.

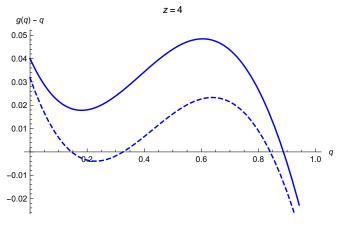


Figure 1: g(q) - q for parameters z = 4; solid is r = 0.35, dashed is r = 0.371

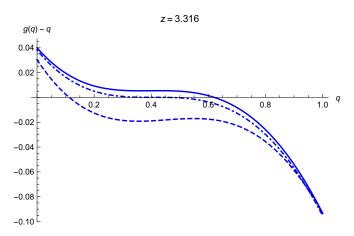


Figure 2: g(q) - q for parameters z = 4; solid is r = 0.35, dot-dashed is r = 0.3543, dashed is r = 0.375

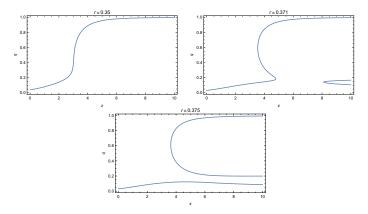


Figure 3: Fixed point plot for increasing r. Note Goldilocks zone in middle figure.

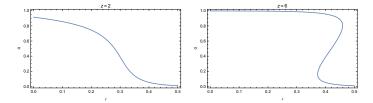
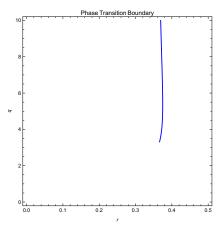


Figure 4: Fixed point plot for increasing z.

Phase transition boundary: discontinuity in $q_{\infty}(z, r)$.



Technically need two q probabilities for the two different layers respectively; however agent randomly samples both thresholds independently from the same normal distribution, so symmetry reduces the system to one variable.

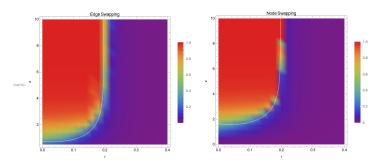


Figure 5: Final fraction size vs (z, r); white curve second order cascade condition

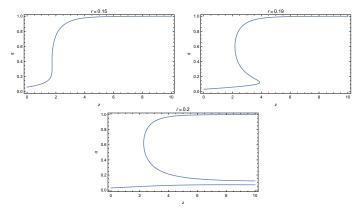


Figure 6: Edge swap: Fixed points cross section across different r values. Note that the Goldilocks zone disappears.

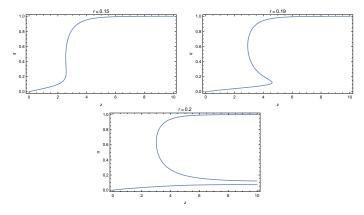


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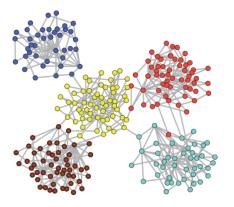


Figure 8: Image due to Preston Engstrom

In our simulations we run the infection on networks of 100 communities, each with 100 members. Average number of edges outside the community is kept at 0.5 per node.

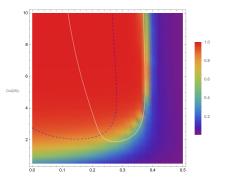


Figure 9: Single Layer simulation. Curve overlays are approximate theoretical thresholds of single layer ER graphs.

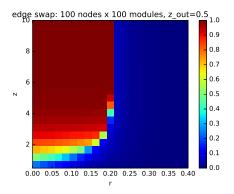


Figure 10: Phase plane with dependency network = edge swap, randomised thresholds.

N.B. in random edge swap, the dependency layer the network looses memory of its community structure and becomes an ER graph.

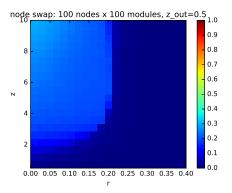


Figure 11: Phase plane with dependency network = node swap, randomised thresholds.

This scenario represents communities of agents trying to solve a *structured* problem which does not correspond to the same social community structure.

Contact Awareness

If the threshold to be socially aware is for only one friend to innovate...

Contact Awareness: Edge swap ER

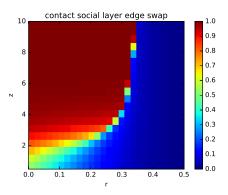


Figure 12: Phase plane with dependency network = edge swap of ER, contact threshold.

Very similar to situation without social network N.B. no Goldilocks.

Contact Awareness: Node swap ER

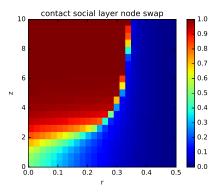


Figure 13: Phase plane with dependency network = node swap of ER, contact threshold.

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Contact Awareness: ER Dependency, PP Social

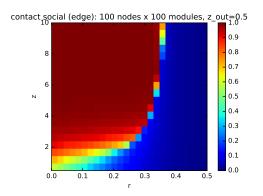


Figure 14: Dependency = ER, Social = PP, contact threshold.

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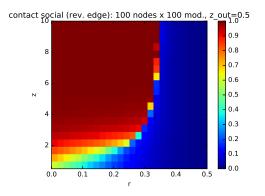


Figure 15: Dependency = PP, Social = ER, contact threshold.

Contact Awareness: node swap, PP Dependency, PP Social

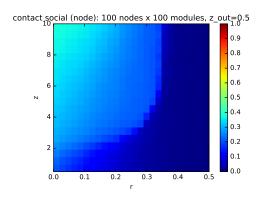


Figure 16: Dependency = PP, Social = ER, contact threshold.

Limited extent of infection appears again when PP swaps nodes.

Summary

- There is a phase transition between very little coordination to extensive, population wide coordination at critical values of social awareness and degree of dependency
- Highly socially aware agents are safe against decoupling of social and dependency network
- Community networks are inflexible against solving structured problems whose dependencies do not correspond to the same community structure. (Limited horizon of infection)
- Node swap is more dangerous because it disconnects agents with a high degree of dependency.
- Future extensions: different threshold distributions, different network topologies, adaptive social networks.

Major thanks to my supervisors Stephen

Cassidy and Stephen Brewis @BT and Renaud Lambiotte and Andrew Mellor @Oxford.

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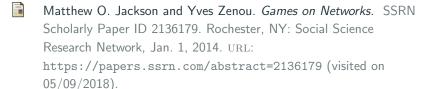


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