

Q1) Row wise & col wise matrix

Adv DSA arrays 2 first ques

1	2	3
4	5	6
7	8	9

$k=2$

Start at

10:35

Start at $i=0$ $j=n-1$

i, j set $i \times 1009 + j$

assume 1 index

$i=17$ $j=21$
 $\rightarrow 18$ $\rightarrow 22$

1	2	2	3
4	5	6	7
8	9	10	11

$k=2$

$A =$

	2	1	0
	1	1	1

conv $A =$

	0	0	0
--	---	---	---

Q2 A, B, C.

- 1) unset C bits from right.
- 2) Restore B bits from right to original value

Eg1 A = 12 B = 2 C = 3

3	2	1	0	
1	1	0	0	\Rightarrow ans = 8
	0			

Eg2 A = 7 B = 1 C = 3

2	1	0	
1	1	1	ans = 1
0	0	1	

Let A =

31	C	C-1	B+1	B	B-1	2	1	0
			0	0	0	0	1	1	1	1	1
			og	og	og	og	og	og	og	og	og

In A set [B, C-1] = 0

```

for (bit = B ; bit ≤ C-1 ; bit++) {
    unset (A, bit)
}

```

```

unset (int A, int bit) {
    if (A & 1 << bit != 0) {
        A = A ^ 1 << bit
    }
}

```

Approach 2

```

original_a = A
for (bit = 0 ; bit < C ; bit++) {
    unset (A, bit)
}

```

} Step 1

```

for (bit = 0 ; bit < B ; bit++) {
    if (original_a & 1 << bit != 0) {
        set (A, bit)
    }
}

```

Q3 Max possible subarray sum with decreasing weight

eg1 3 2 1 ans = 6

eg2 3 3 5 0 1 ans = 5

have we seen some thing like this before ?

Yes

Kadane's Algorithm

⇒ Keep taking until love ≥ 0

Here we have same idea but different condition.

Keep taking while elems are decreasing.

	0	1	2	3	4	5	6
	7	9	8	4	1	3	10
sum = 7	9	17	21	22	3	10	
ans = 7	9	17	21	22	22	22	22

Code

```
int solve ( int [] arr, int n) {  
    int ans = arr[0]  
    int cur_sum = arr[0]  
    for ( i=1 ; i < n ; i++ ) {  
        if ( arr[i] < arr[i-1] ) {  
            cur_sum += arr[i] ; ans = max (ans ,  
                                            cur_sum)  
        }  
        else {  
            cur_sum = arr[i]  
            ans = max (ans , cur_sum )  
        }  
    }  
    return ans  
}
```

Repeat / Reattempt

Sat / Sun