

# C501: Computer Architecture

## Assessed Coursework 1

### Question 1:

a) i)  $E = A \cdot B + B \cdot (A' + A \cdot B)$

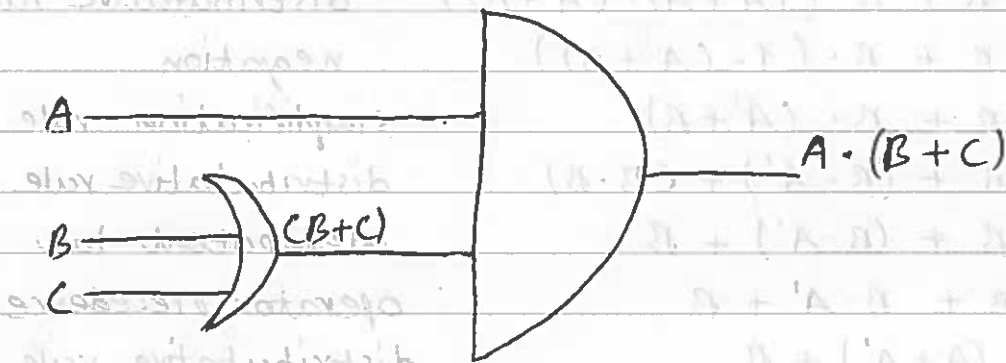
$$\begin{aligned} &\equiv A \cdot B + B \cdot (A' + (A \cdot B)) && \text{operator precedence} \\ &\equiv A \cdot B + B \cdot ((A' + A) \cdot (A' + B)) && \text{distributative rule} \\ &\equiv A \cdot B + B \cdot (1 \cdot (A' + B)) && \text{negation} \\ &\equiv A \cdot B + B \cdot (A' + B) && \text{simplification rule} \\ &\equiv A \cdot B + (B \cdot A') + (B \cdot B) && \text{distributative rule} \\ &\equiv A \cdot B + (B \cdot A') + B && \text{idempotent law} \\ &\equiv A \cdot B + B \cdot A' + B && \text{operator precedence} \\ &\equiv B \cdot (A + A') + B && \text{distributative rule} \\ &\equiv B \cdot 1 + B && \text{negation} \\ &\equiv B + B && \text{simplification rule} \\ &\equiv B && \text{idempotent law} \end{aligned}$$

ii)  $E = (A+B)' \cdot (C+D+F)' + (A' \cdot B')$

$$\begin{aligned} &\equiv (A' \cdot B') \cdot (C' \cdot D' \cdot F') + (A' \cdot B') && \text{De Morgan's rule} \\ &\equiv ((A' \cdot B') \cdot (C' \cdot D' \cdot F')) + (A' \cdot B') && \text{operator precedence} \\ &\equiv ((A' \cdot B') + (A' \cdot B')) \cdot ((C' \cdot D' \cdot F') + (C' \cdot D' \cdot F')) && \text{distributative rule} \\ &\equiv (A' \cdot B') \cdot ((C' \cdot D' \cdot F') + (C' \cdot D' \cdot F')) && \text{idempotent law} \\ &\equiv (A' \cdot B') + (0 \cdot (C' \cdot D' \cdot F')) && \text{distributative rule} \\ &\equiv (A' \cdot B') + 0 && \text{simplification rule} \\ &\equiv A' \cdot B' && \text{simplification rule} \end{aligned}$$

b)  $E = A \cdot (B + A \cdot B) + A \cdot C$

$$\begin{aligned} &\equiv A \cdot (B \cdot (A + 1)) + A \cdot C && \text{distributative rule} \\ &\equiv A \cdot (B \cdot A) + A \cdot C && \text{simplification rule} \\ &\equiv A \cdot (A \cdot B) + A \cdot C && \text{commutative rule} \\ &\equiv (A \cdot A) \cdot B + A \cdot C && \text{associative rule} \\ &\equiv A \cdot B + A \cdot C && \text{idempotent law} \\ &\equiv A \cdot (B + C) && \text{distributative rule} \end{aligned}$$



c)  $E = A \cdot B$

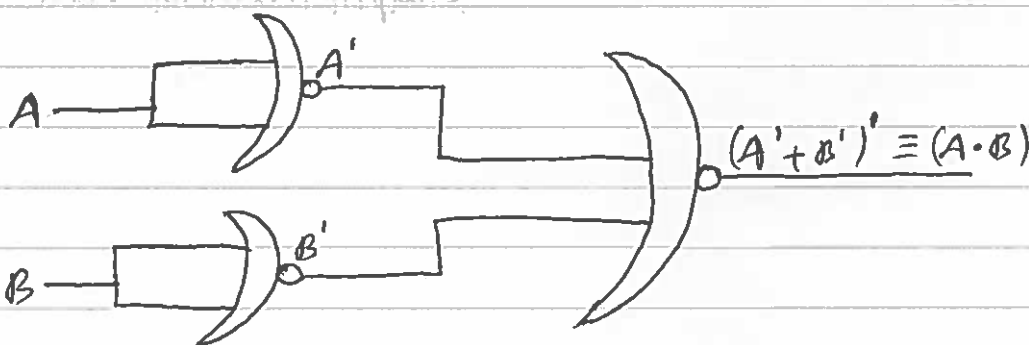
We can show that:

$$(A' + B')' \equiv A'' \cdot B'' \equiv A \cdot B$$

De Morgan's rule      negation rule

or:

A	B	A'	B'	A' + B'	(A' + B')'	A · B
T	T	F	F	F	T	T
T	F	F	T	T	F	F
F	T	T	F	T	F	F
F	F	T	T	T	F	F



## Question 2:

a) 5-bits

i) 9-12

For  $9_{10}$ :

	Quotient	Remainder
9:2	4	1
4:2	2	0
2:2	1	0
1:2	0	1

$\rightarrow 1001_2$

so  $9_{10}$  is  $01001_2$ For  $12_{10}$ :

	Quotient	Remainder
12:2	6	0
6:2	3	0
3:2	1	1
1:2	0	1

$\rightarrow 1100_2$

so  $12_{10}$  is  $01100_2$ .

Two's Complement Subtraction: "negate the subtrahend and add".

$$9 - 12 = 9 + (-12)$$

$$-12: 12 \rightarrow \text{invert} + 1 \rightarrow 10011 + 1 \rightarrow 10100_2$$

$$01001$$

$$+ 10100$$

$$\text{sum } 11101$$

$$\text{carry } 00000$$

So the result is  $11101_2$ .

We don't have overflow since we add operands of different signs.

We can check the result in decimal notation:

We need to turn  $11101_2$  from two's complement into decimal notation.

$11101_2$ :

1. it is negative

2. magnitude  $11101 - 1 = 11100$ , invert  $= 00011_2$

Therefore  $11101_2$  is  $-3$  in decimal notation.

$9 - 12 = -3$ , correct

(i)  $-11-8$

Two's Complement Subtraction: "negate the subtrahend and add"

$-11 + (-8)$

$-11_{10}$ :

	Quotient	Remainder
$11:2$	5	1
$5:2$	2	1
$2:2$	1	0
$1:2$	0	1 $\rightarrow 1011_2$

so  $-11_{10}$ :  $01011 (11_{10}) \rightarrow \text{invert} + 1 \rightarrow 10100 + 1 = 10101_2$

$-8_{10}$ :

	Quotient	Remainder
$8:2$	4	0
$4:2$	2	0
$2:2$	1	0
$1:2$	0	1 $\rightarrow 1000_2$

so  $-8_{10}$ :  $01000 (8_{10}) \rightarrow \text{invert} + 1 \rightarrow 10111 + 1 \rightarrow 11000_2$

$$\begin{array}{r}
 10101 \\
 + 11000 \\
 \hline
 \text{sum } 01101 \\
 \text{carry } 10000
 \end{array}$$

→ we discard carry-out bit

So the result is  $01101_2$

Overflow occurs in this operation, since we add two numbers of the same sign (negative) and the result has the opposite sign (positive).

We can check with decimal notation too:

$01101_2$  (two's complement) is  $2^3 + 2^2 + 2^0 = 8 + 4 + 1 = 13_{10}$  in decimal notation.

$$-11 - 8 = -19 \neq 13.$$

b)  $\frac{189}{27}$

$189_{10}$  is :

Quotient      Remainder

$189:2$

$94$

$1$

$94:2$

$47$

$0$

$47:2$

$23$

$1$

$23:2$

$11$

$1$

$11:2$

$5$

$1$

$5:2$

$2$

$1$

$2:2$

$1$

$0$

$1:2$

$0$

$1 \rightarrow 10111101_2$

So  $189_{10}$  is  $10111101_2$  in 8-bit unsigned binary notation.

$27_{10}$ 

is:

Quotient

Remainder

27:2

13

1

13:2

6

1

6:2

3

0

3:2

1

1

1:2

0

1  $\rightarrow$  11011<sub>2</sub>

so  $27_{10}$  is 11011 in 5-bit unsigned binary notation.

$$\begin{array}{r} 0000111 \\ 11011 \overline{) 10111101} \\ \underline{-0} \end{array}$$

-0

10

-00

101

-000

1011

-0000

10111

-00000

101111

$$(1): - 11011$$

$$(2): - 101000$$

$$(2): - 11011$$

$$0011011$$

$$- 11011$$

$$0000000$$

Therefore the result of this operation is:

quotient: 111 (3-bit unsigned binary)

remainder: 0

$$A' 010$$

$$(1): A 1011111$$

$$- 11011$$

$$010100$$

$$A'' 0110$$

$$A' 1001110$$

$$(2): A 101000$$

$$- 11011$$

$$001101$$

Question 3:

a)  $-31.01$  $31_{10}$  is:

Quotient      Remainder

 $31:2$ 

15

1

 $15:2$ 

7

1

 $7:2$ 

3

1

 $3:2$ 

1

1

 $1:2$ 

0

1

 $\rightarrow 11111_2$ 

$$.01 = \frac{0.02}{2} = \frac{0.04}{4} = \frac{0.08}{8} = \frac{0.16}{16} = \frac{0.32}{2^5} = \frac{0.64}{2^6} = \frac{1.28}{2^7}$$

$$= \frac{1}{2^7} + \frac{0.28}{2^7} = \frac{1}{2^7} + \frac{0.56}{2^8} = \frac{1}{2^7} + \frac{1.12}{2^9} = \frac{1}{2^7} + \frac{1}{2^9} + \frac{0.12}{2^9}$$

$$= \frac{1}{2^7} + \frac{1}{2^9} + \frac{0.24}{2^{10}} = \frac{1}{2^7} + \frac{1}{2^9} + \frac{0.48}{2^{11}} = \frac{1}{2^7} + \frac{1}{2^9} + \frac{0.96}{2^{12}}$$

$$= \frac{1}{2^7} + \frac{1}{2^9} + \frac{1.92}{2^{13}} = \frac{1}{2^7} + \frac{1}{2^9} + \frac{1}{2^{13}} + \frac{0.92}{2^{13}} =$$

$$= \frac{1}{2^7} + \frac{1}{2^9} + \frac{1}{2^{13}} + \frac{1.84}{2^{14}} = \frac{1}{2^7} + \frac{1}{2^9} + \frac{1}{2^{13}} + \frac{1}{2^{14}} + \frac{0.84}{2^{14}} =$$

$$= \frac{1}{2^7} + \frac{1}{2^9} + \frac{1}{2^{13}} + \frac{1}{2^{14}} + \frac{1.68}{2^{15}} = \frac{1}{2^7} + \frac{1}{2^9} + \frac{1}{2^{13}} + \frac{1}{2^{14}} + \frac{1}{2^{15}} + \frac{0.68}{2^{15}}$$

$$= \frac{1}{2^7} + \frac{1}{2^9} + \frac{1}{2^{13}} + \frac{1}{2^{14}} + \frac{1}{2^{15}} + \frac{1.36}{2^{16}} = \frac{1}{2^7} + \frac{1}{2^9} + \frac{1}{2^{13}} + \frac{1}{2^{14}} + \frac{1}{2^{15}} + \frac{1}{2^{16}} + \frac{0.36}{2^{16}}$$

$$= \frac{1}{2^7} + \frac{1}{2^9} + \frac{1}{2^{13}} + \frac{1}{2^{14}} + \frac{1}{2^{15}} + \frac{1}{2^{16}} + \frac{0.72}{2^{17}} = \frac{1}{2^7} + \frac{1}{2^9} + \frac{1}{2^{13}} + \frac{1}{2^{14}} + \frac{1}{2^{15}} + \frac{1}{2^{16}} + \frac{1}{2^{17}} + \frac{0.44}{2^{18}}$$

$$= \frac{1}{2^7} + \frac{1}{2^9} + \frac{1}{2^{13}} + \frac{1}{2^{14}} + \frac{1}{2^{15}} + \frac{1}{2^{16}} + \frac{1}{2^{18}} + \frac{0.44}{2^{18}}$$

$$= \frac{1}{2^7} + \frac{1}{2^9} + \frac{1}{2^{13}} + \frac{1}{2^{14}} + \frac{1}{2^{15}} + \frac{1}{2^{16}} + \frac{1}{2^{18}} + \frac{0.88}{2^{19}}$$

$$\text{So } 0.01 \approx \frac{1}{2^7} + \frac{1}{2^9} + \frac{1}{2^{13}} + \frac{1}{2^{14}} + \frac{1}{2^{15}} + \frac{1}{2^{16}} + \frac{1}{2^{18}}, \text{ or } .0000001010001111010$$

so 31.01 is  $1111.0000001010001111010$   
 $= 1.11110000001010001111010 \times 2^4$

### IEEE Single Precision Format (32-bit)

1 bit is the sign bit  
 8 bit is the exponent  
 23 bit is the coefficient

Exponent is Excess-127

so actual value of 4 is  $127+4=131$  stored value

	Quotient	Remainder
$131:2$	65	1
$65:2$	32	1
$32:2$	16	0
$16:2$	8	0
$8:2$	4	0
$4:2$	2	0
$2:2$	1	0
$1:2$	0	1

$\rightarrow 10000011_2$

Significant:  $11110000001010001111010$

Exponent:  $10000011$

Sign: 1 (negative number)

Sign Exponent Significant

1 10000011 11110000001010001111010

Hexadecimal Value:

1100 0001 1111 1000 0001 0100 0111 1010

C 1 F 8 1 4 7 A

so the Hexadecimal value is:  $0x C1F8147A$



b) 0x 40F108D4

Binary:

4	0	F	1	0	8	D	4
0100	0000	1111	0001	0000	1000	1101	0100

IEEE Single Precision Format

Sign Exponent Significand

0 10000001 11100010000100011010100

Decimal:

Exponent:  $10000001_2$  is 129 stored value $129 - 127 = 2$  actual value in Excess-127.

Significand field + hidden bit: 1.11100010000100011010100

Sign 0, so positive number.

Therefore:  $1.11100010000100011010100 \times 2^2$   
 $= 111.100010000100011010100$

which is:  $2^2 + 2^1 + 2^0 + 2^{-1} + 2^{-5} + 2^{-10} + 2^{-14} + 2^{-15} + 2^{-17} + 2^{-19}$

$$= 4 + 2 + 1 + \frac{1}{2} + \frac{1}{2^5} + \frac{1}{2^{10}} + \frac{1}{2^{14}} + \frac{1}{2^{15}} + \frac{1}{2^{17}} + \frac{1}{2^{19}}$$

$$= 7 + \frac{2^{18} + 2^{14} + 2^9 + 2^5 + 2^4 + 2^2 + 1}{2^{19}}$$

$$= 7 + 0.532327652 = 7.532327652$$