Queue → First in First out

Operations:

(insert) enqueue (x)

(remove) deque (x)

front ()

isEmpty ()

rear ()

5 4 3

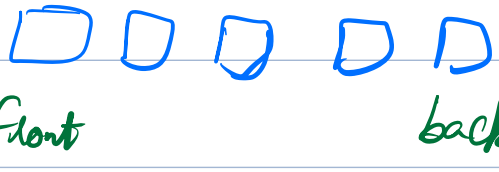
counter

back

0	1	2	3	4	5
10	20	30	40	50	60
	↑				↑
	front				rear

Queue using array

```
class Queue {
    int size
    int arr[size]
    int front = 
    int rear = -1
```



starting idx of queue

ending idx of queue

```
    enque(int x){
        if (rear == size-1) sout("Queue Full")
        else
            rear++          arr[rear] = x

    }

                                    → size of the queue
        deque( ){
    if ( rear-front+1 == 0 )
            sout("Empty Queue")
    else
        front ++

    }
```

Q1 Given K. Series is made up of 1, 2. Return K^{th} num.

	1	2	3	4	5	6	7	8
Series ⇒	1	2	11	12	21	22	111	112

k=7 ans = 111

k=4 ans = 12

Brute Iterate from 1. check if num has only 1, 2. If yes, count++. when count = k, end.

Idea

First 2 no.s are 1, 2

From 1, what all num we can get directly ⇒ 1 → 11 → 111, 112; 12 → 121, 122

From 2, ⇒ 2 → 21, 22;

---

21 22 111 112 121 122

From x ⇒ 1) 10*x + 1

2) 10*x + 2

Code

queue <int> q

q.enque (1)

q.enque (2)

1 2 3 4

21 22 111 112 121 122

cnt = 1

while (cnt != k) {

int x = q.front ()

q.deque ()

cnt ++

q.enque (10x + 1)

q.enque (10x + 2)

}

return q.front ()

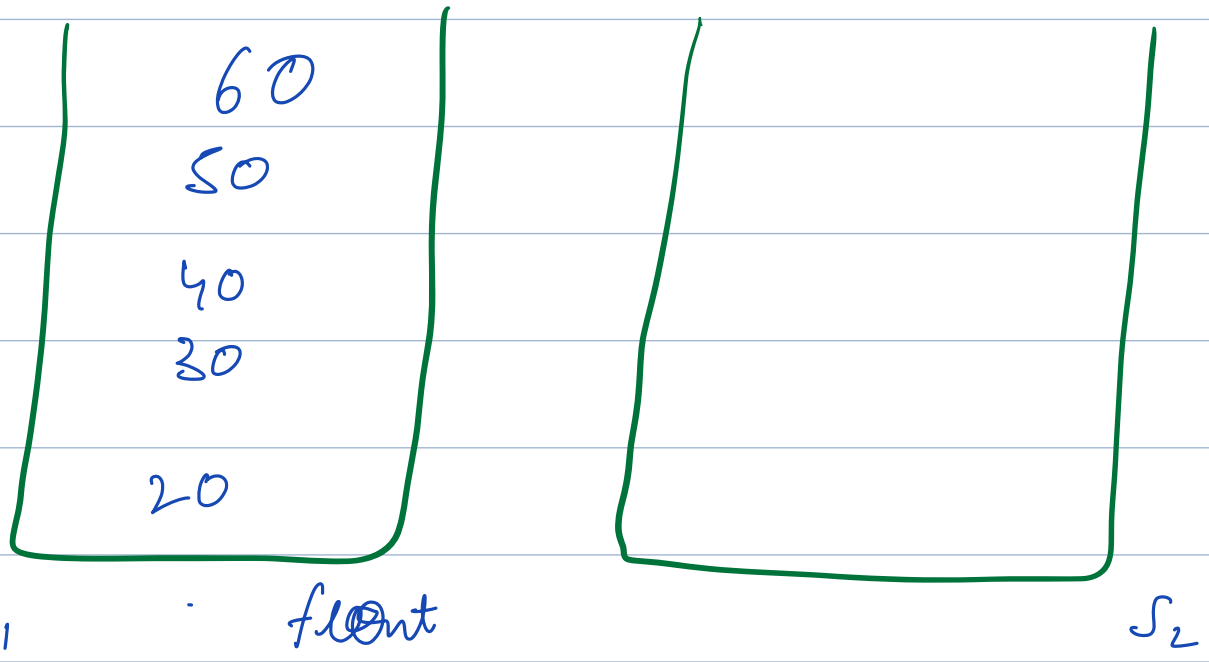
TC: } O(k)

SC: }

MMT

Q3 Implement Queues using stacks (only)

Idea: Use 2 stacks

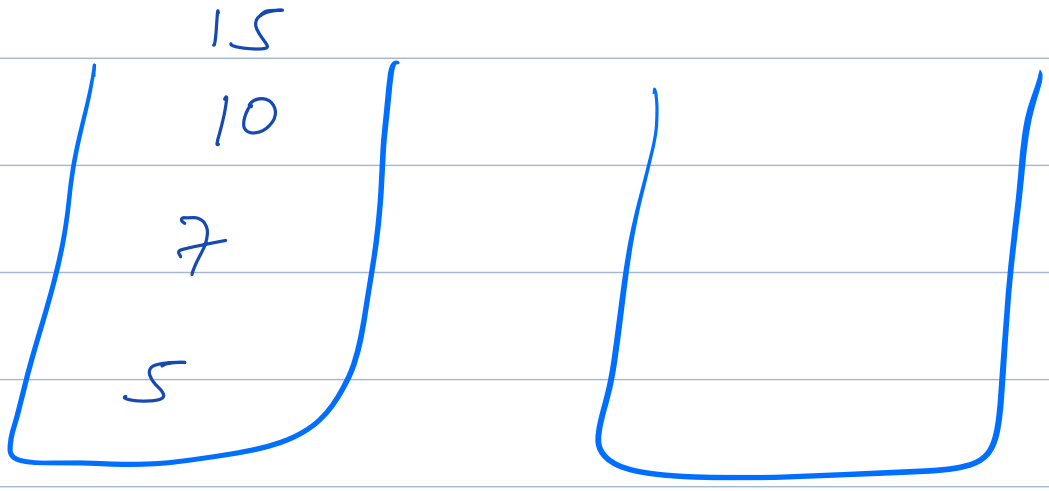


enqueue → Just put on top of S_1

deque → Need to remove bottom elem.

1) Remove all elem & put in S_2
2) S_2.pop()
3) Put all contents of S_2 back in S_1

dry run.



Code

```
void enque (x) {
  s1. push (x)
}

void deque () {
    if ( s1. empty () )
        print ("Error")
  while ( ! s1. empty () )
    s2. push (s1. top ())
    s1. pop ()
}
```

```
    s2.pop()
    while(!s2.empty()){
        s1.push(s2.top())
        s2.pop()
    }
}
```

TC

enque → O(1)

deque → O(n)

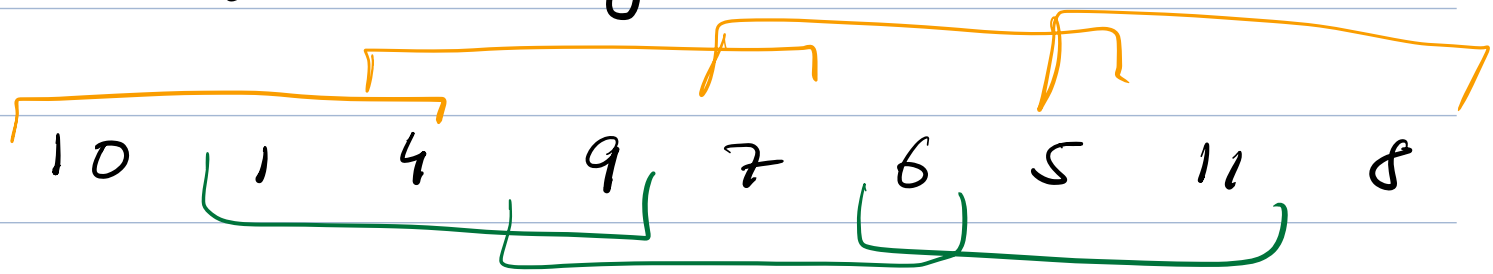
Doubly ended queue Deque

Sliding window maximum

Deque

VVV Imp Adobe GS Amazon Microsoft CRED

Array of size N. Find max elem of every subarray of size K.



k=3 ans ⇒ {10 9 9 9 7 11 11}

Brute: Check all subarrays of size K & find max. TC: $O(n^2)$

Idea k=3

10 1 8 9 7 6 5 11 3

11 3

front

→

access last elem stack

access first elem queue

Deque (Doubly ended queue)

k=4

3 15 6 15 12 4 2 10 9 18

15 15 15 15 12 10 18

18

First prepare the deque for first window

new_elem

rear < new_elem

remove rear

repeat

rear ≥ new_elem

insert at rear

remove elem

check with front

if equal

remove from front

if not

do nothing

Max ⇒ front element

Code

```
Deque <int> dq
for (i=0; i<k; i++){
    while (!dq.empty() && 
           dq.rear() < arr[i])
        dq.pop_rear()
    dq.push_rear(arr[i])
}

ans.insert(dq.front())
int s=1          e=k
while(e<n){
    while (!dq.empty() && 
           dq.rear() < arr[e])
        dq.pop_rear()
    dq.push_rear(arr[e])

    if (dq.front == arr[s-1])
        dq.pop_front()
    ans.insert(dq.front())
```

TC: O(n)
SC: O(n)

Stt ett

y

1 elem → 1 insert, 1 delete } 2

1 → 2 op

N → 2N op

{done}

Idea k=3

10 1 8 9 7 6 5 11 3

{

front