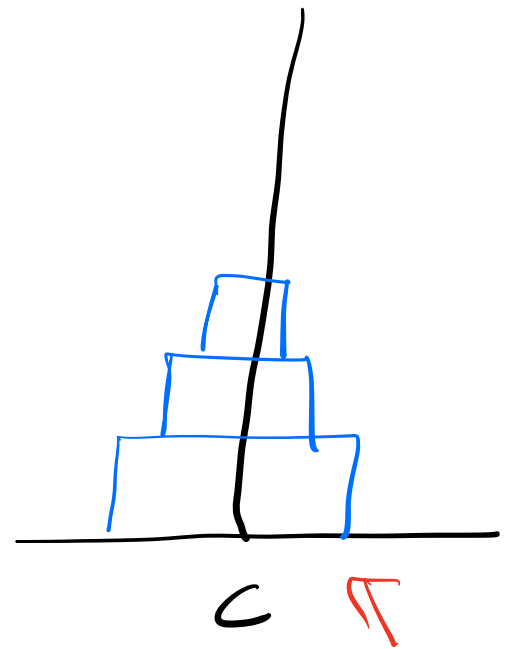
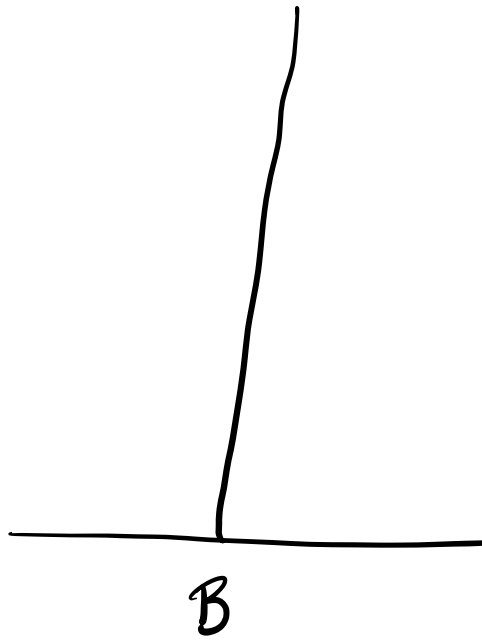
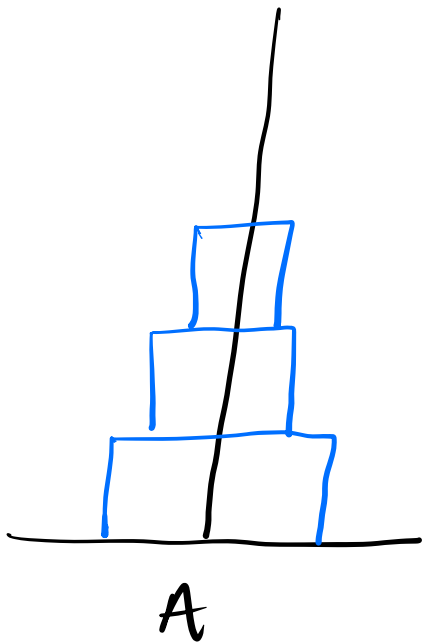


● Tower of Hanoi Given 3 towers A, B, C
N discs are on Tower A
move all discs from A \rightarrow C

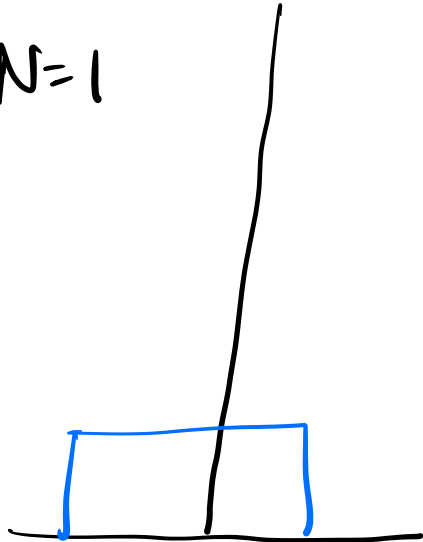
- Only 1 disc moves at a time.
- Larger disc cant be placed on smaller disc.



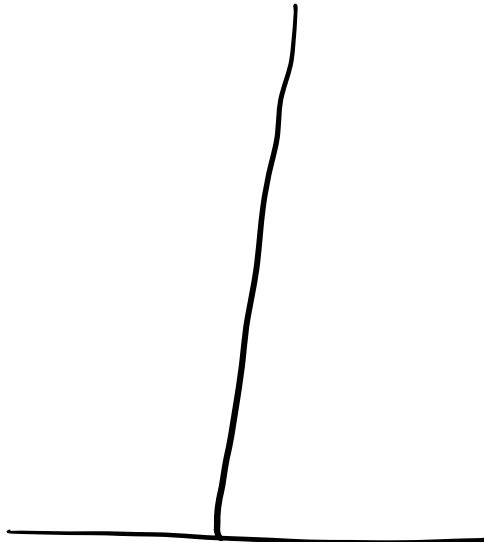
Final

Print the movement of discs

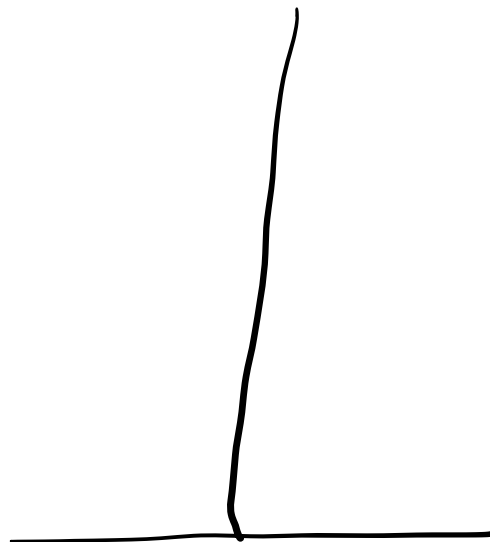
$N=1$



A

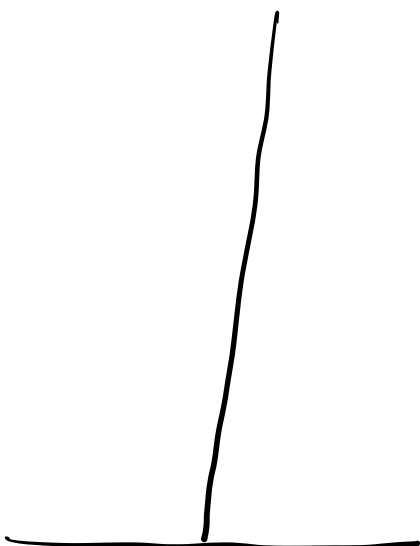


B

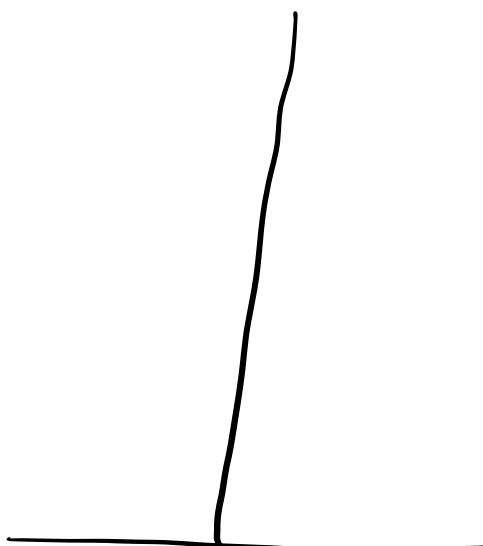


C

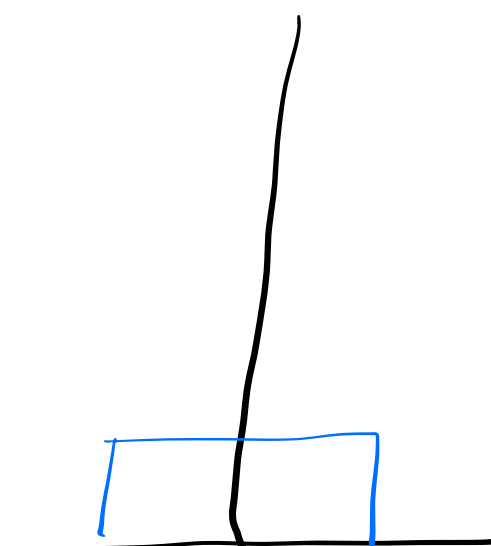
$A \rightarrow C$



A

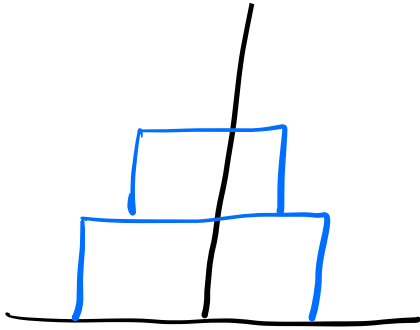


B



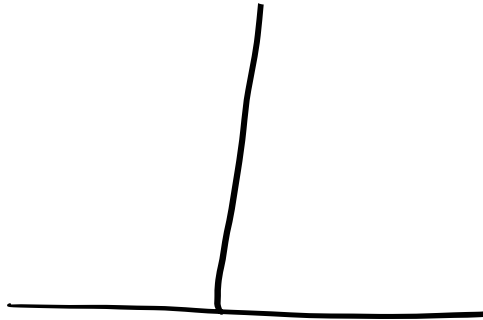
C

$N=2$



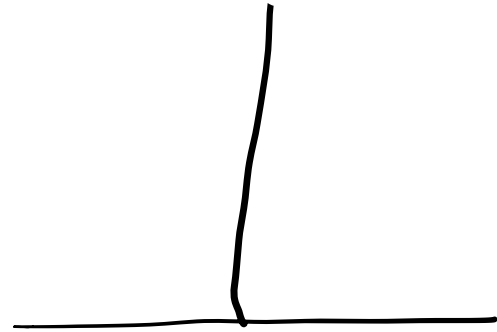
A

src



B

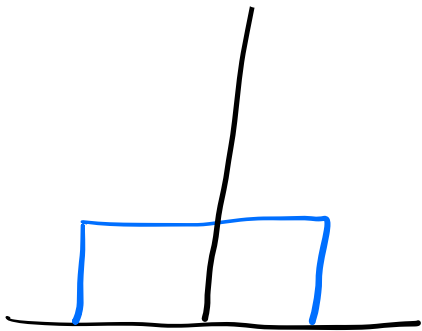
temp



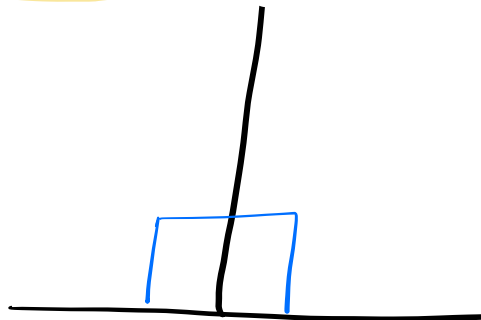
C

dest

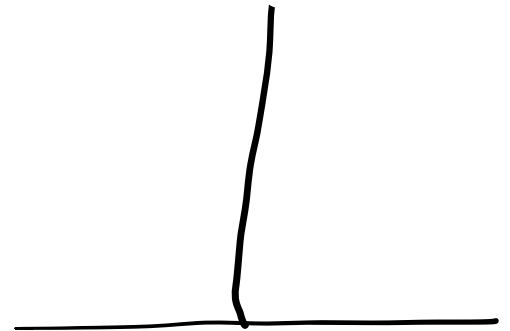
$A \rightarrow B$



A

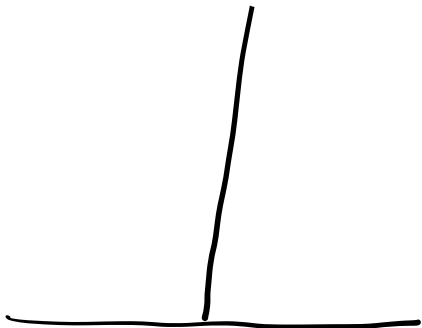


B

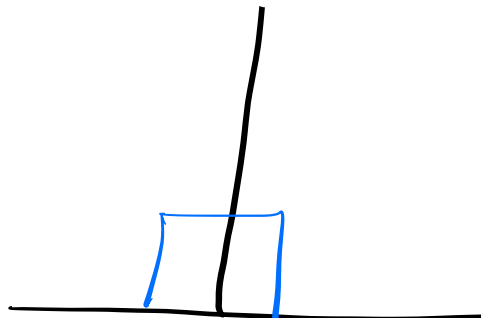


C

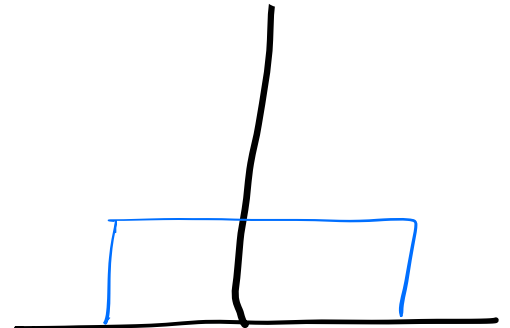
$A \rightarrow C$



A

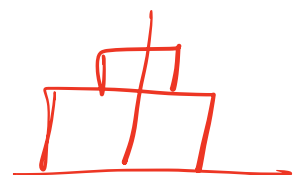


B

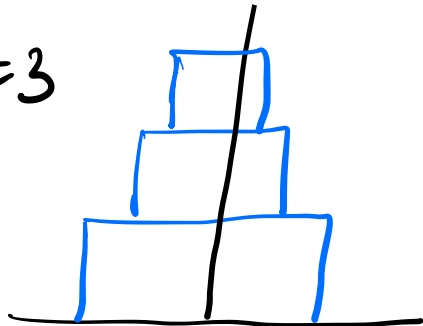


C

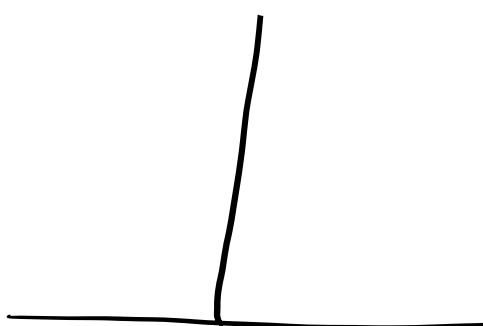
$B \rightarrow C$



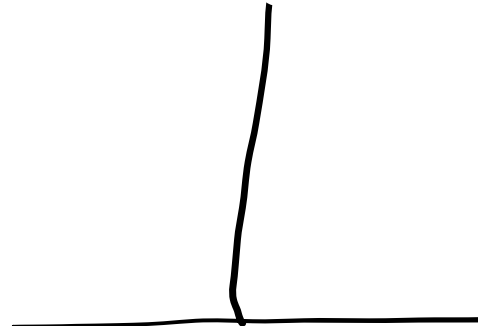
N=3



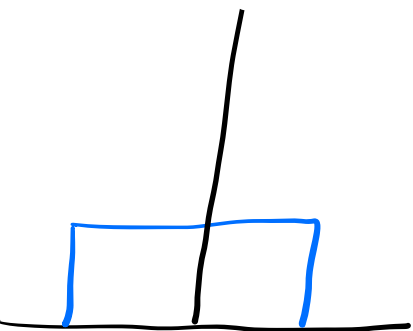
A



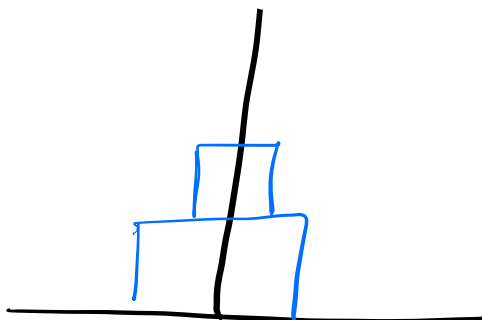
B



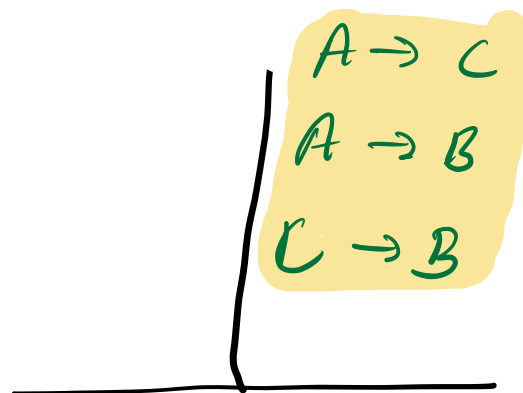
C



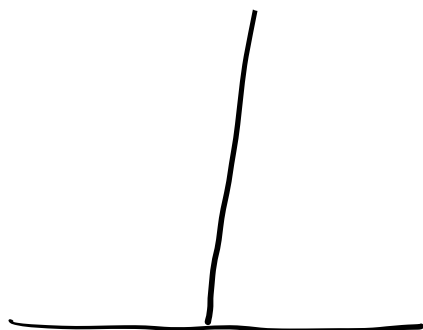
A



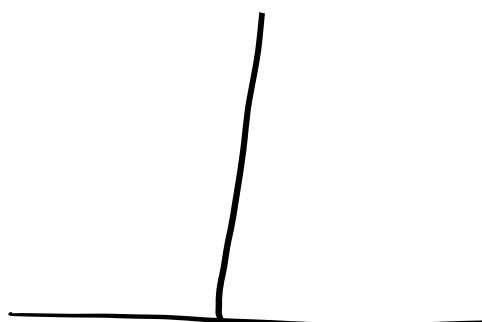
B



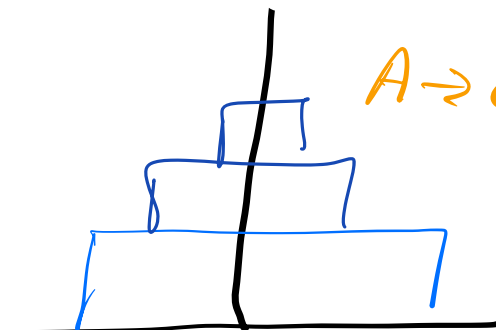
C



A



B



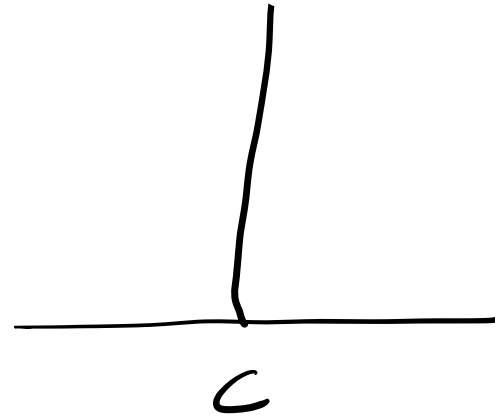
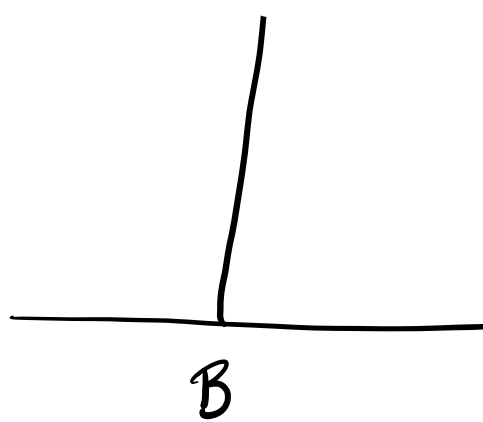
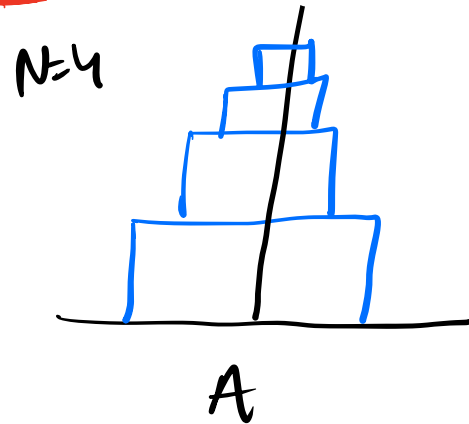
C

Now move $B \rightarrow C$ via A

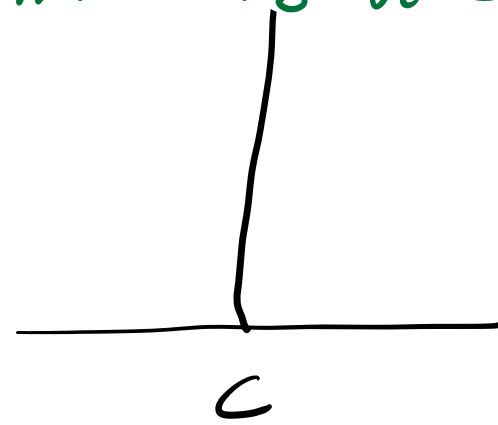
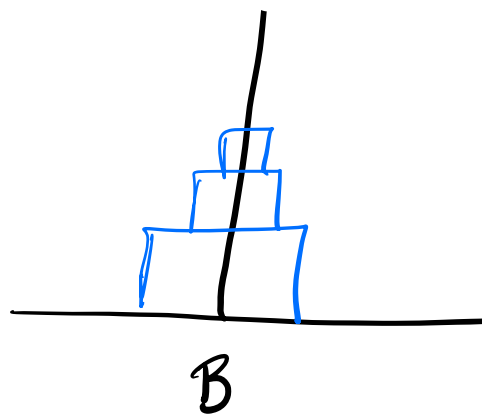
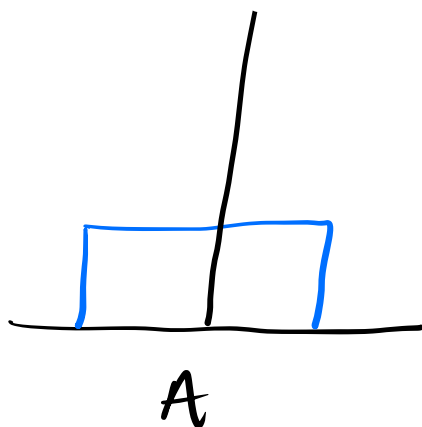
$B \rightarrow A$

$B \rightarrow C$

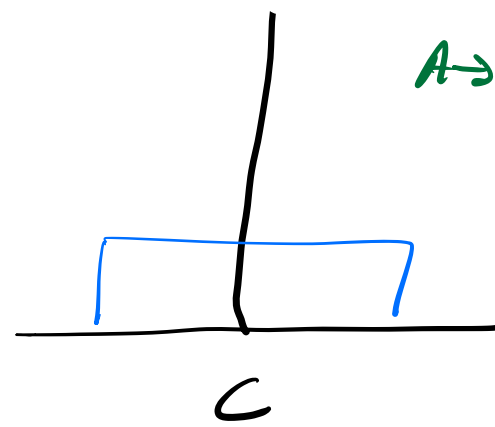
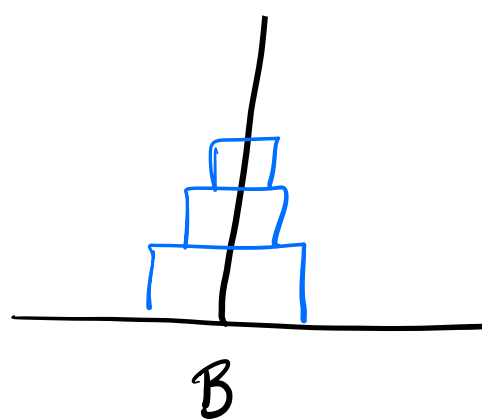
$A \rightarrow C$



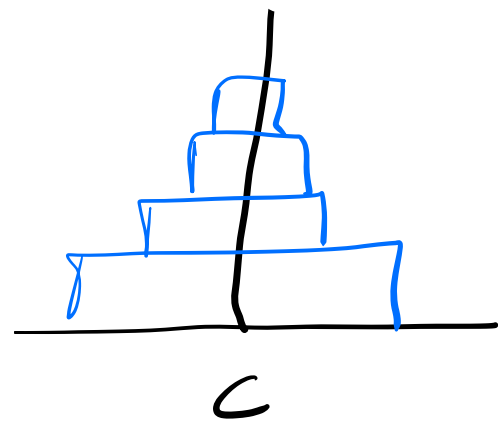
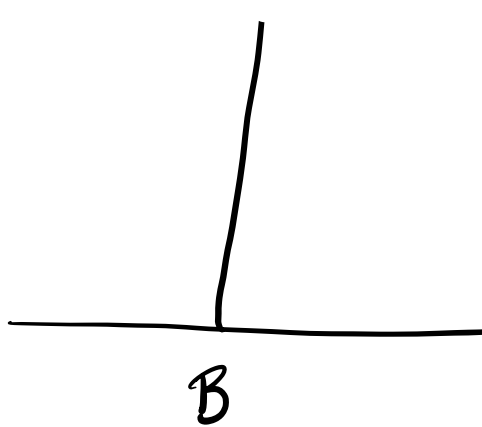
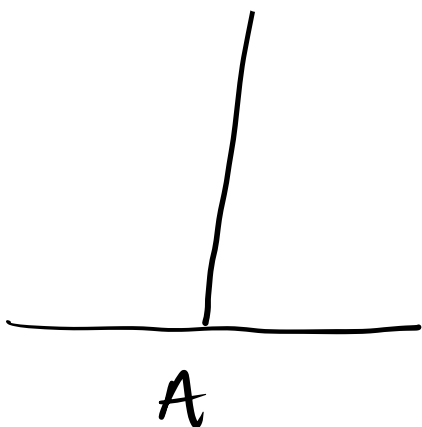
Put $n-1$ $A \rightarrow B$ via C



$A \rightarrow C$



Now move $n-1$ $B \rightarrow C$ via A



● Algorithm

- 1) Move $n-1$ from $A \rightarrow B$ via C
- 2) Move from $A \rightarrow C$
- 3) Move $n-1$ from $B \rightarrow C$ via A

```
void toh (int N, int src, int dest, int temp) {  
    if (N == 1)  
        print (src + " - " + dest)  
    toh (n-1, src, temp, dest)  
    print (src + " - " + dest)  
    toh (n-1, temp, dest, src)  
}
```

Time Complexity : $2^n - 1 \Rightarrow O(2^n)$

Moves $N=1$ 1 $2^1 - 1$

$N=2$ 3 $2^2 - 1$

$N=3$ 7 $2^3 - 1$

Generalize $N=4$ $2^4 - 1$

toh (3, ^{source} A, ^{dest} C, ^{temp} B)

toh (2, A, B, C)

↳ toh (1, A, C, B)

$A \rightarrow C$

$A \rightarrow B$

toh (1, C, B, A)

$C \rightarrow B$

• $A \rightarrow C$

toh (2, B, C, A)

toh (1, B, A, C)

$B \rightarrow A$

$B \rightarrow C$

toh (1, A, C, B)

$A \rightarrow C$

$A \rightarrow C$

$A \rightarrow B$

$C \rightarrow B$

$A \rightarrow C$

$B \rightarrow A$

$B \rightarrow C$

$A \rightarrow C$

$A \rightarrow C$

$A \rightarrow B$

$C \rightarrow B$

$A \rightarrow C$

$B \rightarrow A$

$B \rightarrow C$

$A \rightarrow C$

Q Print all valid parantheses of $len = 2N$

Eg $\rightarrow N = 1$ $()$

$N = 2$ $(())$ $()()$

$N = 3$ $((()))$ $() (())$
 $(()) ()$ $() () ()$ $(() ())$

Idea: try to create all possible parantheses recursively

↑
 $i = 2$

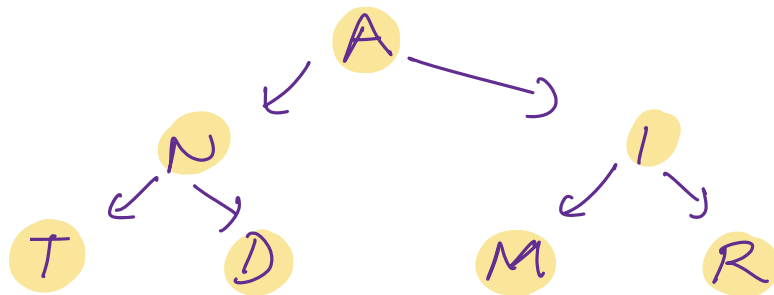
$len = 12$

at all times open \geq close

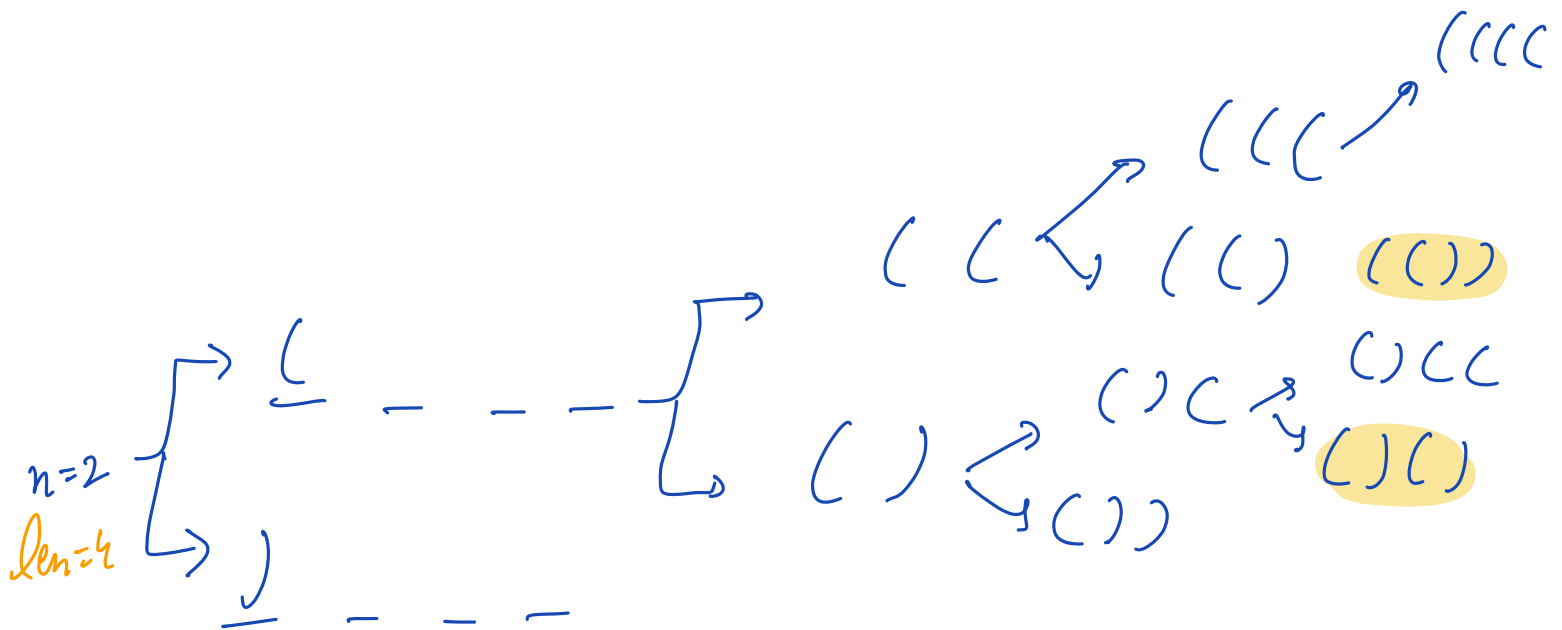
Obs: something like $(()) - - -$ is always invalid. How to use this to better the approach.

Backtracking

Find AIM



Parenthesis via backtracking



obs : open \geq close at all times.

at the end open equal to close

Code

```
void solve (string s, N, open, close) {  
    if (s.length() == 2*N) {  
        if (open == close) {  
            print(s)  
            return;  
        }  
    }  
    if (open < N)  
        solve (s + "(", N, open + 1, close)  
    if (close < open)  
        solve (s + ")", N, open, close + 1)  
}
```

TC: $O(2^n)$

$N=2$ $len=4$

"", 0, 0



(, 1, 0



(, 2, 0



(, 1, 1



((, 3, 0



((), 2, 1



()(, 2, 1



((((, 4, 0

((())

3, 1

(()((, 3, 1

((()))

2, 2

()()((, 3, 1

()()()

2, 2



(

)



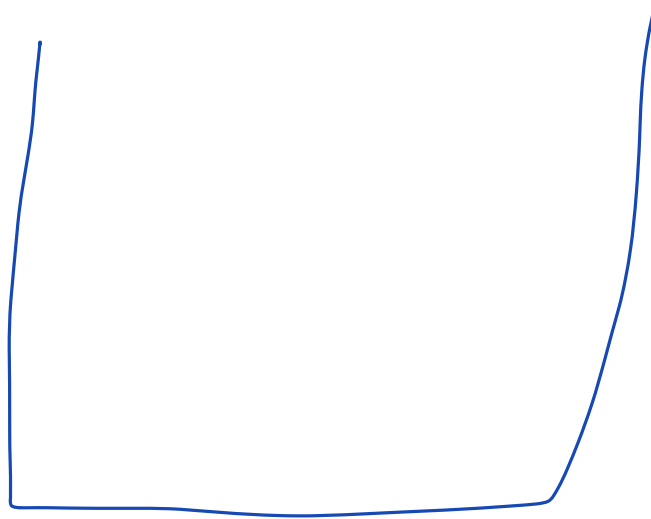
((

(()

```

void solve (n) {
    if (n == 0)
        return
    solve (n-1)
    print (n)
}

```



3 2 1

```

void solve (int n) {

```

mysolver