Q1) Given N array elements, check if there exists a pair (i,j) such that ar[i] + ar[j] = K && i != j

index values

a + b = K

0	1	2	3	4	5	6	7	8	9
8	9	1	-2	4	5	11	-6	7	5

K=11 i=4 j=8

K=6 i=2 j=5

K=22 i=6 j=6 ✗

Brute force: Check each pair

TC: $O(n^2)$ SC: $O(1)$

$a + b = K \implies b = K - a$

K=11 a=5 b=6

a	b (k-a)
8	3
9	2
1	10
-2	13
4	7

K = 11

found it!!!

Hashset

Allows me to check if b exists or not. in O(1)

0	1	2	3	4	5	6	7	8	9
8	9	1	-2	4	5	11	-6	7	5

K=22

HS = { 8, 9, 1, -2, 4, 5, 11, -6, 7 }

a	b(k-a)
8	14
9	13
1	21
-2	24
4	18
5	17
11	11

K = 22

Yes

Learning: We need to maintain freq.

Freq Hashmap

0	1	2	3	4	5	6	7	8	9
8	9	1	-2	4	5	11	-6	7	5

K=22

HM: 8:1, 9:1, 1:1, -2:1, 4:1, 5:2, 11:1, -6:1, 7:1

a	b
8	14
9	13
1	21
-2	24
4	18
5	17
11	11
-6	28
7	15
5	17

freq(11) != 2

return false

Pseudo Code

1) Create the frequency hm.

```
for(i=0; i<n; i++) {
    a = ar[i]      b = k-a
    if(a==b) {
        if( hm.get(a) >= 2 )
            return true.
    }
    else {
        if( hm.containskey(b) )
            return true.
    }
}
return false
```

TC: O(n) SC: O(n)

Q Count no of pairs $a[i] + a[j] = K$ $i != j$

Eg 2 5 2 5 8 5 2 8 $K = 10$

b = 8 5 8 5 2 5 8 2 HM 2 : ✗ ✗ 3

+1 +2 +2 +1 F3 5 : ✗ ✗ 3

8 : ✗ 2

```
hashmap<int,int> hm
ans = 0
for( i=0 ; i<n ; i++) {
    b = K-a
    if (hm.contains(b))
        ans += hm[b]
    hm[arr[i]] ++
}
return ans
```

TC } O(n)
SC }

Q Calc the number of distinct elements in each subarray of size k.

	0	1	2	3	4	5	6	7	8	9
Eg: ar[10] =	2	4	3	8	3	9	4	9	4	10

k = 4

0, 3	4
1, 4	3
2, 5	3
3, 6	4
4, 7	3
5, 8	2
6, 9	3

Idea: Calc using hashset for every window

TC: $O(n^2)$

Idea: Optimize using hashmap

0 : 3 <2,1> <4,1> <3,1> <8,1>

1 : 4 remove a_0 add a_4

add ⇒ freq ++

sub ⇒ freq -- but if freq = 0 remove

<1,4>

<4,1> <3,2> <8,1>

ans = 3

<2,5>

<9,1> <3,2> <8,1>

ans = 3

done

0 k-1

1 k

Code

```
hashmap<int, int> hm
for (i=0 ; i<k ; i++) {
    hm[ a[i] ] ++
}
print (hm.size())
s=1        e=k
while (e<n) {
    hm[ar[ s-1 ]] --
    if ( hm[ar[s-1]] ==0 )
        hm.remove ( ar[s-1] )

    hm[ar[e]] ++
    print (hm.size)
    s++        e++
}
```

TC O(n)
SC O(k)

Q Check if subarray with sum = K exist

Eg - 2 3 9 -4 1 5 6 2 5 K = 11

K = 10

Idea: Similar to prev class where sum = 0

Sum [s:e] ⇒ pf[e] - pf[s-1] s ≠ 0

pf[e] s = 0

Create pf array.

	0	1	2	3	4	5	6	7	8
arr	2	3	9	-4	1	5	6	2	5
pf	2	5	14	10	11	16	22	24	29

Now, pf[e] - pf[s-1] = K

pf[s-1] = pf[e] - K

pf[s-1] : target

1) Create pf sum

2) hashset <int> hs

```
for ( i=0 ; i<n ; i++ ){
    target = pf[i] - k
    if ( hs.contains ( target ) )
        return true
    else {
        hs.add (pf[i])
    }
}
if ( hs.contains(K))
    return true
else
    return false
```

TC } O(n)
SC }

sum += a[i]

if(sum == 0 || set.contains(sum))

set = 1

set.add sum

	1	3	-1	-2	4
sum	1	4	3	1	5